

# Tournament Sort

## PSEUDOCODE

### Tournament Sorting

(a1 ,a2.... an :list of elements with  $n \geq 1$  )  
 We are going to build a binary tree of height  $K=\lceil \log n \rceil$ .  
 The first  $n$  leaves are assigned to be the  $n$  elements in the given list,while the remaining leaves are assigned to be  $-\infty$ .  
 $K= \lceil \log n \rceil$   
**T:**Complete binary tree  
 for  $i:=1$  to  $n$   
     Set the value of the  $i$ th the leaf in  $T$  as  $a_i$   
 for  $i:=n+1$  to  $2^k$   
     Set the value of the  $i$ th leaf in  $T$  as  $-\infty$ .  
 for  $i:=1$  to  $k$   
     for every vertex  $V$  at level  $k-i$   
         The value of  $v$  in  $T$  is the largest value of its two children.  
 for  $i:=1$  to  $n$   
      $C_i$ =value of the root of  $T$   
      $V$ :value of the root of  $T$   
     Change a value of a leaf with value  $v$  to  $-\infty$ .  
     while value of the root is  $v$   
          $V$ =parent( $v$ )  
         The value of the vertex  $V$  in  $T$  is set to the largest value of its two children.  
 List of values  $C_n,C_{n-1},...,C_2,C_1$   
 (list of integers sorted in nondecreasing order)  
 return  $C_n,C_{n-1},...,C_2,C_1$

## DEFINITION

Tournament is a sorting algorithm it improves upon the naive selection sort by using a priority Queue to find the next element in the sort. In the native selection sort, it takes  $O(n)$  operation to select the next element of “ $n$ ” elements. In a tournament sort it takes  $O(\log n)$  operations.

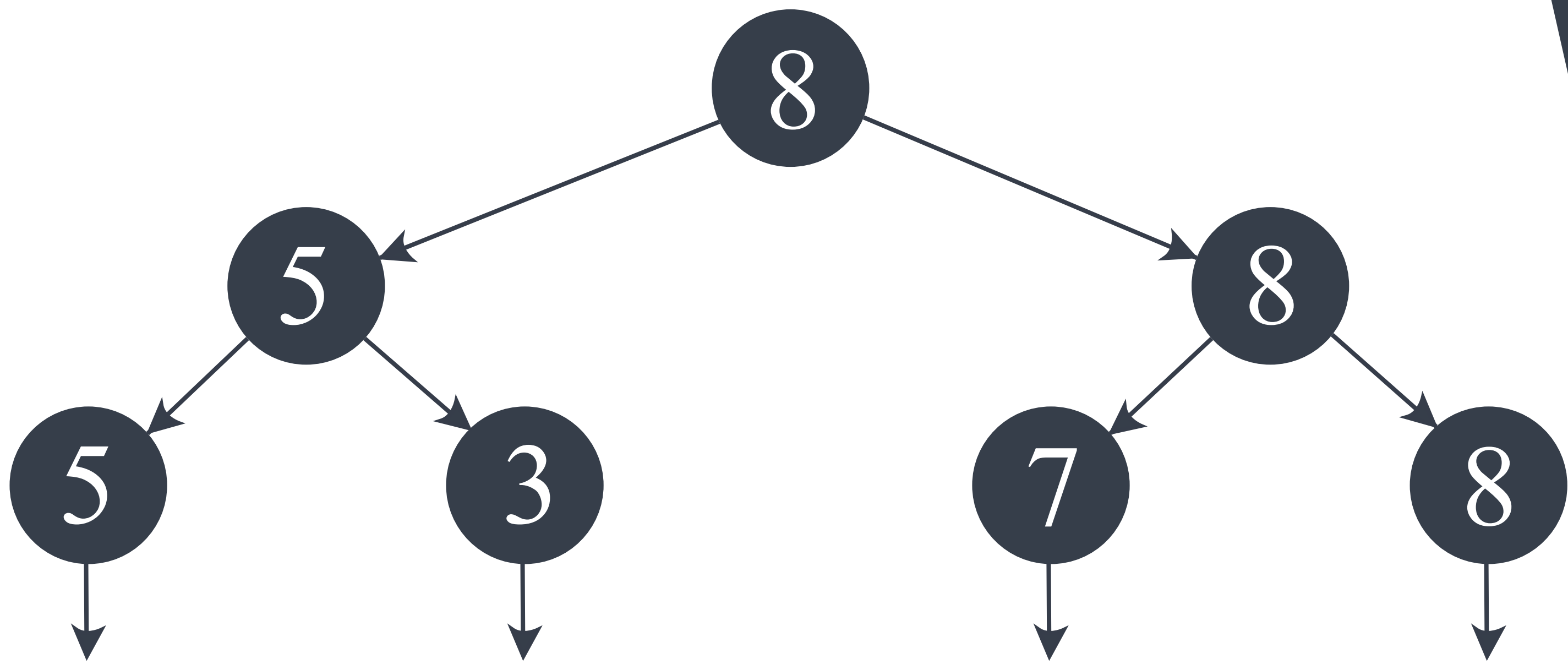
## CLASS

Sorting Algorithm.

## DATA STRUCTURE

Array.

TIME COMPLEXITY		
Best Case	Average Case	Worst Case
$O(\log n)$	$O(\log n)$	$O(\log n)$



Ali Ibrar	17B-092-SE
Ahsan Rasheed	17B-089-SE
Abeer Khan	17B-133-SE

### Result

(a1, a2 ,....an: list of elements with  $n > 1$ )  
 $K:= \lceil \log n \rceil$   
**T:**=complete binary tree  
 for  $i:=1$  to  $n$   
     Set the value of the  $i$ th leaf in  $T$  as  $a_i$   
 for  $i:=n + 1$  to  $2^k$   
     Set the value of the  $i$ th leaf in  $T$  as  $-\infty$   
 for  $i:=1$  to  $k$   
     for every vertex  $v$  at level  $k - i$   
         The value of  $v$  in  $T$  is the largest value of its two children  
 for  $i:=1$  to  $n$   
      $c:=$  value of the root of  $T$   
      $v:=$  value of the root of  $T$   
     Change a value of a leaf with value  $v$  to  $-\infty$   
     while value of the root is  $v$   
          $v:=$ parent( $r$ )  
     The value of vertex  $v$  in  $T$  is set to the largest value of its two children  
 return  $c_n,c_{n-1},...,c_2,c_1$