

BEQ $r1, r2, \text{offset} \rightarrow PC \leftarrow PC + \text{offset}$

THIS INSTRUCTION NEEDS TO :

- compare the content of 2 registers
- add or not a certain offset to PC

we can use the ALU. We set it to perform a subtraction. If it's zero, then the content of 2 registers are equal.

Usually, ALU's have secondary outputs called "Z".
If the result of the sub is zero, then the output of Z will be 1

Since we need to add a certain offset to the PC, we will need another adder. It will have in input the current PC and the immediate.

We add a multiplexer in order to select whether doing $PC + \text{offset}$ if registers are equal otherwise $PC + 4$

We need another control for "beq" because what if we wanted to do the "sub" instruction between 2 regs.? What if it gives 0? It would jump instead of storing the result of sub. That's why a **BRANCH** control (which will be **put at 1** for beq) It will be put in **AND** with z.

HOW WILL BE THE CONTROLS SET FOR BEQ?

ALUSRC - 0 (we need to do subtraction in ALU)

ALUOP → it has to do subtraction (we didn't study yet)

MEMREAD - 0 (we don't have to read memory, only used with lw)

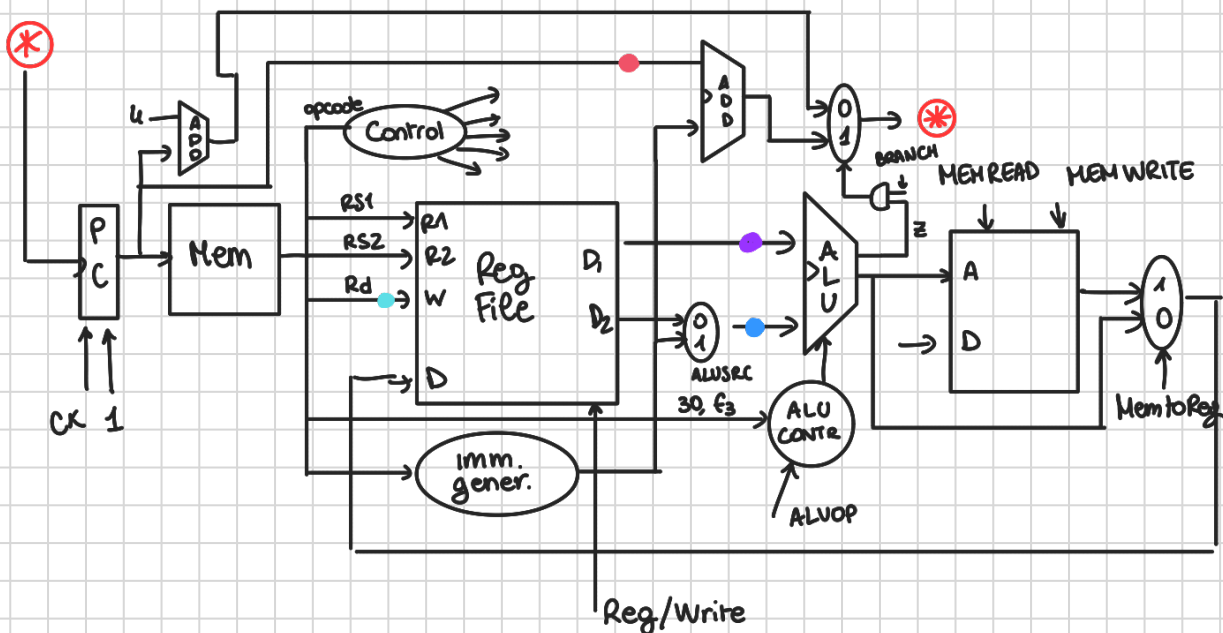
MEMWRITE - 0 (used for sw only)

MEMTOREG - DON'T CARE

REGWRITE - 0

BRANCH - 1

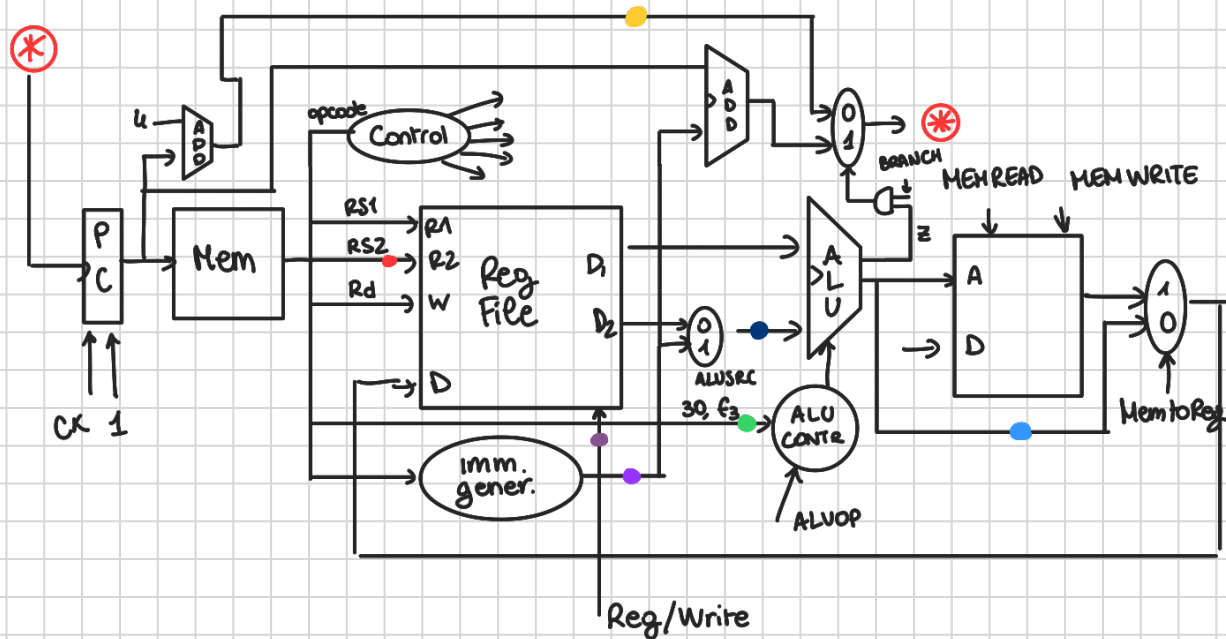
LET'S ASSUME WE'RE EXECUTING SW x6, 8(x12)



- here, we will find the immediate 8
- here, we have the content of x12
- here, we have the PC
- here, we have part of the immediate (it's technically wrong but REGWRITE = 0 so all good)

NOTE: DON'T CARE isn't a value that you can read when an architecture runs.
"NOTHING" also doesn't exist

LET'S ASSUME WE'RE EXECUTING **ADD 5, 6, 7**



- here, we have the content of Register 7
- here, we have the result of $R6 + R7$
- here, we have garbage
- here, we have $PC + 4$ (next instruction)
- here, we will have $f3$ (000) and 30th bit (1)
- here, we have content of $R7$.
- here, we have 1 (REGWRITE)

NOTE: better to say garbage. "Don't care" would be error

This is a slow architecture because it does one instruction in one clock cycle

The slowest instruction is lw , and then sw .

The clock cycle should be very long and the frequency very slow.