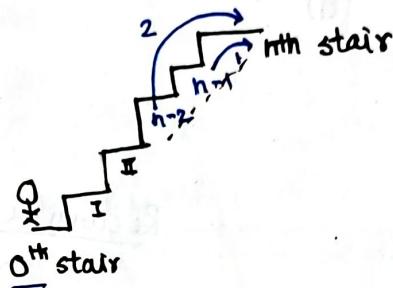


11th March

Week 7 Lee-2

Recursion Level-2

1 Climbing stairs :-



Total ways to reach n^{th}

steps allowed

- 1 stair at a time
- 2 stairs at a time

$$\text{stair} \Rightarrow f(n) = f(n-1) + f(n-2)$$

Find no. of ways to reach n^{th} stair.

```

int climbstairs(int n){
    // base case
    if (n==0) return 1; // 0th stair ki Pahuchne ke tarika
    if (n==1) return 1; // 1st stair ki Pahuchne ke tarika
    int ans = climbstairs(n-1) + climbstairs(n-2);
    return ans;
}
if (n==1) return 1
if (n==2) return 2
    
```

int main(){

int n;

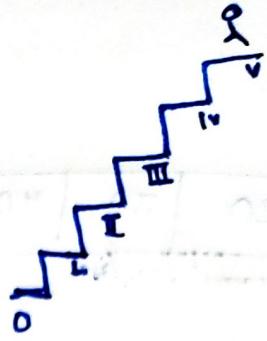
cin >> n;

int ans = climbstairs(n);

cout << ans << endl;

return 0;

$$f(5) = f(4) + f(3)$$



$$f(4) = f(3) + f(2)$$

$$= 3 + 2$$

$$f(3) = f(2) + f(1)$$

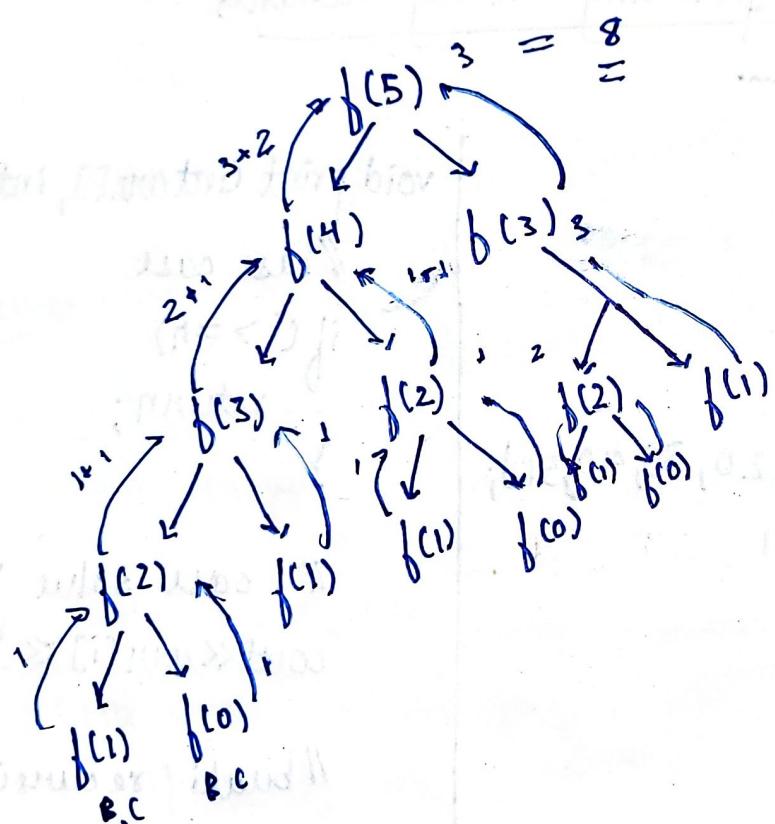
$$= 2 + 1$$

$$f(2) = f(1) + f(0)$$

$$= 1 + 1$$

$$f(1) = 1$$

$$f(0) = 1$$



Array

10 20 30 40 50

```
for (int i=0; i<n; i++) {  
    cout << arr[i];  
}
```

Loop

Using recursion printing Array elements :-

```
#include <iostream>
using namespace std;

int main() {
    int arr[5] = {10, 20, 30, 40, 50};
    int n=5;
    int i=0;
    print(arr, n, i);

    return 0;
}
```

```
void printArr(int arr[], int n, int i)
{
    //base case
    if (i >= n)
        return;
    }

    // i case solve krdia
    cout << arr[i] << " ";
}

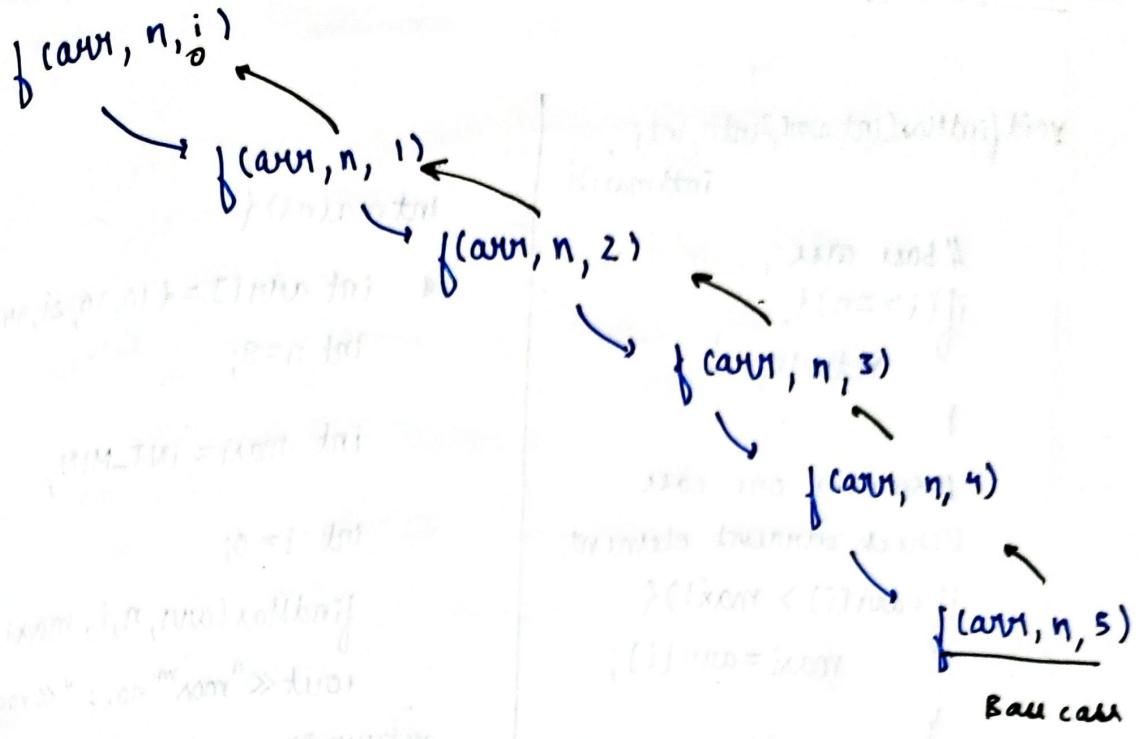
// baaki recursion करेंगा
printArr(arr, n, i+1);
}
```

```

graph TD
    A[f(arr, n, i)] --> B[p = 0]
    B --> C[print arr[0]]
    C --> D[f(arr, n, i + 1)]
    D --> E[print arr[i]]
    E --> F[f(arr, n, i + 1)]
    F --> G[1 + 1 = 2]
    G --> H[p]
  
```

The diagram shows a recursion tree for the `printArr` function. The root node is `f(arr, n, i)`. It branches into two cases: `p=0` which leads to `print arr[0]`, and `f(arr, n, i+1)`. This pattern repeats, with each `f(arr, n, i+1)` node leading to a `print arr[i]` node, which then leads to another `f(arr, n, i+1)` node. The final step is $1 + 1 = 2$, which then leads to the variable `p`.

10	20	30	40	50
0	1	2	3	4



3. Maximum

W1: - Using loops

```

int maxi = INT_MIN
for (int i=0; i<n; i++) {
    if (arr[i] > maxi)
        maxi = arr[i];
}
  
```

OR

```

for (int i=0; i<n; i++) {
    maxi = max(maxi, arr[i]);
}
  
```

1	2	3	4	5	6	7	8
5	3	8	1	7	6	4	2

M2- Recursion

```
void findMax(int arr[], int n, int i,
             int maxi){  
    // base case  
    if (i >= n) {  
        return;  
    }  
    // solving one case  
    // check current element  
    if (arr[i] > maxi) {  
        maxi = arr[i];  
    }  
    // backtracking step  
    findMax(arr, n, i + 1, maxi);  
}
```

O/P - 66

```
int main() {
```

```
    int arr[] = {10, 30, 21, 44, 32, 17, 13, 66};  
    int n = 8;  
    int maxi = INT_MIN;  
    int i = 0;  
    findMax(arr, n, i, maxi);  
    cout << "max no. : " << maxi << endl;  
    return 0;
```

Note - Don't use $i++$ in this. (less confusion)

Dry run :-

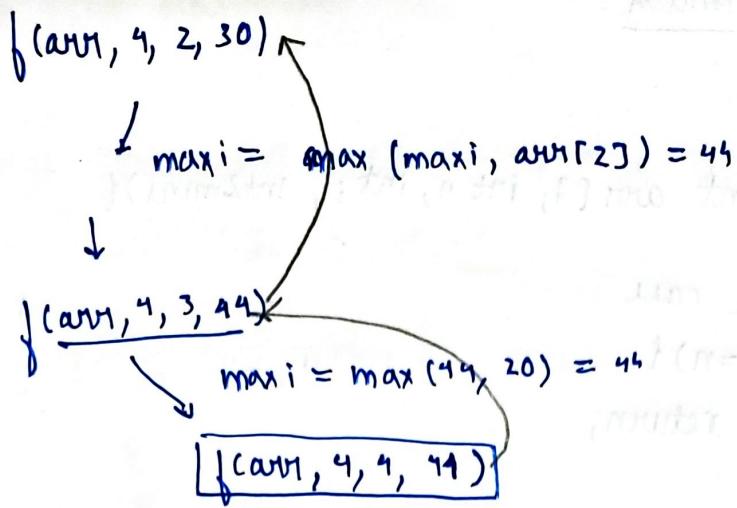
arr,	10	30	41	20
	0	1	2	3

f (arr, 4, 0, INT_MIN)

maxi update

f (arr, 4, 1, 10)

$$\max = \max(\maxi, arr[i])$$



Base case

Dry Run

20	80	10	100
0	1	2	3

arr

$f(\text{arr}, 4, 0, \max_i)$

\downarrow maxi update

$f(\text{arr}, 1, 1, 20)$ \longrightarrow main() - mei phuch k INT_MIN
print hojaega

\downarrow
 $\max_i = \max(\max_i, \text{arr}[i])$
 $= 80$

$f(\text{arr}, 1, 2, 80)$

\downarrow
 $\max_i = \max(\max_i, \text{arr}[z])$

$f(\text{arr}, 1, 3, 80)$

\downarrow
 $\max_i = (\max(\max_i, \text{arr}[3]))$

$f(\text{arr}, 1, 4, 100)$

Base case

4. To find Maximum :-

```
void findMax(int arr[], int n, int i, int& maxi){
```

//base case

```
if(i >= n){
```

```
    return;
```

```
}
```

//i case solve kuna pdega

```
maxi = max(maxi, arr[i]);
```

//baaki recursion samjhali lega

```
findMax(arr, n, i+1, maxi);
```

```
}
```

```
int main(){
```

```
int arr[] = {10, 30, 50, 15, 12, 60};
```

```
int n = 6;
```

```
int maxi = INT_MIN;
```

```
int i = 0;
```

```
findMax(arr, n, i, maxi);
```

```
cout << maxi << endl;
```

```
return 0;
```

```
}
```

3. To find Maximum

```
void findMax(int arr[], int n, int i, int maxi){  
    // base case  
    if (i >= n){  
        return;  
    }  
    // i case solve karna hai  
    maxi = max(maxi, arr[i]);  
    // buaki recursion smhal lega  
    findMax(arr, n, i+1, maxi);  
}
```

}

```
int main(){
```

```
    int arr[] = {10, 20, 90, 70, 20};
```

```
    int n = 5;
```

```
    int maxi = INT_MIN;
```

```
    int i = 0;
```

```
    findMax(arr, n, i, maxi);
```

```
    cout << maxi << endl;
```

```
    return 0;
```

}

Q1P

30

5. To check if in i/p - एक string hai
 string str = "love akshat"
 ex: "love akshat"
- To find: Key → 't'
 t is present in str or not
 (using Recursion)

Ans:

l	o	v	e	a	k	s	h	a	t
0	1	2	3	4	5	6	7	8	9

bool checkKey(string str, int i, int n, char key) {

// base case

if ($i \geq n$) {

// key not found

return false;

}

// I can solve kndo

if ($\text{str}[i] == \text{key}$)

return true;

// baaki recursion samhal lega

return checkKey(str, $i+1$, n, key);

}

Q7.

Note - By reference send kyun se space complexity कम हो जाएगी.

TLE नहीं हो सकता है!

```
int main() {
    string str = "louwakshat";
    int n = str.length();
    char key = 't';
    int i = 0;
    bool ans = checkKey(str, i, n, key);
    cout << ans << endl;
    return 0;
}
```

O/P

1

Now, To return index of key if found, else -1.

- Instead of bool, use int
 - Use -1, where false is ~~returned~~.
 - and 'return i' instead of return true.
 - in main()
 - Nothing would be changed.
-

Making it more correct:-

अगर akshat — a is coming 2 times तो
print both loc"

eg: Found at : 0
 Found at : 4

Code: void checkKey(string& str, int i, int & n, char & key){
 // base case
 if (i >= n){
 // key not found
 return;
 }
 // I case solve kro
 if (str[i] == key){
 cout << "Found at:" << i << endl;
 }
 // backi recur"
 return checkKey(str, i + 1, n, key);
}

```
int main(){
    string str = "lou akshat";
    int n = str.length();
    char key = 'a';
    int i = 0;
    checkKey(str, i, n, key);
    return 0;
}
```

6.

ProblemtatmentIP \rightarrow 647OP \rightarrow print all digit of this no.

void printDigits(int n) {

// base case

if (n == 0) {

return;

}

// I case mai solve karunga

int digit = n % 10;

cout << digit << " ";

int newValueofN = n / 10;

// baaki recursion smhalega })

printDigits(newValueofN);

}

Tail Recursion

int main() {

int n = 647;

printDigits(n);

return 0;

}

OP

7 4 6

For seedha print

Do, head recursion :-

if ($n == 0$) {

 return;

 int newValueOfN = $n / 10$;

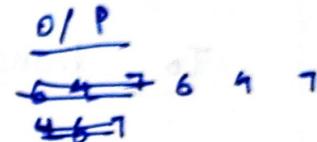
 // baaki recu

 printDigits(newValueOfN);

 // l case mai solve krenge

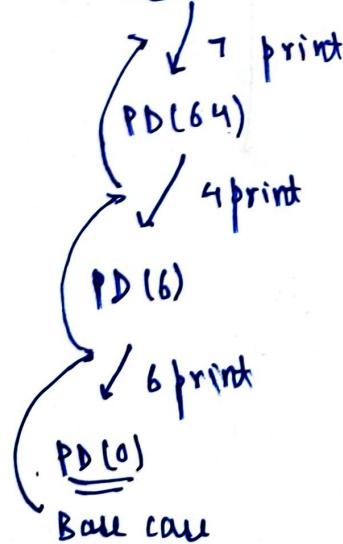
 int digit = $n \% 10$;

 cout << digit << " ";

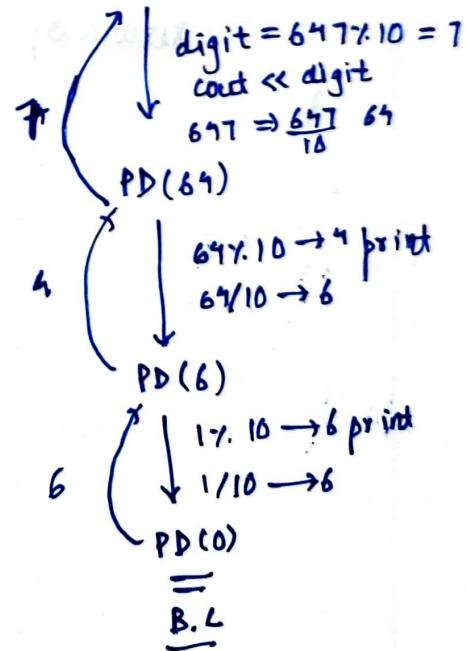


Dry Run

PrintDigits (647)



PD (647)



HW Revise — 5 Eg. Pe Dry Run

(Video ek baar aur
dekh ho)

will fail for '0'

```
int main() {  
    int n = 0;  
    printDigits(n);  
    return 0;  
}
```

O/P

aega hi nhi

To handle this :-

```
int main() {  
    int n = 2000;  
    if (n == 0)  
        cout << 0 << endl;  
    printDigits(n);  
    return 0;  
}
```

2000