§1 POLYSLAVE INTRODUCTION 1

(Downloaded from https://cs.stanford.edu/~knuth/programs.html and typeset on September 17, 2017)

1. Introduction. This program obediently carries out the wishes of POLYNUM, which has compiled a set of one-byte and four-byte instructions for us to interpret.

But instead of producing high-precision output, it does all its arithmetic modulo a given number $m \leq 256$. That trick keeps memory usage small; it allows the user to reconstruct true answers of almost unlimited size just by trying sufficiently many different values of m. And if anybody is concerned about bits getting clobbered by cosmic radiation, they can gain additional confidence in the accuracy by running the calculations for more moduli than strictly necessary.

```
#include <stdio.h>
#include <setjmp.h>
jmp_buf restart_point;

⟨Type definitions 2⟩
⟨Global variables 5⟩
⟨Subroutines 3⟩

main(int argc, char *argv[])
{

⟨Local variables 6⟩;
⟨Scan the command line 4⟩;

setjmp(restart_point); /* longjmp will return here if necessary */
⟨Initialize 15⟩;
⟨Interpret the instructions in the input 17⟩;
⟨Print statistics 25⟩;

exit(0);
}
```

2. It is easy to adapt this program to work with counters that occupy either one byte (unsigned char), two bytes (unsigned short), or four bytes (unsigned int), depending on how much memory is available.

Even if we limit ourselves to one-byte counters, exact results of up to 362 bits can be determined. For example, the eleven moduli 256, 253, 251, 247, 245, 243, 241, 239, 233, 229, 227 will suffice to enumerate n-ominoes for $n \le 46$; and the additional modulus 223 will carry us through $n \le 50$. (If some day we have the resources to go even higher, the next moduli to try would be 211, 199, and 197.)

However, the author's experience with the case n=47 showed that the memory space needed for counters in this program was not as precious as the memory space needed for configurations in POLYNUM. Therefore the four-byte moduli $2^{31}=2147483648$, $2^{31}-1$ (which is prime), and $2^{31}-3$ (which equals $5\cdot 19\cdot 22605091$) worked out best. Together they reach nearly to 10^{28} , which would actually be large enough to count 49-ominoes.

With a little extra work I could have allowed moduli up to 2^{32} . But I didn't bother, because 2^{31} turned out to be plenty big.

```
#define maxm (1 \ll 31) /* the modulus m must not exceed this */ \langle Type definitions 2 \rangle \equiv typedef unsigned int counter; /* the main data type in our arrays */ See also section 11.
This code is used in section 1.
```

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3. The program checks frequently that everything in the input file is legitimate. If not, too bad; the run is terminated (although a debugger can help diagnose nonobvious problems). Extensive checks like this have helped the author to detect errors in the program as well as errors in the input.

```
\langle Subroutines 3 \rangle \equiv
void panic(\mathbf{char} * mess)
\{
fprintf(stderr, "%s! \n", mess);
exit(-1);
\}
See also sections 8, 9, 10, and 12.
This code is used in section 1.
```

4. Several gigabytes of input might be needed, so the input file name will be extended by .0, .1, ..., just as in POLYNUM.

Output data suitable for processing by *Mathematica* will be written on a file with the same name as the input but extended by the modulus and .m.

```
\langle Scan \text{ the command line 4} \rangle \equiv
  if (argc \neq 3 \lor sscanf(argv[2], "%u", \& modulus) \neq 1) {
    exit(-2);
  base\_name = argv[1];
  if (modulus < 2 \lor modulus > maxm) \ panic("Improper_modulus");
  m = modulus;
  sprintf(filename, "%.90s-%u.m", base_name, modulus);
  math\_file = fopen(filename, "w");
  if (\neg math\_file) panic("I_{\sqcup}can't_{\sqcup}open_{\sqcup}the_{\sqcup}output_{\sqcup}file");
This code is used in section 1.
5. \langle \text{Global variables 5} \rangle \equiv
                                 /* results will discard multiples of this number */
  unsigned int modulus;
  char *base_name, filename[100];
  FILE *math\_file;
                          /* the output file */
See also sections 7, 14, 16, and 28.
This code is used in section 1.
6. \langle \text{Local variables 6} \rangle \equiv
  register int k;
                        /* all-purpose index register */
  register unsigned int m;
                                   /* register copy of modulus */
See also section 18.
This code is used in section 1.
```

§7 POLYSLAVE INPUT 3

```
Input. Let's start with the basic routines that are needed to read instructions from the input file(s).
As soon as 2<sup>30</sup> bytes of data have been read from file foo.0, we'll turn to file foo.1, etc.
#define filelength_threshold (1 \ll 30)
                                                /* should match the corresponding number in POLYNUM */
#define buf\_size (1 \ll 16)
                                    /* should be a divisor of filelength_threshold */
\langle Global variables 5\rangle + \equiv
  FILE *in_{-}file;
                      /* the input file */
  union {
    unsigned char buf[buf\_size + 10];
                                               /* place for binary input */
    unsigned int foo;
                             /* force in.buf to be aligned somewhat sensibly */
  } in;
  unsigned char *buf_ptr;
                                  /* our current place in the buffer */
                     /* the number of bytes seen so far in the current input file */
  int bytes_in:
  unsigned int checksum;
                                  /* a way to help identify bad I/O */
  FILE *ck_{-}file;
                       /* the checksum file */
  unsigned int checkbuf;
                                 /* a check sum for comparison */
                          /* the number of GGbytes input */
  int file_extension;
8. \langle \text{Subroutines } 3 \rangle + \equiv
  void open_it()
    sprintf(filename, "%.90s.%d", base_name, file_extension);
    in_file = fopen(filename, "rb");
    if (\neg in\_file) {
       fprintf(stderr, "I_can't_open_file_%s", filename);
       panic("□for□input");
     bytes\_in = checksum = 0;
9. If the check sum is bad, we go back to the beginning. Some incorrect definitions may have been output
to math_file, but we'll append new definitions that override them.
\langle Subroutines 3\rangle + \equiv
  void close_it()
  {
    if (fread \& checkbuf, size of (unsigned int), 1, ck_-file) \neq 1)
       panic("I_{\sqcup}couldn't_{\sqcup}read_{\sqcup}the_{\sqcup}check_{\sqcup}sum");
    if (fclose(in\_file) \neq 0) panic("I_{\sqcup}couldn't_{\sqcup}close_{\sqcup}the_{\sqcup}input_{\sqcup}file");
    printf("[%d_ibytes_iread_ifrom_ifile,|%s,_ichecksum_i%u.]\n",bytes_in,filename,checksum);
    if (checkbuf \neq checksum) {
       printf("Checksum_mismatch!_Restarting...\n");
       longjmp(restart\_point, 1);
    fflush(stdout);
```

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10. My first draft of this program simply used *fread* to input one or three bytes at a time. But that turned out to be incredibly slow on my system, so now I'm doing my own buffering.

The program here uses the fact that six consecutive zero bytes cannot be present in a valid input; thus we need not make a special check for premature end-of-file.

```
#define end\_of\_buffer & in.buf[buf\_size + 4]
\langle \text{Subroutines } 3 \rangle + \equiv
  void read_it()
    register int t, k;
    register unsigned int s;
    if (bytes\_in \geq filelength\_threshold) {
      if (bytes_in ≠ filelength_threshold) panic("Improper_buffer_size");
       close_{-}it();
       file\_extension ++;
       open_it();
    t = fread(in.buf + 4, sizeof(unsigned char), buf\_size, in\_file);
    if (t < buf\_size) in.buf[t+4] = in.buf[t+5] = in.buf[t+6] = in.buf[t+7] = in.buf[t+8] = #81;
         /* will cause sync 1 error if read */
    bytes_in += t;
    for (k = s = 0; k < t; k++) s = (s \ll 1) + in.buf[k+4];
    checksum += s;
  }
```

11. A four-byte instruction has the binary form $(0xaaaaaa)_2$, $(bbbbbbb)_2$, $(cccccc)_2$, $(ddddddd)_2$, where $(aaaaaabbbbbbbbcccccccdddddddd)_2$ is a 30-bit address specified in big-endian fashion. If x=0 it means, "This is the new source address s." If x=1 it means, "This is the new target address t."

A one-byte instruction has the binary form $(1000pppp)_2$, with a 3-bit opcode $(000)_2$ and a 4-bit parameter $(pppp)_2$. If the parameter is zero, the following byte is regarded as an 8-bit parameter $(pppppppp)_2$, and it should not be zero. (In that case the "one-byte instruction" actually occupies two bytes.)

In the instruction definitions below, p stands for the parameter, s stands for the current source address, and t stands for the current target address. The slave processor operates on a large array called count.

 $\S12$ Polyslave input 5

12. The get_inst routine reads the next instruction from the input and returns the value of its parameter, also storing the opcode in the global variable op. Changes to s and t are taken care of automatically, so that op is reduced to either sync, clear, copy, or add.

```
\#define advance_b
         if (++b \equiv end\_of\_buffer) { read\_it(); b = \&in.buf[4]; }
\langle \text{Subroutines } 3 \rangle + \equiv
  opcode get_inst()
     register unsigned char *b = buf_ptr;
     register opcode o;
     register int p;
  restart: advance_b;
     if (\neg(*b \& #80)) (Change the source or target address and goto restart 13);
    o = (*b \gg 4) \& 7;
     p = *b \& #f;
    if (\neg p) {
       advance\_b;
       p = *b;
       if (\neg p) panic("Parameter_is_zero");
     \mathbf{switch} (o) {
     case inc\_src: cur\_src += p; goto restart;
     case dec\_src: cur\_src -= p; goto restart;
     case inc\_trg: cur\_trg += p; goto restart;
     case dec\_trg: cur\_trg -= p; goto restart;
     default: op = o;
     if (verbose) {
       if (op \equiv clear) \ printf("{clear_{\square}%d_{\square}->%d}\n", p, cur_trg);
       else if (op > clear) printf("{%s_{\perp}%d_{\perp}%d->%d}\n", sym[op], p, cur\_src, cur\_trg);
     buf_{-}ptr = b;
    return p;
13. (Change the source or target address and goto restart 13) \equiv
    if (b+3 > end\_of\_buffer) {
       *(b - buf\_size) = *b, *(b + 1 - buf\_size) = *(b + 1), *(b + 2 - buf\_size) = *(b + 2);
       read_it();
       b -= buf\_size;
     }
     p = ((*b \& #3f) \ll 24) + (*(b+1) \ll 16) + (*(b+2) \ll 8) + *(b+3);
     if (*b \& #40) cur_trg = p;
     else cur\_src = p;
     b += 3;
     goto restart;
This code is used in section 12.
```

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```
14. \langle \text{Global variables 5} \rangle + \equiv
                      /* operation code found by get_inst */
  opcode op:
  int verbose = 0; /* set nonzero when debugging */
  \mathbf{char} * sym[4] = \{ \texttt{"sync"}, \texttt{"clear"}, \texttt{"copy"}, \texttt{"add"} \};
  int cur_src, cur_trg;
                              /* current source and target addresses, s and t */
```

15. The first six bytes of the instruction file are, however, special. Byte 0 is the number n of cells in the largest polyominoes being enumerated. When a sync is interpreted, POLYSLAVE outputs the current values

Byte 1 is the number of the final row. If this number is r, POLYSLAVE will terminate after interpreting the instruction sync r.

```
of count[j] for 1 \le j \le n.
  Bytes 2–5 specify the (big-endian) number of elements in the count array.
  Initially s = t = 0, count[0] = 1, and count[j] is assumed to be zero for 1 \le j \le n.
\langle \text{Initialize } 15 \rangle \equiv
  sprintf(filename, "%.90s.ck", base_name);
  ck_{-}file = fopen(filename, "rb");
  if (\neg ck\_file) \ panic("I_{\sqcup}can't_{\sqcup}open_{\sqcup}the_{\sqcup}checksum_{\sqcup}file");
  open_it();
  read_it();
  n = in.buf[4];
  last\_row = in.buf[5];
  prev\_row = 0;
  slave\_size = (in.buf[6] \ll 24) + (in.buf[7] \ll 16) + (in.buf[8] \ll 8) + in.buf[9];
  buf_ptr = \&in.buf[9];
  w = n + 2 - last\_row;
  if (w < 2 \lor n < w + w - 1 \lor n > w + w + 126) panic("Badubytes_at_the_beginning");
  count = (\mathbf{counter} *) \ calloc(slave\_size, \mathbf{sizeof}(\mathbf{counter}));
  if (\neg count) panic("I_{\square}couldn't_{\square}allocate_{\square}the_{\square}counter_{\square}array");
  count[0] = 1;
                       /* prime the pump */
  cur\_src = cur\_trg = 0;
  scount = (\mathbf{counter} *) \ calloc(n+1, \mathbf{sizeof}(\mathbf{counter}));
  if (\neg scount) panic("I_{\square}couldn't_{\square}allocate_{\square}the_{\square}array_{\square}of_{\square}totals");
This code is used in section 1.
16. \langle \text{Global variables 5} \rangle + \equiv
              /* the maximum polyomino size of interest */
  int n;
                       /* the row whose end will complete our mission */
  int last_row;
                       /* the row whose end we've most recently seen */
  int prev_row;
               /* width of polynominoes being counted (deduced from n and last\_row) */
                        /* the number of counters in memory */
  int slave_size;
                           /* base address of The Big Table */
  counter *count;
                            /* base address of totals captured at sync commands */
  counter *scount;
```

§17 POLYSLAVE SERVITUDE 7

Servitude. This program is so easy to write, I could even have done it without the use of literate programming. (But of course it wouldn't be nearly as much fun without CWEB.) \langle Interpret the instructions in the input $17\rangle \equiv$ **while** (1) { $p = get_inst();$ if $(cur_trg + p > slave_size \land op \ge clear)$ $panic("Target_address_out_of_range");$ if $(cur_src + p > slave_size \land op \ge copy)$ $panic("Source_address_out_of_range");$ case sync: (Finish a row; goto done if it was the last 22); break; case *clear*: $\langle \text{Clear } p \text{ counters } 19 \rangle$; **break**: **case** copy: $\langle Copy \ p \ counters \ 20 \rangle$; **break**; case add: $\langle Add p counters 21 \rangle$; break; done:This code is used in section 1. 18. $\langle \text{Local variables } 6 \rangle + \equiv$ /* parameter of the current instruction */ register int p; register unsigned int a; /* an accumulator for arithmetic */ **19.** $\langle \text{ Clear } p \text{ counters } 19 \rangle \equiv$ for $(k = 0; k < p; k++) count[cur_trg + k] = 0;$ This code is used in section 17. **20.** $\langle \text{Copy } p \text{ counters } 20 \rangle \equiv$ for (k = 0; k < p; k++) $count[cur_trg + k] = count[cur_src + k];$ This code is used in section 17. I wonder what kind of machine language code my C compiler is giving me here, but I'm afraid to look. $\langle \text{Add } p \text{ counters } 21 \rangle \equiv$ for (k = 0; k < p; k++) { $a = count[cur_trg + k] + count[cur_src + k];$ **if** (a > m) a -= m; $count[cur_trg + k] = a;$ This code is used in section 17. 22. The sync instruction, at least, gives me a little chance to be creative, especially with respect to checking the sanity of the source file. $\langle \text{Finish a row; goto } done \text{ if it was the last } 22 \rangle \equiv$ \langle Check that p has the correct value 23 \rangle ; (Output the relevant counters for completed polyominoes 24); for $(k = 1; k \le n; k ++)$ scount[k] = count[k];

if $(p \equiv last_row)$ goto done;

This code is used in section 17.

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```
23. \langle Check that p has the correct value 23 \rangle \equiv
  if (p \equiv 255) (Go into special shutdown mode 26);
  if (p \equiv 1) \ panic("File_lread_lerror");
                                                      /* see read_it */
  if (\neg prev\_row) {
     if (p \neq w + 1) panic("Bad_{\sqcup}first_{\sqcup}sync");
  } else if (p \neq prev\_row + 1) panic("Out \cup of \cup sync");
  prev\_row = p;
This code is used in section 22.
24. (Output the relevant counters for completed polyominoes 24) \equiv
  printf("Polyominoes_{\sqcup}that_{\sqcup}span_{\sqcup}%dx%d_{\sqcup}rectangles_{\sqcup}(mod_{\sqcup}%u): \n", p-1, w, m);
  fprintf(math\_file, "p[%d,%d,%u]={0", p-1, w, m};
  for (k = 2; k < w + p - 2; k++) fprintf (math\_file, ", 0");
  for ( ; k \le n; k++)  {
     if (count[k] \ge scount[k]) a = count[k] - scount[k];
     else a = count[k] + m - scount[k];
     printf(" \sqcup %d: %d", a, k);
     fprintf(math\_file, ", %d", a);
  printf("\n");
  fflush(stdout);
  fprintf(math\_file, "}\n");
This code is used in section 22.
25. \langle \text{ Print statistics } 25 \rangle \equiv
  printf("All_{\square}done!_{\square}Final_{\square}totals_{\square}(mod_{\square}%u): \n", m);
  for (k = w + w - 1; k \le n; k++) {
     printf(" \sqcup %d: %d", count[k], k);
  printf("\n");
  close\_it();
This code is used in section 1.
```

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26. Checkpointing. POLYNUM issues the special command *sync* 255 when it wants to pause for breath and shore up its knowledge. Therefore, if we see that instruction, we must immediately dump all the counters into a temporary file. A special variant of this program is able to read that file and reconstitute all the data, as if there had been no break in the action. (See the change file polyslave-restart.ch for details.)

```
\langle Go into special shutdown mode 26 \rangle \equiv
     close_it();
     printf("Checkpoint\_stop:\_After\_processing\_with\_all\_desired\_moduli,\n");
     printf("_please_resume_with_polynum-restart_and_polyslave-restart.\n");
     sprintf (filename, "%.90s-%u.dump", base_name, m);
     out\_file = fopen(filename, "wb");
     if (\neg out\_file) panic("I_{\sqcup}can't_{\sqcup}open_{\sqcup}the_{\sqcup}dump_{\sqcup}file");
     ⟨Dump all information needed to restart 27⟩;
     exit(1);
This code is used in section 23.
27. \langle \text{Dump all information needed to restart } 27 \rangle \equiv
  dump\_data[0] = n;
  dump\_data[1] = w;
  dump\_data[2] = m;
  dump\_data[3] = slave\_size;
  dump\_data[4] = prev\_row;
  if (fwrite(dump\_data, sizeof(unsigned int), 5, out\_file) \neq 5)
     panic("Bad_write_at_beginning_of_dump");
  if (fwrite(scount, sizeof(counter), n+1, out\_file) \neq n+1) panic("Couldn't_udump_uthe_usubtotals");
  if (fwrite(count, sizeof(counter), slave\_size, out\_file) \neq slave\_size)
     panic("Couldn't dump the counters");
  printf("[\%u_{\sqcup}bytes_{\sqcup}written_{\sqcup}on_{\sqcup}file_{\sqcup}\%s.]\n", ftell(out\_file), filename);
This code is used in section 26.
28.
     \langle \text{Global variables 5} \rangle + \equiv
  unsigned int dump\_data[5];
                                        /* parameters needed to restart */
  FILE *out_file;
```

10 CHECKPOINTING POLYSLAVE §29

29. For the record, here are three shell scripts called nums, slaves, and slaves-restart, which were used to run POLYNUM and POLYSLAVE when n = 47:

```
#!/bin/sh
nums
                  if [ $# -ne 3 ]; then
                    echo "Usage: nums width configs counts"
                    exit 255
                  time polynum 47 $1 $2 $3 /home/tmp/poly47-$1
                  slaves $1
                  while [ \$? = 1 ]; do
                    mv /home/tmp/poly47-$1.dump /home/tmp/poly47-$1.dump~
                    time polynum-restart 47 $1 $2 $3 /home/tmp/poly47-$1
                    slaves-restart $1
                  done
                  #!/bin/sh
slaves
                  for m in 2147483648 2147483647 2147483645; do
                    time polyslave /home/tmp/poly47-$1 $m
                  done
                  #!/bin/sh
slaves-restart
                  for m in 2147483648 2147483647 2147483645; do
                    cp /home/tmp/poly47-$1-$m.dump /home/tmp/poly47-$1-$m.dump~
                    time polyslave-restart /home/tmp/poly47-$1 $m
                  done
```

And here is the *Mathematica* script used to convert modular numbers to multiprecise integers:

```
(* for Chinese Remainders, say for example
    chinese[{13,17,19}]
    x=cdecode[{1,2,3}]
and x (= 4031) will satisfy Mod[x,13]=1, Mod[x,17]=2, Mod[x,19]=3 *)
chinese[1_]:=Block[{},chinmod=Apply[Times,1];
    chinlist=Table[(chinmod/l[[k]])PowerMod[chinmod/l[[k]],-1,1[[k]]],
        {k,Length[1]}]]
cdecode[1_]:=Mod[chinlist.1,chinmod]
m=2^31
chinese[{m,m-1,m-3}]
fn[a_,b_]:="poly47-"<>a<>"-"<>ToString[m-b]<>".m"
squash[a_,w_]:=Block[{},Get[fn[a,0]];Get[fn[a,1]];Get[fn[a,3]];
        Do[q[h,w]= cdecode[{p[h,w,m],p[h,w,m-1],p[h,w,m-3]}],{h,w,48-w}];
        Save["poly47-"<>a<>".m",q];
        Clear[q]]
```

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30. Index.

 $a: \underline{18}.$ add: 11, 12, 17. $advance_b$: $\underline{12}$. $argc: \underline{1}, \underline{4}.$ $argv: \underline{1}, \underline{4}.$ *b*: <u>12</u>. base_name: 4, 5, 8, 15, 26. buf: $\underline{7}$, 10, 12, 15. $buf_size: \underline{7}, 10, 13.$ $bytes_in\colon \ \ \underline{7},\ 8,\ 9,\ 10.$ calloc: 15. $checkbuf: \underline{7}, \underline{9}.$ checksum: 7, 8, 9, 10. $ck_file: \underline{7}, 9, 15.$ $clear \colon \ \underline{11}, \ 12, \ 17.$ $close_{-}it: \ \underline{9}, \ 10, \ 25, \ 26.$ copy: 11, 12, 17.count: 11, 15, <u>16</u>, 19, 20, 21, 22, 24, 25, 27. counter: 2, 15, 16, 27. cur_src: 12, 13, <u>14</u>, 15, 17, 20, 21. cur_trg: 12, 13, <u>14</u>, 15, 17, 19, 20, 21. dec_src : 11, 12. dec_trg : 11, 12. done: $\underline{17}$, $\underline{22}$. $dump_data$: 27, 28. $end_of_buffer: \underline{10}, 12, 13.$ exit: 1, 3, 4, 26. fclose: 9. fflush: 9, 24. $filelength_threshold$: 7, 10. filename: $4, \underline{5}, 8, 9, 15, 26, 27$. $foo: \underline{7}.$ fopen: 4, 8, 15, 26. fprintf: 3, 4, 8, 24. fread: 9, 10. ftell: 27.fwrite: 27. $get_inst: 12, 14, 17.$ in: $\underline{7}$, 10, 12, 15. *in_file*: $\underline{7}$, 8, 9, 10. inc_src : 11, 12. $inc_trg: \underline{11}, \underline{12}.$ k: <u>6</u>, <u>10</u>. *last_row*: 15, <u>16</u>, 22. longjmp: 1, 9. $m: \underline{6}.$ $main: \underline{1}.$ $math_file$: 4, $\underline{5}$, 9, 24. maxm: 2, 4.

 $mess: \underline{3}.$ modulus: 4, 5, 6. $n: \underline{16}.$ $o: \ \ \underline{12}.$ $op: 12, \underline{14}, 17.$ opcode: <u>11</u>, 12, 14. $open_{-}it: 8, 10, 15.$ $out_file\colon \ \ 26,\ 27,\ \underline{28}.$ p: <u>12</u>, <u>18</u>. panic: 3, 4, 8, 9, 10, 12, 15, 17, 23, 26, 27. $prev_row\colon \ 15,\ \underline{16},\ 23,\ 27.$ printf: 9, 12, 24, 25, 26, 27. $read_it: 10, 12, 13, 15, 23.$ restart: 12, 13. $restart_point$: $\underline{1}$, $\underline{9}$. $s: \underline{10}$. scount: 15, <u>16</u>, 22, 24, 27. setjmp: 1.slave_size: 15, 16, 17, 27. sprintf: 4, 8, 15, 26. sscanf: 4.stderr: 3, 4, 8. stdout: 9, 24.sym: 12, 14. $sync: 10, \underline{11}, 12, 15, 16, 17, 22, 26.$ t: $\underline{10}$. $targ_bit$: $\underline{11}$. $verbose: 12, \underline{14}.$ $w: \underline{16}$.

12 NAMES OF THE SECTIONS POLYSLAVE

```
\langle \text{Add } p \text{ counters } 21 \rangle Used in section 17.
(Change the source or target address and goto restart 13) Used in section 12.
(Check that p has the correct value 23) Used in section 22.
Clear p counters 19 \ Used in section 17.
 Copy p counters 20 \rangle Used in section 17.
(Dump all information needed to restart 27) Used in section 26.
(Finish a row; goto done if it was the last 22) Used in section 17.
 Global variables 5, 7, 14, 16, 28 \ Used in section 1.
(Go into special shutdown mode 26) Used in section 23.
\langle \text{Initialize } 15 \rangle \quad \text{Used in section } 1.
\langle Interpret the instructions in the input 17 \rangle Used in section 1.
(Local variables 6, 18) Used in section 1.
\langle Output the relevant counters for completed polyominoes 24\rangle Used in section 22.
(Print statistics 25) Used in section 1.
\langle Scan the command line 4 \rangle Used in section 1.
\langle Subroutines 3, 8, 9, 10, 12\rangle Used in section 1.
\langle Type definitions 2, 11 \rangle Used in section 1.
```

POLYSLAVE

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