INDIAN INSTITUTE OF TECHNOLOGY KHARAGPUR

DEPARTMENT OF COMPUTER SCIENCE & ENGINEERING



COMPUTER ORGANIZATION LABORATORY (CS39001)

Abhay Kumar Keshari (20CS10001) Hardik Pravin Soni (20CS30023)

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1 Instruction Set Architecture

Table 1: Instruction Set Architecture

Class	Instruction	$_{ m Usage}$	Meaning
	Add	add rs,rt	$rs \leftarrow rs + rt$
Arithmetic	Comp	$_{ m comp\ rs,rt}$	$rs \leftarrow 2$'s Complement (rt)
Armmenc	Add Immediate	addi rs,imm	$rs \leftarrow (rs) + imm$
	Complement Immediate	compi rs,imm	$rs \leftarrow 2$'s Complement (imm)
Logic	AND	and rs,rt	$rs \leftarrow rs\&rt$
Logic	XOR	xor rs,rt	$rs \leftarrow rs \oplus rt$
	Shift Left Logical	shll rs, sh	$rs \leftarrow (rt)$ left-shifted by sh
	Shift Right Logical	$_{ m shrl}$ $_{ m rs,sh}$	$rs \leftarrow (rt)$ right-shifted by sh
Shift	Shift Left Logical Variable	shllv rs, rt	$rs \leftarrow (rt)$ left-shifted by (rt)
SIIII	Shift Right Logical Variable	shrlv rs, rt	$rs \leftarrow (rt)$ right-shifted by (rt)
	Shift Right Arithmetic	shra rs, sh	$rs \leftarrow (rs)$ arithmetic right-shifted by sh
	Shift Right Arithmatic Variable	shrav rs, sh	$rs \leftarrow (rs)$ arithmetic right-shifted by (rt)
Memory	Load Word	lw rt, imm, (rs)	$rt \leftarrow mem[(rs) + imm]$
Memory	Store Word	sw rt, imm, (rs)	$mem[(rs) + imm] \leftarrow (rt)$
	Unconditional Branch	b L	goto L
	Branch Register	br rs	goto (rs)
	Branch on less than 0	bltz rs,L	if(rs) < 0 then go to L
Branch	Branch on flag zero	$_{ m bz\ rs,L}$	$\mathit{if}(rs) = 0 \ then \ goto \ L$
Dianch	Branch on flag not zero	bnz rs,L	$if(rs) \neq 0 then go to L$
	Branch and link	bl L	goto $L; 31 \leftarrow (PC) + 4$
	Branch on Carry	bcy L	$goto\ L\ if\ Carry=1$
	Branch on No Carry	bncy L	$goto\ L\ if\ Carry=0$
Complex	Diff	diff rs, rt	$rs \leftarrow the \ LSB \ bit \ at \ which \ rs \ and \ rt \ differ$

2 Instruction Format and Encoding

2.1 R-Format Instruction:

add, comp, and, xor, shll, shllv, shrl, shrlv, shra, shrav, diff

Table 2: R-Format Encoding

opcode	rs	rt	shamt	Don't Care	funccode
6 bits	5bits	5 bits	5 bits	6 bits	5 bits

2.2 I-Format Instruction: lw, sw (Memory Access Instructions)

Table 3: I-Format Encoding for lw and sw

opcode	rs	rt	imm
6 bits	5 bits	5 bits	16 bits

2.3 I-format: addi, compi

Table 4: I-Format Encoding for addi and compi

opcode	rs	Don't Care	imm
6 bits	5 bits	5 bits	16 bits

2.4 J1-Format:- bltz, bz, bnz, br

Table 5: J1-Format Encoding

opcode	rs	Don't Care	label
6 bits	5 bits	5 bits	16 bits

2.5 J2-Format:- b,bl,bcy,bncy

Table 6: J2-Format Encoding

opcode	L
6 bits	26 bits

3 Control Signals

Table 7: Control Signal Configuration

Instruction	op-code	func-code	RegDst	RegWrite	MemToReg	MemRead	MemWrite	ALUSel	ALUsrc	ALUop	LblSel	BranchAddrSel	isBranch
add	000000	00000	00	1	10	0	0	0	0	00001	X	X	0
comp	000000	00001	00	1	10	0	0	1	0	00101	X	X	0
AND	000001	00000	00	1	10	0	0	0	0	00010	X	X	0
XOR	000001	00001	00	1	10	0	0	0	0	00011	X	X	0
shll	000010	00000	00	1	10	0	0	0	1	01010	X	X	0
shllv	000010	00010	00	1	10	0	0	0	0	01010	X	X	0
shrl	000010	00001	00	1	10	0	0	0	1	01000	X	X	0
shrlv	000010	00011	00	1	10	0	0	0	0	01000	X	X	0
shra	000010	00100	00	1	10	0	0	0	1	01001	X	X	0
shrav	000010	00101	00	1	10	0	0	0	0	01001	X	X	0
lw	000101	XXXXX	01	1	01	1	0	0	1	10101	X	X	0
SW	000110	XXXXX	XX	0	XX	1	1	0	1	10101	X	X	0
b	001011	XXXXX	XX	0	XX	0	0	X	X	00000	0	0	1
bl	001100	XXXXX	10	1	00	0	0	X	X	00000	0	0	1
bey	001101	XXXXX	XX	0	XX	0	0	X	X	00000	0	0	1
bncy	001110	XXXXX	XX	0	XX	0	0	X	X	00000	0	0	1
br	001010	XXXXX	XX	0	XX	0	0	X	X	00000	X	1	1
bltz	000111	XXXXX	XX	0	XX	0	0	X	X	00000	1	0	1
bz	001000	XXXXX	XX	0	XX	0	0	X	X	00000	1	0	1
bnz	001001	XXXXX	XX	0	XX	0	0	X	X	00000	1	0	1
addi	000011	XXXXX	00	1	10	0	0	0	1	00001	X	X	0
compi	000100	XXXXX	00	1	10	0	0	1	1	00101	X	X	0
diff	001111	00000	00	1	10	0	0	0	0	00100	X	X	0

3.1 Description of Control Signals for KGP-RISC

- RegDst:- Flag Bits to make the choice for the register to which data will be written. Used to make the choice between rs, rt and \$ra.
- RegWrite:- The Flag that controls whether or NOT data is to be written to a register.
- MemToReg:- Selects the suitable line to write to the register file. Used to make a choice between (PC) + 4 (for a bl instruction) or the ALU result (for an R-type instruction), data from the data memory (for a sw instruction)
- MemRead:- The Flag that controls whether or not data is to be read from a given data-Memory address or NOT.
- MemWrite:- The Flag that controls whether or not data is to be written to a given data-memory address or NOT.

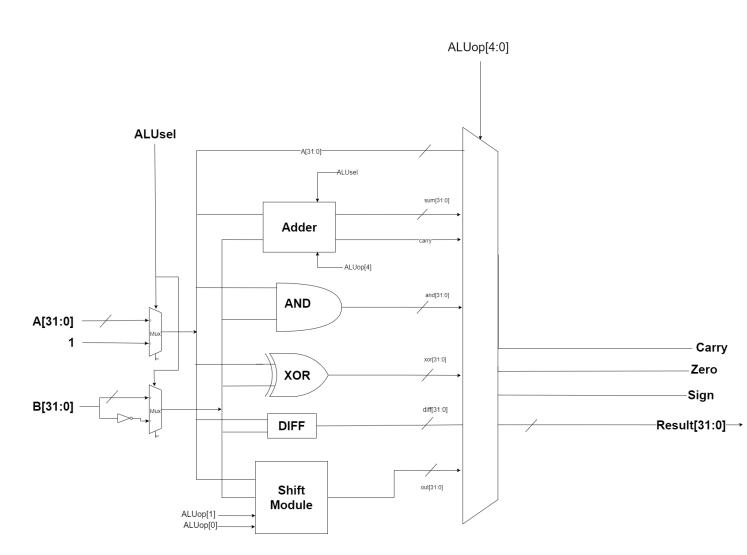


Figure 1: ALU for KGP-RISC

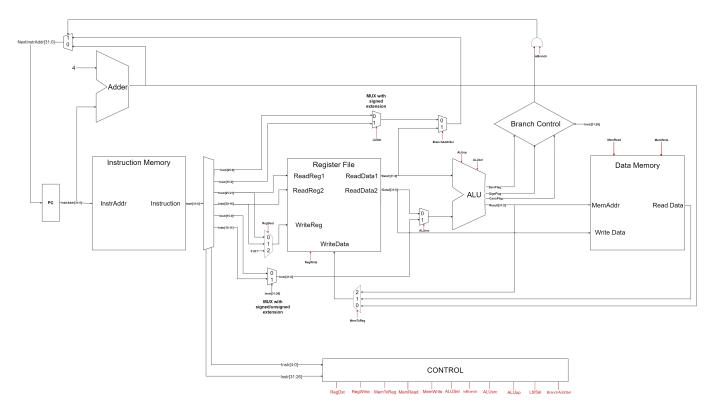


Figure 2: Datapath for KGP-RISC

- ALUsrc:- The Flag bit to control the choice between choose between the sign-extended result (for immediate operations) or the contents of a register (for R-type instructions), this acts as the second input to the ALU.
- **ALUop:-** Determines the type of operation to be carried out in the ALU.
- LblSel:- Controls the length of the label (address) selected for the branch operation. It is 0 for the instructions with a 26-bit label (b, bl, bcy, bncy) and 1 for the instructions with a 16-bit label (bltz, bz, bnz), and
- BranchAddrSel:- Controls the kind of label (address) to be selected during a jump operation. It is 1 only for the br instruction, where the address is the content of a register, otherwise it is 0.
- **isBranch:** Determines whether an instruction is a branch instruction or not(a valid jump is performed or NOT).

Table 8: Instructions in Detail

Instruction	\mathbf{ALUop}	a	b	result	Carry	Sign	Zero
add	add a, b	R[rs]	R[rt]	a+b	c_{out}	X	X
comp	comp b	R[rs]	R[rt]	$\neg b + 1$	no change	X	X
AND	and a, b	R[rs]	R[rt]	$a \wedge b$	no change	X	X
XOR	xor a, b	R[rs]	R[rt]	$a \oplus b$	no change	X	X
diff	diff a, b	R[rs]	R[rt]	$(a \oplus b) \wedge (-(a \oplus b))$	no change	X	X
shll	sll a, b	R[rs]	imm	$a \ll b$	no change	X	X
shllv	sll a, b	R[rs]	R[rt]	$a \ll b$	no change	X	X
shrl	srl a, b	R[rs]	imm	$a \gg b$	no change	X	X
shrlv	srl a, b	R[rs]	R[rt]	$a \gg b$	no change	X	X
shra	sra a, b	R[rs]	imm	$\{b^*(a[31]), a \gg b\}$	no change	X	X
shrav	sra a, b	R[rs]	R[rt]	$\{b^*(a[31]), a \gg b\}$	no change	X	X
lw	add a, b	R[rs]	imm	a+b	no change	X	X
sw	add a, b	R[rs]	imm	a+b	no change	X	X
b	none	X	X	X	no change	X	X
bl	none	X	X	X	no change	X	X
bcy	none	X	X	X	no change	X	X
bncy	none	X	X	X	no change	X	X
br	none	R[rs]	X	a	no change	X	X
bltz	update(sign)	R[rs]	X	a	no change	a[31]	X
bz	update(zero)	R[rs]	X	a	no change	X	a==0
bnz	update(zero)	R[rs]	X	a	no change	X	a == 0
addi	add a, b	R[rs]	imm	a+b	c_{out}	X	X
compi	comp b	R[rs]	imm	$\neg b + 1$	no change	X	X

3.2 Branch Unit

All other input combinations lead to is Branch = 0. The is Branch signal is always 1 in case of an unconditional branch instruction (like Branch branch). For conditional jumps based on some flags or the contents of a register it is 1 only when the condition is satisfied, otherwise it is 0.

Table 9: ALUop for Each Instruction

Operation	ALUop
add a, b	00001
comp b	00101
and a, b	00010
xor a, b	00011
diff a, b	00100
shll a, b	01010
shrl a, b	01000
shra a, b	01001
update(a,s)	00000
update(a,z)	00000
none	00000

Table 10: Truth Table for Jump Control

Instr	opcode	Zero	Sign	Carry	isBranch
bltz	000111	0	1	X	1
bz	001000	1	0	X	1
br	001010	X	X	X	1
bnz	001001	0	X	X	1
b	001011	X	X	X	1
bcy	001101	X	X	1	1
bncy	001110	X	X	0	1
bl	001100	X	X	X	1