

Name(s) _____

CSE Login _____

Instructions Same partner/group policy as the project applies to this assignment. Answer each question as completely as possible. You are encouraged to typeset your assignment using L^AT_EX or some typesetting system. Figures may be hand drawn and scanned/included as an image.

Question	Points	Score
1	15	
2	15	
3	5	
4	5	
5	24	
6	4	
7	9	
8	8	
9	8	
10	2	
11	30	
12	0	
Total:	125	

Algorithm Analysis

When asked to design and analyze an algorithm, be sure to provide the following:

1. Complete pseudocode
2. Identify the input and the input size, n
3. Identify the elementary operation
4. Compute how many times the elementary operation is executed with respect to the input size n
5. Provide a Big-O asymptotic characterization for the algorithm's complexity

1. 15 points Let P be an image represented as an $(n \times m)$ 2-dimensional array of pixels. Design and analyze an algorithm that given an image P will rotate it counter clockwise by 90 degrees.
2. 15 points Let R be a set of rectangles in the 2-D plane represented as a pair of points, $((x, y), (x', y'))$ where (x, y) is the lower left point of the rectangle and (x', y') is the upper right point of the rectangle. Design and analyze an algorithm that given a set R of n rectangles finds the pair of rectangles whose intersection has the greatest area.

For the next few questions, consider defining a *time function* with respect to n . That is, if the algorithm is linear, we could write the time that it takes as a function of n :

$$t = f(n) = cn$$

If it is quadratic, we could write it as a quadratic function:

$$t = f(n) = cn^2$$

Both of these instances ignore lower order terms. If we know the time it takes for a particular value of n then we can compute the constant c and consequently predict the time t it takes for other values of n or vice versa.

3. 5 points An algorithm takes 2.5 ms for an input size 100. How long will it take for input size 2000 (assuming that low-order terms are negligible) if the running time is
 - (a) linear
 - (b) $O(n \log n)$
 - (c) quadratic
 - (d) cubic
 - (e) exponential
4. 5 points An algorithm takes 1 ms for input size 100. How large can an input size be if a problem can be solved in 1 minute (assuming that low-order terms are negligible) if the running time is:
 - (a) linear
 - (b) $O(n \log n)$
 - (c) quadratic
 - (d) cubic
 - (e) exponential
5. 24 points Prove each of the following statements by applying the definition of Big-O. That is, derive an inequality (show your work) and clearly identify the c, n_0 constants you derive as per the definition of Big-O.

- (a) $3n = O(n)$
- (b) $500\sqrt{n} = O(n)$
- (c) $n^2 + 2n + 1 = O(n^2)$
- (d) $2048n + 1234 = O(n^2)$
- (e) $n \log(32n) = O(n \log(n))$
- (f) $12n^3 + 50n^2 - 12n - 60 = O(n^3)$
- (g) $n2^n = O(3^n)$
- (h) $\log(n!) = O(n \log n)$

Recursion

6. 4 points (Weiss 7.17) Give a Big-O characterization for the following recurrences (with initial conditions $T(0) = T(1) = 1$).
- (a) $T(n) = T(n/2) + 1$
 - (b) $T(n) = T(n/2) + n$
 - (c) $T(n) = T(n/2) + n^2$
 - (d) $T(n) = 3T(n/2) + n$
 - (e) $T(n) = 3T(n/2) + n^2$
 - (f) $T(n) = 4T(n/2) + n$
 - (g) $T(n) = 4T(n/2) + n^2$
 - (h) $T(n) = 4T(n/2) + n^3$

Trees

7. 9 points Perform a pre-order, in-order, and post-order traversal on the tree in Figure 1 and give the resulting node sequences for each.

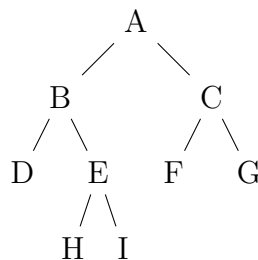


Figure 1: A Tree

8. 8 points Consider a binary search tree T with 5 nodes.
 - (a) What is the *maximum* possible depth of T ?
 - (b) What is the *minimum* possible depth of T ?
 - (c) What are the minimum number of leaves T can have?
 - (d) What are the maximum number of leaves T can have?
9. 8 points Consider a binary search tree T with 15 nodes.
 - (a) What is the *maximum* possible depth of T ?
 - (b) What is the *minimum* possible depth of T ?
 - (c) What are the minimum number of leaves T can have?
 - (d) What are the maximum number of leaves T can have?
10. 2 points Consider a binary search tree T with n nodes.
 - (a) What is the *maximum* possible depth of T ?
 - (b) What is the *minimum* possible depth of T ?
11. When dealing with trees, algorithm complexity is usually measured in terms of the number of nodes in the tree while the elementary operation is usually a node traversal. Let T be a tree such that each node u has a `parent`, `rightChild` and `leftChild`.
 - (a) 15 points Design and analyze an algorithm (provide pseudocode) that, given a node u in T determines its depth d . Analyze your algorithm.
 - (b) 15 points Design and analyze an algorithm (provide pseudocode) that, given an integer d and a tree T with nodes u_1, \dots, u_n determines the *number* of nodes in T at depth d .

Honors/Bonus

These are required for students in the honor(s) section, they are bonus for those in the main section.

12. 0 points When dealing with trees, algorithm complexity is usually measured in terms of the number of nodes in the tree while the elementary operation is usually a node traversal. Let T be a binary tree such that each node u has a `parent`, `rightChild` and `leftChild` and a *key* element, `key`. The left/right child may be `null` indicating it does not exist.

Given a binary tree, we wish to determine if it is a binary search tree or not.

- (a) (5 points) Gomer thinks we can solve this problem by using a preorder traversal. When each node u is processed, it simply verifies that the left child's key is strictly less than u 's key and u 's right child's key is strictly greater than u 's key. Gomer's solution, however, will not work. Provide a counter example to demonstrate this and explain how it shows Gomer is wrong.
- (b) (20 points) Design a *correct* algorithm that, given a root node in a binary tree (with `parent`, `rightChild`, `leftChild` and `key` elements) determines if it is a binary search tree or not. Fully analyze your algorithm.