

Problem 2.2

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a) $T(n) = 36T(n/6) + 2n$

By Master Method we know $\left. \begin{array}{l} a = 36 \\ b = 6 \\ f(n) = 2n \end{array} \right\} T(n) = \Theta(n^{\log_b a})$

if $f(n)$ polynomially smaller than $n^{\log_b a}$

$$\Rightarrow \left. \begin{array}{l} n^{\log_6 36} = n^2 \\ 2n = n^1 \end{array} \right\} n^{\log_b a} > f(n) \Rightarrow \Theta(n^2)$$

b) $T(n) = 5T(n/3) + 17n^{1.2}$

By Master Method we know $\left. \begin{array}{l} a = 5 \\ b = 3 \\ f(n) = 17n^{1.2} \end{array} \right\} T(n) = \Theta(n^{\log_b a})$

if $f(n)$ polynomially smaller than $n^{\log_b a}$

$$\Rightarrow \left. \begin{array}{l} n^{\log_3 5} = n^{1.4} \\ 17n^{1.2} = n^{1.2} \end{array} \right\} n^{\log_b a} > f(n) = \Theta(n^{1.4})$$

c) $T(n) = 12T(n/2) + n^2 \lg n$

We get $\left. \begin{array}{l} a = 12 \\ b = 2 \\ f(n) = n^2 \lg n \end{array} \right\}$ where

$$\begin{aligned} f(n) &= n^2 \lg n \\ n^{\log_b a} &= n^{\log_2 12} = n^{3.5} \\ \Rightarrow n^{3.5} &> n^2 \lg n \end{aligned}$$

regularity condition

$$af(n/b) \leq cf(n) \quad c < 1$$

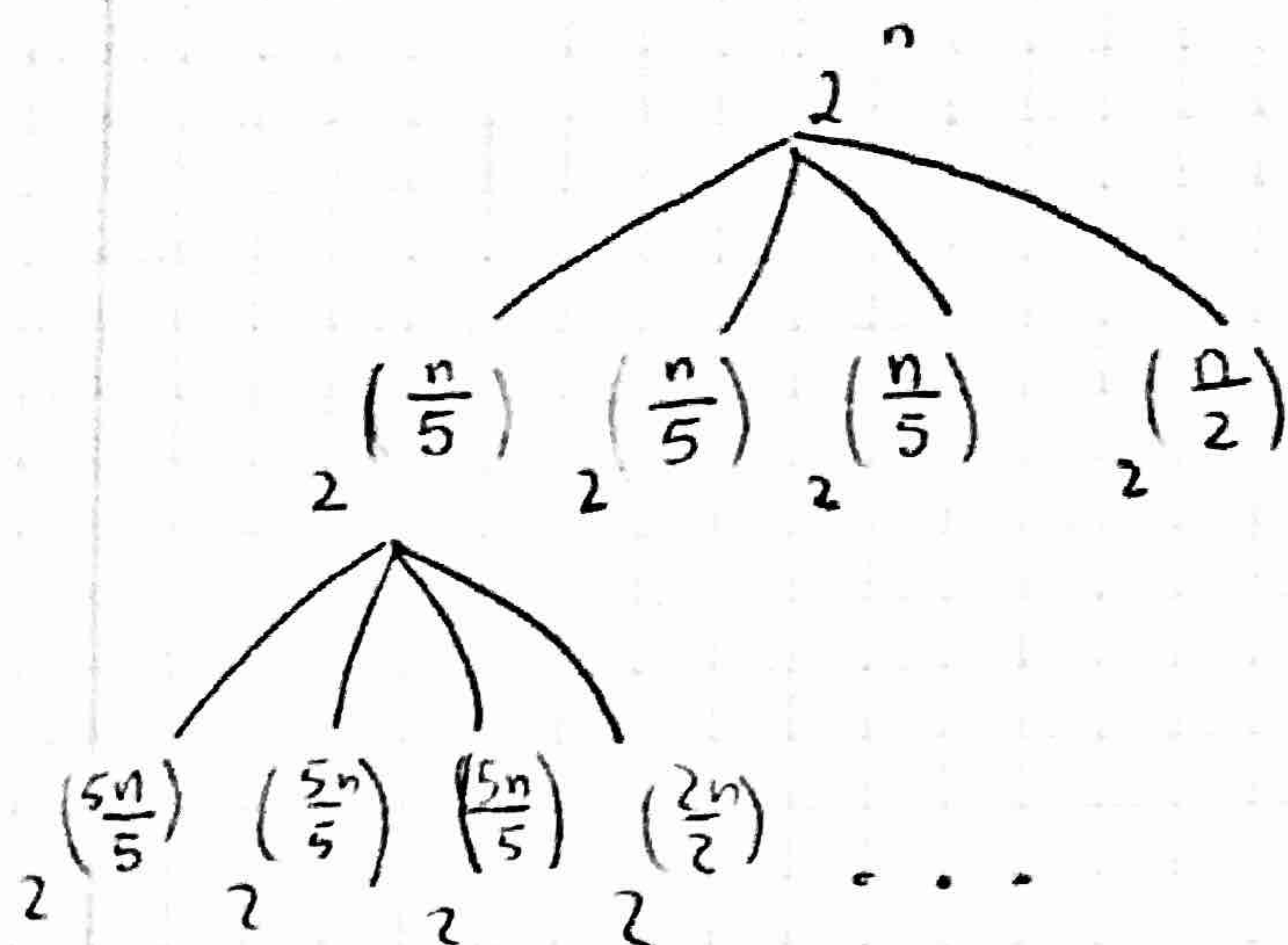
$$a \cdot (n/b)^2 \log_{10}(n/b) \leq cf(n)$$

$$12(n/2)^2 \log_{10}(n/2) \leq cf(n)$$

$$\frac{12}{4} = 3 \quad \text{True for } c = \frac{12}{4}$$

So $T(n) = \Theta(f(n)) = \Theta(n^2 \lg n)$

d) $T(n) = 3T(n/5) + T(n/2) + 2^n$ - Using Recursion Tree



We get:

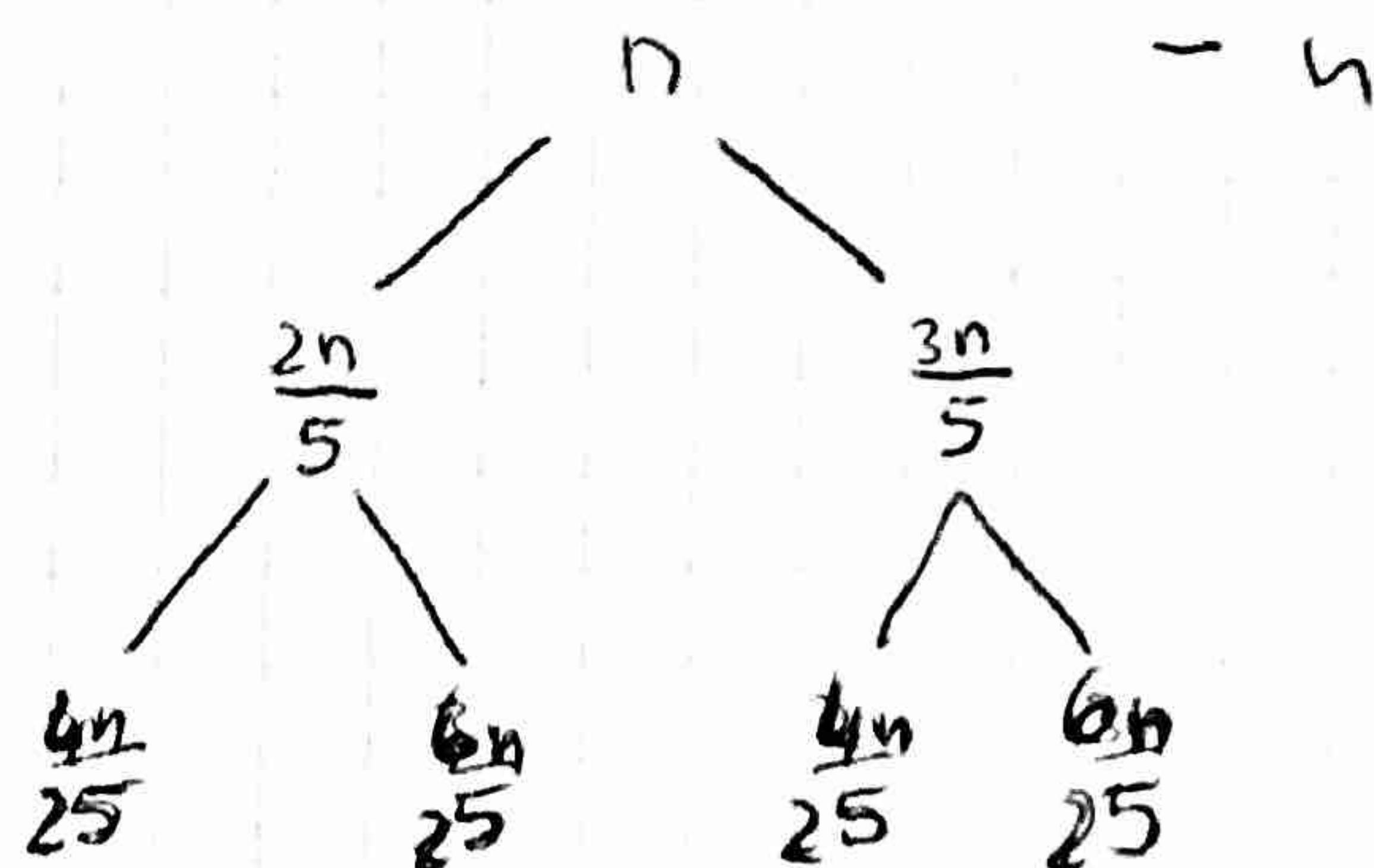
$$2^n + 3 \times 2^{(n/5)} + 2^{(n/5)} + 2^{(n/5)} + \dots$$

We see that the common base is 2^n as:

$$2^n (1 + \dots) \quad \text{geometric series}$$

Therefore $\Theta(2^n)$

e) $T(n) = T(2n/5) + T(3n/5) + \Theta(n)$ - Using Recursion Tree



$$T(n) = n + n + \dots$$

So $T(n) \in \Theta(\log n)$
 $= \Theta(\log n)$