

Quiz 1 Solution.

Q1

$$f(n) = n^{\frac{1}{2}} \log n^{\frac{1}{2}}$$
$$g(n) = 30000 n^{\frac{1}{2}} \log n^{\frac{1}{4}}$$

$$\boxed{f(n) = \Theta(g(n))}$$

$$\Rightarrow f(n) = \frac{1}{2} \sqrt{n} \log n \quad \Rightarrow \quad g(n) = \frac{30000}{4} \sqrt{n} \log n$$

$f(n)$ and $g(n)$ are both of same growth rate ($\sqrt{n} \log n$). So, so particular c_1, c_2 and n_0 , $g(n)$ can be upper as well as lower bound of $f(n)$.

$$f(n) = \Omega(g(n))$$

$$f(n) = O(g(n))$$

$$\text{So, } f(n) = \Theta(g(n))$$

Q2 $f(n) = \frac{\sqrt{n}}{18} - 19 \log n + 20$

$$k_1 g(n) \leq f(n) \leq k_2 g(n)$$

$$\boxed{g(n) = \sqrt{n}}$$

$$k_1 \sqrt{n} \leq \frac{\sqrt{n}}{18} - 19 \log n + 20 \leq k_2 \sqrt{n}$$

for Upper bound:

let $k_2 = 50$, then

$f(n) \leq 50\sqrt{n}$ holds for all values of n , so $n_0 = 1$.

for lower bound:

$$\text{let } k_1 = \frac{1}{1000}$$

$$\frac{\sqrt{n}}{1000} \leq \frac{\sqrt{n}}{18} - 19 \log n + 20$$

Now, we can't simply find n_0 , as for some range of n , the factor $19 \log n$ would be large and make the R.H.S. -ve. And since, we can't solve this eq. by calculator to find roots, we'll find n_0 by hit and trial.

$$\text{for } n_0 = 5000000 \quad k_1 g(n) \leq f(n)$$

Ans

$$k_1 = \frac{1}{1000}, \quad k_2 = 50, \quad n_0 = 5000000$$

pro tip

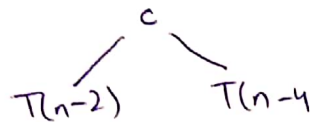
In log changing value of n from 1000 to 10000 will show no difference.

Change it by a factor of 10 i.e 1000 to 10000

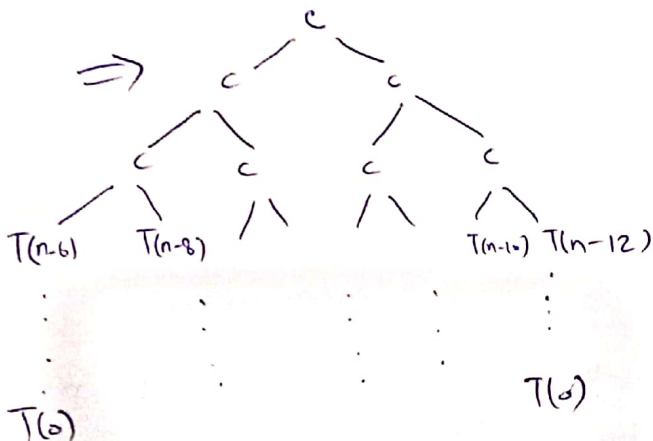
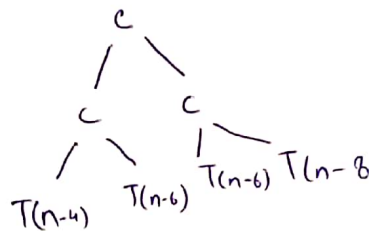
Q3

$$T(n) = T(n-2) + T(n-4) + c$$

By Tree Method



\Rightarrow



$T(n-2)$ will give the longest branch. Size of subproblems at different levels:

$$n, n-2, n-4, n-6, \dots, 0$$

$$\boxed{\text{Total levels} = \frac{n+1}{2}}$$

Work/cost at different levels

$$\begin{array}{l} c \\ 2c \\ 4c \\ 8c \\ \vdots \end{array}$$

$$= (c + 2c + 4c + 8c + \dots + 2^k c)$$

$$= c(1 + 2^1 + 2^2 + \dots + 2^k)$$

Total terms of this geometric series will be $\frac{n}{2} + 1$.

$$= \frac{c(2^{\frac{n}{2}+1} - 1)}{2-1} = c \cdot 2^{\frac{n}{2}+1} - c = c_1 2^{\frac{n}{2}} - c$$

$$\boxed{T(n) = O(2^{\frac{n}{2}})}$$

OR $2^{\frac{n}{2}} = (2^n)^{\frac{1}{2}} = (2^{\frac{1}{2}})^n = (\sqrt{2})^n = (1.414)^n$

$$\boxed{T(n) = O(1.4^n)}$$

At some point when the smallest branch ends, cost will not follow this series form, but for simplicity we assume that it does.