CSE 548: Analysis of Algorithms

Lecture 8
(Amortized Analysis)

Rezaul A. Chowdhury
Department of Computer Science
SUNY Stony Brook
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A Binary Counter

counter value	counter	#bit flips	#bit resets (1→ 0)	#bit sets (0 → 1)
0	0 0 0 0 0 0 0 0			
1	0 0 0 0 0 0 0 1	1	0	1
2	0 0 0 0 0 0 1 0	2	1	1
3	0 0 0 0 0 0 1 1	1	0	1
4	0 0 0 0 0 1 0 0	3	2	1
5	0 0 0 0 0 1 0 1	1	0	1
6	0 0 0 0 0 1 1 0	2	1	1
7	0 0 0 0 0 1 1 1	1	0	1
8	0 0 0 0 1 0 0 0	4	3	1
9	0 0 0 0 1 0 0 1	1	0	1
10	0 0 0 0 1 0 1 0	2	1	1
11	0 0 0 0 1 0 1 1	1	0	1
12	0 0 0 0 1 1 0 0	3	2	1
13	0 0 0 0 1 1 0 1	1	0	1
14	0 0 0 0 1 1 1 0	2	1	1
15	0 0 0 0 1 1 1 1	1	0	1
16	0 0 0 1 0 0 0 0	5	4	1

A Binary Counter

Consider a k-bit counter initialized to 0 (i.e., all bits are 0's).
Suppose we increment the counter n times.
and cost of an increment = #bits flipped

Question: What is the worst-case total cost of n increments?

Worst-case cost of a single increment:

#bit sets ($0 \rightarrow 1$), $b_1 \leq 1$
#bit resets ($1 \rightarrow 0$), $b_0 \leq k - b_1$
#bit flips $= b_1 + b_0 \leq k$

Worst-case cost of n increments:

#bit flips $\leq nk$

This turns out to be a very loose upper bound!

Aggregate Analysis

A better upper bound can be obtained as follows.

Each increment sets ($0 \rightarrow 1$) at most one bit, i.e., $b_1 \leq 1$
So, total number of bits set by n increments, $B_1 = b_1 n \leq n$

Since at most n bits are set, there cannot be more than n bit resets ($1 \rightarrow 0$), i.e., $B_0 \leq B_1 \leq n$

So, total number of bit flips $= B_1 + B_0 \leq n + n = 2n$

Thus worst-case cost of a sequence of n increments, $T(n) \leq 2n$

Hence, in the worst case, average cost of an increment $= \frac{T(n)}{n} \leq 2$

This *worst-case average cost* is called the *amortized cost* of an increment in a sequence of n increments.

(Lecture 8) Amortized Analysis

A Binary Counter

counter value	counter	#bit flips	#bit resets (1→ 0)	#bit sets (0 → 1)	total #bit flips
0	0 0 0 0 0 0 0 0				
1	0 0 0 0 0 0 0 1	1	0	1	1
2	0 0 0 0 0 0 1 0	2	1	1	3
3	0 0 0 0 0 0 1 1	1	0	1	4
4	0 0 0 0 0 1 0 0	3	2	1	7
5	0 0 0 0 0 1 0 1	1	0	1	8
6	0 0 0 0 0 1 1 0	2	1	1	10
7	0 0 0 0 0 1 1 1	1	0	1	11
8	0 0 0 0 1 0 0 0	4	3	1	15
9	0 0 0 0 1 0 0 1	1	0	1	16
10	0 0 0 0 1 0 1 0	2	1	1	18
11	0 0 0 0 1 0 1 1	1	0	1	19
12	0 0 0 0 1 1 0 0	3	2	1	22
13	0 0 0 0 1 1 0 1	1	0	1	23
14	0 0 0 0 1 1 1 0	2	1	1	25
15	0 0 0 0 1 1 1 1	1	0	1	26
16	0 0 0 1 0 0 0 0	5	4	1	31

Amortized Analysis

– often obtains a tighter worst-case upper bound on the cost of a sequence of operations on a data structure by reasoning about the interactions among those operations

– the actual cost of any given operation may be very high, but that operation may change the state of the data structure in such a way that similar high-cost operations cannot appear for a while

– tries to show that there must be enough low-cost operations in the sequence to average out the impact of high-cost operations

– unlike average case analysis proves a worst-case upper bound on the total cost of the sequence of operations

– unlike expected case analysis no probabilities are involved

Accounting Method (Banker's View)

Consider a k-bit counter initialized to 0 (i.e., all bits are 0's).

Worst-case cost of a single increment:

#bit sets ($0 \rightarrow 1$), $b_1 \leq 1$

#bit resets ($1 \rightarrow 0$), $b_0 \leq k - b_1$

#bit flips $= b_1 + b_0 \leq k$

Thus each increment is paying for the bit it sets (fair).
But also paying for resetting bits set by prior increments (unfair)!

A fairer cost accounting for each increment:

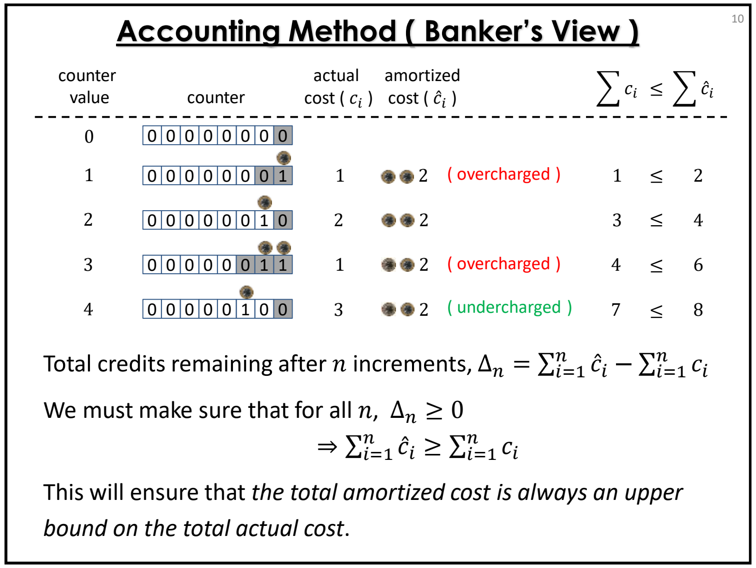
(1) Pay for the bit it sets.

(2) Pay in advance for resetting this bit (by some other increment) in the future. Store this advanced payment as a *credit* associated with that bit position.

(3) When resetting a bit use the credit stored in that bit position.

Accounting Method (Banker's View)

counter value	counter	actual cost (c_i)	amortized cost ($\hat{c}_i$)		$\sum c_i$	$\leq$	$\sum \hat{c}_i$
0	0 0 0 0 0 0 0 0						
1	0 0 0 0 0 0 0 1	1	●● 2	(overcharged)	1	$\leq$	2
2	0 0 0 0 0 0 1 0	2	●● 2		3	$\leq$	4
3	0 0 0 0 0 0 1 1	1	●● 2	(overcharged)	4	$\leq$	6
4	0 0 0 0 0 1 0 0	3	●● 2	(undercharged)	7	$\leq$	8
5	0 0 0 0 0 1 0 1	1	●● 2	(overcharged)	8	$\leq$	10
6	0 0 0 0 0 1 1 0	2	●● 2		10	$\leq$	12
7	0 0 0 0 0 1 1 1	1	●● 2	(overcharged)	11	$\leq$	14
8	0 0 0 0 1 0 0 0	4	●● 2	(undercharged)	15	$\leq$	16
9	0 0 0 0 1 0 0 1	1	●● 2	(overcharged)	16	$\leq$	18



Accounting Method (Banker's View)

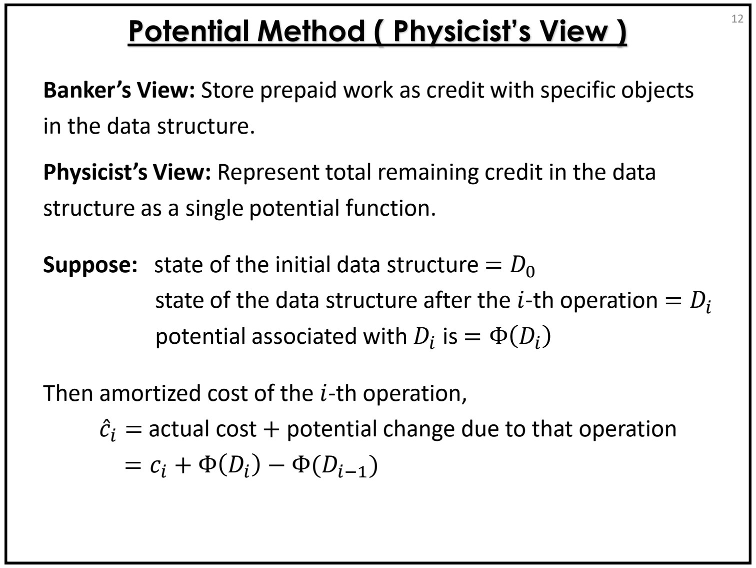
counter value	counter	actual cost (c_i)	amortized cost ($\hat{c}_i$)		$\sum c_i$	$\leq$	$\sum \hat{c}_i$
0	0 0 0 0 0 0 0 0						
1	0 0 0 0 0 0 0 1	1	●● 2	(overcharged)	1	$\leq$	2
2	0 0 0 0 0 0 1 0	2	●● 2		3	$\leq$	4
3	0 0 0 0 0 0 1 1	1	●● 2	(overcharged)	4	$\leq$	6
4	0 0 0 0 0 1 0 0	3	●● 2	(undercharged)	7	$\leq$	8

Total credits remaining after n increments, $\Delta_n = \sum_{i=1}^{n} \hat{c}_i - \sum_{i=1}^{n} c_i$

We must make sure that for all n, $\Delta_n \geq 0$

$$\Rightarrow \sum_{i=1}^{n} \hat{c}_i \geq \sum_{i=1}^{n} c_i$$

This will ensure that *the total amortized cost is always an upper bound on the total actual cost*.



Potential Method (Physicist's View)

Banker's View: Store prepaid work as credit with specific objects in the data structure.

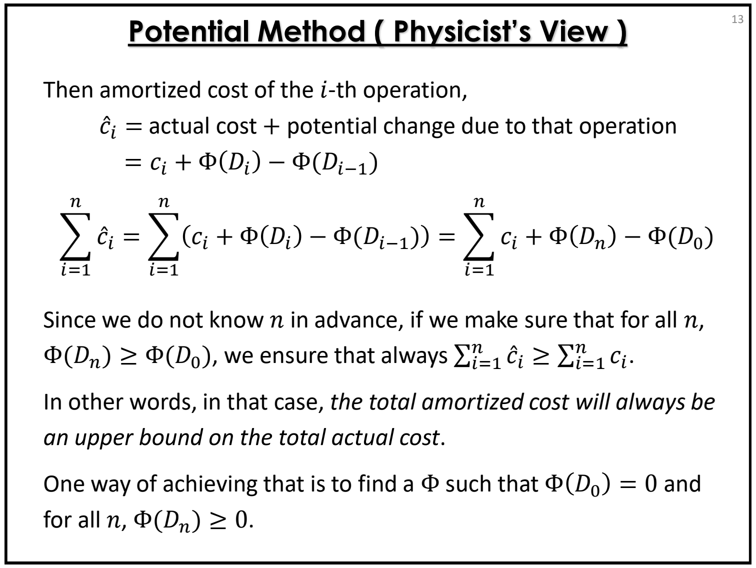
Physicist's View: Represent total remaining credit in the data structure as a single potential function.

Suppose: state of the initial data structure $= D_0$
state of the data structure after the i-th operation $= D_i$
potential associated with D_i is $= \Phi(D_i)$

Then amortized cost of the i-th operation,

$\hat{c}_i =$ actual cost $+$ potential change due to that operation

$= c_i + \Phi(D_i) - \Phi(D_{i-1})$



Potential Method (Physicist's View)

Then amortized cost of the i-th operation,

$\hat{c}_i =$ actual cost $+$ potential change due to that operation

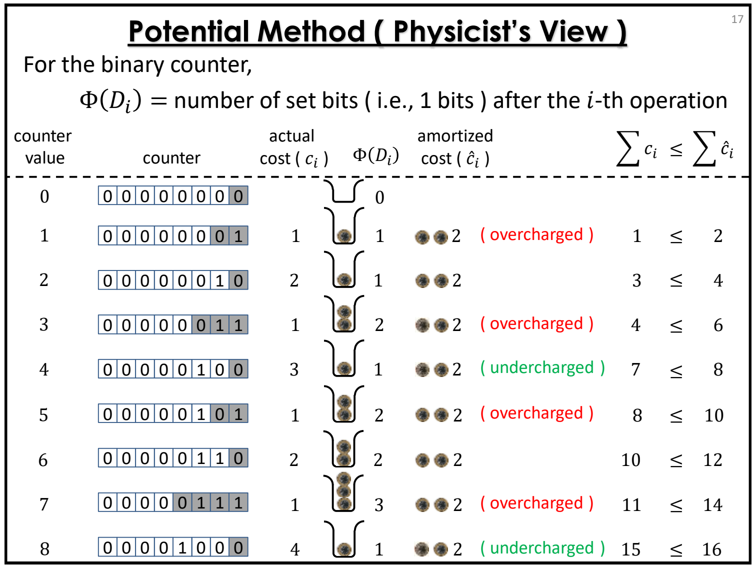
$= c_i + \Phi(D_i) - \Phi(D_{i-1})$

$$\sum_{i=1}^{n} \hat{c}_i = \sum_{i=1}^{n} (c_i + \Phi(D_i) - \Phi(D_{i-1})) = \sum_{i=1}^{n} c_i + \Phi(D_n) - \Phi(D_0)$$

Since we do not know n in advance, if we make sure that for all n, $\Phi(D_n) \geq \Phi(D_0)$, we ensure that always $\sum_{i=1}^{n} \hat{c}_i \geq \sum_{i=1}^{n} c_i$.

In other words, in that case, *the total amortized cost will always be an upper bound on the total actual cost*.

One way of achieving that is to find a Φ such that $\Phi(D_0) = 0$ and for all n, $\Phi(D_n) \geq 0$.



Potential Method (Physicist's View)

For the binary counter,

$\Phi(D_i) =$ number of set bits (i.e., 1 bits) after the i-th operation

counter value	counter	actual cost (c_i)	$\Phi(D_i)$	amortized cost ($\hat{c}_i$)		$\sum c_i$	$\leq$	$\sum \hat{c}_i$
0	0 0 0 0 0 0 0 0		0					
1	0 0 0 0 0 0 0 1	1	1	●● 2	(overcharged)	1	$\leq$	2
2	0 0 0 0 0 0 1 0	2	1	●● 2		3	$\leq$	4
3	0 0 0 0 0 0 1 1	1	2	●● 2	(overcharged)	4	$\leq$	6
4	0 0 0 0 0 1 0 0	3	1	●● 2	(undercharged)	7	$\leq$	8
5	0 0 0 0 0 1 0 1	1	2	●● 2	(overcharged)	8	$\leq$	10
6	0 0 0 0 0 1 1 0	2	2	●● 2		10	$\leq$	12
7	0 0 0 0 0 1 1 1	1	3	●● 2	(overcharged)	11	$\leq$	14
8	0 0 0 0 1 0 0 0	4	1	●● 2	(undercharged)	15	$\leq$	16