

DATA STRUCTURE & ALGORITHM

Chapter 09

Queue

Part 1 -Queue Implementation Linked List

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Objectives

At the end of the lesson students are expected to be able to:

- Understand queue concepts and applications.
- Understand queue structure and operations that can be done on queue.
- Understand and know how to implement queue using array and linked list : linear array, circular array, linear link list and circular list.

Introduction to Queue

New items enter at the back, or rear, of the queue

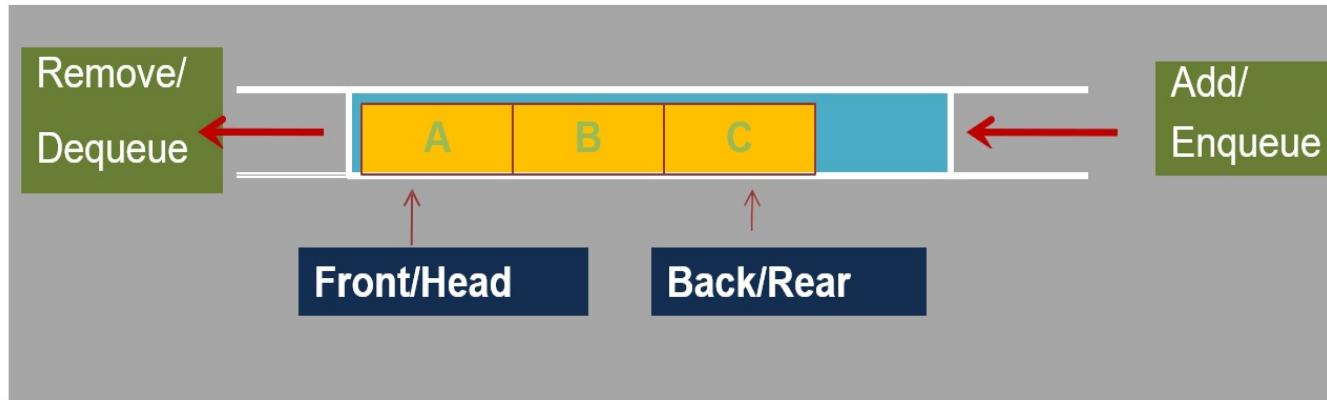
- Items leave from the front of the queue
- First-in, first-out (FIFO) property
 - The first item inserted into a queue is the first item to leave
 - Middle elements are logically inaccessible
- Important in simulation & analyzing the behavior of complex systems

Queue Applications

Real-World Applications

- Cashier lines in any store
- Check out at a bookstore
 - Bank / ATM
 - Call an airline
- Computer Science Applications
 - Print lines of a document
 - Printer sharing between computers
 - Recognizing palindromes
 - Shared resource usage (CPU, memory access, ...)
- Simulation
 - A study to see how to reduce the wait involved in an application

Queue Implementation

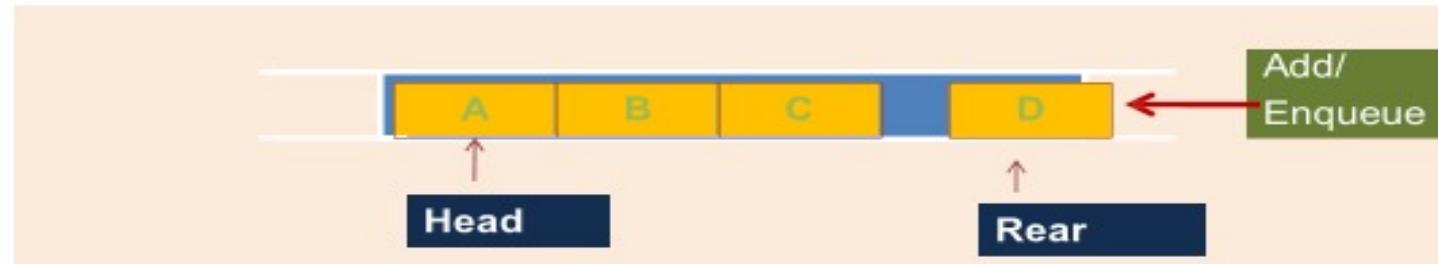


Basic Structure of a Queue:

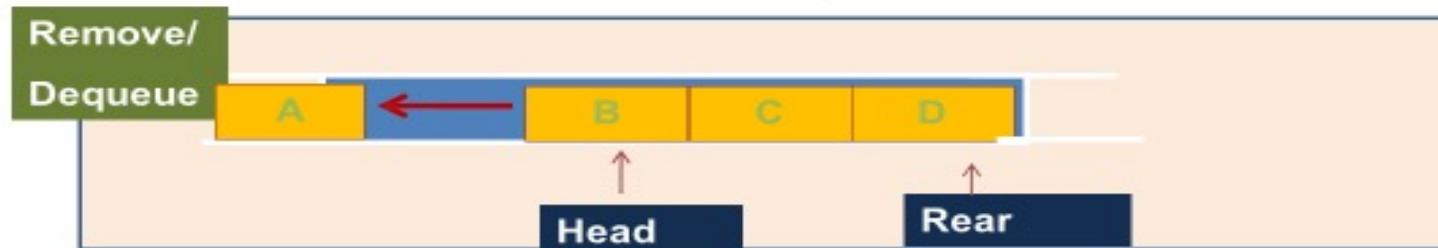
- Data structure that hold the queue
- head
- rear

Queue Implementation

Queue implementation



Insert D into Queue (enQueue) : D is inserted at rear



Delete from Queue (deQueue) : A is removed

Abstract Data Type Queue

Implementation:

- Array-based (Linear or Circular)
- Pointer-based : Link list (Linear or Circular)

Abstract Data Type Queue

- ADT queue operations
 - Create an empty queue
 - Destroy a queue
 - Determine whether a queue is empty
 - Add a new item to the queue
 - Remove the item that was added earliest
 - Retrieve at Front
 - Retrieve at Back the item that was added earliest

Queue: Linear Array Implementation

- Number of elements in Queue are fixed during declaration.
- Need ***isFull ()*** operation to determine whether a queue is full or not.
- Queue structure need at least 3 elements:
 - 1) Element to store items in Queue
 - 2) Element to store index at head
 - 3) Element to store index at rear

Queue

front

rear

items

createQueue ()

destroyQueue ()

isEmpty ()

isFull ()

enQueue ()

deQueue ()

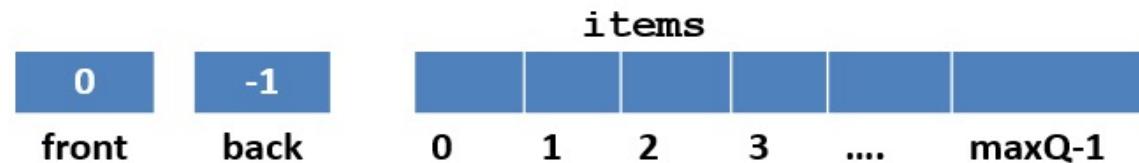
getFront ()

getRear ()

CreateQueue() operation

- Linear Array implementation
 - *front* & *back* are indexes in the array
 - Initial condition: *front* = 0 & *back* = -1

```
Queue::Queue()  
{ front = 0;  
  back = -1;  
}
```



Initial state for a queue linear array

Queue operations

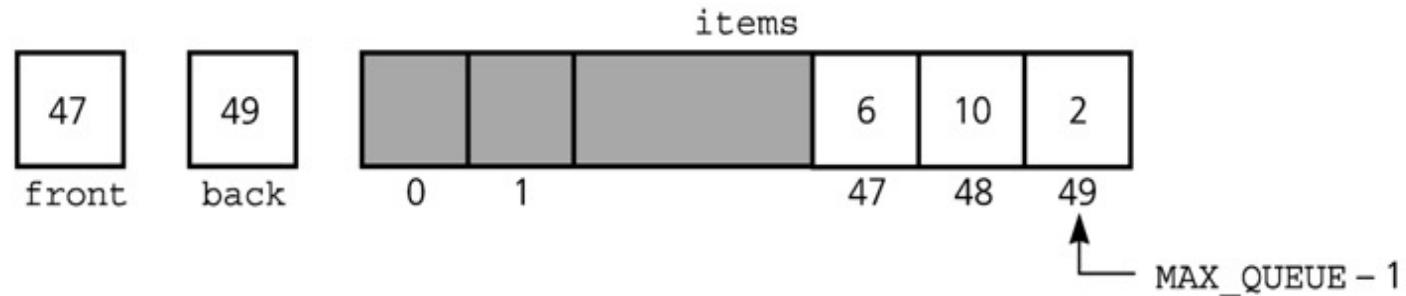
- **Destroy Queue distructor : All elements in the queue will be disposed.**

```
queue::~queue()
{ delete [ ] items; }
```

- **Check whether a queue is empty**
 - Queue Empty Condition : back < front

```
bool queue::isEmpty()
{ return bool(back < front); }
```

Queue operations



Check whether a queue is Full

- Queue Full Condition : $\text{back} = \text{size} - 1$

```
bool queue::isFull()
{ return bool(back == size - 1); }
```

- No more item can be insert into a queue, when a queue is full.

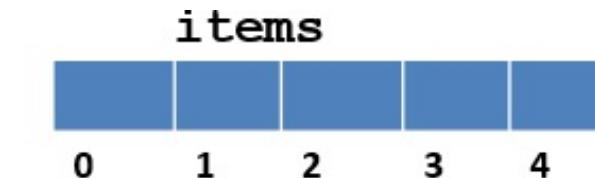
Queue operations

- Insert into a queue (enqueue)
 - Increment back
 - Insert item in `items[back]`

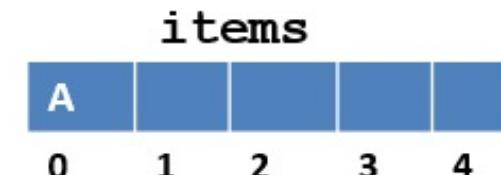
```
void queue::enqueue(char insertItem)
{ if (isFull())
    cout<< "\nCannot Insert. Queue is full!";
else
{ //insert at back
    back++;
    items[back] = insertItem;
} // end else if
}
```

enQueue operations for a queue with size = 5

Queue myQueue;



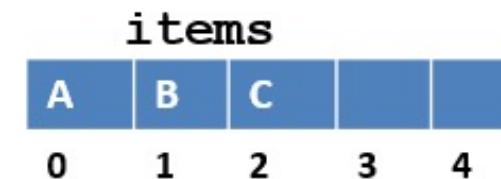
myQueue.enQueue('A');



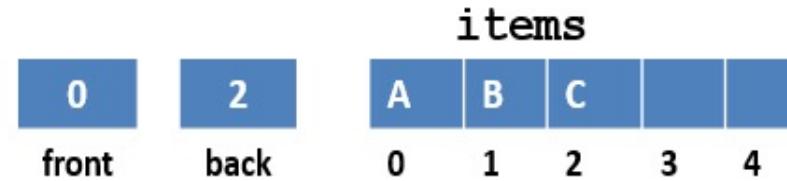
myQueue.enQueue('B');



myQueue.enQueue('C');



Queue operations



- Item at front and back can be retrieved from queue

```
char queue:: getFront() // get item at Front  
{ return items[front] ; }
```

```
char queue::getRear() // get item at Back  
{ return items[back] ; }
```

```
cout << myQueue.getFront(); //output is A  
cout << myQueue.getRear(); // output is C
```

Queue operations

- Delete from a queue (deQueue)
 - Increment front

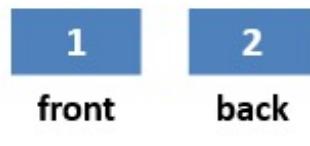
```
void queue::deQueue()
{ if (isEmpty())
    cout<< "\nCannot remove item. Empty Queue!";
else
{ //retrieve item at front
 deletedItem = items[front];
 front++;
 } // end else if
}
```

deQueue operations

`myQueue.deQueue();`

`deletedItem`

A



`items`



`myQueue.deQueue();`

`deletedItem`

B



`items`



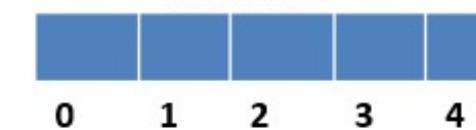
`myQueue.deQueue();`

`deletedItem`

C



`items`



`myQueue.deQueue();`

Cannot remove item.

Queue is Empty with back < front

Queue operations - enQueue

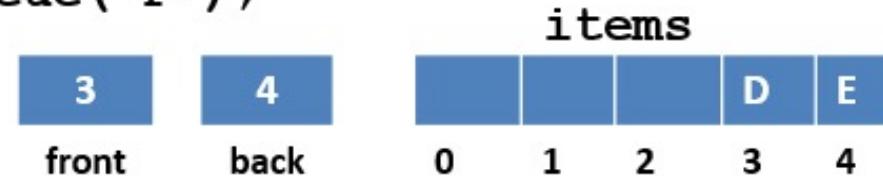
`myQueue.enQueue('D') ;`



`myQueue.enQueue('E') ;`



`myQueue.enQueue('F') ;`



Cannot insert F, even though there are empty spaces
in front of the queue array.

Currently, Queue is FULL with `back == size - 1`.

Problem: Rightward-Drifting:

- After a sequence of additions & removals, items will drift towards the end of the array
- Even though, there are empty spaces in front of the queue array, enQueue operation cannot be performed on the queue, since $\text{back} = \text{max_queue} - 1$.

Rightward
drifting



Rightward drifting solutions

To optimize space and to solve rightward drifting:

- Shift array elements after each deletion.

`myQueue . deQueue () ;`

`deletedItem`

A



`items`



Shift array elements
to front array

`items`



However, shifting is not effective and dominates the cost of the implementation.

Queue circular array

- Problem:
 - front & back no longer can be used to distinguish between queue-full & queue-empty conditions
- Solution:
 - Use a counter
 - count == 0 means empty queue
 - count == MAX_QUEUE means full queue
- Disadvantage
 - Overhead of maintaining a counter or flag

Circular Array Implementation

- Queue declarations

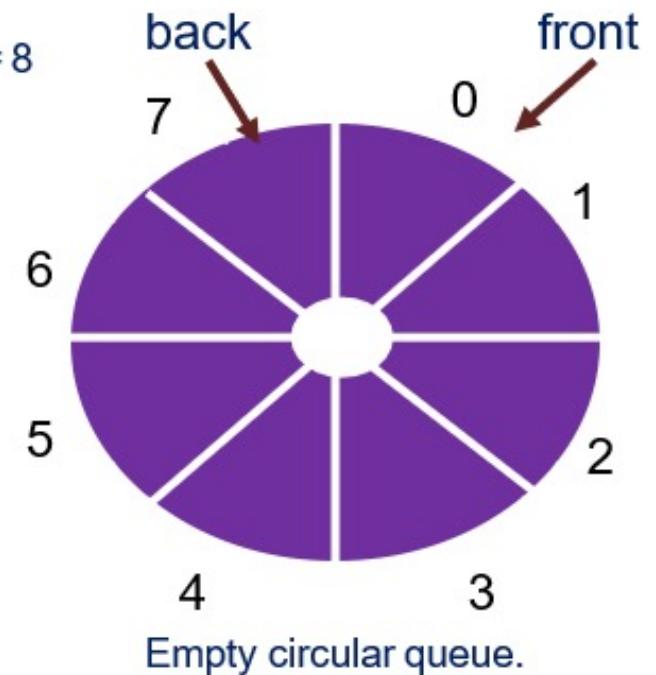
```
const int MAX_QUEUE = maximum-size-of-queue;
QueueItemType items [MAX_QUEUE];
int front;
int back;
int count
```

MAX_QUEUE = 8
count = 0

- Initial condition:

- count = 0, front = 0,
- back = MAX_QUEUE - 1

- The Wrap-around effect is obtained by using arithmetic (%-operator)



Circular Array Implementation

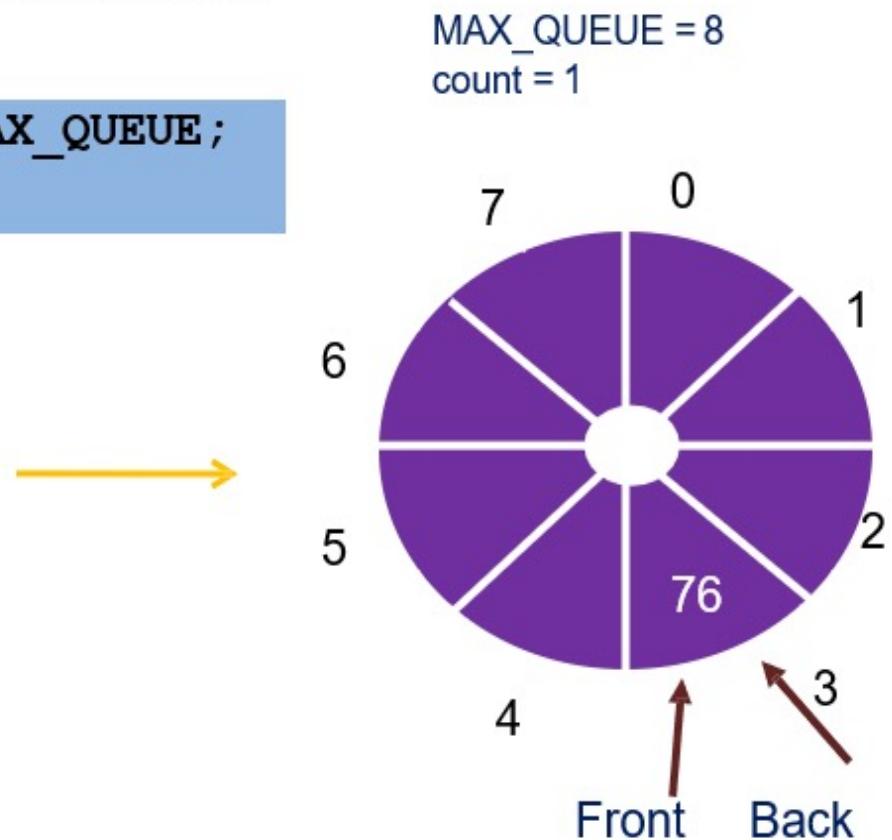
Deletion

Increment *front* using modulo arithmetic

Decrement *count*

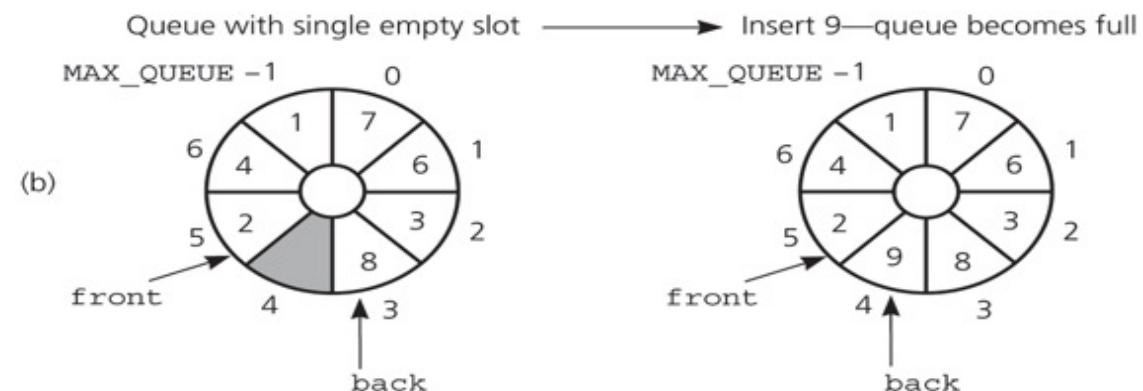
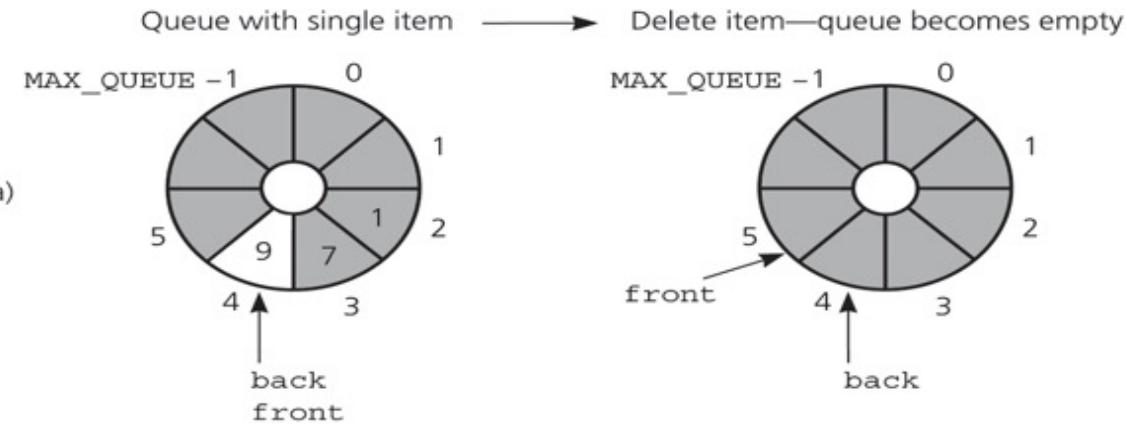
```
front = ( front + 1 ) % MAX_QUEUE;
--count;
```

After delete 20, 45 and 51 sequentially from circular queue



Circular Array Implementation

- (a) *front* passes *back* when the queue becomes empty;
- (b) *back* catches up to *front* when the queue becomes full





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Next....

Queue implementation Linked List