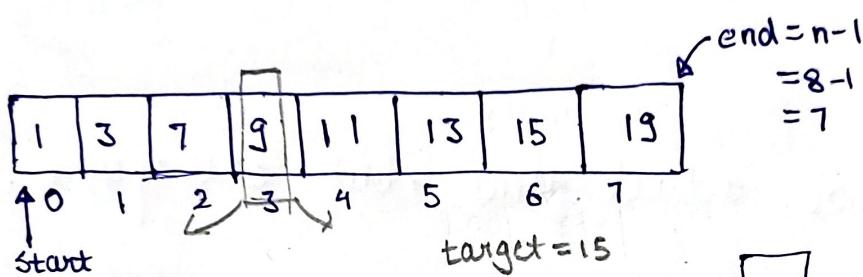
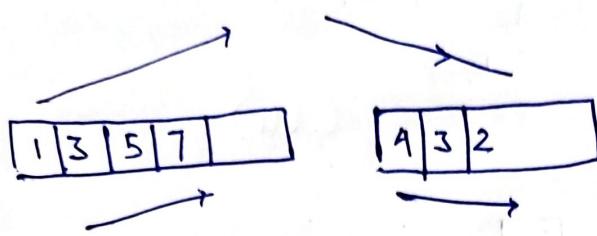


Lec1:Binary Search :-

cond: element should be in monotonic order.

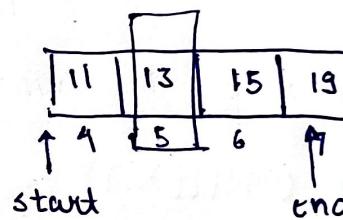


① start  $\rightarrow$  0 index

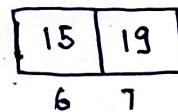
② end  $\rightarrow$  last index

$$n-1 = 7$$

$$\begin{aligned} \textcircled{3} \quad \text{mid} &= \left( \frac{\text{start} + \text{end}}{2} \right) \\ &= \frac{0+7}{2} = \textcircled{3} \end{aligned}$$



$$\text{mid} = \frac{4+7}{2} = \frac{11}{2} = 5$$



$$\text{mid} = \frac{6+7}{2} = \frac{13}{2} = 6$$

if 15 on 6 is target?  
Yes

```
#include <iostream>
using namespace std;
```

```
int binarySearch(int arr[], int size, int target){
    int start = 0;
    int end = size - 1;
```

```

int mid = (start + end) / 2;

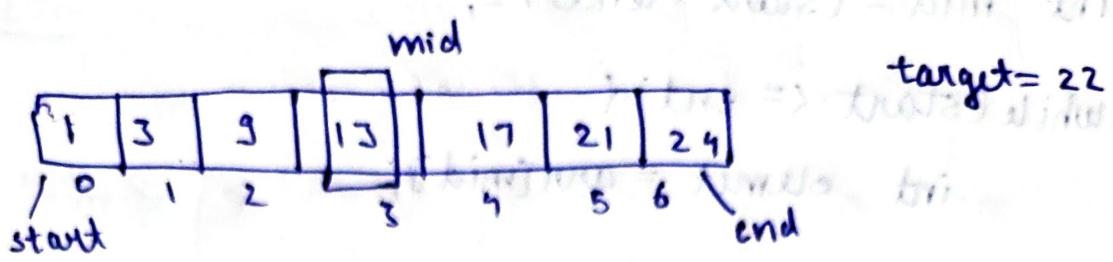
while (start <= end) {
    int element = arr[mid];
    if (element == target) {
        return mid;
    } else if (target < element) {
        // search in left
        end = mid - 1;
    } else {
        start = mid + 1; // search in right
    }
    mid = (start + end) / 2;
}
// element not found
return -1;
}

```

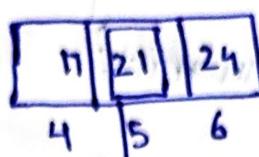
```

int main() {
    int arr[] = {2, 4, 6, 8, 10, 12, 16};
    int size = 7;
    int target = 2;
    int indexOfTarget = binarySearch(arr, size, target);
    if (indexOfTarget == -1) {
        cout << "target not found" << endl;
    } else {
        cout << "target found at " << indexOfTarget << endl;
    }
    return 0;
}

```



if target < arr[mid]:



if target > arr[mid]:

$$22 > 21$$



Start > e

Loop stop

$(s)(\text{arr}[s] + \text{arr}[e]) = \text{target}$

Here element not found

break for break

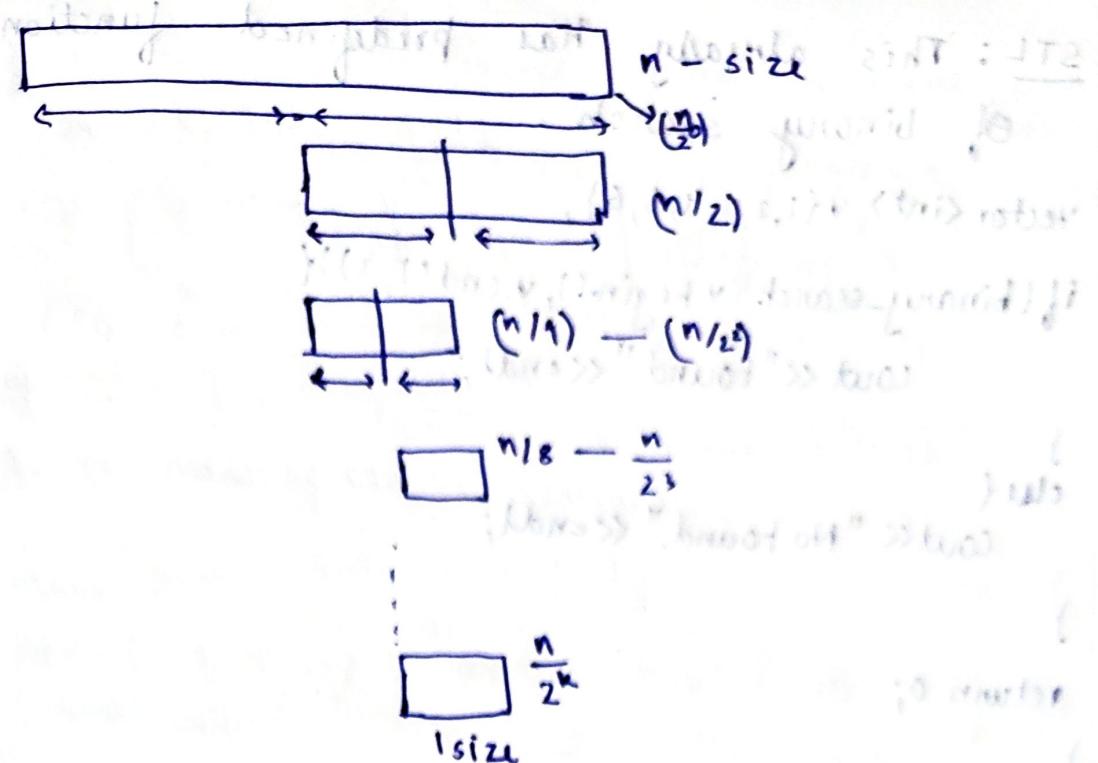
Note:-

$\text{mid} = \frac{s+e}{2}$  has a issue.

Integer overflow हो सकता है।

so, we use:

$$\boxed{\text{mid} = s + \frac{(e-s)}{2}}$$



### Time complexity (Derivation)

n size array, two at (middle) width

1st iteration  $\rightarrow \frac{n}{2^1}$

2nd "  $\rightarrow \frac{n}{2^2}$

kth "  $\rightarrow \frac{n}{2^k}$

$\frac{n}{2^k} = 1$ ; as if "only one element is now there"

$$n = 2^k$$

$$k = \log_2 n$$

So, Time.C of Binary search is  $O(\log n)$ .

- Binary Search is faster than linear search.

STL: This already has predefined function of binary search.

```
vector<int> v{1,2,3,4,5,6};  
if(binary-search(v.begin(), v.end(), 3)){  
    cout << "Found" << endl;  
}  
else{  
    cout << "No Found." << endl;  
}  
return 0;
```

To use binary-search, we need #include <algorithm> in our program.

Also, For array:

```
int arr[] = {1, 2, 3, 4};  
int size = 4;  
if (binary-search(arr, arr + size, 3)){  
    cout << "Found" << endl;  
}  
else{  
    cout << "Not Found" << endl;  
}
```

① find the first occurrence of an element, using

Binary Search

mid										target = 4
1	3	4	4	4	4	4	4	6	7	9
0	1	2	3	4	5	6	7	8	9	

(target is occurring in the given list)

Ans -

It is monotonic.

start = 0

end = 9

$$\text{mid} = \frac{\text{start} + \text{end}}{2} = 4$$

i) arr[mid] == target

$4 == 4 \rightarrow \text{yes/true} \rightarrow \text{element stored} \rightarrow \text{index } 4$   
↳ search in left

1	3	4	4
0	1	2	3

$$\text{mid} = 3$$

$$4 > 3$$

1	3
2	3

↳ (bin < mid) & (arr[mid] == target) if yes

$$\text{mid} = 2$$

arr[mid] == target  $\rightarrow$  arr[2]  $\rightarrow$  index  $\rightarrow 2$

$$\rightarrow L-S \xrightarrow{\text{each}} e = \text{mid} - 1$$

↳ (bin > target) if no

Search failed

$$1 - bin = 0$$

$$16 + 0 = bin$$

16 is result

```

int main(){
    vector<int> v{1, 3, 4, 4, 4, 4, 6, 7};
    int target=4;
    int indexOfFirstOcc = firstOccurrence(v, target);
    cout << "ans is " << ans << endl;
    return 0;
}

```

```

int firstOccurrence(vector<int> arr, int target){
    int s=0;
    int e = v.size() - 1;
    int mid = (s+e)/2;
    int ans = -1;

    while(s <= e){
        if(arr[mid] == target){
            // ans store and then left search
            ans = mid;
            e = mid - 1;
        }
        else if(target > arr[mid]){
            // right search
            s = mid + 1;
        }
        else if(target < arr[mid]){
            // left search
            e = mid - 1;
        }
        mid = (s+e)/2;
    }
    return ans;
}

```

## Dry Run

① if ( $\text{arr}[\text{mid}] == \text{target}$ )

    ↳  $\text{ans} = \text{mid};$

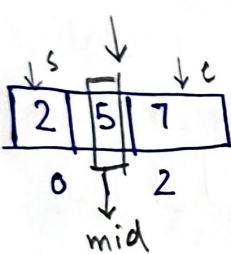
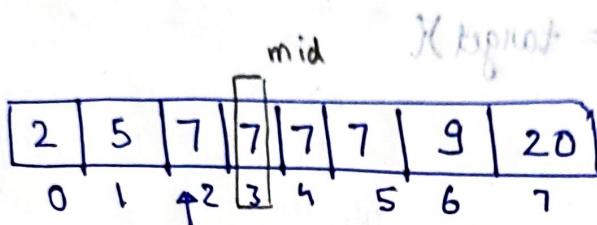
    ↳  $e = \text{mid} - 1;$

② if ( $\text{target} < \text{arr}[\text{mid}]$ )

    ↳  $e = \text{mid} - 1;$

③ if ( $\text{target} > \text{arr}[\text{mid}]$ )

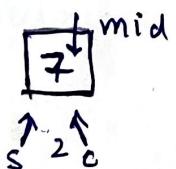
    ↳  $s = \text{mid} + 1;$



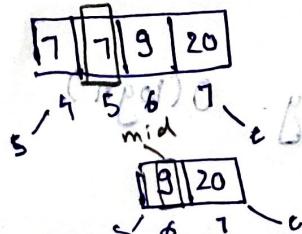
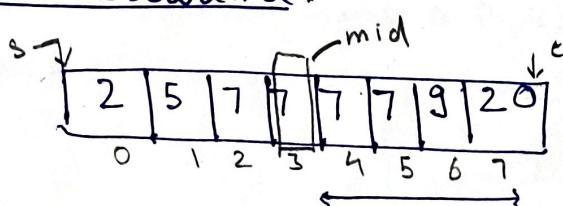
~~target < arr[mid]~~ X

mid target > arr[mid]

    ↳  $s = \text{mid} + 1$



Last occurrence :-



if ( $\text{arr}[\text{mid}] == \text{target}$ ) {

    ↳  $\text{ans} = \text{mid}$

    ↳  $s = \text{mid} + 1$

if ( $\text{target} < \text{arr}[\text{mid}]$ ) {

    ↳  $e = \text{mid} - 1$

if ( $\text{target} > \text{arr}[\text{mid}]$ ) {

    ↳  $s = \text{mid} + 1$

```

#include <iostream>
#include <vector>
using namespace std;

int lastOcc(vector<int> arr, int target) {
    int s = 0;
    int e = arr.size() - 1;
    int mid = s + (e - s) / 2;
    int ans = -1;

    while (s <= e) {
        if (arr[mid] == target) {
            // ans store
            ans = mid;
            // right search
            s = mid + 1;
        } else if (target < arr[mid]) {
            // left search
            e = mid - 1;
        } else if (target > arr[mid]) {
            // right search
            s = mid + 1;
        }
        mid = s + (e - s) / 2;
    }
    return ans;
}

int main() {
    vector<int> arr{1, 2, 2, 2, 2, 2, 3};
    int target = 2;
}

```

Time complexity  $O(\log n)$

STL functions for first and last occurrence:-

lower\_bound(v.begin(), v.end(), 20);  
                        └→ target

Lower bound is used to find the first occurrence and it returns an iterator.

Similarly, upper-bound is used to find last occurrence.

upper-bound (v.begin(), v.end(), 20)  
                        └→ target element .

2) Total number of occurrences of element.

i/p - 2, 4, 4, 4, 4, 5, 6      ~~8 5 3 1 4 3 1~~  
                        └→ 2 2 2 2 3 1 0

o/p → 4

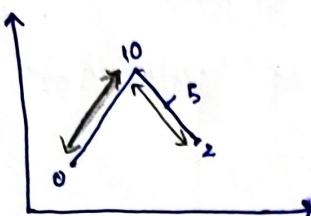
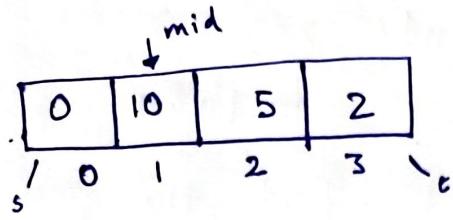
Aus:

Total occurrence = lastOccurrence - firstOccurrence + 1.

Code:

```
int first = firstOccurrence(v, 4);
int last = lastOccurrence(v, 4)
cout << (last - first + 1) << endl;
```

## Leetcode: Peak Index in a Mountain Array:

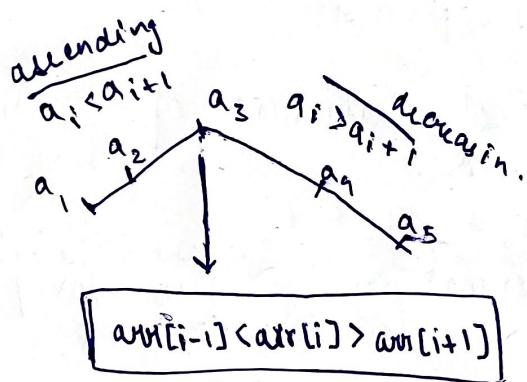


Brute force →

LS → max element

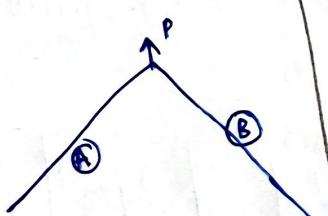
↓  
 $O(n)$

BS →  $O(\log n)$   
optimal



if ( $arr[mid] > arr[mid+1]$ ) → mid is Line A

if ( $arr[mid] < arr[mid+1]$ ) → mid element may be a peak element  
→ mid element in Line B



Line A →

$arr[i] < arr[i+1]$

P → peak element

$arr[i-1] < arr[i] > arr[i+1]$

Line B →  $arr[i] > arr[i+1]$

```
if (arr[mid] < arr[mid + 1])
```

```
s = mid + 1
```

```
else
```

```
{
```

```
e = mid;
```

```
}
```

Code:-

```
class Solution{
```

```
public:
```

```
int findPeakIndex(vector<int> arr){
```

```
int s = 0;
```

```
int e = arr.size() - 1;
```

```
int mid = s + (e - s) / 2;
```

```
while (s <= e) {
```

```
if (arr[mid] < arr[mid + 1]) {
```

// right search

```
s = mid + 1;
```

```
} else {
```

```
e = mid;
```

```
}
```

```
mid = s + (e - s) / 2;
```

```
}
```

```
return s;
```

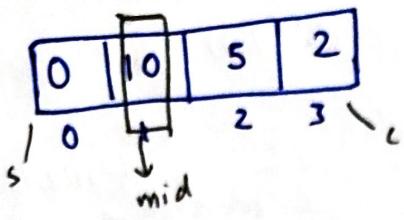
```
}
```

```
int peakIndexInMountainArray(vector<int>& arr){
```

```
return findPeakIndex(arr);
```

```
}
```

```
};
```



scc

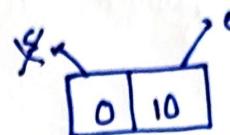
↙  
go in loop

if (arr[mid] < arr[mid+1])  
    ↳ R-S      FALSE

↓  
else {

e = mid;

}

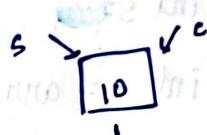


' returns;

OR

return e;

arr[mid] < arr[mid+1]

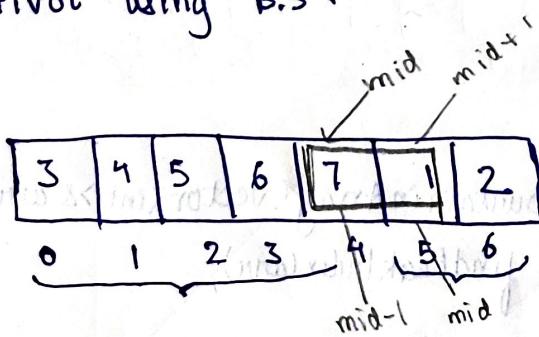


① find pivot using Binary search

②

Lec 2:-

① Pivot using B.S.



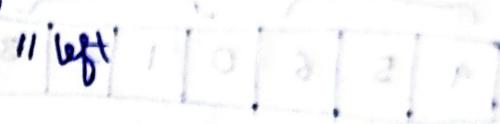
if ( $\text{arr}[\text{mid}] > \text{arr}[\text{mid} + 1]$ )  
    return mid;

→ Jiske bade monotonic fun break krrha hai  
use 'pivot'.

if ( $\text{arr}[\text{mid} - 1] > \text{arr}[\text{mid}]$ )  
    return mid - 1

if ( $\text{arr}[s] > \text{arr}[\text{mid}]$ )

s = mid - 1



else if ( $\text{arr}[s] < \text{arr}[\text{mid}]$ )

s = mid + 1

    // right

### CODE:

```
int findPivot(vector<int> arr){  
    int s = 0;  
    int e = arr.size() - 1;  
    int mid = s + (e - s) / 2;  
  
    while (s <= e){  
        if ( $\text{arr}[\text{mid}] > \text{arr}[\text{mid} + 1]$ )  
            return mid;  
        if ( $\text{arr}[\text{mid} - 1] > \text{arr}[\text{mid}]$ )  
            return mid - 1;  
  
        else  
            s = mid + 1;
```

$mid = s + (e-s)/2;$

( $b1 \text{ and } b2 < b3$ )  $\Rightarrow$   $b1 \text{ and } b2$  are in left half

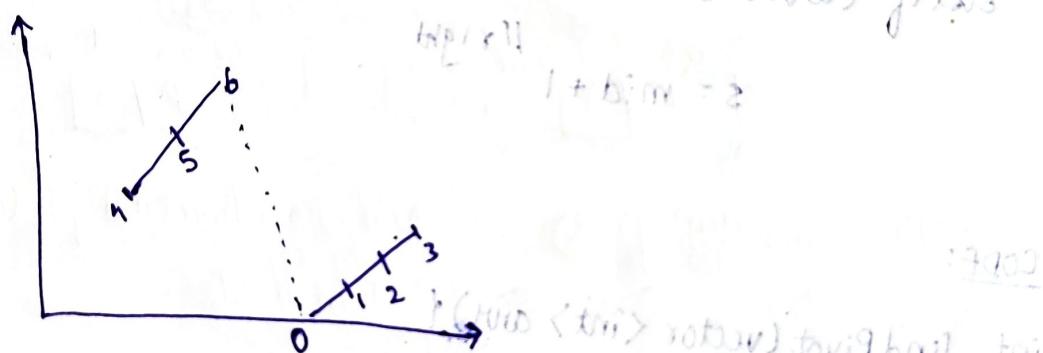
}  
return -1;

and current does not lie in some good block  
 (if  $b1 \text{ and } b2 > b3$ )  $\Rightarrow$   $b1 \text{ and } b2$  are in right half  
 (if  $b1 \text{ and } b2 < b3$ )  $\Rightarrow$   $b1 \text{ and } b2$  are in left half

## ② Search in rotated sorted array.

4	5	6	0	1	2	3
---	---	---	---	---	---	---

( $b1 \text{ and } b2 < b3$ )  $\Rightarrow$   $b1 \text{ and } b2$  are in left half  
 $l + bim = 3$



sorted  $\rightarrow$   $0 \rightarrow$  pivot

sorted  $\rightarrow$  pivot  $+ 1 \rightarrow n-1$

$$l + (2-3) + 3 = bim \text{ tri}$$

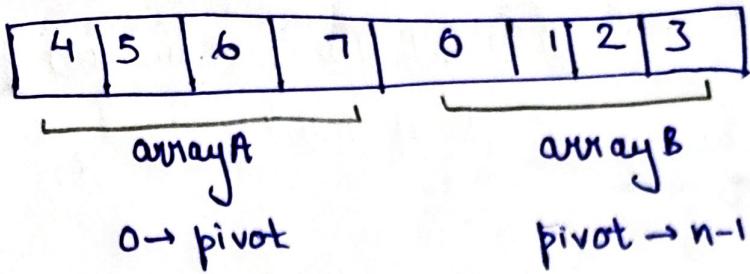
( $b1 \text{ and } b2 < b3$ )  $\Rightarrow$   $b1 \text{ and } b2$  are in left half

if  $b1 \text{ and } b2$  are in left half

( $b1 \text{ and } b2 > b3$ )  $\Rightarrow$   $b1 \text{ and } b2$  are in right half

if  $b1 \text{ and } b2$  are in right half

$$l + bim = 3$$



```

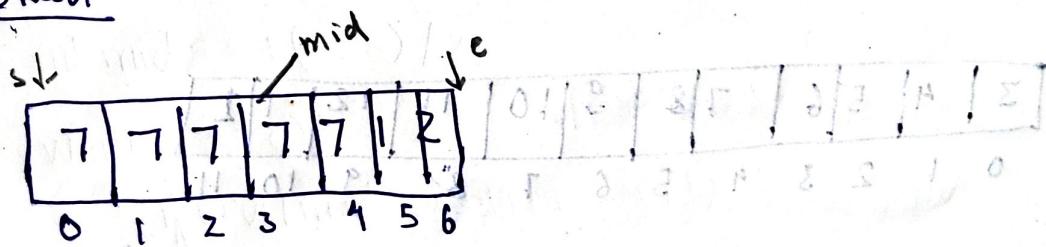
int searchVector<int> & nums, int target) {
    int pivotIndex = findPivot(nums);
    if (target >= nums[0] && target <= nums[pivot]) {
        // search A
        int ans = binarySearch(nums, target, 0, pivot);
        return ans;
    }
    if (target >= nums[pivot+1] && target <=
        // search B
        int ans = binarySearch(nums, target, pivot+1,
                               nums.size() - 1);
    }
}

```

([2, 3, 4, 5, 6, 7]) 10 < [1 - bim] return ans;

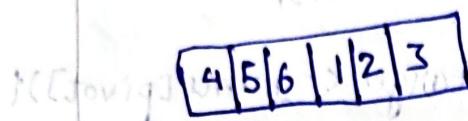
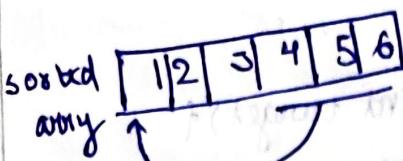
(1 - bim) mit target  
return -1;

D Run

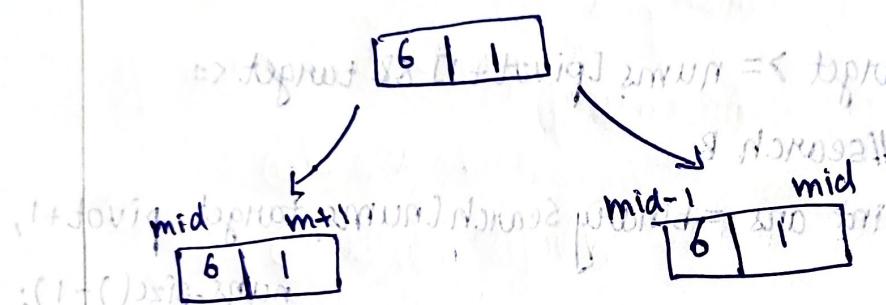
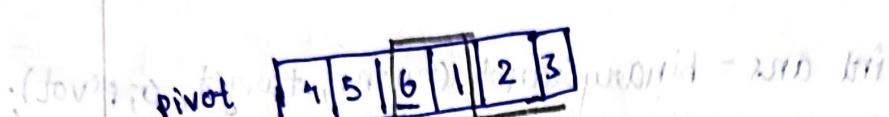


([2, 3, 4, 5, 6, 7]) 10 > [1 - bim] 4 + bim

4 + bim - 2



sorted &  
Rotated



① if arr[mid] > arr[mid+1]

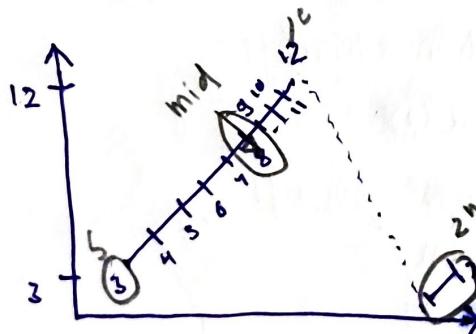
return mid

② if arr[mid-1] > arr[mid]

return (mid-1);

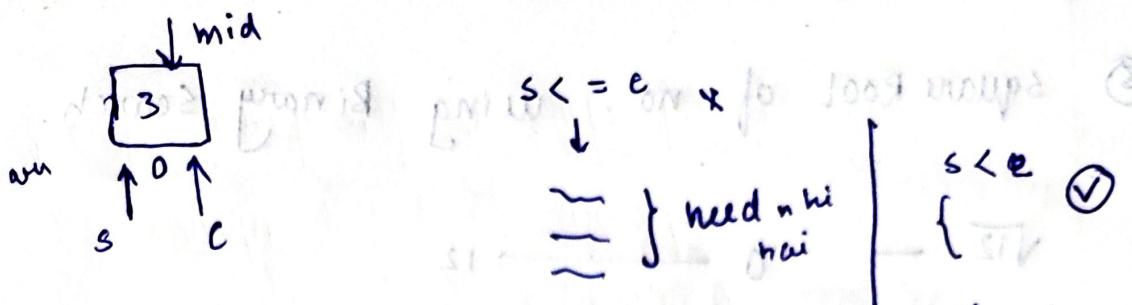
Now, check left search or right search:

3	4	5	6	7	8	9	10	11	12	13
0	1	2	3	4	5	6	7	8	9	10, 11,



③ if arr[mid] < arr[s]  
s = mid + 1;

else  
e = mid - 1;



$s < e$  for loop range  $\checkmark$   
 $\{ \}$  need  $n^{hi}$  mai  
 $\{ \}$   $s < e$   $\checkmark$   
 returns;  
 $s = mid$   
 or  
 $e = mid$

3	4	5	6	7	8	9	10	11	12	1	2
$s$		*					*	$e$	$s$	$B$	$c$

rotated &  
sorted  
Arr mai  
search.

$A \rightarrow (0 \text{ to pivot})$

$B \rightarrow (\text{pivot} + 1 \rightarrow n - 1)$

$3 - 12$



2 exist  $\times$

$3 - 2 \checkmark$

2 exist  $\checkmark$

$(\text{pivot} = \text{bin} * \text{bin}) \checkmark$

bin marker

$\beta(n \log^2 \log n)$

$O = \delta \cdot \log$

$m = \delta \cdot \log$

$\{ s \}(2 - 2) + 2 = \log \log$

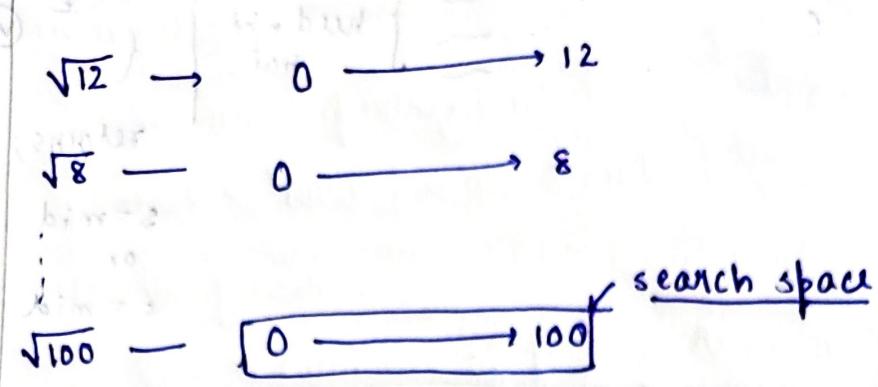
$\log \log \log \log$

$(\text{pivot} = \text{bin} * \text{bin}) \checkmark$

bin marker

$(\text{pivot} < \text{bin} * \text{bin}) \checkmark$

### ③ square root of no. using Binary Search.



Sol<sup>n</sup>

Algo:  
if ( $mid * mid > target$ )

    ↳ L.S

if ( $mid * mid < target$ )

    ① L.s store

    ② Right search

if ( $mid * mid == target$ )

    return mid

int findSqrt (int n) {

    int s = 0;

    int e = n;

    int mid = s + (e - s) / 2;

    while (s <= e) {

        if ( $mid * mid == target$ )

            return mid;

        if ( $mid * mid > target$ )

            // LS

            e = mid - 1;

}

else {

// ans store

ans = mid;

// R.S

s = mid + 1;

}

mid = s + (e - s) / 2;

}

return ans;

}

int main() {

int n;

cin >> n;

int ans = findSqrt(n);

cout << "Ans is " << ans << endl;

return 0;

0	1	2	3	4	5	6
0	1	2	3	4	5	6

mid



$\sqrt{5}$

```

int main() {
    int ans = findSqrt(n);
    cout << "Ans is " << ans << endl;
    int precision;
    cout << "Enter no. of floating digit in precision" << endl;
    cin >> precision;
    double step = 0.1;
    double finalAns = ans;
    for (int i = 0; i < precision; i++) {
        for (double j = finalAns; j * j <= n; j += step) {
            finalAns = j;
        }
        step = step / 10;
    }
    cout << "Final ans is." << finalAns << endl;
    return 0;
}

```

Dry Run:-

```

for (int i = 0; i < precision; i++) // adding new digit
    for (double j = finalAns; j * j <= n; j = j + step) {
        finalAns = j; // ans store
    }
    step = step / 10;
}

```

$$J = 3 \\ 3 \times 3 = 9 < 10 \rightarrow T$$

$$J = 3 + 0.1 = 3.1$$

$$3.1 \times 3.1 = 9.61 < 10 \checkmark$$

$$\therefore 3.1 + 0.1 = 3.2$$

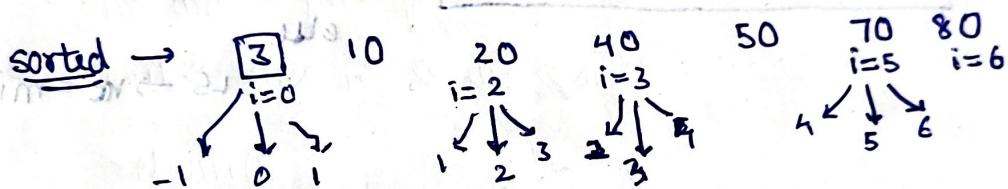
$$3.2 \times 3.2 = 10.24 < 10 \checkmark$$

- ① Search in a nearly sorted array in log time.
- ② Division of 2 no. using B.S.
- ③ Find odd occurring element using Binary Search.

Lec 3:

### Binary Search in nearly sorted Array:-

10	3	40	20	50	80	70
5 - bim 0	1	2	3	4	5	6



Nearly sorted array:-

Element  $\downarrow$  j (topnot  $\Rightarrow$  target)  $\rightarrow$  increasing order  $\Rightarrow$  i

sorted array  $\Rightarrow$   $i = 0$

$\downarrow$   $i^{th}$  index pe hai

$\downarrow$  then, in Nearly sorted array  $\Rightarrow$   $i \geq 2$  width

$\downarrow$   $(topnot = [bim] now)$

$(i-1)^{th}, i^{th}, (i+1)^{th}$  indexes  $\Rightarrow$  bim number

$(topnot = [target] = [i = bim] now)$

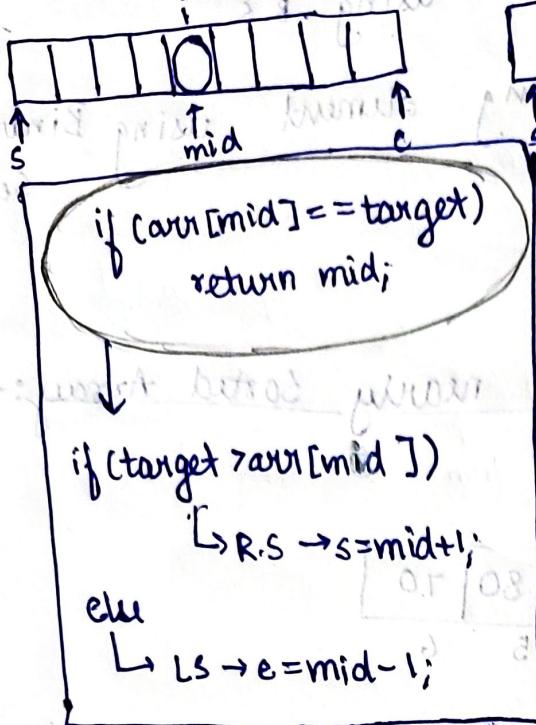
$\downarrow$  bim number

$(topnot = (1 + bim)(topnot) \geq [1 + bim] now)$

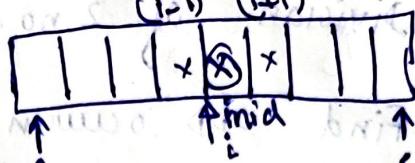
$\downarrow$  bim number

$2.1 \Rightarrow ([bim] now < topnot)$

sorted Array



Nearly Sorted



if (arr[mid] == target)  
    return mid;  
if (arr[mid - 1] == target)  
    return mid - 1;  
if (arr[mid + 1] == target)  
    return mid + 1;  
  
if (target > arr[mid])  
    R.S → s = mid + 2;  
else  
    L.S → e = mid - 2;

CODE:

```
int binarySearch(vector<int> arr, int target){  
    int s=0;  
    int e = arr.size() - 1;  
    int mid = (s + (e - s)) / 2;  
  
    while (s <= e){  
        if (arr[mid] == target)  
            return mid;  
        if (mid - 1 >= 0 && arr[mid - 1] == target)  
            return mid - 1;  
        if (mid + 1 < arr.size() && arr[mid + 1] == target)  
            return mid + 1;  
  
        if (target > arr[mid]){ // R.S  
            s = mid + 2;  
        }  
        else  
            e = mid - 2;  
    }  
}
```

```

        else { // L.S
            e = mid - 2;
        }
        mid = s + (e - s) / 2;
    }
    return n - 1;
}

int main() {
    vector<int> arr{10, 3, 40, 20, 50, 80, 70},
    int target = 40;
    int ans = binarySearch(arr, target);
    cout << "index of " << target << " is " << ans << endl;
    return 0;
}

```

T.C  $\rightarrow O(\log n)$

## (2) Divide two no.s using Binary Search.

dividend = 10

divisor = 2

quotient = ?

$$\begin{array}{r}
 2 \overline{) 10} \quad | 5 \text{ Quotient} \\
 \underline{10} \\
 0 \text{ remainder}
 \end{array}$$

$$\cancel{\text{quotient} * \text{divisor} + \text{remainder}} = \text{dividend}$$

$\text{quotient} * \text{divisor} \leq \text{dividend}$	(CONDITION)
---	-------------

DRY  
Run:-

$$\frac{22}{7} \leftarrow \text{dividend}$$

$$7 \leftarrow \text{divisor}$$



$\text{mid} * \text{divisor} > \text{dividend}$

$$11 * 7 = 77 > 22$$

L.S.

$$0 \quad 5 \quad 10$$

$\uparrow \quad \uparrow \quad \uparrow$

$s \quad \text{mid}, \text{ele}$

101 = right part

mid \* divisor

$$5 * 7 = 35 > 22$$

Ans  $\rightarrow$  store & right search

$$0 \quad 2 \quad 4$$

$\uparrow \quad \uparrow \quad \uparrow$

$e \quad \text{mid}, \text{ele}$

mid \* divisor

$$2 * 7 = 14 \leq 22$$

Ans  $\rightarrow$  store

right search

$$3 \quad 4$$

$$s \quad e$$

$$= 21 \leq 22$$

Ans  $\rightarrow$  store & RS

$$4 * 7 = 28 > 22$$

L.S

$$s = 9 \\ e = 3$$

Ans  $\rightarrow$  store & RS

code:

```
int solve(int dividend, int divisor) {
    int s=0; // part of ans
    int e = dividend; // also dividend
    int ans = 0;
    int mid = s + (e-s)/2;

    while (s<=e) {
        // perfect solution
        if (mid * divisor == dividend) {
            return mid;
        }
        // not perfect answer
        if (mid * divisor > dividend) {
            // left search
            e = mid - 1;
        } else { // ans store
            ans = mid;
            // right search
            mid = s + (e-s)/2;
        }
    }
    return ans;
}

int main() {
    // code
}
```

```
int ans = solve();
cout << "Ans is " << ans << endl;
int precision;
cout << "Enter no. of floating digits: " << endl;
cin >> precision;
double step = 0.1;
```

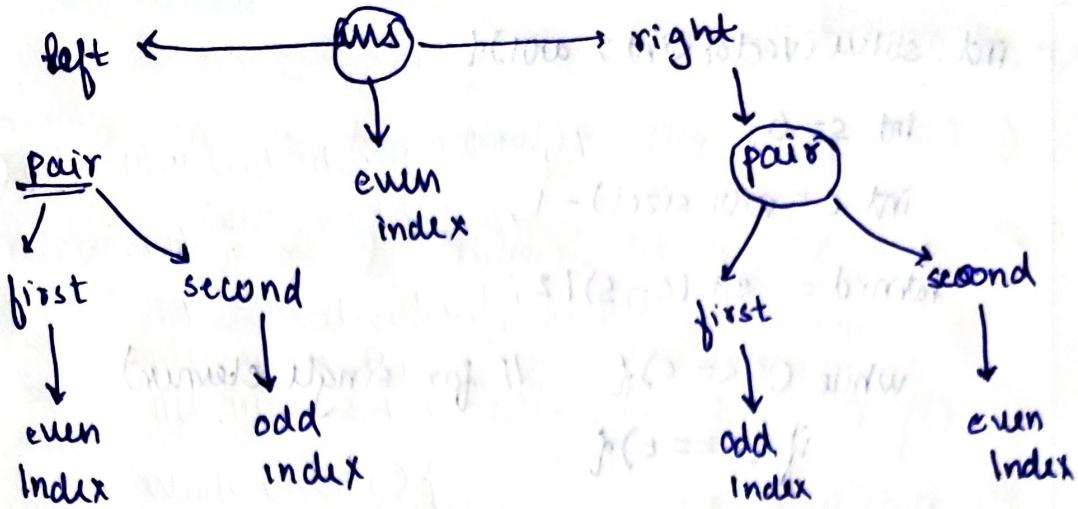
③ Find odd occurring element in array in B.S.,

1	1	2	2	3	3	4	4	3	6	0	0	6	0	4	4
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

odd odd even odd even odd even odd even odd even even  
 $bim = 210$

→ all repeating occurrences of element appear in pairs, are not adjacent. (There cannot be more than two consecutive occurrence of any element).

→ find element that appears odd no. of times.



Conditions:-

$s = 0;$

$c = n - 1;$

$mid = s + (c-s)/2$

while ( $s \leq c$ ) {

    if ( $s == c$ )

        return  $s;$

①

    if ( $mid \% 2 == 0$ ) // even

        // even

        | Go to Left | I am on right side

        if ( $arr[mid] == arr[mid+1]$ )

$\rightarrow R.S \rightarrow s = mid + 2$

        else

$e = mid$

        | Go to R.S | I am on LS

        if ( $arr[mid-1] == arr[mid]$ )

$\rightarrow s = mid + 1;$

$\square ! = \square \rightarrow \text{else}$

        mid-1   mid  
        ↓        ↓  
        even    odd

$(s-odd+2) = bin$

```

int solve(vector<int> arr) {
    int s=0;
    int e=arr.size()-1;
    int mid = s + (e-s)/2;
    while(s <= e){ // for single element
        if(s == e){
            return s;
        }
    }
    // 2 cases → mid even or odd
    if(mid % 2 == 0){
        if(arr[mid] == arr[mid+1]){
            s=mid+2;
        } else {
            e=mid;
        }
    } else {
        if(arr[mid] == arr[mid-1]){
            s=mid+1;
        } else {
            e=mid-1;
        }
    }
    mid = s + (e-s)/2;
}
return -1;
}

```

```

int main() {
    vector<int> arr {1, 1, 2, 2, 3, 3, 4, 4, 3, 600, 600, 4, 4};
    int ans = solve(arr);
    cout << arr[ans] << endl;
    return 0;
}

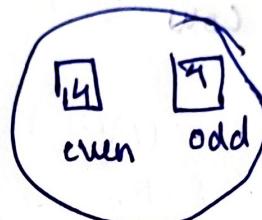
```

Dry Run

1	1	2	2	3	3	4	1	5	4	4	2	2
0	1	2	3	4	5	6	7	8	9	10	11	12

mid

(i)



LS mei kde ho

So, go R.S

5	4	4	2	2
8	9	10	11	12

mid

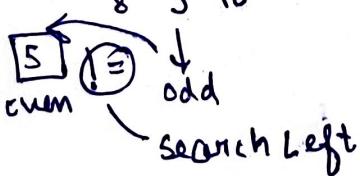
(ii)

5	4	4	2	2
8	9	10	11	12

mid=10

(iii)

5	4	4
8	9	10



$$e = \text{mid} - 1$$

5

### Homework:

- ① Find pairs with difference ' $k$ ' in an array.
- ② Find ' $k$ ' closest element to a given value in an array.
- ③ Exponential search
- ④ Unbounded BS
- ⑤ Advt B.S
  - ↳ Book Allocation
  - ↳ Roti Spoj
  - ↳ Painter Partition
  - ↳ EKO SPOT
  - ↳ Aggressive Cow