

Adaptive Ant Colony Optimization with Node Clustering for the Multi-Depot Mixed Fleet Capacitated Multiple TSP

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1 Variables

- $V = \{D \cup C\}$: total vertices, Depots & Customers
- VT : total Vehicle Types
- D : total Depots
- R : solution
- K : cluster matrix
- τ : pheromone matrix
- n_{ants} : number of ants in colonies
- n_{freq} : frequency of the local optimization
- n_{prim} : number of primary clusters
- n_{size} : number of vertices in clusters
- n_{sect} : number of sectors
- T_{update} : temperature updating coefficient
- α_{update} : temperature cooling coefficient
- ρ_{min}, ρ_{max} : minimum and maximum limits of the pheromone evaporation coefficient
- δ : pheromone updating coefficient
- α : distance probability coefficient
- β : pheromone probability coefficient

At the initial phase of the algorithm, the pheromone matrix τ is initialized using (1). τ is a 4-dimensional matrix.

$$\tau_{ij}^{(k)(h)} = 1 \text{ for all } v_i, v_j \in V, vt_k \in VT \text{ and } d_h \in D \quad (6)$$

Pheromone matrix update, evaporation coefficient ρ and pheromone evaporation follow the same principles with the original AACO-NC found here.

The node clustering technique also follows the original paper, which can be found here.

Although in this algorithm each vertex has more than one set of clusters, which is equal to the number of different vehicle types. This is because of the restrictions which forbid some vehicles from visiting certain customers.

2 Functions

In general, the functions used in this algorithm are slightly altered versions of the ones presented in the AACO-NC for the MDVRP by Stodola, with the purpose of adapting the AACO-NC algorithm to solve the MD-mfcmTSP.

The MD-mfcmTSP differs from the MDVRP in the following ways:

1. Supports multiple vehicle types (different capacities and speeds).
2. Each vehicle type has an assigned speed.
3. Considers realistic scenarios where big vehicles cannot access certain customers and drone safe landing spaces (not all customers can be serviced by drone).
4. Minimizes makespan instead of distance.
5. Vehicle capacity is dictated by the vehicle type and not by the vehicle's depot (depots are considered always stocked and as a reloading station for their vehicles).
6. Not every depot has to have the same number of vehicles or vehicle types assigned (e.g. Depot 1 has 4 trucks, Depot 2 has 10 drones and Depot 3 has 3 trucks, 5 motorcycles and 7 drones).

Algorithm 1: AACONC

Input: $V, n_{\text{ants}}, n_{\text{freq}}, n_{\text{size}}, n_{\text{sect}}, n_{\text{prim}}, T_{\text{update}}, \alpha, \beta, \rho_{\text{min}}, \rho_{\text{max}}, \delta$

```
1  $|R| \leftarrow \infty$ 
2  $iter \leftarrow 0$ 
3 Initialize pheromone matrices  $\tau$ 
4 for each  $t \in VT$ 
5   for each  $v_i \in V^{(t)}$ 
6      $K^{(t)(v_i)} \leftarrow \text{CreateClusters}(V^{(t)}, v_i, n_{\text{size}}, n_{\text{sect}}, n_{\text{prim}})$ 
7   end for
8 end for
9 while not terminated do
10   $|R_{\text{best}}| \leftarrow \infty$ 
11   $iter \leftarrow iter + 1$ 
12  for  $a = 1$  to  $n_{\text{ants}}$  do
13     $R_a \leftarrow \text{AntSolution}(V, K, \tau, \alpha, \beta)$ 
14    if  $|R_a| < |R_{\text{best}}|$  then
15       $R_{\text{best}} \leftarrow R_a$ 
16    end if
17  end for
18  if  $iter \bmod n_{\text{freq}} = 0$  then
19     $R_{\text{best}} \leftarrow \text{LocalOptimization}(V, R_{\text{best}})$ 
20  end if
21  if  $|R_{\text{best}}| < |R|$  then
22     $R \leftarrow R_{\text{best}}$ 
23  end if
24  Update pheromone matrices  $\tau$ 
25  Calculate evaporation coefficient  $\rho$ 
26  Evaporate pheromone matrices  $\tau$  using  $\rho$ 
27 end while
28 return  $R$ 
```

Algorithm 2: antSolution

```
1 Function  $\text{antSolution}(V = \{D, C\}, K, \tau, \alpha, \beta)$ 
2    $V_{\text{free}} = C$ 
3   while  $V_{\text{free}} \neq \emptyset$  do
4      $vt = \text{selectVehicleType}(V_{\text{free}}, K, \tau)$ 
5      $d = \text{selectDepot}(vt, V_{\text{free}}, K^{(vt)}, \tau)$ 
6      $v = \text{selectVehicle}(vt, d, V_{\text{free}}, K^{(vt)}, \tau)$ 
7      $pos \leftarrow$  vehicle's position
8      $k = \text{selectCluster}(vt, d, v, V_{\text{free}}, K^{(pos)(vt)}, \tau, \alpha, \beta)$ 
9      $V_{\text{candidates}} = V_{\text{free}} \cap K_k^{(pos)(vt)}$ 
10     $c = \text{selectCustomer}(vt, d, pos, V_{\text{candidates}}, \tau, \alpha, \beta)$ 
11    if  $v_{\text{load}} < c^{(\text{demand})}$  then
12       $R_d^{(vt)} = R_d^{(vt)} + \{d\}$ 
13       $v_{\text{load}} = vt_{\text{capacity}}$ 
14    end if
15     $R_d^{vt} = R_d^{vt} + \{c\}$ 
16     $v_{\text{load}} = v_{\text{load}} - c^{(\text{demand})}$ 
17     $V_{\text{free}} = V_{\text{free}} - \{c\}$ 
18  end while
19  for each  $d \in D$  and  $vt \in VT$  //Vehicles return to
20    their depots
21     $R_d^{vt} = R_d^{vt} + \{d\}$ 
22  end for
23  return  $R = \{R_1^1, R_2^1, \dots, R_2^3, R_3^3, R_D^{VT}\}$ 
```

Algorithm 3: selectVehicleType

```
1 Function  $\text{selectVehicleType}(V_{\text{free}}, K, \tau)$ 
2   for each vehicle type  $t_i \in VT$  do
3      $V_{\text{cand}} = \emptyset$ 
4     for each vehicle do
5        $pos \leftarrow$  vehicle's current location
6        $d \leftarrow$  vehicle's depot
7       for  $k = 1$  to  $n_{\text{prim}}$  do
8          $V_{\text{cand}} = V_{\text{cand}} + V_{\text{free}} \cap K_k^{(pos)(t_i)}$ 
9       end for
10    end for
11     $p(t_i) = \sum_{v_j \in V_{\text{cand}}} \tau_{v_{\text{pos}}v_j}^{(t_i)(d)}$ 
12  end for
13   $p_{\text{sum}} = \sum_{t_i \in VT} p(t_i)$ 
14  return  $\text{rouletteWheel}(p(VT), p_{\text{sum}})$ 
```

Algorithm 4: selectDepot

```
1 Function  $\text{selectDepot}(vt, V_{\text{free}}, K^{(vt)}, \tau)$ 
2   for each  $d_i \in D^{(vt)}$  do
3      $V_{\text{cand}} = \emptyset$ 
4     for each vehicle do
5        $pos \leftarrow$  vehicle's current location
6       for  $k = 1$  to  $n_{\text{prim}}$  do
7          $V_{\text{cand}} = V_{\text{cand}} + V_{\text{free}} \cap K_k^{(pos)(vt)}$ 
8       end for
9     end for
10     $p(d_i) = \sum_{v_j \in V_{\text{cand}}} \tau_{v_{\text{pos}}v_j}^{(vt)(d_i)}$ 
11  end for
12   $p_{\text{sum}} = \sum_{d_i \in D^{(vt)}} p(d_i)$ 
13  return  $\text{rouletteWheel}(p(D^{(vt)}), p_{\text{sum}})$ 
```

Algorithm 5: selectVehicle

```
1 Function  $\text{selectVehicle}(vt, d, V_{\text{free}}, K^{(vt)}, \tau)$ 
2   for each  $v_i \in V^{(vt)(d)}$  do
3      $V_{\text{cand}} = \emptyset$ 
4      $pos \leftarrow$  vehicle's current location
5     for  $k = 1$  to  $n_{\text{prim}}$  do
6        $V_{\text{cand}} = V_{\text{cand}} + V_{\text{free}} \cap K_k^{(pos)(vt)}$ 
7     end for
8      $p(v_i) = \sum_{v_j \in V_{\text{cand}}} \tau_{v_{\text{pos}}v_j}^{(vt)(d)}$ 
9   end for
10   $p_{\text{sum}} = \sum_{v_i \in V^{(vt)(d)}} p(v_i)$ 
11  return  $\text{rouletteWheel}(p(V^{(vt)(d)}), p_{\text{sum}})$ 
```

Algorithm 6: selectCluster

```
1 Function selectCluster( $vt, d, pos, V_{free}, K, \tau, \alpha, \beta$ )
2   for  $k = 1$  to  $n_{prim}$  do
3      $V_{cand} = \emptyset$ 
4      $V_{cand} = V_{free} \cap K_k^{(pos)(vt)}$ 
5     if  $V_{cand} = \emptyset$  then
6        $\eta_k = \tau_k = 0$ 
7     end if
8     else
9        $\eta_k = |V_{cand}| \cdot \sum_{v_j \in V_{cand}} |v_{pos} - v_j|^{-1}$ 
10       $\tau_k = \frac{1}{|V_{cand}|} \cdot \sum_{v_j \in V_{cand}} \tau_{v_{pos}v_j}^{(vt)(d)}$ 
11    end if
12  end for
13   $\eta_{sum} \leftarrow \sum_{k=1}^{n_{prim}} \eta_k^\alpha$ 
14   $\tau_{sum} \leftarrow \sum_{k=1}^{n_{prim}} \tau_k^\beta$ 
15  if  $\eta_{sum} = 0$  then
16    // return first cluster with a free customer
17    for  $k = n_{prim} + 1$  to  $|K^{(pos)(vt)}|$  do
18       $V_{cand} = V_{free} \cap K_k^{(pos)(vt)}$ 
19      if  $V_{cand} \neq \emptyset$  then
20        return  $k$ 
21      end if
22    end for
23  end if
24  for  $k = 1$  to  $n_{prim}$  do
25     $p(K_k^{(pos)(vt)}) = \frac{\eta_k^\alpha \cdot \tau_k^\beta}{\eta_{sum} \cdot \tau_{sum}}$ 
26  end for
27   $p_{sum} = \sum_{k \in n_{prim}} p(K_k^{(pos)(vt)})$ 
28  return rouletteWheel( $p(K^{(pos)(vt)}), p_{sum}$ )
```

Algorithm 7: selectCustomer

```
1 Function selectCustomer( $vt, d, pos, V_{candidates}, \tau, \alpha, \beta$ )
2   for each  $v_i \in V_{cand}$  do
3      $p(v_i) = |v_{pos} - v_i|^{-\alpha} \cdot (\tau_{v_{pos}v_i}^{(vt)(d)})^\beta$ 
4   end for
5    $p_{sum} = \sum_{v_i \in V_{cand}} p(v_i)$ 
6   return rouletteWheel( $p(v), p_{sum}$ )
```

Algorithm 8: LocalOptimization

```
1 Function LocalOptimization( $V, R_{best}$ )
2   for each  $t_i \in VT$  do
3     for each  $d_i \in D^{(t_i)}$  do
4       singleColonyOpt( $R_{best}^{(t_i)(d_i)}, n_{max} = 1$ )
5       singleColonyOpt( $R_{best}^{(t_i)(d_i)}, n_{max} = 2$ )
6       (Moves  $n_{max}$  successive customer node(s) to
        different positions in the same route)
7     end for
8     mutualColonyOpt( $R_{best}^{(t_i)}, n_{max} = 1$ )
9     if  $t_i \neq Drone$  then
10      mutualColonyOpt( $R_{best}^{(t_i)}, n_{max} = 2$ )
11    end if (Moves  $n_{max}$  successive customer node(s)
12      ) from each  $R_{best}^{(t_i)(d_i)}$ ,  $d_i \in D^{(t_i)}$  to different
13      positions in each  $R_{best}^{(t_i)(d_j)}$ ,  $d_j \neq d_i$ )
14  end for
15  return  $R_{best}$ 
```

Algorithm 9: singleColonyOptimization

```
1 Function singleColonyOptimization( $R_{best}^{(t_i)(d_i)}, n_{max}$ )
2   for  $n = 1$  to  $n_{max}$  do
3     for each combination of  $n$  successive nodes in
       the route
4       move the nodes to a different place on the
       same route
5       evaluate the newly-created solution
6       if this solution is better than the original and
         all constraints are satisfied then
7         replace the original with the new solution
8       end if
9       continue in point 4 unless all possible places
       in the route have been already evaluated
10    end for
11  end for
12  return  $R_{best}$ 
```

Algorithm 10: mutualColonyOptimization

```
1 Function mutualColonyOptimization( $R_{best}^{(t_i)}, n_{max}$ )
2   for  $n = 1$  to  $n_{max}$  do
3     for each possible pair of depots  $d1$  and  $d2$ 
4       for each combination of  $n$  successive nodes in
         the route of  $d1$ 
5         remove the nodes from the route of  $d1$  and
6         insert them into the route of  $d2$ 
7         evaluate the newly-created solution
8         if this solution is better than the original
           and all constraints are satisfied then
9           replace the original with the new
             solution
10          end if
11          continue in point 6 unless all possible
            places in the route of  $d2$  have been
            already evaluated
12        end for
13      end for
14    end for
15  return  $R_{best}$ 
```

1, 2 and 3

"3" means that the customer can only be served by vehicle type 3

3 Other

Because this is a new problem there are no existing instances with which we can test the algorithm, so we will create new random ones with the following structure:

{number of customers} {number of depots} {number of different vehicle types}

{capacity_{vt₁}}, {capacity_{vt₂}}, {...}, {capacity_{vt_n}}

{speed_{vt₁}}, {speed_{vt₂}}, {...}, {speed_{vt_n}}

{customer1_{id}}{customer1_x}{customer1_y}{customer1_{demand}}{customer1_{accessibility}}

{customer2_{id}}{customer2_x}{customer2_y}{customer2_{demand}}{customer2_{accessibility}}

.

.

.

{customerN_{id}}{customerN_x}{customerN_y}{customerN_{demand}}{customerN_{accessibility}}

{depot1_{id}}, {depot1_x}, {depot1_y}, {number of vehicles of type 1}, {n of vehicles of type 2}, ..., {n of vehicles of type N}

{depot2_{id}}, {depot2_x}, {depot2_y}, {number of vehicles of type 1}, {n of vehicles of type 2}, ..., {n of vehicles of type N}

.

.

.

{depotN_{id}}, {depotN_x}, {depotN_y}, {number of vehicles of type 1}, {n of vehicles of type 2}, ..., {n of vehicles of type N}

customer_{accessibility} is an integer ≥ 1 which shows what vehicle type(s) can serve the customer.

Examples:

"123" means that the customer can be served by vehicle types