A Las vegas algorithm to solve Elliptic Curve Discrete Logarithm Problem in public key cryptography

Interdisciplinary School of Scientific Computing - SPPU

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Day Month Year

Dedication

Lorem Ipsum is simply dummy text of the printing and typesetting industry. Lorem Ipsum has been the industry's standard dummy text ever since the 1500s, when an unknown printer took a galley of type and scrambled it to make a type specimen book. It has survived not only five centuries, but also the leap into electronic typesetting, remaining essentially unchanged. It was popularised in the 1960s with the release of Letraset sheets containing Lorem Ipsum passages, and more recently with desktop publishing software like Aldus PageMaker including versions of Lorem Ipsum.

Certificate of the GUIDE

I declare that the thesis titled **"Las Vegas algorithm to solve Elliptic Curve Discrete Logarithm problem in public key cryptography"** submitted by me for the degree of Doctor of Philosophy is the record of word carried out by me during the period from January 2016 to JUNE 2021 under the guidance of Dr. Smita Bedekar (ISSC-SPPU) and Dr. Ayan Mahanalobis (IISER-PUNE) and had not formed the basis for the award of any degree, diploma, associateship, fellowship, titles in this or any other University or the other institution of Higher learning. I further declare that the material obtained from other sources had been duly acknowledged in the thesis.

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Chapter 1

An Introduction

1.1 Introduction

An ability to have a secure communication medium has been of human interest since time immemorial. A medium is secure for communication if it allows information flow between two entities such that, even if any third person gains access to this channel the information flowing is of no use to that person. For the information to be useless for any third person it has to be converted to a form that makes no sense for any third person. This conversion process is known as encryption and the converted message as ciphertext. The process of getting back the original message from ciphertext is decryption. Thus,

$$Cryptography = Encryption + Decryption.$$

The idea of Public Key Cryptosystem was proposed by Diffie and Hellman in 1976. Since then a number of cryptosystems have been proposed. Public key cryptographic protocols used in practical applications are based on some known mathematical hard problems. RSA is based on the **I**nteger **F**actorization **P**roblem (IFP), ElGamal is based on the **D**iscrete **L**og **P**roblem (DLP) and the [1] **E**lliptic **C**urve **C**ryptosystem (ECC) are based over the Discrete Log Problem in a group of points over an elliptic curve.

Recent see 1.1 advances in the index calculus attack for DLP by Joux have shown that the cryptosystems over a finite field of small characteristics are no longer secure. A Quasi-polynomial time algorithm to solve the DLP in a multiplicative group of finite extension field was proposed by Joux in. Since then the DLP is not as hard to solve as it used to be in some setting.

Index calculus algorithms are not suitable for for solving DLP efficiently in an elliptic curve group. Thus only attacks available for the DLP over elliptic curve are the generic algorithms like the Pollard rho or the r-adding walk. The time complexity for these algorithms is about the square root of the size of the group.

Figure 1.1: ASTON MARTIN DB9

As a result of the advancement in the index calculus attack for fields with small characteristics, the set of cryptographic algorithms proposed in "Suite B Cryptography" by the National Security Agency (USA) recommends ECC. Thus currently, there is no other highly recommended public key cryptosystem other than the ECC in use for practical applications. Public key Cryptosystems based on finite fields are over small characteristics, mostly binary field, i.e characteristic 2. As binary fields are similar to the binary nature of a digital computer there are extremely efficient algorithms for performing operations over the field elements. Hence, cryptosystems over fields of characteristic greater than 2 were never needed in practice.

Thus to summarize the present situation in public key cryptography :-

- No recommended Public key cryptosystem other than ECC

- Non-ECC over small characteristics are not secure for practical applications

As no cryptosystem other than Elliptic curve cryptosystem is available for higher characteristic we aim to develop a new public-key cryptosystem based on the Discrete Log Problem for higher characteristics. This system would be based on the DLP over circulant matrices which are defined over a finite field of large characteristic. A cryptosystem based on characteristic 2 fields was proposed in. Cryptographic protocols based over any general or a special type of matrices are not used in everyday cryptography. Generally, it is believed that matrices offer no major security or efficiency advantage. Here we would be working with matrices of a special type. These matrices have a certain structure in them, using this structure matrix operations can be performed in an efficient manner. These matrices can provide security and computational efficiency as compared to systems over finite fields used today in public key cryptography.

1.2 Cryptosystem to be implemented

1.2.1 Circulant matrices

An $n \times n$ matrix C is a circulant matrix if all its rows (or columns) other than the first are a circular shift of the previous row (or column).

$$C = \begin{pmatrix} c_0 & c_1 & \cdots & c_{n-2} & c_{n-1} \\ c_{n-1} & c_0 & \cdots & c_{n-3} & c_{n-2} \\ c_{n-2} & c_{n-1} & \cdots & c_{n-4} & c_{n-3} \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ c_1 & c_2 & \cdots & c_{n-1} & c_0 \end{pmatrix}$$

An $n \times n$ circulant matrix has atmost n different entries and can be determined completly by the first row, i.e. A circulant matrix is a special case of the Toeplitz matrix which is a square matrix with constant diagonals. Even though a matrix is a 2-dimensional entity, due to the structure of a circulant matrix it behaves like a one-dimensional array. Thus allowing operations on the matrix to be performed in an efficient manner. Circulant matrices form a commutative ring under matrix multiplication and addition.

1.2.2 ElGamal cryptosystem

The ElGamal Cryptosystem is a public key cryptosystem based on the discrete logarithm in a group G. Given a primitive element a $\in$ G and another element b, DLP is the computational problem of finding $x = log_a(b)$ such that $b = a^x$.

1.2.3 Discrete Logarithm problem in matrices

The DLP in matrices is to find m, from A and A^m where $A \in GL(d,q)$. Here $GL(d,q)$ is a group of nonsingular matrices of size d over the finite field F

1.2.4 Aim

We will work with circulant matrices and ElGamal cryptosystem, to determine how fast one can do computations with circulant matrices when implemented in a **C**omputer **A**lgebra **S**ystem (CAS) like Victor Shoup's **N**umber **T**heory **L**ibrary (NTL). How this cryptosystem will behave with finite fields of higher characteristics will be interesting as public key cryptosystems have never been implemented over finite fields of large characteristics.

How much secure this cryptosystem is? Is the cryptosystem secure for large characteristic? Is it efficient than the cryptosystem implemented over a finite field for large characteristic? These will be some of the questions to be answered during the course of the work.

Exponentiation is the fundamental operation required for the ElGamal cryptosystem. Developing a cryptosystem over ElGamal with circulant matrices will require implementing exponentiation in circulant matrices. Exponentiation over matrices is available in NTL. However this implementation is for generic matrices, i.e structure of the matrix is not used for exponentiation computation. One of the better-known algorithm for exponentiation computation is the square and multiply algorithm. The two fundamental operations here are squaring and multiplication. It is known that squaring is computationally free for finite fields of characteristic 2 using normal basis. How fast can matrix multiplication be done in this case is a question that we will try to answer. For a finite field of higher characteristic (>2), squaring is not computationally free. How fast one can perform squaring and multiplication of circulant matrices in these fields will be another question to be answered.

Generating secure instances of circulant matrices to support the validity of the cryptosystem will be needed. This is a computationally demanding task. Computational complexity increases exponentially as size of the field increases.

The ElGamal cryptosystem works with any type of matrices. Circulant matrices were chosen because they have a structure which makes them interesting from the cryptographic perspective. Similar generalization of circulant matrices are the Block circulant matrices, which too have some structure. The structure here applies to the blocks rather than the elements of the matrix.

1.3 Study the generalization of circulant matrices

As there is a structure in Block circulant matrices it would be interesting to study them with the objective of using them to develop a public key cryptosystem. Parallel matrix multiplication algorithms operate on matrices by dividing the input matrix into blocks and distributing these blocks onto different processors to get the multiplication done. As block circulant matrices have an internal structure implementing them in a parallel environment can improve the efficiency of multiplication. Thus, block circulant matrices can extract advantage from a parallel implementation, whereas for finite field implementation parallelization is not straightforward. The three generalization of the circulant matrix are:

- Block Circulant matrices with circular blocks

- Block Circulant matrices with non-circular block

- Non-Circulant matrices with circular block

1.3.1 Block Circulant matrices with Circular Blocks (BCCB)

$$C_1 = \left(\begin{array}{ccc|ccc} c_0 & c_1 & c_2 & c_3 & c_4 & c_5 \\ c_2 & c_0 & c_1 & c_5 & c_3 & c_4 \\ c_1 & c_2 & c_0 & c_4 & c_5 & c_3 \\ \hline c_3 & c_4 & c_5 & c_0 & c_1 & c_2 \\ c_5 & c_3 & c_4 & c_2 & c_0 & c_1 \\ c_4 & c_5 & c_3 & c_1 & c_2 & c_0 \end{array} \right)$$

A BCCB is not a circular matrix, the blocks are circular and each block row is a circular shift of the previous block row. In the example above C_1 is a BCCB with two 3×3 circular blocks, each block is a circular matrix in itself.

1.3.2 Block Circulant matrices with non-circular block

$$C_2 = \left(\begin{array}{ccc|ccc} c_0 & c_1 & c_2 & c_3 & c_4 & c_5 \\ c_2 & c_0 & c_1 & c_5 & c_3 & c_4 \\ c_1 & c_2 & c_0 & c_4 & c_5 & c_3 \\ \hline c_6 & c_7 & c_8 & c_9 & c_{10} & c_{11} \\ c_8 & c_6 & c_7 & c_{11} & c_9 & c_{10} \\ c_7 & c_8 & c_6 & c_{10} & c_{11} & c_9 \end{array} \right)$$

In block circulant matrix with non-circular block the blocks themselves are circular. The matrix is not circular. Thus a row from each block is required to represent the entire matrix.

1.3.3 Non-Circulant matrices with circular block

$$C_3 = \left(\begin{array}{ccc|ccc} c_0 & c_1 & c_2 & c_9 & c_{10} & c_{11} \\ c_3 & c_4 & c_5 & c_{12} & c_{13} & c_{14} \\ c_6 & c_7 & c_8 & c_{15} & c_{16} & c_{17} \\ \hline c_9 & c_{10} & c_{11} & c_0 & c_1 & c_2 \\ c_{12} & c_{13} & c_{14} & c_3 & c_4 & c_5 \\ c_{15} & c_{16} & c_{17} & c_6 & c_7 & c_8 \end{array} \right)$$

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Chapter 2

Algorithms to solve ECDLP

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Chapter 3

The LasVegas Algorithm

3.1 Introduction

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Chapter 4

Schur Complement to solve ECDLP

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Extension Minor, Maximal Minor, minus-3 hypothesis

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Conclusion

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typesetting, remaining essentially unchanged. It was popularised in the 1960s with the release of Letraset sheets containing Lorem Ipsum passages, and more recently with desktop publishing software like Aldus PageMaker including versions of Lorem Ipsum.

Appendix A

How to use the code

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Appendix B

Something more

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List of publications