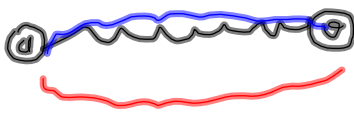


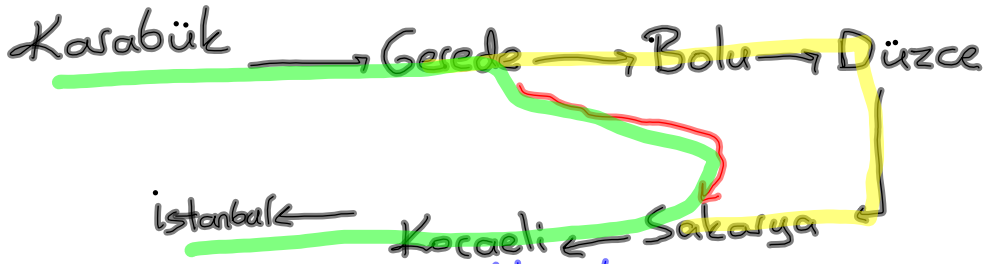
## En kısa Yollar



$\delta(u, v)$ , en kısa yolun ağırlığı

## Optimal Altıyapı

**Teorem:** En kısa yolun bir alt yolu en kısa bir yoldur.



## Dijkstra Algoritması

Bu alg. pozitif ağırlıklı veya negatif ağırlıklı çevrime sahip olmayan graflarda çalışır. Bir kaynaktan diğer tüm tepelelere varsa en kısa yolları bulur.

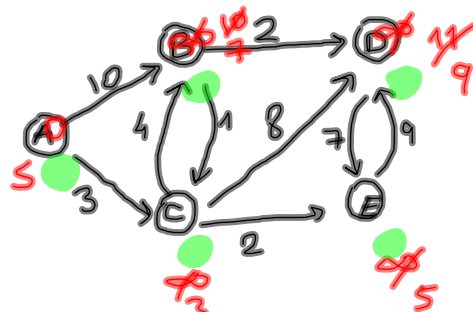
Dijkstra ( $G, s$ )

```

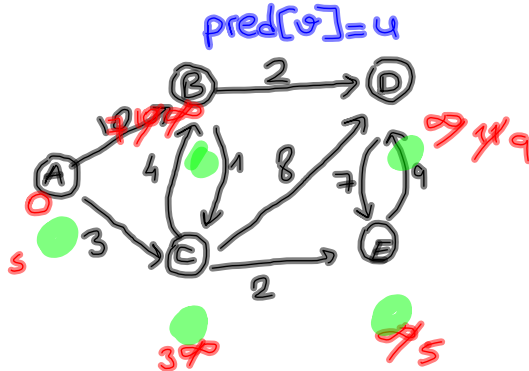
d[s] ← 0
for each v ∈ V - {s}
  do d[v] ← ∞; pred[v] = NULL
S ← ∅
Q ← V
while Q ≠ ∅
  do u ← EXTRACT-MIN(Q)
  S ← S ∪ {u}
  for each v ∈ Adj[u]
    do if d[v] > d[u] + w(u, v)
       then d[v] ← d[u] + w(u, v)
       pred[v] = u

```

*relaxation step*



Q:	A	B	C	D	E
d:	0	∞	∞	∞	∞
		10	3	∞	∞
		7		11	5
		7		11	9



pred =	NULL	C	A	B	C
Q =	A	B	C	D	E
d =	0	7	3	9	5

Dijkstra ( $G, s$ )

```

d[s] ← 0; pred[s] = NULL
for each v ∈ V - {s}
  do d[v] ← ∞
  pred[v] = NULL
Q ← V
while Q ≠ ∅
  do u ← EXTRACT-MIN(Q)

```

## Dijkstra Algoritmasının Analizi

$\Theta(V)$  {  $\Theta(1)$   $d[s] \leftarrow 0$   
 $\Theta(V)$  for each  $v \in V - \{s\}$   
 $\Theta(V)$  do  $d[v] \leftarrow \infty$   
 $\Theta(V)$   ~~$S \leftarrow S \cup \{u\}$~~   
 $\Theta(V)$   $Q \leftarrow V$   $\triangleright Q$  is a priority queue maintaining  $V - S$   
 $\Theta(V)$  while  $Q \neq \emptyset$   
 $\Theta(V)$  do  $u \leftarrow \text{EXTRACT-MIN}(Q)$  ?  
 $\Theta(V)$   ~~$S \leftarrow S \cup \{u\}$~~   
 $\Theta(V)$  for each  $v \in \text{Adj}[u]$   
 $\Theta(V)$  do if  $d[v] > d[u] + w(u, v)$  relaxation step ? decrease key  
 $\Theta(V)$  then  $d[v] \leftarrow d[u] + w(u, v)$  decrease key

$$T(E, V) = \Theta(V) \cdot T_{\text{extract-min}} + \Theta(E) \cdot T_{\text{decrease-key}}$$

$Q$	$T_{\text{extract-min}}$	$T_{\text{decrease-key}}$	$T(E, V)$
dizi	$\Theta(V)$	$\Theta(1)$	$\Theta(V^2)$
ikili heap	$\Theta(\lg V)$	$O(\lg V)$	$\Theta(E \lg V)$