

DATA STRUCTURE & ALGORITHM

Chapter 09

Queue

Part 2 -Queue Implementation Linked List

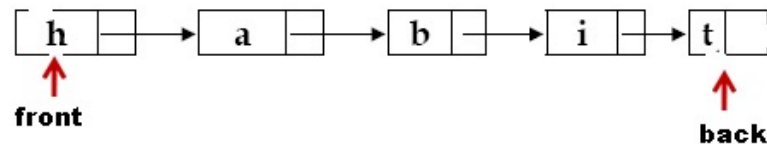
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Queue Implementation Linked List

Pointer-Based Implementation

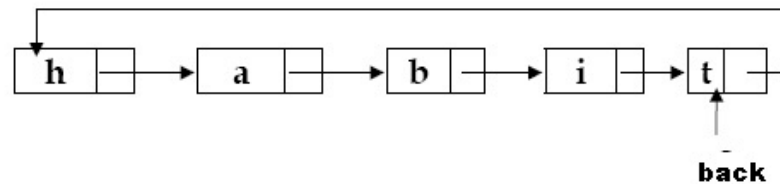
- More straightforward than array-based
- Possible options:
 - Linear linked list

Two external pointer (Front & Back)



- Circular linked list

One pointer will be enough (Back)



Queue Implementation Linked List

Need 2 structure

- Declaration of the node

```
struct nodeQ {  
    char item;  
    nodeQ * next;  
}
```

- Declaration of the queue

```
class queue  
{public:  
    nodeQ *backPtr, * frontPtr;  
    // operations for queue  
};
```

Queue

frontPtr
backPtr

```
createQueue()  
destroyQueue()  
isEmpty();  
enqueue();  
dequeue();  
getFront();  
getRear();
```

Queue Implementation Linked List

createQueue()

```
backPtr = Null; frontPtr = NULL;
```

destroyQueue()

Destroy the whole nodes in the queue

```
nodeQ *temp = frontPtr;  
while (temp){  
    frontPtr = temp->next;  
    delete temp; temp=frontPtr; }
```

isEmpty()

```
backPtr == Null && frontPtr == NULL
```

Insert to a linear queue

- Inserting a new node at the Back needs 3 pointer changes

Change Next pointer in the new node

Change the Next pointer in the Back node

Change the External pointer

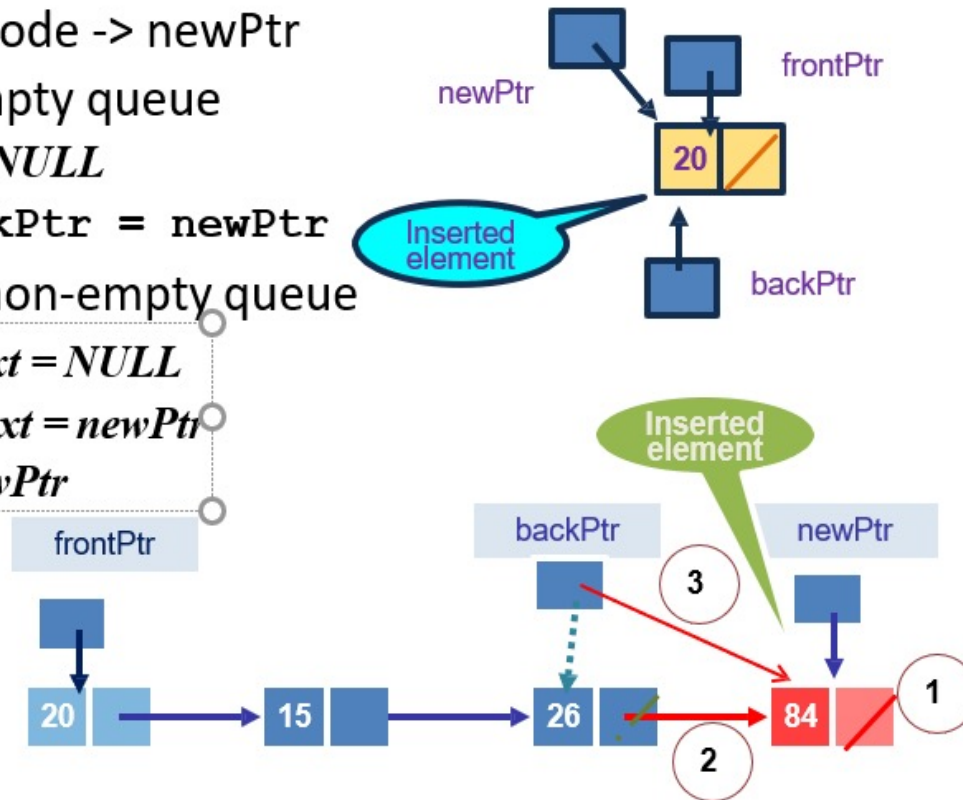
- Special case:
 - If the queue is empty

Queue Implementation Linked List

Linear linked list with 2 external pointers

1. Create a new node -> newPtr
2. Insert to an empty queue
 $newPtr \rightarrow next = NULL$
 $frontPtr = backPtr = newPtr$
3. Insertion to a non-empty queue

- 1 $newPtr \rightarrow next = NULL$
- 2 $backPtr \rightarrow next = newPtr$
- 3 $backPtr = newPtr$



Delete from Linear queue

- Deletion
 - Delete from the Front
 - Only one pointer change is needed
 - Special case:
 - If the queue contains one item only

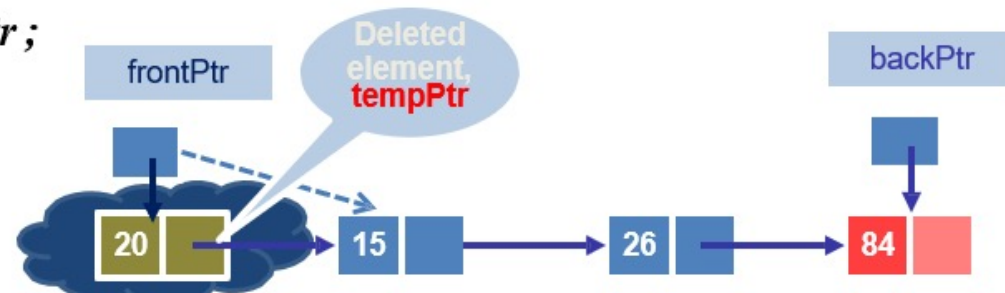
- Deletion Code

tempPtr = frontPtr

frontPtr = frontPtr -> next

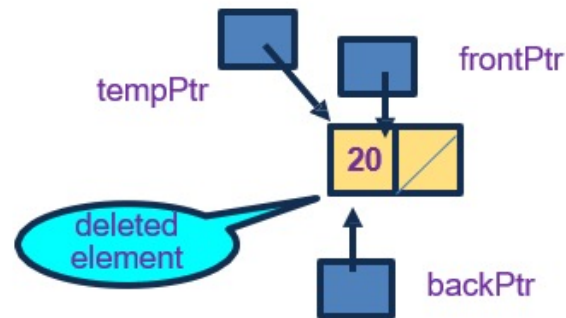
tempPtr -> next = NULL

delete tempPtr ;



Delete from Linear queue

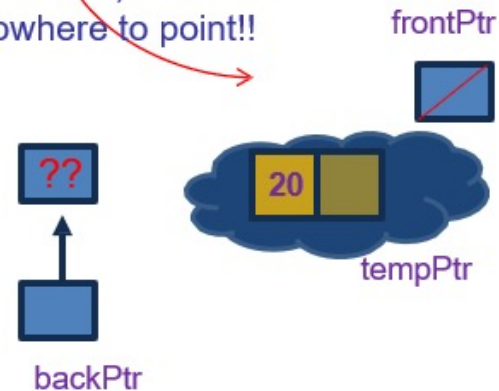
If the queue contains one item only,



Need to add this statement:

If (!frontPtr)
backPtr = NULL;

After deletion, backPtr
has nowhere to point!!

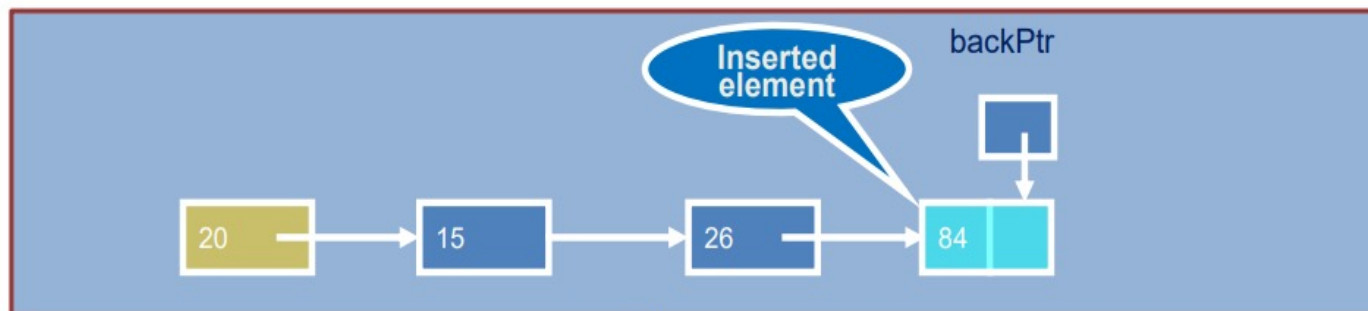
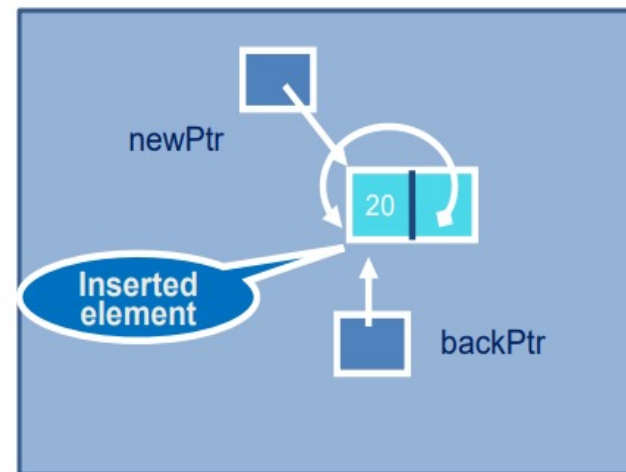


Circular Queue Implementation

Circular linear linked list with one external pointer

– Insertion

- Into an empty queue
 $\text{NewPtr} \rightarrow \text{Next} = \text{NewPtr}$
 $\text{BackPtr} = \text{NewPtr}$
- Into a non-empty queue
 $\text{NewPtr} \rightarrow \text{Next} = \text{BackPtr} \rightarrow \text{Next}$
 $\text{BackPtr} \rightarrow \text{Next} = \text{NewPtr}$
 $\text{BackPtr} = \text{NewPtr}$



Circular Queue Implementation

Deletion

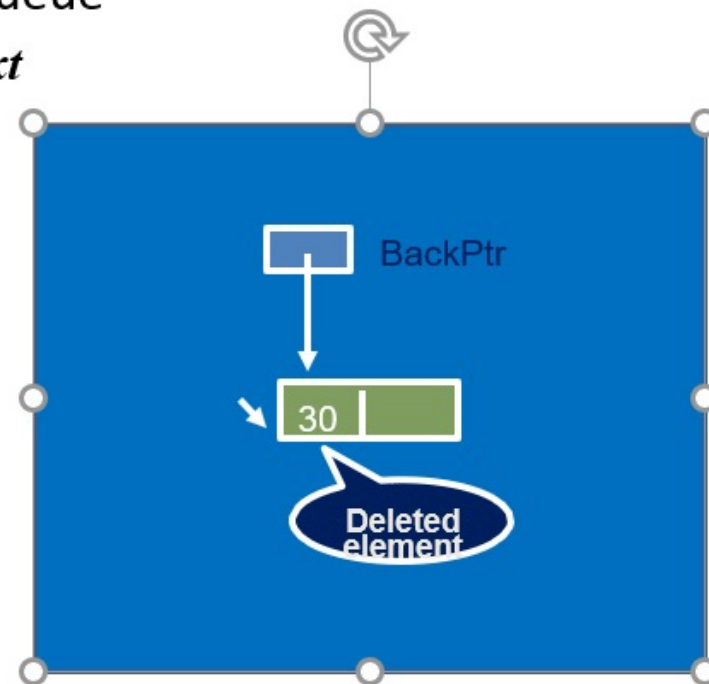
- From a one-node (one item) queue

$deletePtr = BackPtr \rightarrow Next$

If ($deletePtr = BackPtr$)

$BackPtr = NULL$

delete deletePtr



Circular Queue Implementation

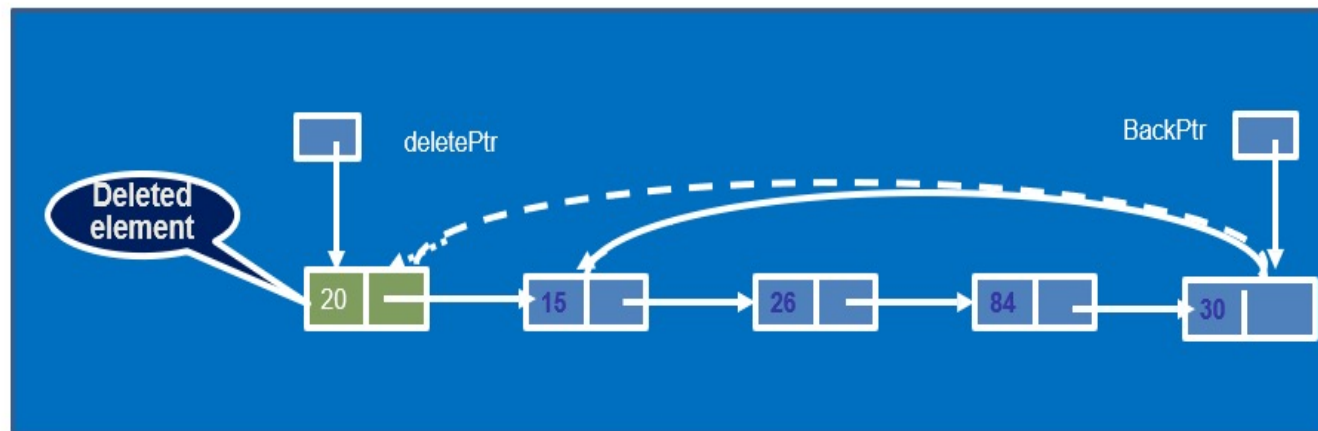
Deletion

- From a non-empty, more than one item queue

deletePtr = BackPtr -> Next

BackPtr -> Next = deletePtr -> Next

delete deletePtr



Comparing Implementations

- Fixed size versus dynamic size
 - A statically allocated array
 - Prevents the *enqueue* operation from adding an item to the queue if the array is full
 - A resizable array or a reference-based implementation
 - Does not impose this restriction on the *enqueue* operation
- Pointer-based implementations
 - A linked list implementation
 - More efficient

Comparing Implementations

- Pointer-Based
 - More complicated than ADT List
 - Most flexible No size restrictions
- Array-Based
 - No overhead of pointer manipulation
 - Prevents adding elements if the array is full

Summary of Position-Oriented ADTs

Stacks

- Operations are defined in terms of position of data items
- Position is restricted to the Top of the stack. Only one end position can be accessed
- Operations:
 - *create*:
 - Creates an empty ADT of the Stack type
 - *isEmpty*:
 - Determines whether an item exists in the ADT
 - *push*:
 - Inserts a new item into the Top position
 - *pop*:
 - Deletes an item from the Top position
 - *peek*:
 - Retrieves the item from the Top position

Summary of Position-Oriented ADTs

Queues

- Operations are defined in terms of position of data items
- Position is restricted to the front & back of the queue. Only the end positions can be accessed
- Operations:
 - *create*:
 - Creates an empty ADT of the Queue type
 - *isEmpty*:
 - Determines whether an item exists in the ADT
 - *enqueue*:
 - Inserts a new item in the Back position
 - *dequeue*:
 - Deletes an item from the Front position
 - *peek*:
 - Retrieves the item from the Front position

Summary of Position-Oriented ADTs

- Stacks and queues are very similar
 - Operations of stacks and queues can be paired off as
 - *createStack* and *createQueue*
 - Stack *isEmpty* and queue *isEmpty*
 - *push* and *enqueue*
 - *pop* and *dequeue*
 - Stack *getTop* and queue *getFront*

References

- Frank M. Carano, Janet J Prichard. *“Data Abstraction and problem solving with C++” Walls and Mirrors*. 5th edition (2007). Addison Wesley.
- Nor Bahiah et al. *“Struktur data & algoritma menggunakan C++”*. Penerbit UTM. 2005.