

# DSA

→ Infinite algorithm's time complexity can't be measured

Efficiency  $\Rightarrow$  time  
memory occupation

## Methods for time complexity

→ Any method which relies on ~~prod~~ the machine is invalid (machines could have varying speeds)

prior method (approximation)

→ Input size matters e.g

```
for(int i=0; i<4; i++)  
{ cout << "Wassup"; }
```

If ~~one~~ <sup>single</sup> cycle execution of command takes 1 sec then  $f(n) = 4$

```
int a = 0;
```

```
cin >> a;
```

```
for (int i = 0; i < n9; i++)
```

```
{ a = a + 1; }
```

$$f(n) = n + 2$$

~~f(n)~~  $f(n) = 11$

```
for (int i = 0; i < 4; i++)
```

```
{ for (int j = 0; j < 4; j++)
```

```
{ cout << "hi"; }
```

$$f(n) = n^2$$

in this case  $f(n) = 16$

Growth rate of  $n$  : (asymptotic complexity)

$$f(n) = 3n^2 + 8n$$



$$\text{If } n=1$$

$$\therefore \text{ due to } 3n^2 = 25 \text{ i.e. app}$$

$$\therefore \text{ due to } 8n = 75 \text{ i.e. app}$$

$$\text{Let } n=10$$

$$\therefore 3n^2 = 79 \text{ i.e.}$$

$$\therefore 8n = 21 \text{ i.e.}$$

$$\text{Let } n=100$$

$$\therefore 3n^2 = 97 \text{ i.e. app}$$

$$\therefore 8n = 3 \text{ i.e. app}$$

$$\text{Big O} \\ O()$$

$$\text{Big omega} \\ \Omega()$$

$$\text{Big theta} \\ \Theta()$$

→ upper bound (least)  
Big O (worst case scenario)

$$f(n) = O(g(n))$$

$$f(n) \leq c \cdot g(n)$$

$$c > 0$$

$$n > k$$

$$k > 1$$

e.g.

$$f(n) = 3n^2 + n$$

$$3n^2 + n \leq c \cdot g(n)$$

$$3n^2 + n \leq 4n^2$$

$$n \geq 4/3$$

