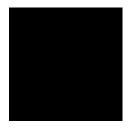
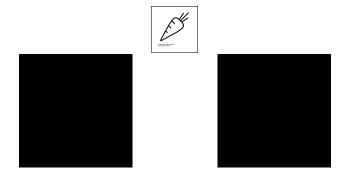
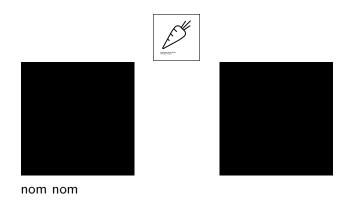
Tamarin: Concolic Disequivalence for MIPS

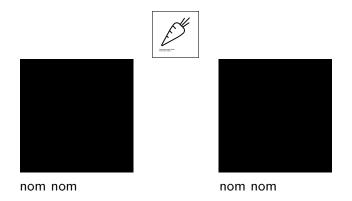
Abel Nieto

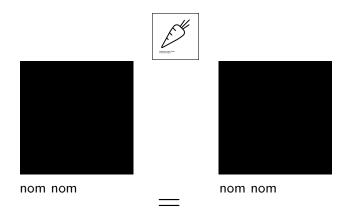


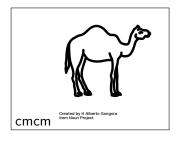
















Given MIPS program P_1 and P_2 , when are they equivalent?

Attempt 1: two programs are equivalent if they give the same output (resp.) for all inputs.

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What's an input?

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What's an input? Register \$1 and \$2.

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What's an input? Register \$1 and \$2. What's an output?

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What's an input? Register \$1 and \$2. What's an output? Register \$3.

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What's an input? Register \$1 and \$2. What's an output? Register \$3.

Don't care about (most) CPU interrupts/IO.

Attempt 1: two programs are equivalent if they give the same output (resp.) for all inputs.

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Problem: undecidable via Rice's theorem.

Attempt 2: two programs are S-equivalent if they cannot be told apart after S steps.

S-equivalent (e.g. for S = 10), but not equivalent:

R_1	R_2	S-equiv
-------	-------	---------

R_1	R_2	S-equiv
V	V	yes

R_1	R_2	S-equiv
V	V	yes
V	$w \neq v$	no

R_1	R_2	S-equiv
V	V	yes
V	$w \neq v$ error	no
V	error	no

R_1	R_2	S-equiv
V	V	yes
V	$w \neq v$ error	no
V	error	no
error	error	yes

R_1	R_2	S-equiv
V	V	yes
V	$w \neq v$	no
V	error	no
error	error	yes
non-termination	???	yes

Lemma

Equivalence implies S-equivalence.

Lemma

Equivalence implies S-equivalence.

Corollary (Soundness)

If two programs are not S-equivalent (for any S), then they are not equivalent.

Attempt 2: two programs are *S*-equivalent if they cannot be told apart after *S* steps.

Which inputs?

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Which inputs?

Try some inputs by hand: low coverage, fast (unit tests)

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Try all 2⁶⁴ values of \$1 and \$2: high coverage, slow (but decidable)

Attempt 2: two programs are *S*-equivalent if they cannot be told apart after *S* steps.

Which inputs?

Try some inputs by hand: low coverage, fast (unit tests)

Try all 2⁶⁴ values of \$1 and \$2: high coverage, slow (but decidable)

Tamarin: use concolic execution: higher coverage(?), not too slow(?)



```
# P_1
    bne $1, 42, end
    add $3, $1, $2

add $3, $3, $0

end:
    add $3, $1, $2

end:
# P_2
add $3, $1, $2

bne $2, 100, end
add $3, $3, $2

end:
```

Alternating concolic execution

```
# P_1
bne $1, 42, end
add $3, $1, $2
add $3, $3, $0
bne $2, 100, end
end:
add $3, $1, $2
end:
```

Run | Driver | Verifier | \$1 | \$2 | Path $|R_D|R_V$

```
# P_1
    bne $1, 42, end
    add $3, $1, $2

add $3, $3, $0

end:
    add $3, $1, $2

end:
# P_2
add $3, $1, $2

bne $2, 100, end
add $3, $3, $2

end:
```

Run	Driver	Verifier	\$1	\$2	Path	R_D	R_V
1	P_1	P_2	1	1	\$1 ≠ 42	2	2

```
# P_1
    bne $1, 42, end
    add $3, $1, $2

add $3, $3, $0

end:
    add $3, $1, $2

end:
# P_2
add $3, $1, $2

bne $2, 100, end
add $3, $3, $2

end:
```

Run	Driver	Verifier	\$1	\$2	Path	R_D	R_V
1	P_1	P_2	1	1	\$1 ≠ 42	2	2
2	P_2	P_1	1	1	\$2 \neq 100	2	2

```
# P_1
    bne $1, 42, end
    add $3, $1, $2

add $3, $3, $0

end:
    add $3, $1, $2

end:
# P_2
add $3, $1, $2

bne $2, 100, end
add $3, $3, $2

end:
```

Run	Driver	Verifier	\$1	\$2	Path	R_D	R_V
1	P_1	P_2	1	1	\$1 \neq 42	2	2
2	P_2	P_1	1	1	\$2 \neq 100	2	2
3	P_1	P_2	42		\$1 = 42	2	2

```
# P_1
    bne $1, 42, end
    add $3, $1, $2

add $3, $3, $0

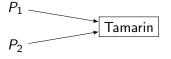
end:
    add $3, $1, $2

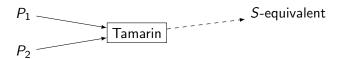
end:
# P_2
add $3, $1, $2

bne $2, 100, end
add $3, $3, $2

end:
```

Run	Driver	Verifier	\$1	\$2	Path	R_D	R_V
1	P_1	P_2	1	1	\$1 \neq 42	2	2
2	P_2	P_1	1	1	$$2 \neq 100$	2	2
3	P_1	P_2	42	1	$\$2 \neq 100$ \$1 = 42	2	2
4	P_2	P_1	1	100	\$2 = 100	201	2







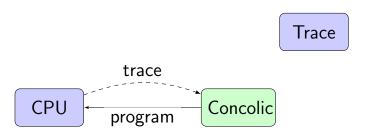
Concolic

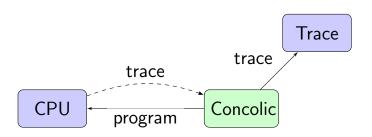
CPU

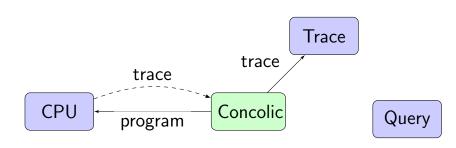
Concolic

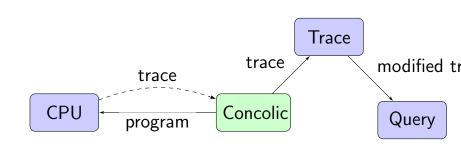


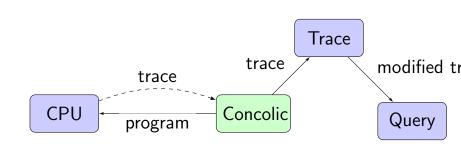


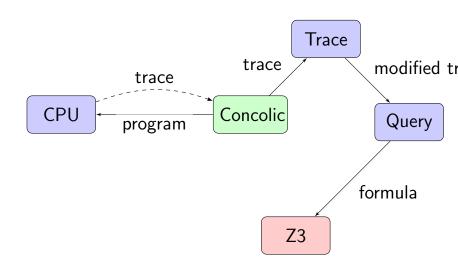


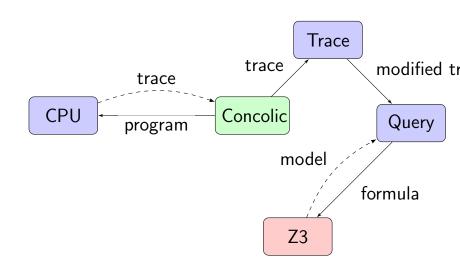


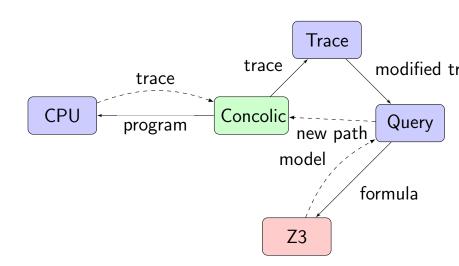


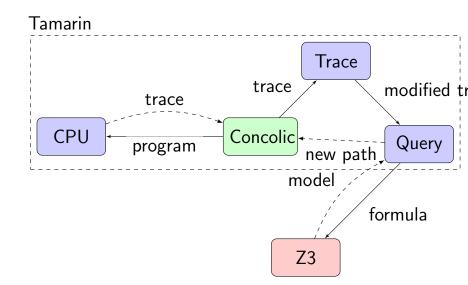




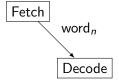


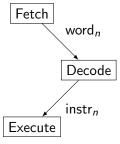


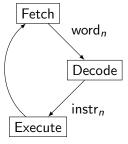


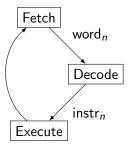


Fetch



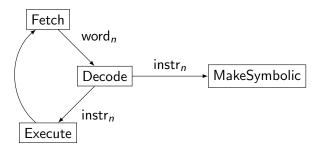




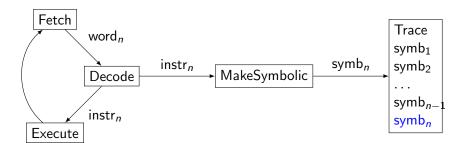


 ${\sf Make Symbolic}$

Trace $symb_1$ $symb_2$... $symb_{n-1}$ $symb_n$



Trace $symb_1$ $symb_2$... $symb_{n-1}$ $symb_n$



CPU (MakeSymbolic)

Instruction Symbolic

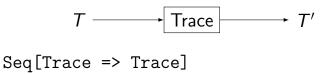
CPU (MakeSymbolic)

Instruction	Symbolic
add \$3, \$1, \$2	$r_3 \leftarrow r_1 + r_2$

CPU (MakeSymbolic)

Instruction	Symbolic
add \$3, \$1, \$2	$r_3 \leftarrow r_1 + r_2$ $r_1 = r_2 \text{ or } r_1 \neq r_2$ $r_3 \leftarrow 0$ x8BADF00D
beq \$1, \$2, label	$r_1 = r_2$ or $r_1 \neq r_2$
add \$3, \$pc, \$0	$r_3 \leftarrow 0$ x8BADF00D
lis \$3; 42	$r_3 \leftarrow 42$







Seq[Trace => Trace]

Desugar

$$T \longrightarrow Trace \longrightarrow T'$$

Seq[Trace => Trace]

- Desugar
- Simplify

$$T \longrightarrow Trace \longrightarrow T'$$

Seq[Trace => Trace]

- Desugar
- Simplify
- ▶ Trim

$$T \longrightarrow Trace \longrightarrow T'$$

Seq[Trace => Trace]

- Desugar
- Simplify
- Trim
- SSA convert

Desugar

Desugar

```
Mult(s, t) \rightarrow Mult64(tmp, s, t); Low32(lo, tmp); High32(hi, tmp)
```

Desugar

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Simplify

Desugar

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Mult(s, t) \rightarrow Mult64(tmp, s, t); Low32(lo, tmp); High32(hi, tmp)
```

Simplify

```
beq $0, $0, label 	o \emptyset
```

Trace

Desugar

```
Mult(s, t) \rightarrow Mult64(tmp, s, t); Low32(lo, tmp); High32(hi, tmp)
```

Simplify

```
beq $0, $0, label 	o \emptyset
```

Trim

Trace

Desugar

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Mult(s, t) \rightarrow Mult64(tmp, s, t); Low32(lo, tmp); High32(hi, tmp)
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Simplify

```
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```

Trim

limit trace to D path conditions

Trace

Desugar

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Mult(s, t) \rightarrow Mult64(tmp, s, t); Low32(lo, tmp); High32(hi, tmp)
```

Simplify

beq \$0, \$0, label
$$ightarrow \emptyset$$

Trim

limit trace to D path conditions

SSA convert

Trace	Incorrect	Correct

Trace	Incorrect	Correct
Add(\$3, \$1, \$2)		

Trace	Incorrect	Correct
Add(\$3, \$1, \$2)	z = x + y	

Trace	Incorrect	Correct
Add(\$3, \$1, \$2)	z = x + y	$z_1 = x_1 + y_1$

Trace	Incorrect	Correct
Add(\$3, \$1, \$2)	z = x + y	$z_1 = x_1 + y_1$
Sub(\$1 , \$1 , \$2)		

Trace	Incorrect	Correct
Add(\$3, \$1, \$2)	z = x + y	$z_1 = x_1 + y_1$
Sub(\$1 , \$1 , \$2)	x = x - y	

Trace	Incorrect	Correct
		$z_1 = x_1 + y_1$
Sub(\$1 , \$1 , \$2)	x = x - y	$x_2 = x_1 - y_1$

Trace	Incorrect	Correct
Add(\$3, \$1, \$2)	z = x + y	$z_1 = x_1 + y_1$
Sub(\$1 , \$1 , \$2)	x = x - y	$x_2 = x_1 - y_1$
Add(\$2, \$1, \$2)		

Trace	Incorrect	Correct
Add(\$3, \$1, \$2)	z = x + y	$z_1 = x_1 + y_1$
Sub(\$1 , \$1 , \$2)	x = x - y	$x_2 = x_1 - y_1$
Add(\$2, \$1, \$2)	y = x + y	

	Incorrect	
Add(\$3, \$1, \$2)	z = x + y	$z_1 = x_1 + y_1 x_2 = x_1 - y_1 y_2 = x_2 + y_1$
Sub(\$1 , \$1 , \$2)	x = x - y	$x_2 = x_1 - y_1$
Add(\$2, \$1, \$2)	y = x + y	$y_2 = x_2 + y_1$

	Incorrect	
Add(\$3, \$1, \$2)	z = x + y	$z_1 = x_1 + y_1 x_2 = x_1 - y_1 y_2 = x_2 + y_1$
Sub(\$1 , \$1 , \$2)	x = x - y	$x_2 = x_1 - y_1$
Add(\$2, \$1, \$2)	y=x+y	$y_2 = x_2 + y_1$

$$$3 = 9 \land $1 = 3$$

	Incorrect	
Add(\$3, \$1, \$2)	z = x + y	$z_1 = x_1 + y_1 x_2 = x_1 - y_1 y_2 = x_2 + y_1$
Sub(\$1 , \$1 , \$2)	x = x - y	$x_2 = x_1 - y_1$
Add(\$2, \$1, \$2)	y = x + y	$y_2 = x_2 + y_1$

$$$3 = 9 \land $1 = 3$$
 $z = 9 \land x = 3$

	Incorrect	
Add(\$3, \$1, \$2)	z = x + y	$z_1 = x_1 + y_1 x_2 = x_1 - y_1 y_2 = x_2 + y_1$
Sub(\$1 , \$1 , \$2)	x = x - y	$x_2 = x_1 - y_1$
Add(\$2, \$1, \$2)	y = x + y	$y_2 = x_2 + y_1$

$$3 = 9 \land 1 = 3$$

 $z = 9 \land x = 3 \text{ (no sol)}$

	Incorrect	
Add(\$3, \$1, \$2)	z = x + y	$z_1 = x_1 + y_1 x_2 = x_1 - y_1 y_2 = x_2 + y_1$
Sub(\$1 , \$1 , \$2)	x = x - y	$x_2 = x_1 - y_1$
Add(\$2, \$1, \$2)	y=x+y	$y_2=x_2+y_1$

$$\$3 = 9 \land \$1 = 3$$

 $z = 9 \land x = 3 \text{ (no sol)}$
 $z_1 = 9 \land x_2 = 3$

	Incorrect	
Add(\$3, \$1, \$2)	z = x + y	$z_1 = x_1 + y_1 x_2 = x_1 - y_1 y_2 = x_2 + y_1$
Sub(\$1 , \$1 , \$2)	x = x - y	$x_2 = x_1 - y_1$
Add(\$2, \$1, \$2)	y = x + y	$y_2 = x_2 + y_1$

$$\$3 = 9 \land \$1 = 3$$

 $z = 9 \land x = 3 \text{ (no sol)}$
 $z_1 = 9 \land x_2 = 3 \text{ } (x_1 = 6 \land y_1 = 3)$