



**FACULTY OF ENGINEERING AND TECHNOLOGY  
ELECTRICAL AND COMPUTER ENGINEERING DEPARTMENT  
ADVANCED DIGITAL DESIGN ENCS3310  
COURSE PROJECT**

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**Section: 2**

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## **Abstract**

The project's goal is to create a Finite State Machine (FSM) is intended to recognize the particular sequence "1011." The FSM functions by transitioning between a number of states in response to input bits. When it starts out, it is in the first state, but as soon as it reads "1," it moves on to the second state. When it comes across "0," it then switches to the third state. The FSM advances to the fourth state, which denotes successful detection of the intended sequence, after receiving a second "1." The FSM goes back to its initial state if any unexpected input is received. The FSM architecture makes sure that input bit streams containing the "1011" sequence are correctly identified.

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## **Brief introduction and background:**

### **- Introduction:**

The main goal of the project is to develop and implement a Moore Finite State Machine (FSM) that can identify the specific binary sequence "1011" in a stream of binary input data. Finite State Machines, which are used to represent and regulate sequential activity, are crucial elements in the design of digital logic. The project builds a Moore FSM specifically designed to detect the target sequence using T-flip-flops as essential building elements.

### **-Background:**

It is frequently necessary for digital systems to be able to identify particular patterns inside binary data streams. A computer model known as a finite state machine (FSM) offers a structured and methodical method for addressing sequential patterns. The Moore FSM is a form of FSM in which inputs cause state transitions and outputs are connected to states.

The T-flip-flop, which changes its state in response to a particular input signal, serves as the project's central component. Its actions are in line with the requirement to switch between states when certain binary patterns are recognized. A Moore FSM can be built to detect complex sequences by interconnecting numerous T-flip-flops and customizing their connections.

The number sequence "1011" serves as a real-world illustration of the project's fundamental ideas. In order to detect this sequence, the FSM must react to particular input combinations and move through many states until arriving at the ultimate acceptance state. As part of the design process, a state diagram that describes the sequence detection procedure is created. T-flip-flops are then given states, state transitions are decided upon based on input conditions, and the necessary output logic is established for each state.

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The capability of accurately detecting sequences and patterns becomes increasingly important as digital systems get more complex. This project offers insight into the actual use of T-flip-flops as memory components inside a sequential logic circuit in addition to practical expertise in developing and implementing finite state machines. The research effectively exhibits the power of combining fundamental digital components to produce complex pattern recognition by successfully developing a Moore FSM that correctly detects the "1011" sequence.

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## Design philosophy:

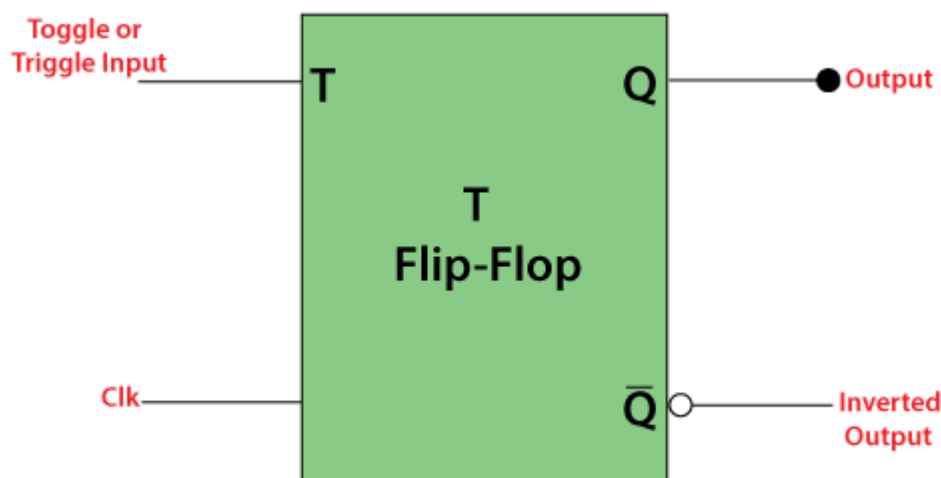
We used in project:

- T-Flip Flops:

In T flip flop, "T" defines the term "Toggle". In SR Flip Flop, we provide only a single input called "Toggle" or "Trigger" input to avoid an intermediate state occurrence. Now, this flip-flop work as a Toggle switch. The next output state is changed with the complement of the present state output. This process is known as "Toggling".

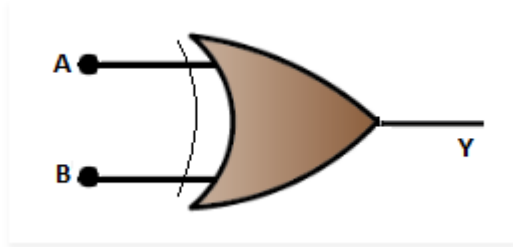
We can construct the "T Flip Flop" by making changes in the "JK Flip Flop". The "T Flip Flop" has only one input, which is constructed by connecting the input of JK flip flop. This single input is called T. In simple words, we can construct the "T Flip Flop" by converting a "JK Flip Flop". Sometimes the "T Flip Flop" is referred to as single input "JK Flip Flop".

Block diagram of the "T-Flip Flop" is given where T defines the "Toggle input", and CLK defines the clock signal input.



### -Xor gate:

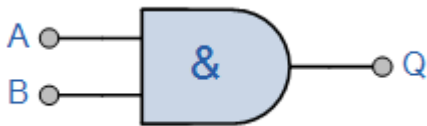
The Exclusive-OR gate or XOR gate is achieved by combining standard logic gates together. XOR gate is used extensively in error detection circuits, computational logic comparators and arithmetic logic circuits. The Exclusive OR gate gives an output only if its two inputs are dissimilar, namely if one of them is high (one) and the other is low (zero).



### -And gate:

The output state of a digital logic AND gate only returns “LOW” again when ANY of its inputs are at a logic level “0”. In other words for a logic AND gate, any LOW input will give a LOW output.

The logic or Boolean expression given for a digital AND gate is that for Logical Multiplication which is denoted by a single dot or full stop symbol, ( . ) giving us the Boolean expression of:  $A.B = Q$ .



## Results:

This is the code for the T-Flip Flops that takes the clock and reset to trigger the output q and the module was build behaviorally.

```
1 //abdelrhman abed ~1193191~
2 module TFF (clk,reset,t,q); //T-Flip Flops
3     input clk,reset,t;
4     output reg q;
5
6     always @(posedge clk or posedge reset)
7     begin
8         if (reset)
9             q <= 1'b0;
10        else if (t)
11            q <= ~q;
12    end
13 endmodule
```

This's the module for the output detector which takes an inputs and only gives the output if it meets the sequence 1011.

```
15 module Det (q,det); //DetectionLogic
16
17     input [0:3]q;
18     output reg det;
19
20
21     always @(*)
22     begin
23         if (q[0] & ~q[1] & q[2] & q[3])
24             det = 1'b1;
25         else
26             det = 1'b0;
27     end
28
29 endmodule
```

This's the module for the structural description of the circuit .

```
31 module SDet(clk,reset,input1,det);    //SequenceDetector
32
33     input clk,reset,input1;
34     output reg det;
35     wire t0, t1, t2, t3;
36     reg [0:3]q;
37
38     TFlipFlop tff0 (.clk(clk) , .reset(reset) , .t(t0) , .q(q[0]));
39     TFlipFlop tff1 (.clk(clk) , .reset(reset) , .t(t1) , .q(q[1]));
40     TFlipFlop tff2 (.clk(clk) , .reset(reset) , .t(t2) , .q(q[2]));
41     TFlipFlop tff3 (.clk(clk) , .reset(reset) , .t(t3) , .q(q[3]));
42
43     assign t0 = (q[0] ^ input1) & (~q[1]);
44     assign t1 = (q[1] ^ ~input1) & (~q[2]);
45     assign t2 = (q[2] ^ input1) & (~q[3]);
46     assign t3 = (q[3] ^ input1) & q[0];
47
48     Det1 (.q(q), .det(det));
49
50
51
52
53 endmodule
```

The code for the behavioral description of the circuit

```
55 module SDetBehavioral (clk, reset, input1, det);    //SequenceDetectorBehavioral
56
57     input clk, reset, input1;
58     output reg det;
59     reg [0:3] q;
60
61     always @(posedge clk or posedge reset) begin
62         if (reset)
63             q <= 4'b0000;
64         else begin
65             q[0] <= q[0] ^ input1;
66             q[1] <= q[1] ^ ~input1;
67             q[2] <= q[2] ^ input1;
68             q[3] <= q[3] ^ input1;
69         end
70     end
71
72     always @(*) begin
73         if (q[0] & ~q[1] & q[2] & q[3])
74             det = 1'b1;
75         else
76             det = 1'b0;
77     end
78
79 endmodule
```



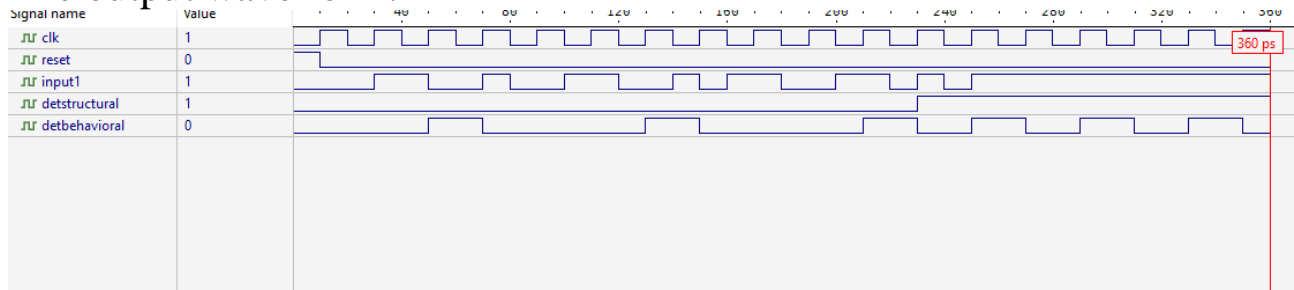
The test bench module for the structural and behavioral circuit modules.

```
81 module SDet_TB;
82
83     reg clk;
84     reg reset;
85     reg input1;
86
87
88     wire detstructural;
89     wire detbehavioral;
90
91
92     SDet structural (.clk(clk),.reset(reset),.input1(input1), .det(detstructural));
93
94     SDetBehavioral behavioral (.clk(clk),.reset(reset),.input1(input1),.det(detbehavioral));
95
96
97     always
98         #10 clk = ~clk;
99
100
```

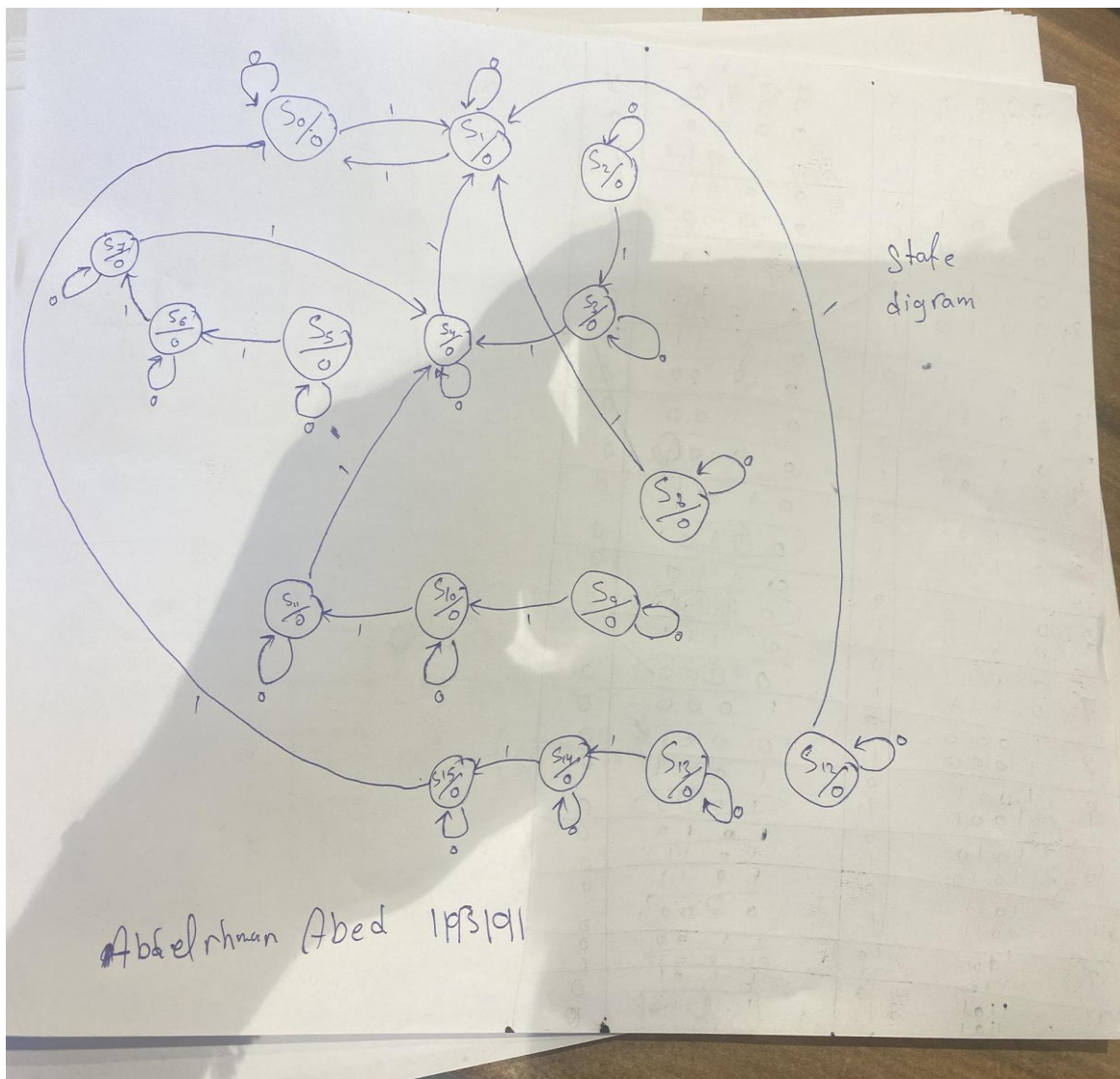
```
103
104     initial
105         begin
106
107             clk = 0;
108             reset = 1;
109             input1 = 0;
110
111             #10 reset = 0;
112
113             for (int i = 0; i < 16; i = i + 1)
114                 begin
115                     input1 = i;
116                     if (detstructural == detbehavioral)
117                         begin
118                             #10;
119                             end
120                             #10;
121                         end
122
123             #100
124             $finish;
125         end
126
127     endmodule
128
```

<

## The output wave form.



## State diagram



## State table

state	Q[3],Q[2],Q[1],Q[0]	input	Next state	output
0	0000	0	0	0
0	0000	1	1	0
1	0001	0	1	0
1	0001	1	0	0
2	0010	0	2	0
2	0010	1	3	0
3	0011	0	3	0
3	0011	1	4	0
4	0100	0	4	0
4	0100	1	1	0
5	0101	0	5	0
5	0101	1	6	0
6	0110	0	6	0
6	0110	1	7	0
7	0111	0	7	0
7	0111	1	4	0
8	1000	0	8	0
8	1000	1	1	0
9	1001	0	9	0
9	1001	1	10	0
10	1010	0	10	0
10	1010	1	11	0
11	1011	0	11	0
11	1011	1	4	0
12	1100	0	12	0
12	1100	1	1	0
13	1101	0	13	0
13	1101	1	14	0
14	1110	0	14	0
14	1110	1	15	0
15	1111	0	15	0
15	1111	1	0	1

## Input for the fourth TFF

For  $X=0$  for all cases

$Q_3 Q_2$ \ $Q_1 Q_0$	00	01	11	10
00	X	X	X	X
01	X	X	X	X
11	X	X	X	X
10	X	X	X	X

When the input  $X=0$   
the machine acts like a memory  
for the state or the sequence

For  $X=1$   $Q_3^+$

$Q_3 Q_2$ \ $Q_1 Q_0$	00	01	11	10
00				
01				
11		1		1
10		1	1	1

$Q_3^+ = Q_3 \cdot Q_1' \cdot Q_0 + Q_3' \cdot Q_1 \cdot Q_0'$   
 by simply using boolean algebra and considering the input for both cases  
 $Q_3^+ = (Q_3 \wedge X) \cdot Q_0$

## The output for the third TFF

for  $x=1$   $Q_2^+$

$Q_3 Q_2$ \ $Q_1 Q_0$	00	01	11	10
00				
01		1	1	
11	1	1		1
10		1	1	

$Q_2^+ = Q_3' \cdot Q_1 \cdot Q_0 + Q_2 \cdot Q_1' \cdot Q_0 + Q_3 \cdot Q_0 \cdot Q_3' \cdot Q_2' + Q_2 \cdot Q_1 \cdot Q_3'$   
 by simplifying using boolean algebra and considering the input for both causes

$Q_2^+ = (Q_2 \cdot 1) \cdot Q_3$

## The output for the second TFF

$Q_1^+$

$Q_3 Q_2$ \ $Q_1 Q_0$	00	01	11	10
00				
01		1	1	1
11				
10	1	1	1	1

$Q_1^+ = Q_1 \cdot Q_0' + Q_2 \cdot Q_1' \cdot Q_0 + Q_3 \cdot Q_1' \cdot Q_0$   
 by simplifying using boolean algebra and considering the input for both cases

$Q_1^+ = (Q_1 \cdot Q_0) + \sim Q_2$

## The output of the first TFF and the output y

$Q_0^+$

$Q_3 Q_2$ \ $Q_1 Q_0$	00	01	11	10
00	1	1	1	1
01				
11				
10	1	1	1	1

$Q_0^+ = Q_1 \cdot Q_0' + Q_1' \cdot Q_0'$

By Simplifying using boolean algebra and considering the input for both cases

$Q_0^+ = (Q_0 \wedge x) \cdot \sim Q_1$

Output(y)

$Q_3 Q_2$ \ $Q_1 Q_0$	00	01	11	10
00	1			
01				
11				
10				

$y = Q_1' \cdot Q_0' \cdot Q_3' \cdot Q_2'$

I used a specific module called Det1 to implement this output.



## Conclusions and Future works:

I concluded that there is a slight error in the modules but both of them detected the sequence as they should. The error in the wave form could be from the clock's alignments between the two modules. I believe that if I separate the inputs and reset and clocks and assign specific inputs for each case they will function as required.

Also, I concluded that we can depend on the more ~fsm~ to design several circuits for any problem we want to obtain a solution for it.

For future works I want to test the theory of why the clocks resulted in change of how the output came and try to fix it.