# National University of Computer and Emerging Sciences, Lahore Campus

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| Course Name:    | Design and Analysis of Algorithms | Course Code: | CS2009      |
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| Degree Program: | BSCS                              | Semester:    | Spring 2023 |
| Due Date:       | 16-05-2023                        | Total Marks: | 70          |
| Section:        | All                               | Page(s):     | 2           |
| Exam Type:      | Assignment 3                      | CLO          | 1           |

| Student: Name: | Roll No | Section: |
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### Instructions/Notes:

- 1. You have to submit this assignment in hand written form in class, no typed/printed/online submission will be accepted.
- 2. No late submission will be entertained.
- 3. Assignment will be collected at the start of class on Tuesday, May 16, 2023.

#### Q1:

Professor Sabatier conjectures the following converse of Theorem 23.1. Let G = (V, E) be a connected, undirected graph with a real-valued weight function w defined on E. Let A be a subset of E that is included in some minimum spanning tree for G, let (S, V - S) be any cut of G that respects A, and let (u, v) be a safe edge for A crossing (S, V - S). Then, (u, v) is a light edge for the cut. Show that the professor's conjecture is incorrect by giving a counterexample.

#### Q2:

What is the running time of BFS if we represent its input graph by an adjacency matrix and modify the algorithm to handle this form of input?

### Q3:

Give an example of a directed graph G=(V,E), a source vertex  $s \in V$ , and a set of tree edges  $E_{\pi} \subseteq E$  such that for each vertex  $v \in V$ , the unique simple path in the graph  $(V, E_{\pi})$  from s to v is a shortest path in G, yet the set of edges  $E_{\pi}$  cannot be produced by running BFS on G, no matter how the vertices are ordered in each adjacency list.

### Q4:

Modify the pseudocode for depth-first search so that it prints out every edge in the directed graph G, together with its type. Show what modifications, if any, you need to make if G is undirected.

#### Q5:

Rewrite the procedure DFS, using a stack to eliminate recursion.

# Q6:

Give a counterexample to the conjecture that if a directed graph G contains a path from u to v, and if u.d < v.d in a depth-first search of G, then v is a descendant of u in the depth-first forest produced.

# Q7:

Give a counterexample to the conjecture that if a directed graph G contains a path from u to v, then any depth-first search must result in  $v \cdot d \leq u \cdot f$ .