1.1. For each algorithm below calculate its time complexity.

```
b)
a)
    for i in range(n):
                                            for i in range(n):
      j = 0
                                               j = i
      while j * j < i:
                                               while j > 0:
                                                 j = j // 2
         i += 1
c)
                                        d)
    \mathbf{def} \ \mathbf{f(n)}:
                                            \mathbf{def} \ \mathbf{f(n)}:
      if n = 0:
                                               if n = 0:
         return 1
                                                 return 1
                                               else
      else
                                                 return f(n // 3) + f(n // 3)
         return 5 * f(n // 3)
```

- 1.2. Prove using mathematical induction that if T(n) = 2T(n/2) + n, then $T(n) = \Omega(n \log n)$ (lower bound).
- 1.3. Prove using mathematical induction that if T(n) = 2T(n/2 + 20) + n, then $T(n) = O(n \log n)$.
- 1.4. Prove using mathematical induction that if $T(n) = \log n \cdot T(n/\log n) + n$, then $T(n) = O(n \log n)$.
- 1.5. Prove using mathematical induction that if $T(n) = 2T(\sqrt{n}) + 1$, then $T(n) = O(\log n)$.
- 1.6. You are given two arrays a and b sorted in non-decreasing order. Determine if there is a number that occurs in both arrays. Time O(n).
- 1.7. You are given two arrays a and b sorted in non-decreasing order. Find i and j such that the difference $|a_i b_j|$ is minimal. Time O(n).
- 1.8. You are given two arrays a and b sorted in non-decreasing order and a number S. Find such i and j such that the sum $a_i + b_j = S$. Time O(n).
- 1.9. You are given two arrays a and b sorted in non-decreasing order. Find the number of pairs (i, j) such that $a_i = b_j$. Time O(n).
- 1.10. You are given two arrays a and b sorted in non-decreasing order. Find the number of pairs (i, j) such that $a_i > b_j$. Time O(n).
- 1.11. Given an array a. The pair (i, j) such that i < j and $a_i > a_j$ is called <u>inversion</u>. Find the number of inversions in array a. Time $O(n \log n)$
- 1.12. Show that, with the correct implementation, merge sort is stable (that is, does not change the order of equal elements).
- 1.13. Show how to implement merge sort with a single additional array of size n (without constructing new arrays in each recursive call).
- 1.14. Show how to implement merge sort without recursion.