Certification of Python Programs on the Basis of Static Information Flow Analysis

A Thesis
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by

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Abstract

In this thesis, we present our work on secure information flow analysis of python programs. We have built a platform that takes source code and labels of all objects used in python program as input for static analysis of information flow throughout the program. We started with Denning's lattice model [1] for verification of secure information flow. In this model, every object is associated with its security class. To prevent unauthorized leak of information, the flow of information should be in one way – from less secure to more secure class. A lattice represents such information model very well, the upward direction in lattice represents secure information flow. Verification of information flow only on the basis of security class is not sufficient to certify the security of system, there is a need to consider the process, user or subject that executes the code. Use of Reader Writer Flow model [1] with subjects makes it possible to do secure information flow analysis. We have developed four type of constraint generators C1, C2, C3 and C4 each implementing different approach. Constraint generator C1 considers fixed label and PC reset (PC reset denotes PC is not retaining information out of scope), C2 considers fixed label and monotonic PC (monotonic PC never lose information), C3 considers dynamic label and PC reset and finally the constraint generator C4 considers dynamic label and monotonic PC. Constraint resolver takes these constraints and RWFM labels [2] for each object defined by the user as input and provides answers to various queries related to information flow security. The report describes, the approach, implementation and case study done so far.

Contents

1	Intr	oduction	1
	1.1	Motivation	3
	1.2	Goals	3
	1.3	Contributions	4
		1.3.1 Implementation Details	4
2	Rela	ated work	7
3	Info	rmation Flow Analysis using Denning's Lattice Model	8
	3.1	Analysis In Local Scope	8
		3.1.1 Benchmarking of Certification Script using Denning's Examples [1]	9
	3.2	Further analysis: considering global influence of while	10
		3.2.1 Benchmarking of Certification Script using Denning's Examples [1]	11
4	Ana	lysis of Multi-threaded Programs	13
	4.1	Handling Information flow due to WAIT and SIGNAL operations	13
		4.1.1 Benchmarking of Certification Script using Denning's Example [1]	14
5	Cat	regory 1 constraint Generator: PC reset and fixed labels.	16
	5.1	Working	16
	5.2	Constraint Rules	17
	5.3	Key Idea	17
	5.4	Limitations	18
6	Cate	egory 2 constraints : PC monotonic and fixed labels.	20
	6.1	Working	20
	6.2	Key Idea	20
	6.3	Limitations	20

ii Contents

7	Cate	gory 3 constraints : PC reset and dynamic labels.	22
	7.1	Constraint Rules	22
	7.2	Key Idea	24
	7.3	Limitations	24
8	Cate	gory 4 constraints: PC monotonic and dynamic label.	25
	8.1	Key Idea	25
	8.2	Limitations	26
9	Com	parison among all categories of constraints.	27
10	Cons	straint Verification with RWFM Labels	31
	10.1	RWFM label [2]	31
		10.1.1 Secure Information Flow	31
		10.1.2 Definition of Join and Meet operations [2]	31
		10.1.3 Create, Read and Write operations on objects with RWFM label [2]	32
	10.2	Constraint Verification	32
11	Con	clusion & Future work	35
12	Pyth	on Syntax	36
Ap	pend	ices	38
A	Pyth	on Script category 1: Constraint Generator	40
В	Pyth	on Script category 2: Constraint Generator	57
C	Pyth	on Script category 3: Constraint Generator	74
D	Pyth	on Script category 4: Constraint Generator	93
E	Pyth	on Script 2: Constraint Checker	111
F	Cop	y Programs	115

List of Figures

1.1	Lattice of Listing 1.5	3
1.2	Block diagram of script1	4
1.3	Block diagram of parsing function	5
3.1	Control flow graph of If else and while loop	8
5.1	Example for PC reset	17
5.2	Example for PC monotonic	18
9.1	Category of constraint generator	29
9.2	Set diagram	29
9.3	Performance test	30
10.1	Block diagram of script2	32
10.2	Lattice for information flow in $a = x$ and $y = b$, (a) without subject (b) with subject.	33

Listings

1.1	Python example	1
1.2	Python example	1
1.3	Python example	2
1.4	Python example	2
1.5	Python example	2
3.1	Python version of copy1 example in [1]. goal: information flow from x to y	9
3.2	Python version of copy2 example in [1]. goal: information flow from x to y	9
3.3	Non terminating while	10
3.4	Python version of copy5 example in [1]. goal: information flow from x to y	11
3.5	Python version of copy4 example in [1]. goal: information flow from x to y	11
4.1	Example of wait() operation on binary semaphore. Info Flow: $s \rightarrow x$	13
4.2	Infinite while loop, Information Flow: $x \rightarrow y$	13
4.3	Python version of copy3 example in [1]. goal: information flow from x to y	14
5.1	Python example	17
5.2	Python version of copy5 example in [1]. goal: information flow from x to y	18
5.3	Python version of dynamic label example in [1]. goal: information flow from x	
	to y	19
6.1	Python version of dynamic label example in [1]. goal: information flow from x	
	to y	20
12.1	if else, indentation is must	36
12.2	while loop	36
12.3	threads	36
12.4	threads	36
p1.py	7	40
p2.py	,	57
р3.ру	7	74
p4.py	7	93

Listings		7

constraint_checker.py	 														111
programs.py	 														115

Introduction

Data Security can be achieved using information flow policies. Execution of any action like copying of file, read operation on file or memory, write operation on file or memory, execution of statement etc may cause flow of information. Usually, information revealed that was unknown before execution of statement termed as information flow otherwise if it was known already then there will be no information flow. This thesis focuses on information flow between variables used in a python program irrespective of the amount of information flowing. There are two kinds of information flow among variables:

1. Explicit Information Flow

```
x = y
x = math.log(y)
```

Listing 1.1: Python example

these assignment operations are example of explicit flow because information related to variable y is flowing into x using data dependency in both cases.

2. Implicit Information Flow

Listing 1.2: Python example

2 Introduction

```
while y < z:
x = 1
y += 1
```

Listing 1.3: Python example

in these examples value of x after execution of statements depends on the control path taken by the program, variable y and z used in choosing between two control path in while program, this involvement of y and z in making decision reduce the uncertainty of variables y and z. Whether y < z or y > z can be observed with help of final value of x after execution of given code in listing 1.3. So there is indirect implicit flow from y to x in the first example and implicit flow from y and z to x in the second example[1].

```
x = 0
while y == 1:
pass
x = 1
```

Listing 1.4: Python example

Here assignment operation on x in Listing 1.4 is outside of while body but still information flowing from y to x because execution of x = 1 statement depends on termination of while loop, so you can know value of y by checking the value of x after execution of program, for example if value of x is 1 and program terminated then y must be other than 1, if while loop goes in infinite loop then y must be 1.

The first chapter describes python program certification with the help of Denning's Lattice Model. The lowest security class in this lattice assumed is Low everyone can read information from this class, the highest class in the lattice is High, information from all security classes can flow into it but no information can flow from this class to others. For Example certification of python code in Listing 1.5 needs to follow Denning's lattice and constraints written in figure 1.1, security class of variable var is denoted by <u>var</u>.

Listing 1.5: Python example

1.1 Motivation 3

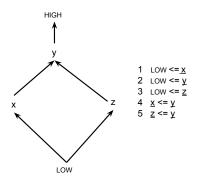


Figure 1.1: Lattice of Listing 1.5

Chapter 3 describes information flow by nonterminating while loop and how to certify program in such situation. Chapter 12 describes how information flows among thread in the multi-threaded program using semaphore. Chapter 10 presents a new approach to certifying a program using more sophisticated labels (RWFM label) and it also considers subjects for certification of the program.

1.1 Motivation

In the field of data security, there are a lot of approaches to prevent leak of information for example cryptography takes care of confidentiality and integrity while data is transmitting through less secure networks, access permission on files prevent unauthorized access to files in a system where users have different privileges. But at the time of execution of program, data used in the program is vulnerable to various attacks so to maintain security at the time of execution of program and processing of data, information flow policies are used. The subject is defined as an executing authority it can be a user or parent process, object can be a file, program variable, memory location etc. In a multilevel security system a subject has permission related to objects. Information flow verification of program only considering objects may seem to be secure but with a particular subject same program may be insecure, so we considered subjects in static analysis of python program.

1.2 Goals

To develop a platform that takes input a python program and labels of each variable used in the program, and provide answers to various queries regarding the security of information flow. 4 Introduction

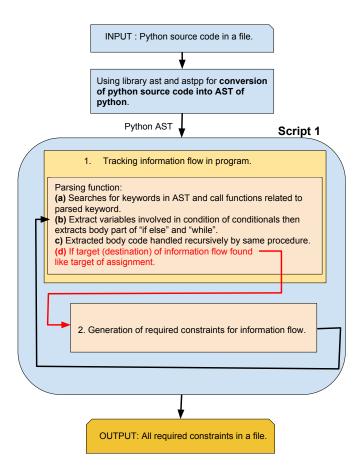


Figure 1.2: Block diagram of script1

1.3 Contributions

Subject considered with objects for certification of python program. Reader Writer Flow Model [2] used to verify information flows in python program.

1.3.1 Implementation Details

Prerequisite Third Party libraries:

- 1. ast python library (for conversion of python source code into abstract syntax tree)
- 2. astpp python library (for readability of of abstract syntax tree of python source code)

Subset of features of python language considered for analysis.

- Assignment operations : x = e (expression)
- Conditional statements: "if else", "elif".
- Iteration: "while".
- Semaphore operations : set(), wait(), clear(), initialization of semaphore.
- Global variables and local variable in a function.

1.3 Contributions 5

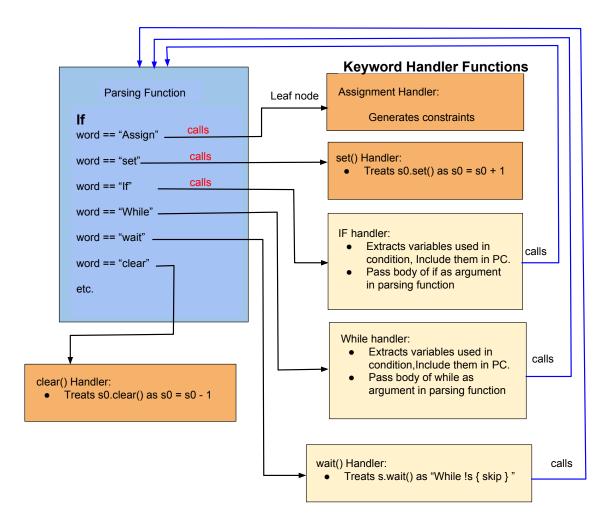


Figure 1.3: Block diagram of parsing function

- Function calls and definitions.
- Return Statement.

I implemented fully automated certification platform for python source code using two python scripts, given in Appendix ?? and E. Block diagram in figure 1.2 shows modules of script1, first python source code is converted into abstract syntax tree (AST) with the help of ast library. The purpose of this step is to avoid tedious work of parsing of source code and comments. Figure 1.3 shows that parsing function reads AST word by word. If function finds any desired word it calls other handler functions to handle code related to particular word, for example: if "While" word is found, parsing function parses body of while and passes it as argument to while_handler(while_code) function. Handler function parses variables used in condition and passes the body part to parsing function again. Whenever parsing function finds assignment operation it generates constraint and goes to next word. The Block diagram for script 2 is given in figure 10.1 of chapter 10. Effectiveness of this platform depends on the phase of constraint

6 Introduction

generation because if constraint generator fails to track information flow properly then verifier can not produce correct results. Chapter 5 to 8 describe about different constraint generators. Chapter 9 shows the comparison among all constraint generator.

Related work

There have been many studies on information flow control and all of them share some basic properties like information flow should be from less secure entity to more secure entity. Denning's book [1] has a chapter on information flow control, this chapter describes lattice model for information flow [3], this makes it easy to track information flow in a program using transitivity property. Analysis has been done on basic operations which involve information flow like assignment operation (explicit flow) based on data flow, conditional operations like if else, while etc. (implicit flows) based on control flow, information flow through covert channel based on traps and exception in programs. Here are some basic rules given in [1], (arrow \rightarrow denotes information flow).

```
• x = y : y \rightarrow x
```

- if e then $x = y : e \rightarrow x$
- while w if e then x = y w = false : $w \oplus e \rightarrow x$
- infinite loop: while w; $x = y :- w \rightarrow x$ etc.

Chen et al. [4](published in 2014) presented work on python byte-code and claimed that there was no work related to python at that time. They implemented information flow checker for python byte-code using static and dynamic analysis but their main focus is on information flow policies related to objects. Kumar et al. [2] introduce a new model to work with subjects and my work will be focused on this. Conti et al. [5] provide library support in python for information flow analysis in explicit flows only.

Information Flow Analysis using Denning's Lattice Model

3.1 Analysis In Local Scope

We started certification of python code at local scope initially. In figure 3.1 we considered

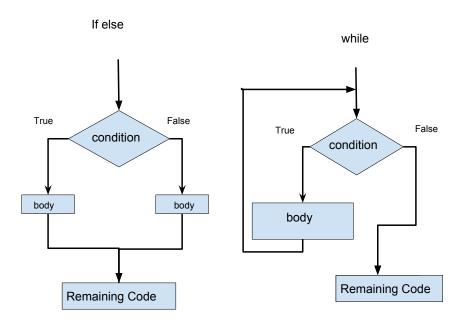


Figure 3.1: Control flow graph of If else and while loop

information flowing from variables used in condition to targets of assignment operation in body code only. Further analysis will be dealing with information flows to targets of assignment in

remaining code. Now a program with if else and terminating loops can be certified to be secure (no information leaks) by using basic rules of information theory implemented in my script. Fifth chapter of Denning's book[1] used six examples to describe how can information flow through different ways in the program. We used python version of these examples as benchmark.

3.1.1 Benchmarking of Certification Script using Denning's Examples [1]

```
#Procedure copy1

2 x = 0

3 z = 0

4 y = 0

5 if x == 0:

6 z=1

7 if z == 0:

8 y=1
```

Listing 3.1: Python version of copy1 example in [1]. goal: information flow from x to y

Constraints for program in 3.1 generated by our script are:

```
1. low \le \underline{x}
2. low \le z
```

3. $low \le y$

4. $low \le y$

5. $\underline{x} \leq \underline{z}$

6. $\underline{z} \leq \underline{y}$

Program in listing 3.1 shows how information can flow indirectly from x to y, for secure information flow $\underline{x} \le \underline{y}$ must be true. Constraint 5 (x \le z) and constraint 6 (z \le y) with transitivity property of lattice model satisfies x \le y.

```
# Procedure copy2

2 x = 0

3 z = 1

4 y = -1

5 while z == 1:

6 y = y + 1

7 if y == 0:

8 z = x

9 else:
```

```
z = 0
```

Listing 3.2: Python version of copy2 example in [1]. goal: information flow from x to y Constraints of program in Listing 3.2 generated by our script are:

- 1. $low \leq \underline{x}$
- 2. $low \leq \underline{z}$
- 3. $low \le y$
- $4. \ \underline{y} \oplus \underline{z} \leqslant \underline{y}$
- 5. $y \oplus \underline{x} \oplus \underline{z} \leq \underline{z}$

Information flows from x to y using iteration of the while loop in listing 3.2. One iteration of while loop changes the value of y to -1 to 0 by incrementing it by 1, and two iteration of same while loop changes the value of y to -1 to 1. Number of iterations of while loop depends on value of x so this program is transmitting information from x to y indirectly using both explicit and implicit information flows.

Proof:

 $\underline{x} \leq \underline{y}$ must be true for secure information flow. This can be proved using constraint 4 ($\underline{y} \oplus \underline{z} \leq \underline{y}$) and constraint 5 ($\underline{y} \oplus \underline{x} \oplus \underline{z} \leq \underline{z}$) in two ways. First: $\underline{z} \leq \underline{y}$ is given in constraint 4 and $\underline{y} \leq \underline{z}$ is given in constraint 5 so \underline{y} must be equal to \underline{z} using this new constraint ($\underline{y} \equiv \underline{z}$) and reduced constraint 5 ($\underline{x} \leq \underline{z}$) desired constraint $\underline{x} \leq y$ is proved.

3.2 Further analysis: considering global influence of while

A program either terminates after a finite period of time or executes for an infinite period of time. The latter can be achieved through while loop and for loop. Such a non-terminating loop influences variables within the body of the loop as well as after the body of loop.

```
x = 0
z = True
y = 0
while z :
y = 5
x = 5
```

Listing 3.3: Non terminating while.

In listing 3.3 execution of last statement x = 5 is conditioned on "while z" statement similar to "if" branching but "if" has limited body but in this case every statement that comes after "while

z" share fate with x = 5, because of this property of "while" it is very crucial to know whether loop is terminating or nonterminating, so either there should be a mechanism to determine to terminating and nonterminating nature of loop or treat all while loop as nonterminating, latter approach is imprecise but secure, for now we are using this approach to handle while loops.

3.2.1 Benchmarking of Certification Script using Denning's Examples [1]

```
#Procedure copy5
y = 0
while x==0:
pass
y = 1
```

Listing 3.4: Python version of copy5 example in [1]. goal: information flow from x to y

Constraints generated by our script:

- 1. $low \le y$
- 2. $\underline{x} \leq y$

so here our script able to track information flow from x to y and generating constraint $\underline{x} \leq \underline{y}$ for verification.

```
#Procedure copy4
2 import thread
3 import time
4 import threading
6 def thread1():
7 global x
8 global e0
9 global e1
10 if x = = 0:
    e0 = False
12 else:
    e1 = False
def thread2():
16 global e0
17 global e1
18 global y
19 while e0
    pass
y = 1
e1 = False
```

```
def thread3():
global e1
global e0
global y
while e1:
pass
y = 0
e0 = False

thread.start_new_thread(thread1,())
thread.start_new_thread(thread2,())
thread.start_new_thread(thread3,())
```

Listing 3.5: Python version of copy4 example in [1]. goal: information flow from x to y

Constraints generated by our script for program in listing 3.5 are:

```
1. \underline{x} \leq \underline{e0}

2. \underline{x} \leq \underline{e1}

3. \underline{e0} \leq \underline{y}

4. \underline{e0} \leq \underline{e1}

5. \underline{e1} \leq \underline{y}

6. \underline{e1} \leq \underline{e0}

(\underline{x} \leq \underline{e0}) and (\underline{e0} \leq \underline{y}) \equiv \underline{x} \leq \underline{y} (transitivity)

(\underline{x} \leq \underline{e1}) and (\underline{e1} \leq \underline{y}) \equiv \underline{x} \leq \underline{y}
```

So script able to track hidden information flow from x to y and generating constraint for this flow $\underline{x} \le y$ by using transitivity.

Analysis of Multi-threaded Programs

In a multi-threaded program, information flows among threads because of communication and synchronization among them. There are two types of semaphores — counting and binary, for synchronization among threads. For now, our script handles binary semaphores only.

4.1 Handling Information flow due to WAIT and SIGNAL operations

Traditional operations related to binary semaphore are WAIT and SIGNAL. SIGNAL operation changes the value of semaphore 0 to 1 and WAIT operation wait for an infinite time if the current value of the semaphore is 0 otherwise it changes value 1 to 0 and allows control flow forward. There are three operations related to the binary semaphore in python language wait(), set() and clear(). Traditional WAIT operation can be simulated using python wait() followed by clear() operation, SIGNAL is equivalent to set().

```
s = threading.Event()
s.wait()
x = 1
y = 0
while x:
pass
y = 1
y = 0
y = 1
```

Listing 4.1: Example of wait() operation Li on binary semaphore. Info Flow: $s \rightarrow x$ In

Listing 4.2: Infinite while loop, Information Flow: $x \rightarrow y$

Listing 4.1 and Listing 4.2 show that control flow of wait() is similar to infinite while loop so we treat wait() in a similar way. All statements which use global variables as a target of assignment and are preceded by wait() may transmit information to other threads. So information flows

from semaphore s_0 to targets of assignment operations which follows s_0 .wait() statement. All semaphore operations simplified into normal operations.

- s.set() treated as s = s + 1.
- s.clear() treated as s = s 1.
- Listing 4.1 and 4.2 shows s.wait() equivalent to while(s == 0) { skip }.

4.1.1 Benchmarking of Certification Script using Denning's Example [1]

```
1 #Procedure copy3
2 import thread
3 import time
4 import threading
s s0 = threading.Event()
6 s1 = threading. Event()
8 def thread1():
    global x
    if x==0:
      s0. set()
    else:
      s1. set()
13
def thread2():
    global y
    s0. wait()
    s0.clear()
    y=1
19
    s1. set()
21
def thread3():
    global y
    s1. wait()
24
    s1.clear()
    y=0
26
    s0. set()
28
29 thread.start_new_thread(thread1,())
thread . start_new_thread (thread2 ,())
thread.start_new_thread(thread3,())
```

Listing 4.3: Python version of copy3 example in [1]. goal: information flow from x to y

To certify the multi-threaded program in Listing 4.3 correctly our script must track information flow from x to y $(x \rightarrow y)$ and must generate constraints accordingly.

Constraints generated by our script for program in Listing 4.3 are:

- 1. $\underline{\mathbf{x}} \oplus \underline{\mathbf{s0}} \leq \underline{\mathbf{s0}}$
- 2. $\underline{\mathbf{x}} \oplus \underline{\mathbf{s}} \mathbf{1} \leq \underline{\mathbf{s}} \mathbf{1}$
- 3. $\underline{s0} \leq \underline{s0}$
- 4. $\underline{s0} \le y$
- 5. $\underline{s1} \oplus \underline{s0} \leq \underline{s1}$
- 6. $s1 \le s1$
- 7. $\underline{s1} \leq y$
- 8. $\underline{s1} \oplus \underline{s0} \leq \underline{s0}$

```
constraint 1 \ (\underline{x} \oplus \underline{s0} \leq \underline{s0}) and constraint 4 \ (\underline{s0} \leq \underline{y}) \equiv \underline{x} \leq \underline{y}.

constraint 2 \ (\underline{x} \oplus \underline{s1} \leq \underline{s1}) and constraint 7 \ (\underline{s1} \leq \underline{y}) \equiv \underline{x} \leq \underline{y}.
```

Hence script is able to generate correct constraints in multi-threaded program too.

Category 1 constraint Generator : PC reset and fixed labels.

This chapter describes working of first algorithm for constraint generation. All four algorithms are different because of different combination of PC label management scheme and use of dynamic labels. In this algorithm we are using fixed label. Fixed label denotes that label will not change throughout the program. PC reset denotes that after completion of scope of a particular conditional/iteration body PC reset back to the PC just before the execution of conditional/iteration statement. PC monotonic denotes a scheme in which PC label never lose any information once it acquires, it just grows monotonically. This algorithm is able to capture basic information flows within a program, so this algorithm will certify a large number of programs as secure. Program certified secure by this algorithm does not mean that its fully secure, it certifies secure because of the limitation in detection of information flows in program. Some information flows which are not captured by this algorithm may violate information security.

5.1 Working

All four algorithm shares same basic structure for parsing input program. Dynamic analysis can not process all branches in a one go but static analysis process all branches in one run. Algorithms generates constraints for all possible control branches in program. PC keeps track of variables used in conditional statement and iteration statement. Assignment operation are responsible for information flows so at each occurrence of assignment operation constraints generated with the help of PC label.

5.2 Constraint Rules 17

```
1.
     x = 1
                PC{Low}
                                      1.
                                           z = 1
                                                           PC{Low}
                                      2.
2.
     y = 0
                PC{Low}
                                           y = -1
                                                           PC{Low}
                             reset
                                                                         reset
                                      3.
     if x == 0: PC\{x\}
                                           while z == 1: PC\{z\}
        y = 0 PC\{x\}
                                      4.
4.
                                               y = y + 1 PC\{z\} \sim
                                                                         reset
5.
               PC{Low}
                                      5.
                                              if y == 0:
    z = 1
                                                           PC\{z,y\}
                                      6.
                                                    z = x PC\{z,y\}
                                      7.
                                              else:
                                      8.
                                                    z = 0 PC\{z,y\}
                                      9.
                                              a = 1
                                           b = 1
                                    10.
                                                            PC{Low}
```

PC reset

Figure 5.1: Example for PC reset

5.2 Constraint Rules

- 1. $\langle x := e \rangle$ generate constraint $[\lambda(e) \oplus \lambda(PC) \leq \lambda(x)]$ and update PC label $\lambda(PC) = \lambda(e) \oplus \lambda(PC)$
- 2. $\langle x := e \rangle$ generate constraint $[\lambda(e) \oplus \lambda(PC) \leq \lambda(x)]$ and update PC label $\lambda(PC) = \lambda(e) \oplus \lambda(PC)$
- 3. < if e then c1 else c2> $\forall x \in (modified_global(\ c1\ and\ c2) \cup \{PC\})$ generate constraints $[\lambda(e) \oplus \lambda(PC) \leq \lambda(x)]$ and update PC label $\lambda(PC) = \lambda(e) \oplus \lambda(PC)$
- 4. < while e do c > $\forall x \in (modified_global(c) \cup \{PC\})$ generate constraints $[\lambda(e) \oplus \lambda(PC) \le \lambda(x)]$ and update PC label $\lambda(PC) = \lambda(e) \oplus \lambda(PC)$

5.3 Key Idea

Direct information flow happens because of copying or assigning values. Implicit information flow happens because of control dependency, this algorithm focuses on information flow from variables used in condition statement of if else(conditional) and while(iteration) to all variables modified in the body of iteration or conditional.

```
def(x,y): #copy x to y
y = 0
z = 0
if x == 0: # implicit flow x -> z PC{x}
z = 1
if z == 0: # implicit flow z-> y PC{z}
y = 1
```

```
PC{Low}
                            1.
                                z = 1
1.
    x = 1
            PC{Low}
                                y = -1 PC{Low}
   y = 0 PC\{Low\}
                            2.
2.
                                while z == 1: PC\{z\}
                            3.
    if x == 0: PC\{x\}
3.
                                    y = y + 1 PC\{z,y\}
                            4.
       y = 0 PC\{x\}
4.
                                   if y == 0: PC\{z,y\}
                            5.
    z = 1
5.
                                        z = x PC\{z,y,x\}
                            6.
                            7.
                                   else:
                                        z = 0 PC\{z,y,x\}
                            8.
                                          PC\{z,y,x\}
                            9.
                                   a = 1
                           10.
                                b = 1
```

PC monotonic

Figure 5.2: Example for PC monotonic

```
p = q # direct flow q \rightarrow p PC\{q\}
```

Listing 5.1: Python example

Constraints generated by this algorithm for listing 5.1 are given below.

- x <= z</p>
- z <= y

5.4 Limitations

For listing 5.1 this algorithm is able to capture all information flows, but in more complex programs it may declare falsely a program secure.

```
# Procedure copy5
2 y = 0
3 while x==0:
4 pass
5 y = 1
```

Listing 5.2: Python version of copy5 example in [1]. goal: information flow from x to y

For listing 5.2 this algorithm generated only one constraint Low <= y, that shows this algorithm will certify listing 5.2 secure always irrespective of information flow x to y is secure or not. All these limitation in capturing information flow raised because of PC label management in this algorithm is only focus on local information flow. Next algorithm will try to remove these

5.4 Limitations 19

limitation using monotonic PC label management. Appendix A shows the implementation of this algorithm.

```
def fun(x, y, z):
    a = x
    y = a
    a = z
    fun(x, y, z)
```

Listing 5.3: Python version of dynamic label example in [1]. goal: information flow from x to y

Another limitation of this algorithm is related to use of fixed label. In listing 5.3 this algorithm detect false information flow z to y. Category 3 uses both dynamic and fixed label to remove this limitation.

Category 2 constraints: PC monotonic and fixed labels.

This chapter describes working of category 2 constraint generator. This algorithm is a improved version of previous category 1 algorithm. This algorithm using monotonic PC label instead of PC reset. By using monotonic PC this algorithm is able to detect additional information flows in program.

6.1 Working

rules

6.2 Key Idea

This analysis extension of previous algorithm, non terminating loops create a control dependency between variables used in condition of loop and the rest of the code where control can go subsequently on termination of loop, because of this behavior of non terminating loop PC storing all the dependencies.

6.3 Limitations

In static analysis if constraint resolver ignores the order of generated constraints then it may show some additional false information flow in program, this is responsible for overhead and imprecision in certification process. Use of dynamic label solves this problem easily on the cost of more complex analysis.

```
def fun(x, y, z):

2 a = x

3 y = a

4 a = z
```

6.3 Limitations 21

fun(x, y, z)

Listing 6.1: Python version of dynamic label example in [1]. goal: information flow from x to y

constraints generated for listing 6.1 are given below:

- x <= a
- $a + x \le y$
- $a + x + z \le a$

these constraints shows that there is information flow z to y (using a $\leq y$ and $z \leq a$) but in program there is no such flow exist. Because of such information flow constraint resolver may certify a secure program as not secure and it also create extra overhead on resolver. This example shows flaw in approach of using fixed label everywhere. Next algorithm will try to remove this limitation.

Category 3 constraints: PC reset and dynamic labels.

This chapter describes category 3 constraint generator. This algorithm introduces use of dynamic label. Global variable are using fixed label and all local variables are assigned dynamic label. Information can flow outside only because of modification of global variable, modification of local variable doe not cause information because information remains in program itself.

7.1 Constraint Rules

- $< x := e > generate constraint [\lambda(e) \oplus \lambda(PC) \le \lambda(x)]$ and update PC label $\lambda(PC) = \lambda(e) \oplus \lambda(PC)$
- < if e then c1 else c2>
 - 1. $\forall x \in (modified_global(\ c1\ and\ c2) \cup \{PC\})$ generate constraints $[\lambda(e) \oplus \lambda(PC) \le \lambda(x)]$ and update PC label $\lambda(PC) = \lambda(e) \oplus \lambda(PC)$
 - 2. $\forall x \in (modified_local(c1 \ and c2) \cup \{PC\}) \text{ update PC label } \lambda(PC) = \lambda(e) \oplus \lambda(PC)$
- < while e do c >
 - 1. $\forall x \in (modified_global(c) \cup \{PC\})$ generate constraints $[\lambda(e) \oplus \lambda(PC) \leq \lambda(x)]$ and update PC label $\lambda(PC) = \lambda(e) \oplus \lambda(PC)$
 - 2. $\forall x \in (modified_local(c) \cup \{PC\})$ update PC label $\lambda(PC) = \lambda(e) \oplus \lambda(PC)$

In static analysis if constraint resolver ignores the order of generated constraints then it may show some additional false information flow in program, this is responsible for overhead and imprecision in certification process. Use of dynamic label solves this problem easily on the cost of more complex analysis.

^{&#}x27;a' is a local variable in function defined below.

7.1 Constraint Rules 23

def function(x,y,z):

a = x

y = a

a = z

static analysis will generate constraints 1. $x \le a$, 2. $a \le y$, 3. $z \le a$.

Last two constraints shows false information flow from z to y $(z\rightarrow y)$.

Dynamic label analysis

- 1. $\lambda(a) := \mathbf{x}$
- 2. $y \le \lambda(a)\{x\}$
- 3. $\lambda(a) := z$

This analysis treats global and local variable differently so it avoids false constraints successfully without tracking order of constraints explicitly.

a = x

while 1:

y = a

a = z

Dynamic label Analysis:

First iteration of while:

- 1. $\lambda(a) := \mathbf{x}$
- 2. $y \le \lambda(a)\{x\}$
- 3. $\lambda(a) := z$

Second Iteration:

- 2. $y \le \lambda(a)\{z\}$
- 3. $\lambda(a) := z$

Dynamic label analysis generating different constraints for first iteration and second iteration but static analysis is not able to distinguish between information flow in first iteration and second iteration of while loop.

7.2 Key Idea

Definition of information flow among objects: If any data can be guessed by using given objects which was unknown previously, by using this idea information can flow outside only because of modification of global objects, so any information flow from local objects to global objects must be checked for security breach. Local variable plays a role of temporary in flow of information from one global to another, so local variable must keep track of information they hold, dynamic label is a good technique to keep track of history of information stored in a local variable.

7.3 Limitations

This constraint generator again using PC reset label scheme. We created this category for thorough analysis and comparison among all categories. So it shares the first limitation of category 1. It fails to capture global information flows created by non terminating loops.

Category 4 constraints: PC monotonic and dynamic label.

This algorithm is best among all four algorithm in terms information security. Dynamic label processing increases time complexity of this algorithm but we used a property of PC label to make optimization. Constraint generation rules are same as category 3.

8.1 Key Idea

In this analysis PC never gets reset because we want to track all possible information flows including information flows from a nonterminating loop to rest of code.

```
'a' is a local var

a = x

while w:

y = a

a = z

z = y
```

Dynamic label Analysis:

```
1. \lambda(a) = x

2.PC{w}

3.PC{w, \lambda(a){x}} w \oplus x \le y

4.\lambda(a) = z

5.PC{w, x, y} w \oplus x \oplus y \le z
```

So this algorithm is capable to track information flow $w\rightarrow z$ in last statement z=y by using

monotonic PC without generating additional false constraints.

8.2 Limitations

Use of dynamic label with monotonic creates challenge for processing of large number of labels.

Comparison among all categories of constraints.

Example given in table 9.1 is suitable to differentiate between all category. First algorithm fails to track information flow $w\rightarrow z$ in last statement z=y. Second algorithm is able to track information flow $w\rightarrow z$ in last statement z=y but it will show additional false information flow $z\rightarrow y$ too. Third algorithm avoids tracking of additional false information flow $z\rightarrow y$ but it fails to show information flow $w\rightarrow z$ because of PC reset. Fourth analysis avoids tracking of false information flow as well as tracks information flow caused by nonterminating loop($w\rightarrow z$). Table ?? shows constraints generated for copy program given in Denning,s book by all constraint generator. Figure 9.2 shows the relationship between set of programs declared secure by all constraint generator. C1-C4 are abbreviation for category 1 - category 4. Number of constraints is inversely proportional to size of set of program declared secure, because more constraints means high probability of violation of security. In category 1 generator generates many false constraints because of absence of dynamic label, but category 3 uses dynamic labels with fixed so it reduces number of constraints. Constraints generated by category 1 are superset

Input Program	C1	C2	C3	C4
#'a' is a local var	1	x <= a		
a = x	x <= a	a + x + w <= y	$x + w \le y$	x + w <= y
while w:	a + w <= y	$a + x + z + w \ll a$	$x + w \le y$ $x + z + w \le y$	x + z + w <= y
y = a	$z + w \le a$	$a + x + z + w \le y$	$y \le z$	y + x + w <= z
a = z	y <= z	$a + x + z + y + w \le z$	y \- L	$y + x + z + w \ll z$
z = y		$y + x + w \ll z$		

Table 9.1: Example for comparison

Programs	C1	C2	C3	C4
Copy1	x <= z z <= y	$ \begin{array}{cccccccccccccccccccccccccccccccccccc$	x <= y	x <= y
Copy2	Low <= z Low <= y y + z <= y y + x + z <= z y + z <= z	Low <= z Low <= y y + z <= y y + x + z <= z y + z <= z y + z <= z y + x + z <= y	Low <= y y <= y y + x <= y	Low <= y y <= y y + x <= y
Copy3	$x + s0 \le s0$ $x + s1 \le s1$ $s0 \le s0$ $Low \le y$ $s1 \le s1$	$x + s0 \le s0$ $x + s1 \le s1$ $s0 \le s0$ $s0 \le y$ $s1 + s0 \le s1$ $s1 \le s1$ $s1 \le s1$ $s1 \le s1$	x + s0 <= s0 x + s1 <= s1 s0 <= s0 Low <= y s1 <= s1	$x + s0 \le s0$ $x + s1 \le s1$ $s0 \le s0$ $s0 \le y$ $s1 + s0 \le s1$ $s1 \le s1$ $s1 \le y$ $s1 + s0 \le s0$
Copy4	x <= e0 x <= e1 Low <= y Low <= e1 Low <= e0	$x \le e0$ $x \le e1$ $e0 \le y$ $e0 \le e1$ $e1 \le e0$	x <= e0 x <= e1 Low <= y Low <= e1 Low <= e0	x <= e0 x <= e1 e0 <= y e0 <= e1 e1 <= y e1 <= e0
Copy5	Low <= y	Low <= y x <= y	Low <= y	Low <= y x <= y
Сору6	Low <= z Low <= sum Low <= y x + sum + z <= sum y + z <= y	Low <= z Low <= sum Low <= y x + sum + z <= sum y + x + sum + z <= y y + x + sum + z <= sum	Low <= y y <= y	Low <= y y + x <= y
Dynamic label	$x \le a$ $a \le y$ $z \le a$	x <= a $a + x <= y$ $a + x + z <= a$	x <= y	x <= y

Table 9.2: Constraints generated by all four algorithm for copy programs given Denning [1]

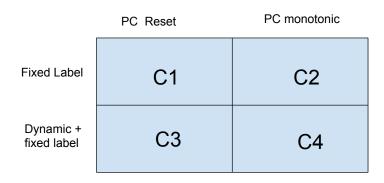


Figure 9.1: Category of constraint generator

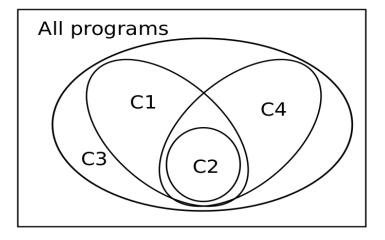


Figure 9.2: Set diagram

of constraints generated by category 3 this relationship shows that set of program declared secure by C1 must be subset of C3. Similarly C2 and C4 differ by use of dynamic labels so set of accepted program by C2 is a subset of set of programs accepted by C4. Use of monotonic PC label helps to capture global information flows so use of monotonic PC label increases the number of constraints. C3 and C4 differ by use of PC label scheme, C4 using monotonic PC and C3 using PC reset so constraints generated by C4 are superset of constraints generated by C3 so set of accepted programs of C4 must be subset of C3. Similarly C2 and C1 differ by PC label scheme so set of accepted programs is a subset of set of programs accepted by C1. Figure

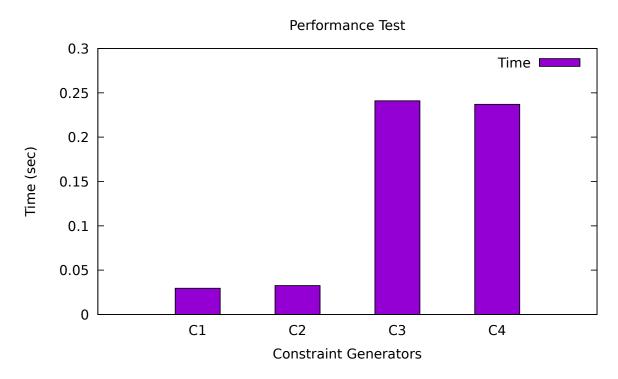


Figure 9.3: Performance test

9.3 shows the average time taken in processing one copy program by all four generator. Time taken in order C1 < C2 < C4 < C3. C4 taking little less time than C3 because of optimization in label generation, this optimization uses property of monotonic PC so it can not applied in C3.

Chapter 10

Constraint Verification with RWFM Labels

Information flow policies require some data related to objects and subjects to enforce any security protocol on them. Each object and subject needs to specify information related to permission like who can read, who can write etc. Label or security class is a way to represent this information. Information flow policies add some prerequisite condition with each operation on objects based on labels of objects and subject.

10.1 RWFM label [2]

Reader Writer Flow Model label has three tuple (s,R,W) the first s is a subject, R is a set of subjects allowed to read object, W is a set of subjects allowed to write and subjects who influenced this object so far.

10.1.1 Secure Information Flow

All information flow security models follow property of lattice because, for secure flow, information must flow from less secure class or label to more secure class or label, any reverse flow is a violation of security. So each information flow model redefines \leq operator to check validity of flow.

Label1 = (s_1, R_1, W_2) , **Label2** = (s_2, R_2, W_2)

information can flow from Label1 to Label2 if and only if $R_1 \supseteq R_2$ and $W_1 \subseteq W_2$.

10.1.2 Definition of Join and Meet operations [2]

Join: Label $1 \oplus \text{Label } 2 = (s_3, R_1 \cap R_2, W_1 \cup W_2)$

Meet: Label1 \otimes Label2 = $(s_3, R_1 \cup R_2, W_1 \cap W_2)$

10.1.3 Create, Read and Write operations on objects with RWFM label [2]

If a subject $s^{(s,R_1,W_1)}$ creates object o, o will be assigned a default label derived from subject is $(s,R_1,W_1\cup\{s\})$. Clearance of subject s is assumed (s,R_s,W_s) , clearance is used to set an upper bound on labels of a subject.

Read: A subject $s^{(s_1,R_1,W_1)}$ can read object $o^{(s_2,R_2,W_2)}$ if and only if $s \in R_2$ and $(R_1 \cap R_2) \supseteq R_s \land (W_1 \cup W_2) \subseteq W_s$. after read operation label of subject s changes into $(s_1,R_1 \cap R_2,W_1 \cup W_2)$ **Write:** A subject $s^{(s_1,R_1,W_1)}$ can read object $s^{(s_2,R_2,W_2)}$ if and only if $s \in W_2$ and $s^{(s_1,R_1,W_1)}$ can read object $s^{(s_2,R_2,W_2)}$ if and only if $s \in W_2$ and $s^{(s_1,R_1,W_1)}$ can read object $s^{(s_2,R_2,W_2)}$ if and only if $s \in W_2$ and $s^{(s_1,R_1,W_1)}$ can read object $s^{(s_1,R_1,W_1)}$ can read object $s^{(s_1,R_1,W_1)}$ if and only if $s \in W_2$ and $s^{(s_1,R_1,W_1)}$ can read object $s^{(s_1,R_1,W_1)}$ if and only if $s \in W_2$ and $s^{(s_1,R_1,W_1)}$ can read object $s^{(s_1,R_1,W_1)}$ if and only if $s \in W_2$ and $s^{(s_1,R_1,W_1)}$ of $s^{(s_1,R_1,W_1)}$ if $s^{(s_1,R_1,W_1)}$ if $s^{(s_1,R_1,W_1)}$ and $s^{(s_1,R_1,W_1)}$ and $s^{(s_1,R_1,W_1)}$ if $s^{(s_1,R_1,W_1)}$ and $s^{(s_1,R_1,$

10.2 Constraint Verification

First of all constraint checker script processes inputs and populate data structures. Inputs are two files, one has all constraints generated by our constraint generator script, other has labels of each object provided by user. Figure 10.1 shows modules of script2 with flow of control. After preprocessing of given constraints and labels, script provides answer to following queries.

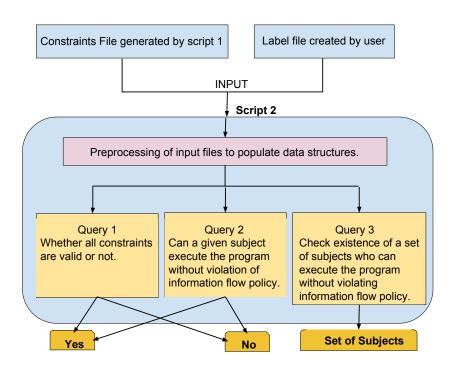


Figure 10.1: Block diagram of script2

1. Whether all constraints are valid or not.

Constraint format in input file : $\underline{x} \le y$.

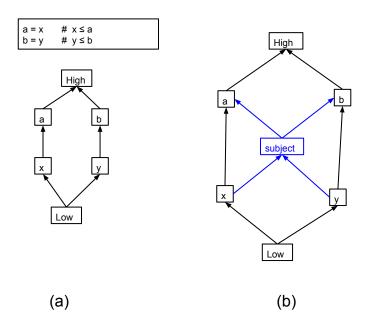


Figure 10.2: Lattice for information flow in a =x and y=b,(a)without subject (b) with subject.

Format of RWFM labels : label(x) (s_1, R_1, W_1) label(y) (s_2, R_2, W_2)

Definition of \leq_{RWFM} operator in RWFM model : label(x) \leq_{RWFM} label(y) if only if $R_1 \supseteq R_2$ and $W_1 \subseteq W_2$ [2]

Script reads all constraints from input file one by one then converts all objects involved in constraint into RWFM label and checks whether they follow \leq_{RWFM} , if any constraint fails to follow \leq_{RWFM} then returns false otherwise true.

2. Can a given subject execute the program without violation of information flow policy.

In the previous case, we are checking constraints without considering the subject, but now a subject is given as input and we have to check whether given subject has all required permissions to execute all statements without violation of information flow policy. Figure 10.2 shows that subject needs to read x and write a in order to execute a = x, so it must follow label(x) \leq_{RWFM} label(subject) and label(subject) \leq_{RWFM} label(a) to maintain secure information flow. If the subject does not have required permissions for any statement then script returns false otherwise true.

3. Check existence of a set of subjects who can execute the program without violating

information flow policy.

In this case, we need to find a maximal set of those subjects who can pass 2nd case. A naive way to implement it: for each subject run the 2nd test and if it passes, add it in a set; after checking all subjects this set will be desired output. There is an efficient way to do it, any subject can pass 2nd test if and only if it is present in reader set of all objects which lie on the left side of \leq operator in constraints and it is present in writer set of all objects which lie on the right side of \leq operator in constraints. So we take the stepwise intersection of all reader sets of objects that lie in left side of \leq operator in constraints and writer sets of all objects that lie on the right side of \leq operator in constraints.

Example: constraint is given in equation 10.1

Conversion of security class of objects into RWFM labels: $x^{(s,R_1,W_1)}$, $y^{(s,R_2,W_2)}$, $z^{(s,R_3,W_3)}$.

Output of script: $R_1 \cap R_2 \cap W_3$

$$\underline{x} \oplus \underline{y} \le \underline{z} \tag{10.1}$$

Chapter 11

Conclusion & Future work

We have implemented various algorithms for constraint generation for capturing information flow in program, category 4 algorithm represents intuitive notion of correct security. Second part is verification of constraints using RWFM [2] is also implemented. We considered the subject in information flow analysis, it helps to deal with real world problems related with information flow.

Future work

- Information flow analysis on python data structures list, dictionary etc.
- Information flow analysis related to dynamic types and objects in python
- Implementation for all features of python.

Chapter 12

Python Syntax

1. if e == 1 then x = y else x = z:

```
if e == 1:
    x = y
selse:
    x = z
```

Listing 12.1: if else, indentation is must

2. while $(x < y) \{ x = y; z += 1 \}$

```
while x < y:
x = y
z += 1
```

Listing 12.2: while loop

3. **function**: (x is a global variable) foo(int z, int y) x+=1; int y; y=1; return y;

```
def foo(z , y):
    global x
    x +=1
    y = 1
    return y
```

Listing 12.3: threads

4. threads

```
import thread
import time
import threading
4
```

```
s s0 = threading.Event()

def thread1():
    global x
    s0.set()
    s0.clear()
    s0.wait()

thread.start_new_thread(thread1,())
```

Listing 12.4: threads

Appendices

Appendix A

Python Script category 1: Constraint Generator

```
#INPUT P - the set of principals that have a stake in the computation.
       p - computing athority
        S - set of all principals in system.
4 from more_itertools import unique_everseen
5 import sys
6 import re, pdb
7 import copy
8 import itertools
iteration = 20
debug = 0
def debugPrint(x):
      if debug == 1:
          print x
 class const:
      otime = "*" # u"\u2295"
      oplus = "+" # u "\u2297"
18
      1t = "<=" # u" \ u2264"
 def extract_Globals(fun_str):
      global_index = [m.start() for m in re.finditer("Global\(", fun_str)]
22
      globals = \{\}
      for it in global_index:
          global_str = parse_parenthesis(it + 6, fun_str)[0]
25
          sq_str = parse_square_br(global_str.find("["), global_str)[0]
```

```
ss = sq_str.strip("["].strip("]")
           sslist = ss.split(",")
28
           for it in sslist:
29
               if it == '':
30
                   continue
31
               globals [(it.strip("'"))] = 1
      return globals
34
  def SemanticsOfProgram(P,p,c,PC,S):
      if p not in P:
36
           print "MISSUSE";
      for x in AccessedGlobal(c):
          if p not in R(lamda(x)):
               print "MISSUSE";
40
      #intialization
      for x in Global(c):
          M[x] = Md[x]
43
          lamda[x] = lamdad[x]
44
      for x in ((VA(c) - Global(c)) | set(PC)):
45
          M[x] = 0
          lamda[x] = (p, S, set([p]))
47
  def VA(data):
49
      return set(parse_variables(data))
51
52 def Global(data): #discard all whiles, ifs and functions -> then remaining
      code will have only globals.
      str = ""
53
      length = len(data)
54
      i = 0
55
      while i < length - 1:
          # checking for keyword
          if parse_keyword(i, data) == "FunctionDef":
               i += 11
59
               i = parse_parenthesis(i, data)[1]
           elif parse_keyword(i, data) == "Expr(":
61
               i += 4
62
               i = parse_parenthesis(i, data)[1]
63
           elif parse_keyword(i, data) == "AugAssign":
64
               i += 9
65
               i = parse_parenthesis(i, data)[1]
66
           elif parse_keyword(i, data) == "If":
67
```

```
i += 2
68
                i = parse_if(i, data)[1]
69
70
           elif parse_keyword(i, data) == "While":
71
                i += 5
72
                i = parse_parenthesis(i, data)[1]
           else:
                str += data[i]
75
           i += 1
       return set(parse_variables(str))
       parseTestVariables (data):
79
       test_index = [ m. start() for m in re.finditer('test=', data) ]
80
       11 = []
81
       for it in test index:
82
           11 += parse_variables(parse_next_parenthesis(it, str)[0])
       return 11
84
85
  def AccessedGlobal(data):
86
       #(i)right-hand side of assignment
       #(ii) condition of branching/iteration
88
       #(iii) return
       11 = []
90
       length = len(data)
       i = 0
92
       while i < length - 1:
93
           # checking for keyword
94
           if parse_keyword(i, data) == "AugAssign":
95
                i += 9
               tmp = parse_parenthesis(i, data)
                11 += parse_variables(tmp[0])
98
                i = tmp[1]
99
           elif parse_keyword(i, data) == "Assign":
                i += 6
101
               tmp = parse_parenthesis(i, data)
                ryt = tmp[0]. split("value=")[1]
103
                11 += parse_variables(ryt)
                i = tmp[1]
105
           elif parse_keyword(i, data) == "If":
106
                i += 2
107
                tmp = parse parenthesis(i, data)
108
                11 += parseTestVariables(tmp[0])
109
```

```
i = tmp[1]
110
           elif parse_keyword(i, data) == "While":
111
                i += 5
                tmp = parse_parenthesis(i, data)
113
                11 += parseTestVariables(tmp[0])
114
                i = tmp[1]
           elif parse_keyword(i, data) == "Return":
                i += 6
                tmp = parse_parenthesis(i, data)
118
                11 += parse_variables(tmp[0])
119
                i = tmp[1]
120
           i += 1
       return set(11)
123
  def ModifiedGlobal(data):
124
       ss = target_of_assignment(data)
125
       modifiedVarList = parse_variables(ss)
126
       return Global(data) & set(modifiedVarList)
127
128
129
  def make_lub_string(llist): # assumption list containing string elemnts
130
       if len(1list) == 0:
           return "Low"
       if len(1list) == 1:
133
           return llist[0]
134
       tmp = set(llist)
135
       uniq_list = list(tmp)
136
       if len(uniq_list) == 1:
137
           return str(uniq_list[0])
138
       ret = ""
139
       ret += uniq_list[0]
140
       i = 1
141
       while i < len(uniq_list):</pre>
142
           ret += " " + const.oplus + " "
143
           ret += uniq_list[i]
           i += 1
145
       return ret
147
148
  def make_glb_string(llist): # assumption list containing string elemnts
149
       # type: (list) -> string
150
       if len(1list) == 0:
151
```

```
return "High"
152
       if len(1list) == 1:
153
           return llist[0]
154
       uniq_list = list(set(llist))
155
       if len(uniq_list) == 1:
156
           return uniq_list[0]
       ret = ""
       ret += uniq_list[0]
159
       i = 1
       while i < len(uniq_list):</pre>
161
           ret += " " + const.otime + " "
           ret += uniq_list[i]
163
           i += 1
164
       return ret
165
166
167
  def split_through_orelse(if_str):
168
       # find first body word
169
       i = if_str.find("body=[")
       i = parse\_square\_br(i + 5, if\_str)[1] + 1
       return ["{" + if_str[1:i] + "}", if_str[i + 7:]]
172
173
174
  def parse_keyword(i, data):
175
       # checking for
176
       funLen = len("Expr(value=Call(func=Name(id='")
177
       attrLen = len("Expr(value=Call(func=Attribute(value=Name(id='")
178
179
       if i + 6 < len(data) - 1 and data[i:i + 6] == 'Assign':
180
           return "Assign"
181
       if i + 9 < len(data) - 1 and data[i:i + 9] == 'AugAssign':
182
           return "AugAssign"
183
       if i + 2 < len(data) - 1 and data[i:i + 2] == 'If':
           return "If"
185
       if i + 5 < len(data) - 1 and data[i:i + 5] == 'While':
           return "While"
187
       if i + 11 < len(data) - 1 and data[i:i + 11] == 'FunctionDef':
           return "FunctionDef"
189
       if i + 6 < len(data) - 1 and data[i:i+6] == "Return":
190
           return "Return"
191
       if i + funLen < len(data) - 1 and data[i:i + funLen] == "Expr(value=
192
      Call(func=Name(id='":
```

```
return "fun_call"
193
       if i + attrLen < len(data) - 1 and data[i:i + attrLen] == "Expr(value
      =Call(func=Attribute(value=Name(id='":
           if extract_variavle_name(i+attrLen, data) == 'thread':
195
                return "thread_fun_call"
196
197
         return "set_clear_wait"
198
       return "none"
199
  def parse_square_br(i, data):
201
       if data[i] != '[':
202
           print "Error: [ is missing"
203
           return []
204
       ret = "["]
205
       count = 1
206
       i += 1
207
       while count > 0 and i < len(data) - 1:
208
           if data[i] == '[':
209
                count += 1
210
           if data[i] == ']':
                count -= 1
           ret += data[i]
           i += 1
214
       return [ret, i]
215
216
217
  def parse_parenthesis(i, data):
218
       # type: (int , string) -> string
219
       if data[i] != '(':
220
           print data [i-4:i+4], data [i]
           print "Error: ( is missing"
222
           return []
       ret = "("
224
       count = 1
225
       i += 1
       while count > 0 and i < len(data) - 1:
           if data[i] == '(':
                count += 1
229
           if data[i] == ')':
230
                count -= 1
           ret += data[i]
232
           i += 1
```

```
return [ret, i]
234
235
  def parse_next_parenthesis(i, data):
236
       while i < len(data)-1 and data[i] != '(':
            i += 1
238
       if i == len(data)-1:
239
            print "No parenthesis in string"
            return ["",i]
241
       ret = "("
       count = 1
243
       i += 1
       while count > 0 and i < len(data) - 1:
245
           if data[i] == '(':
246
                count += 1
247
            if data[i] == ')':
248
                count -= 1
249
            ret += data[i]
250
            i += 1
251
       return [ret, i]
252
253
  def extract_variavle_name(startpos, line):
254
       # string -> string
255
       var = ""
256
       while line[startpos] != "'":
            var += line[startpos]
258
            startpos += 1
259
       return var
260
261
  def target_of_assignment(str): # find all targets
262
       # string -> list
263
       targets\_ptrn = r"targets = \setminus [.*?\setminus]"
264
       ctargets_ptrn = re.compile(targets_ptrn)
265
       temp_list = ctargets_ptrn.findall(str)
       ret = ''.join(temp_list) # converting to string
267
       return ret
269
  def parse_variables(line):
271
       # type: (string) -> list
272
       # type: (str) -> object
273
       id_index = [m. start() for m in re.finditer('id=', line)]
274
       var_list = []
275
```

```
for it in id_index:
276
           vname = extract_variavle_name(it + 4, line)
277
           if vname == "False" or vname == "True":
278
                continue
279
           var_list.append(vname)
280
       return var_list
281
283
  def multiple_assign(assign_str, target_id_index, PC):
       global line
285
       tmp = assign_str.split("value", 1)
286
       rvalue = parse_variables(tmp[0])
287
       lvalue = parse_variables(tmp[1])
288
289
       # printing denning's rule
290
       for it in rvalue:
291
           # left = make_lub_string(dict[key])
292
           if len(lvalue) == 0:
293
                print "low " + const.lt + " " + it
294
                line += 1
           else:
296
                print make_lub_string(lvalue), const.lt, it
                line += 1
298
       return PC[:]
299
300
  def pc_update(PC, list):
301
       #print "pre update", PC
302
       for pc in PC:
303
           pc += 1ist
304
       #print "post update",PC
305
306
def assign_denning(assign_str, PC): # applying dennig's model on
      assignments
       #pdb.set_trace()
308
       global line
       global output
310
       ss = assign_str.split("value")
       target_id_index = [m. start() for m in re.finditer('id=', ss[0])]
312
313
       if len(target_id_index) > 1:
314
           return multiple_assign(assign_str, target_id_index, PC)
315
316
```

```
if "id='" in ss[0]:
317
           left = extract_variavle_name(0, ss[0].split("id='")[1])
318
       else:
319
           left = ['const']
       id_index = [m. start() for m in
321
                re.finditer('id=', ss[1])] # list of starting index of
      variables in right part of string
       rvalue = []
323
       if len(id_index) == 0:
           # rvalue.append("low")
325
           pass
326
       else:
327
           for it in id_index:
328
                startpos = it + 4
329
                vname = extract_variavle_name(startpos, ss[1])
330
                """Exclusion of False keyword"""
                if vname == "False" or vname == "True":
332
                    continue
                rvalue.append(vname)
334
       #pc_update(PC, rvalue)
336
       ret = ""
       1 = rvalue
338
       #11 = 1 + lambda(pc)
       #updating PC
340
       for pc in PC:
341
                output.append(make_lub_string(pc+rvalue)+ " " + const.lt + " "
342
      + make_glb_string([left]))
       line += 1
343
       return PC[:]
344
345
346
  def augAssign_denning(called_by_fun, fun_global, augAssign_str, PC):
       \# i = augAssign_str.find("id=")
348
       # Su = [extract_variavle_name(i + 4, augAssign_str)]
       global output
350
       Sr = parse_variables(augAssign_str)
       for pc in PC:
352
           output.append( make_lub_string(pc) + " " + const.lt + " " + Sr[0])
353
       return PC[:]
354
^{356} PCA = []
```

```
357
  def if_denning(called_by_fun, fun_global, if_str, rest, PC):
358
       # type: (list, list, string, dict, string) -> print rules
359
       #pdb.set trace()
360
       #print "begin IF",PC
361
       if "orelse=" not in if_str:
362
           # print "termination", if_str
           if if_str[0:2] == "[]": # absence of else part
364
                return []
365
           else: # handling else part
366
                else\_str = if\_str
                continuous_parse([], ccalled_by_fun, fun_global, else_str, PC
368
      [:])
                return []
369
370
       tmp = split_through_orelse(if_str)
371
       if_half = tmp[0]
372
       ladder = tmp[1]
373
374
       if if_str[1:5] != "test":
           print "Error test not found in if"
376
       """ extract test = ...() from if half """
378
       i = if_half.find("(")
379
       tmp = parse_parenthesis(i, if_half)
380
       test_str = tmp[0]
381
       #parent_list += parse_variables(test_str)
382
       i = tmp[1]
383
       pc_update(PC, parse_variables(test_str))
384
385
       """then extract body part and process like normal AST text """
386
       # body processing
387
       body_onward_str = if_half[i:] ### Asumption : Compare string always
      followed by body = [...] imediatly
       # setting i to location of [ in body_str: ,body = [...
       i = body_onward_str.find("[")
390
       body_str = parse_square_br(i, body_onward_str)[0]
      memo = \{ \}
392
       memo2 = \{\}
393
      PC1 = copy._deepcopy_list(PC, memo)
394
       PC2 = copy. deepcopy list(PC, memo2)
395
      PCA = list(continuous_parse([], called_by_fun, fun_global, body_str,
396
```

```
PC1[:]))
      PCB = list(continuous_parse([], called_by_fun, fun_global, ladder, PC2
397
      [:]))
      #debugPrint (PCA)
398
      #debugPrint(PCB)
399
      #print "end IF", PCB + PCA
400
       return PCB + PCA
402
  def while_denning(iteration, called_by_fun, fun_global, while_str, PC):
       # debug print "printing while_str", while_str[6:10]
404
       compare = "()"
       if while_str[6:10] == "Name":
406
           tmp = parse_parenthesis(10, while_str)
407
           compare = tmp[0]
408
           i = tmp[1]
409
       elif while_str[6:10] == "Comp":
410
           tmp = parse_parenthesis(13, while_str)
411
           compare = tmp[0]
412
           i = tmp[1]
413
       pc_update(PC, parse_variables(compare))
      # body processing
415
       body_onward_str = while_str[i:] ### Asumption : Compare string always
      followed by body = [...] imediatly
      # setting i to location of [ in body_str: ,body = [...
       i = body_onward_str.find("[")
418
       body_str = parse_square_br(i, body_onward_str)[0]
419
      memo = \{ \}
420
      memo2 = \{ \}
421
      PC1 = copy._deepcopy_list(PC, memo)
422
      PC2 = copy._deepcopy_list(PC, memo2)
423
424
      PCB = list(continuous_parse([], called_by_fun, fun_global, body_str,
425
      PC2[:]))
       return PC1 + PCB
426
  def set_clear_denning( called_by_fun, fun_global, expr_str, PC):
428
      #ASSUMPTION SEMAPHORE VAR IS ALWAYS GLOBAL
       global output
430
       i = expr_str.find("Call(")
431
       i += 4
432
       call_str = parse_parenthesis(i, expr_str)[0]
433
       i = call_str.find("Attribute(")
434
```

```
i += 9
435
       attribute_str = parse_parenthesis(i, call_str)[0]
436
       i = attribute_str.find("Name(")
437
       i += 4
438
       # name_str = parse_parenthesis(i, attribute_str)[0]
439
       i = attribute_str.find("id=")
440
       var_name = extract_variavle_name(i + 4, attribute_str)
       i = attribute_str.find("attr=")
442
       attr = extract_variavle_name(i + 6, attribute_str)
       if attr == "set" or attr == "clear":
444
           # treat it like AugAssign s0 += 1
           # treat it like AugAssign s0 -= 1
446
           #print "set clear -> ",PC
447
           for pc in PC:
448
                output.append( make_lub_string(pc+[var_name])+ " " + const.lt +
449
       " " + make_glb_string([var_name])) #label[left]
       #elif attr == "wait":
450
           # global_while_list.append(var_name)
451
           #print "wait -> ",PC
452
453
454
455
  def fun_denning(fun_str,PC):
456
       fun_name = extract_variavle_name(fun_str.find("name=") + 6, fun_str)
       fun_globals = extract_Globals(fun_str)
458
       funPC = []
459
       PC = continuous_parse(funPC, fun_name, fun_globals, fun_str,PC)
460
       return PC + funPC
461
462
463 # global var for counting
464 \text{ ww} = \text{ww}1 = \text{ww}2 = 1
465
466 # global while list
467 global_while_list = []
469
  def parse_if(i, data):
470
       tmp = parse_parenthesis(i, data)
471
       i = tmp[1]
472
       if_str = tmp[0]
473
       rest = data[i:]
474
       return [if_str , i , rest]
475
```

```
476
  def uniq(1):
477
       11 = []
478
       for it in 1:
479
           11 . append(list(set(it)))
480
       return 11
481
  fun_hash = \{\}
483
  def duplicateRemoval(PC):
485
       tmpPC = []
486
       for pc in PC:
487
           tmppc = list(set(pc))
488
           tmppc.sort()
489
           tmpPC.append(tmppc)
490
       tmpPC.sort()
491
       return list(tmpPC for tmpPC, _ in itertools.groupby(tmpPC))
492
493
  recCount =0
  def continuous_parse( funPC, called_by_fun, fun_global, data, PC):
       #pdb.set_trace()
496
       # type: (object, object) -> object
       global recCount
498
       length = len(data)
       i = 0
500
       recCount += 1
501
       debugPrint(recCount)
502
       debugPrint("Continuous_parse before loop:")
503
       while i < length - 1:
504
           #checking for keyword
505
           if parse_keyword(i,data) == "FunctionDef": #skipping all function
506
      definition
                tmp = parse_parenthesis(i+11, data)
                i = tmp[1]
508
           if parse_keyword(i,data) == "thread_fun_call": #parsing function
      call used in threads
                tmp = parse_parenthesis (i+4, data)
                i = tmp[1]
511
                if 'start_new_thread' in tmp[0]:
512
                    #print "got thread call"
513
                    fi = tmp[0]. find("args=[")
514
                    funName = extract_variavle_name(fi+15,tmp[0]) #len(args=[
515
```

```
Name (id = ') = 15
                    if funName in fun_hash:
516
                        findex = fun_hash[funName]
517
                    else:
518
                        print "Function not found but prgram called function"
519
                    findex += 11
520
                    tmp = parse_parenthesis(findex, data)
                    fun_str = tmp[0]
522
                    tmp = fun_denning(fun_str,[[]])
                    memo = \{ \}
524
                    1s = duplicateRemoval(tmp)
                    PC = copy._deepcopy_list(ls, memo)
526
527
           if parse_keyword(i,data) == "fun_call": #parsing functioncalls
528
               tmp = parse_parenthesis(i+4,data)
529
               #print "got fun call"
530
               funName = extract_variavle_name(i + len("Expr(value=Call(func=
531
      Name(id='"), data)
               i = tmp[1]
532
               if funName in fun_hash:
533
                    findex = fun_hash[funName]
534
               else:
                    print "Function not found but prgram called function"
536
               findex += 11
               tmp = parse_parenthesis(findex, data)
538
               fun_str = tmp[0]
539
               tmp = fun_denning(fun_str,PC)
540
               memo = \{ \}
541
               1s = duplicateRemoval(tmp)
542
               PC = copy._deepcopy_list(ls, memo)
543
544
           if parse_keyword(i,data) == "Return":
545
               debugPrint("Return stmt:")
               funPC += PC
547
           if parse_keyword(i, data) == "set_clear_wait":
549
               tmp = parse_parenthesis(i, data)
               expr_str = tmp[0]
551
               if expr_str.find("'set'") != -1 or expr_str.find("'clear'") !=
552
      -1 or expr_str.find("'wait'") != -1 :
                    i = tmp[1]
553
                    set_clear_denning( called_by_fun, fun_global, expr_str, PC)
554
```

```
if parse_keyword(i, data) == "AugAssign":
555
                i += 9
556
                tmp = parse_parenthesis(i, data)
557
                augAssign_str = tmp[0]
558
                i = tmp[1]
559
                tmp = augAssign_denning(parent_list[:], global_while_list,
560
      called_by_fun , fun_global , augAssign_str , PC[:])
                memo = \{ \}
561
                1s = duplicateRemoval(tmp)
                PC = copy._deepcopy_list(ls, memo)
563
           if parse_keyword(i, data) == "Assign":
565
                global ww
566
                memo = \{ \}
567
                1s = duplicateRemoval(PC)
568
                PC_reset = copy._deepcopy_list(ls,memo)
                ww += 1
570
                i += 6
571
                tmp = parse_parenthesis(i, data)
572
                assign_str = tmp[0]
                i = tmp[1]
574
                if "value=Name(id='threading'" in assign_str:
                    continue
576
                assign_denning(assign_str, PC[:])
                memo = \{ \}
578
                PC = copy._deepcopy_list(PC_reset, memo)
579
           elif parse_keyword(i, data) == "If":
580
                debugPrint("If stmt:")
581
                memo = \{ \}
582
                1s = duplicateRemoval(PC)
583
                PC_reset = copy._deepcopy_list(ls,memo)
584
                global ww1
585
                ww1 += 1
                i += 2
587
                tmp = parse_if(i, data)
                if_str = tmp[0]
589
                i = tmp[1]
                rest = tmp[2]
591
                if_denning(called_by_fun, fun_global, if_str, rest, PC[:])
592
                memo = \{\}
593
                PC = copy._deepcopy_list(PC_reset, memo)
594
           elif parse_keyword(i, data) == "While":
595
```

```
debugPrint("while stmt:")
596
               global ww2
597
              ww2 += 1
598
               i += 5
599
               tmp = parse_parenthesis(i, data)
600
               while_str = tmp[0]
601
               i = tmp[1]
               it = 1
603
              memo = \{ \}
               1s = duplicateRemoval(PC)
605
               PC_reset = copy._deepcopy_list(ls,memo)
606
               while( it <= iteration ):</pre>
607
                  memo = \{ \}
608
                   lastPC = copy._deepcopy_list(PC, memo)
609
                   tmp = duplicateRemoval(PC)
610
                  memo = \{ \}
611
                  PC = copy._deepcopy_list(tmp, memo)
612
                   #print "### While Iteration:", it
613
                   tmp = while_denning(it, called_by_fun, fun_global,
614
      while_str, PC[:])
                   1s = duplicateRemoval(tmp)
615
                  memo = \{\}
                  PC = copy._deepcopy_list(ls, memo)
617
                   #print lastPC,"|--|", PC
618
                   #print "-> PC:", PC
619
                   if lastPC == PC:
620
                       #print "Saturation point of loop!"
621
                       break
622
                   it += 1
623
              memo = \{ \}
624
              PC = copy._deepcopy_list(PC_reset, memo)
625
           i += 1
626
      debugPrint(recCount)
      recCount -= 1;
628
      debugPrint("END continuous parse:")
      return PC[:]
630
632
with open(sys.argv[1], "r") as inputfile:
      # data = inputfile.read().replace('\n', '').replace(' ','')
```

```
data = "".join(inputfile.read().split())
637 \ 11 i s t = []
638 dummy = []
su_sr_list = []
1ine = 0
PC = [[]]
642 gg = []
output = []
  fun_call_index = [m. start() for m in re.finditer("FunctionDef", data)]
  for index in fun_call_index:
      varName = extract_variavle_name(index+len("FunctionDef(name='"), data)
      fun_hash[varName] = index
648
649
  \#G = Global(data)
  continuous_parse("", [], dummy, data, PC)
652
for it in list(unique_everseen(output)):
   print it
```

Appendix B

Python Script category 2: Constraint Generator

```
1 #INPUT P - the set of principals that have a stake in the computation.
       p - computing athority
        S - set of all principals in system.
4 from more_itertools import unique_everseen
5 import sys
6 import re, pdb
7 import copy
8 import itertools
iteration = 20
debug = 0
def debugPrint(x):
      if debug == 1:
          print x
16 class const:
      otime = "*" # u"\u2295"
      oplus = "+" # u "\u2297"
18
      1t = "<=" # u" \ u2264"
20
 def SemanticsOfProgram(P,p,c,PC,S):
      if p not in P:
          print "MISSUSE";
      for x in AccessedGlobal(c):
          if p not in R(lamda(x)):
```

```
print "MISSUSE";
      #intialization
28
      for x in Global(c):
29
          M[x] = Md[x]
30
          lamda[x] = lamdad[x]
31
      for x in ((VA(c) - Global(c)) | set(PC)):
          M[x] = 0
          lamda[x] = (p,S, set([p]))
34
  def VA(data):
      return set(parse_variables(data))
 def Global (data): #discard all whiles, ifs and functions -> then remaining
     code will have only globals.
      str = ""
40
      length = len(data)
41
      i = 0
42
      while i < length - 1:
43
          # checking for keyword
44
          if parse_keyword(i, data) == "FunctionDef":
               i += 11
46
               i = parse_parenthesis(i, data)[1]
          elif parse_keyword(i, data) == "Expr(":
48
               i += 4
               i = parse_parenthesis(i, data)[1]
          elif parse_keyword(i, data) == "AugAssign":
51
               i += 9
52
               i = parse_parenthesis(i, data)[1]
          elif parse_keyword(i, data) == "If":
               i += 2
               i = parse_if(i, data)[1]
          elif parse_keyword(i, data) == "While":
               i += 5
59
               i = parse_parenthesis(i, data)[1]
          else:
               str += data[i]
          i += 1
      return set(parse_variables(str))
64
65
 def parseTestVariables (data):
      test_index = [ m. start() for m in re.finditer('test=', data) ]
```

```
11 = []
68
       for it in test_index:
69
           11 += parse_variables(parse_next_parenthesis(it, str)[0])
70
       return 11
71
72
  def AccessedGlobal(data):
73
       #(i)right-hand side of assignment
       #(ii) condition of branching/iteration
75
       #(iii) return
       11 = []
       length = len(data)
       i = 0
       while i < length - 1:
80
           # checking for keyword
81
           if parse_keyword(i, data) == "AugAssign":
82
                i += 9
               tmp = parse_parenthesis(i, data)
84
                11 += parse_variables(tmp[0])
85
                i = tmp[1]
86
           elif parse_keyword(i, data) == "Assign":
                i += 6
88
               tmp = parse_parenthesis(i, data)
                ryt = tmp[0]. split("value=")[1]
90
                11 += parse_variables(ryt)
                i = tmp[1]
92
           elif parse_keyword(i, data) == "If":
93
                i += 2
94
                tmp = parse_parenthesis(i, data)
95
                11 += parseTestVariables(tmp[0])
                i = tmp[1]
           elif parse_keyword(i, data) == "While":
98
                i += 5
99
                tmp = parse_parenthesis(i, data)
                11 += parseTestVariables(tmp[0])
101
                i = tmp[1]
102
           elif parse_keyword(i, data) == "Return":
103
                i += 6
               tmp = parse_parenthesis(i,data)
105
                11 += parse_variables(tmp[0])
106
                i = tmp[1]
107
           i += 1
108
       return set(11)
109
```

```
110
  def ModifiedGlobal(data):
111
       ss = target_of_assignment(data)
       modifiedVarList = parse_variables(ss)
113
       return Global (data) & set (modified Var List)
114
  def make_lub_string(llist): # assumption list containing string elemnts
117
       if len(1list) == 0:
118
            return "Low"
119
       if len(1list) == 1:
120
           return llist[0]
121
       tmp = set(1list)
122
       uniq_list = list(tmp)
123
       if len(uniq_list) == 1:
124
            return str(uniq_list[0])
125
       ret = ""
126
       ret += uniq_list[0]
127
       i = 1
128
       while i < len(uniq_list):</pre>
            ret += " " + const.oplus + " "
130
            ret += uniq_list[i]
            i += 1
       return ret
133
134
135
  def make_glb_string(llist): # assumption list containing string elemnts
136
137
       # type: (list) -> string
       if len(1list) == 0:
138
            return "High"
139
       if len(1list) == 1:
140
            return llist[0]
141
       uniq_list = list(set(llist))
142
       if len(uniq_list) == 1:
143
            return uniq_list[0]
       ret = ""
145
       ret += uniq_list[0]
146
       i = 1
147
       while i < len(uniq_list):</pre>
148
            ret += " " + const.otime + " "
149
            ret += uniq_list[i]
150
            i += 1
151
```

```
return ret
152
153
154
  def split_through_orelse(if_str):
155
       # find first body word
156
       i = if_str.find("body=[")
       i = parse\_square\_br(i + 5, if\_str)[1] + 1
       return ["{" + if_str[1:i] + "}", if_str[i + 7:]]
159
161
  def parse_keyword(i, data):
162
       # checking for
163
       funLen = len("Expr(value=Call(func=Name(id='")
164
       attrLen = len("Expr(value=Call(func=Attribute(value=Name(id='")
165
166
       if i + 6 < len(data) - 1 and data[i:i + 6] == 'Assign':
167
           return "Assign"
168
       if i + 9 < len(data) - 1 and data[i:i + 9] == 'AugAssign':
169
           return "AugAssign"
       if i + 2 < len(data) - 1 and data[i:i + 2] == 'If':
171
           return "If"
172
       if i + 5 < len(data) - 1 and data[i:i + 5] == 'While':
173
           return "While"
174
       if i + 11 < len(data) - 1 and data[i:i + 11] == 'FunctionDef':
           return "FunctionDef"
176
       if i + 6 < len(data) - 1 and data[i:i+6] == "Return":
           return "Return"
178
       if i + funLen < len(data) - 1 and data[i:i + funLen] == "Expr(value=
179
      Call(func=Name(id='":
     return "fun call"
180
       if i + attrLen < len(data) - 1 and data[i:i + attrLen] == "Expr(value
181
      =Call(func=Attribute(value=Name(id='":
           if extract_variavle_name(i+attrLen, data) == 'thread':
               return "thread_fun_call"
183
           else:
         return "set clear wait"
185
       return "none"
187
  def parse_square_br(i, data):
188
       if data[i] != '[':
189
           print "Error: [ is missing"
190
           return []
191
```

```
ret = "["
192
       count = 1
193
       i += 1
194
       while count > 0 and i < len(data) - 1:
195
            if data[i] == '[':
196
                count += 1
197
           if data[i] == ']':
                count -= 1
199
            ret += data[i]
            i += 1
201
       return [ret, i]
202
203
204
   def parse_parenthesis(i, data):
205
       # type: (int , string) -> string
206
       if data[i] != '(':
207
            print data [i-4:i+4], data [i]
208
            print "Error: ( is missing"
209
            return []
       ret = "("
211
       count = 1
       i += 1
213
       while count > 0 and i < len(data) - 1:
214
           if data[i] == '(':
215
                count += 1
216
            if data[i] == ')':
217
                count -= 1
218
            ret += data[i]
219
            i += 1
220
       return [ret, i]
222
  def parse_next_parenthesis(i, data):
223
       while i < len(data)-1 and data[i] != '(':
224
            i += 1
225
       if i == len(data)-1:
            print "No parenthesis in string"
            return ["",i]
       ret = "("
229
       count = 1
230
       i += 1
       while count > 0 and i < len(data) - 1:
232
            if data[i] == '(':
```

```
count += 1
234
           if data[i] == ')':
235
                count -= 1
236
           ret += data[i]
           i += 1
238
       return [ret, i]
239
  def extract_variavle_name(startpos, line):
241
       # string -> string
242
       var = ""
243
       while line[startpos] != "'":
           var += line[startpos]
245
           startpos += 1
246
       return var
247
248
  def target_of_assignment(str): # find all targets
249
       # string -> list
250
       targets_ptrn = r"targets = \[.*?\]"
251
       ctargets_ptrn = re.compile(targets_ptrn)
252
       temp_list = ctargets_ptrn.findall(str)
253
       ret = ''.join(temp_list) # converting to string
254
       return ret
255
256
257
  def parse_variables(line):
258
       # type: (string) -> list
259
       # type: (str) -> object
260
       id_index = [m. start() for m in re.finditer('id=', line)]
261
       var_list = []
262
       for it in id index:
263
           vname = extract_variavle_name(it + 4, line)
           if vname == "False" or vname == "True":
265
                continue
           var_list.append(vname)
267
       return var_list
268
269
270
  def multiple_assign(assign_str, target_id_index, PC):
271
       global line
272
       tmp = assign_str.split("value", 1)
273
       rvalue = parse variables (tmp[0])
274
       lvalue = parse_variables(tmp[1])
275
```

```
pc_update(PC, lvalue)
276
277
       # printing denning's rule
278
       for it in rvalue:
279
           # left = make_lub_string(dict[key])
280
           if len(lvalue) == 0:
281
                print "low " + const.lt + " " + it
                line += 1
283
           else:
                print make_lub_string(lvalue), const.lt, it
285
                line += 1
       return PC[:]
287
288
  def pc_update(PC, list):
289
       #print "pre update", PC
290
       for pc in PC:
291
           pc += list
292
       #print "post update",PC
293
294
  def assign_denning(assign_str, PC): # applying dennig's model on
      assignments
       #pdb.set_trace()
       global line
297
       global output
       ss = assign_str.split("value")
299
       target_id_index = [m. start() for m in re.finditer('id=', ss[0])]
300
301
       if len(target_id_index) > 1:
302
           return multiple_assign(assign_str, target_id_index, PC)
303
304
       if "id='" in ss[0]:
305
           left = extract_variavle_name(0, ss[0].split("id='")[1])
306
       else:
           left = ['const']
308
       id_index = [m. start() for m in
                re.finditer('id=', ss[1])] # list of starting index of
      variables in right part of string
       rvalue = []
311
       if len(id_index) == 0:
312
           # rvalue.append("low")
313
           pass
314
       else:
315
```

```
for it in id_index:
316
                startpos = it + 4
317
                vname = extract_variavle_name(startpos, ss[1])
318
                """ Exclusion of False keyword """
319
                if vname == "False" or vname == "True":
320
                    continue
                rvalue.append(vname)
       #pc_update(PC, rvalue)
323
324
       ret = ""
325
       1 = rvalue
326
       #11 = 1 + lambda(pc)
327
       #updating PC
328
       for pc in PC:
329
           pc += 1
330
       for pc in PC:
                output.append(make_lub_string(pc)+ " " + const.lt + " " +
332
      make_glb_string([left]))
       line += 1
       return PC[:]
335
  def augAssign_denning(called_by_fun, fun_global, augAssign_str, PC):
337
       \# i = augAssign_str.find("id=")
338
       # Su = [extract_variavle_name(i + 4, augAssign_str)]
339
       global output
340
       Sr = parse_variables(augAssign_str)
341
       pc_update(PC, Sr)
342
       for pc in PC:
343
           output.append( make_lub_string(pc) + " " + const.lt + " " + Sr[0])
344
       return PC[:]
345
346
PCA = []
348
  def if_denning(called_by_fun, fun_global, if_str, rest, PC):
       # type: (list, list, string, dict, string) -> print rules
350
       #pdb.set_trace()
351
       #print "begin IF",PC
352
       if "orelse=" not in if_str:
353
           # print "termination", if_str
354
           if if_str[0:2] == "[]": # absence of else part
355
                return []
356
```

```
else: # handling else part
357
               else\_str = if\_str
358
               continuous_parse([], ccalled_by_fun, fun_global, else_str, PC
359
      [:])
               return []
360
361
      tmp = split_through_orelse(if_str)
       if_half = tmp[0]
363
       ladder = tmp[1]
365
       if if_str[1:5] != "test":
           print "Error test not found in if"
367
368
       """ extract test = ...() from if half """
369
       i = if half.find("(")
       tmp = parse_parenthesis(i, if_half)
371
       test str = tmp[0]
372
      #parent_list += parse_variables(test_str)
373
       i = tmp[1]
374
       pc_update(PC, parse_variables(test_str))
376
       """then extract body part and process like normal AST text """
      # body processing
378
       body_onward_str = if_half[i:] ### Asumption : Compare string always
      followed by body = [...] imediatly
      # setting i to location of [ in body_str: ,body = [...
380
       i = body_onward_str.find("[")
381
       body_str = parse_square_br(i, body_onward_str)[0]
382
      memo = \{ \}
383
      memo2 = \{\}
384
      PC1 = copy._deepcopy_list(PC, memo)
385
      PC2 = copy._deepcopy_list(PC, memo2)
386
      PCA = list(continuous_parse([], called_by_fun, fun_global, body_str,
      PC1[:]))
      PCB = list(continuous_parse([], called_by_fun, fun_global, ladder, PC2
      [:])
      #debugPrint (PCA)
      #debugPrint (PCB)
390
      #print "end IF", PCB + PCA
391
       return PCB + PCA
392
def while_denning(iteration, called_by_fun, fun_global, while_str, PC):
```

```
# debug print "printing while_str", while_str[6:10]
395
       compare = "()"
396
       if while_str[6:10] == "Name":
397
           tmp = parse_parenthesis(10, while_str)
398
           compare = tmp[0]
399
           i = tmp[1]
400
       elif while_str[6:10] == "Comp":
           tmp = parse_parenthesis(13, while_str)
402
           compare = tmp[0]
           i = tmp[1]
404
       pc_update(PC, parse_variables(compare))
      # body processing
406
       body_onward_str = while_str[i:] ### Asumption : Compare string always
407
      followed by body = [...] imediatly
      # setting i to location of [ in body_str: ,body = [...
408
       i = body_onward_str.find("[")
409
       body_str = parse_square_br(i, body_onward_str)[0]
410
      memo = \{ \}
411
      memo2 = \{\}
412
      PC1 = copy._deepcopy_list(PC, memo)
      PC2 = copy._deepcopy_list(PC, memo2)
414
      PCB = list(continuous_parse([], called_by_fun, fun_global, body_str,
416
      PC2[:]))
       return PC1 + PCB
417
418
  def set_clear_denning( called_by_fun, fun_global, expr_str, PC):
419
       #ASSUMPTION SEMAPHORE VAR IS ALWAYS GLOBAL
420
       global output
421
       i = expr_str.find("Call(")
422
       i += 4
423
       call_str = parse_parenthesis(i, expr_str)[0]
424
       i = call_str.find("Attribute(")
       i += 9
426
       attribute_str = parse_parenthesis(i, call_str)[0]
       i = attribute str.find("Name(")
428
       i += 4
       # name_str = parse_parenthesis(i, attribute_str)[0]
430
       i = attribute_str.find("id=")
431
       var_name = extract_variavle_name(i + 4, attribute_str)
432
       i = attribute str.find("attr=")
433
       attr = extract_variavle_name(i + 6, attribute_str)
434
```

```
if attr == "set" or attr == "clear":
435
           # treat it like AugAssign s0 += 1
436
           # treat it like AugAssign s0 -= 1
437
           pc_update(PC,[var_name])
438
           #print "set clear -> ",PC
439
           for pc in PC:
440
                output.append( make_lub_string(pc)+ " " + const.lt + " " +
      make_glb_string([var_name])) #label[left]
       elif attr == "wait":
           #global_while_list.append(var_name)
443
           pc_update(PC,[var_name])
           #print "wait -> ",PC
445
446
  def extract_Globals(fun_str):
447
       global_index = [m. start() for m in re.finditer("Global\(", fun_str))]
448
       globals = \{\}
449
       for it in global_index:
450
           global_str = parse_parenthesis(it + 6, fun_str)[0]
451
           sq_str = parse_square_br(global_str.find("["), global_str)[0]
452
           ss = sq_str.strip("["].strip("]")
           sslist = ss.split(",")
454
           for it in sslist:
                if it == '':
456
                    continue
                globals [(it.strip("'"))] = 1
458
       return globals
459
460
461
  def fun_denning(fun_str,PC):
462
       fun_name = extract_variavle_name(fun_str.find("name=") + 6, fun_str)
463
       fun_globals = extract_Globals(fun_str)
464
       funPC = []
465
       PC = continuous_parse(funPC, fun_name, fun_globals, fun_str,PC)
       return PC + funPC
467
469 # global var for counting
470 \text{ ww} = \text{ww}1 = \text{ww}2 = 1
471
472 # global while list
global_while_list = []
474
475
```

```
def parse_if(i, data):
476
       tmp = parse_parenthesis(i, data)
477
       i = tmp[1]
478
       if_str = tmp[0]
479
       rest = data[i:]
480
       return [if_str , i , rest]
481
  def uniq(1):
483
       11 = []
484
       for it in 1:
485
            11 . append(list(set(it)))
       return 11
487
488
  fun_hash = \{\}
489
490
  def duplicateRemoval(PC):
491
       tmpPC = []
492
       for pc in PC:
493
           tmppc = list(set(pc))
494
           tmppc.sort()
           tmpPC.append(tmppc)
496
       tmpPC.sort()
       return list(tmpPC for tmpPC, _ in itertools.groupby(tmpPC))
498
  recCount = 0
500
  def continuous_parse( funPC, called_by_fun, fun_global, data, PC):
       #pdb.set_trace()
502
       # type: (object, object) -> object
503
       global recCount
504
       length = len(data)
505
       i = 0
506
       recCount += 1
507
       debugPrint(recCount)
       debugPrint("Continuous_parse before loop:")
509
       while i < length - 1:
510
           #checking for keyword
511
            if parse_keyword(i,data) == "FunctionDef": #skipping all function
512
      definition
                tmp = parse_parenthesis(i+11, data)
513
                i = tmp[1]
514
            if parse_keyword(i,data) == "thread_fun_call": #parsing function
515
      call used in threads
```

```
tmp = parse_parenthesis(i+4,data)
516
                i = tmp[1]
517
                if 'start_new_thread' in tmp[0]:
518
                    #print "got thread call"
519
                    fi = tmp[0]. find("args=[")
520
                    funName = extract_variavle_name(fi+15,tmp[0]) #len(args=[
521
      Name (id = ') = 15
                    if funName in fun_hash:
522
                         findex = fun_hash[funName]
                    else:
524
                         print "Function not found but prgram called function"
                    findex += 11
526
                    tmp = parse_parenthesis(findex, data)
527
                    fun_str = tmp[0]
528
                    tmp = fun_denning(fun_str,[[]])
529
                    memo = \{ \}
530
                    1s = duplicateRemoval(tmp)
531
                    PC = copy._deepcopy_list(ls, memo)
532
533
           if parse_keyword(i,data) == "fun_call": #parsing functioncalls
                tmp = parse_parenthesis(i+4,data)
535
                #print "got fun call"
                funName = extract_variavle_name(i + len("Expr(value=Call(func=
537
      Name ( id = '" ) , data )
                i = tmp[1]
538
                if funName in fun_hash:
539
                    findex = fun_hash[funName]
540
                else:
541
                    print "Function not found but prgram called function"
542
                findex += 11
543
                tmp = parse_parenthesis(findex, data)
544
                fun_str = tmp[0]
545
                tmp = fun_denning(fun_str,PC)
               memo = \{ \}
547
                1s = duplicateRemoval(tmp)
               PC = copy._deepcopy_list(ls, memo)
549
           if parse_keyword(i,data) == "Return":
551
                debugPrint("Return stmt:")
552
                funPC += PC
553
           if parse keyword(i, data) == "set clear wait":
554
                i += 4
555
```

```
tmp = parse_parenthesis(i, data)
556
               expr_str = tmp[0]
557
               if expr_str.find("'set'") != -1 or expr_str.find("'clear'") !=
558
      -1 or expr_str.find("'wait'") != -1:
                    i = tmp[1]
559
                    set_clear_denning( called_by_fun, fun_global, expr_str, PC)
560
           if parse_keyword(i, data) == "AugAssign":
               i += 9
562
               tmp = parse_parenthesis(i, data)
               augAssign_str = tmp[0]
564
               i = tmp[1]
               tmp = augAssign_denning(parent_list[:], global_while_list,
566
      called_by_fun, fun_global, augAssign_str, PC[:])
               memo = \{\}
567
               1s = duplicateRemoval(tmp)
568
               PC = copy._deepcopy_list(ls, memo)
570
           if parse_keyword(i, data) == "Assign":
571
               global ww
572
               ww += 1
               i += 6
574
               tmp = parse_parenthesis(i, data)
               assign_str = tmp[0]
576
               i = tmp[1]
               if "value=Name(id='threading'" in assign_str:
578
579
               tmp = assign_denning(assign_str, PC[:])
580
               memo = \{ \}
581
               1s = duplicateRemoval(tmp)
582
               PC = copy._deepcopy_list(ls,memo)
583
           elif parse_keyword(i, data) == "If":
584
               debugPrint("If stmt:")
585
               global ww1
               ww1 += 1
587
               i += 2
               tmp = parse_if(i, data)
589
               if_str = tmp[0]
               i = tmp[1]
591
               rest = tmp[2]
592
               tmp = if_denning(called_by_fun, fun_global, if_str, rest, PC
593
      [:]
               1s = duplicateRemoval(tmp)
594
```

```
memo = \{ \}
595
              PC = copy._deepcopy_list(ls, memo)
596
           elif parse_keyword(i, data) == "While":
597
               debugPrint("while stmt:")
598
               global ww2
599
              ww2 += 1
600
               i += 5
               tmp = parse_parenthesis(i, data)
602
               while_str = tmp[0]
               i = tmp[1]
604
               it = 1
605
               while( it <= iteration ):</pre>
606
                  memo = \{ \}
607
                   lastPC = copy._deepcopy_list(PC, memo)
608
                   tmp = duplicateRemoval(PC)
609
                   memo = \{ \}
610
                   PC = copy._deepcopy_list(tmp, memo)
611
                   #print "### While Iteration:", it
612
                   tmp = while_denning(it, called_by_fun, fun_global,
613
      while_str, PC[:])
                   1s = duplicateRemoval(tmp)
614
                  memo = \{\}
                   PC = copy._deepcopy_list(ls, memo)
616
                   #print lastPC,"|--|", PC
617
                   #print "-> PC:", PC
618
                   if lastPC == PC:
619
                       #print "Saturation point of loop!"
620
                       break
621
                   it += 1
622
           i += 1
623
      debugPrint (recCount)
624
      recCount -= 1;
625
      debugPrint("END continuous parse:")
      return PC[:]
627
629
  with open(sys.argv[1], "r") as inputfile:
      # data = inputfile.read().replace('\n', '').replace(' ','')
632
      data = "".join(inputfile.read().split())
634 \ 11ist = []
```

```
635 dummy = []
su_sr_list = []
637 \ 1ine = 0
^{638} PC = [[]]
gg = []
output = []
642 fun_call_index = [m. start() for m in re.finditer("FunctionDef", data)]
  for index in fun_call_index:
       varName = extract_variavle_name(index+len("FunctionDef(name='"), data)
644
       fun_hash[varName] = index
646
^{647} #G = Global (data)
  continuous_parse("", [], dummy, data, PC)
649
  for it in list(unique_everseen(output)):
650
   print it
```

Appendix C

Python Script category 3: Constraint Generator

```
1 #INPUT P - the set of principals that have a stake in the computation.
       p - computing athority
        S - set of all principals in system.
        label file
from more_itertools import unique_everseen
6 import sys
7 import re, pdb
8 import copy
9 import itertools
iteration = 20
debug = 0
def debugPrint(x):
      if debug == 1:
          print x
17 class const:
     otime = "*" # u"\u2295"
      oplus = "+" # u" \setminus u2297"
     1t = "<=" # u" \ u 2264"
def SemanticsOfProgram(P,p,c,PC,S):
      if p not in P:
          print "MISSUSE";
    for x in AccessedGlobal(c):
```

```
if p not in R(lamda(x)):
               print "MISSUSE";
28
      #intialization
29
      for x in Global(c):
30
          M[x] = Md[x]
31
          lamda[x] = lamdad[x]
      for x in ((VA(c) - Global(c)) | set(PC)):
          M[x] = 0
34
          lamda[x] = (p, S, set([p]))
36
  def VA(data):
      return set(parse_variables(data))
38
40 def Global(data): #discard all whiles, ifs and functions -> then remaining
     code will have only globals.
      str = ""
41
      length = len(data)
42
      i = 0
43
      while i < length - 1:
44
          # checking for keyword
          if parse_keyword(i, data) == "FunctionDef":
46
               i += 11
               i = parse_parenthesis(i, data)[1]
48
           elif parse_keyword(i, data) == "Expr(":
               i += 4
50
               i = parse_parenthesis(i, data)[1]
51
           elif parse_keyword(i, data) == "AugAssign":
52
               i += 9
53
               i = parse_parenthesis(i, data)[1]
           elif parse_keyword(i, data) == "If":
55
               i += 2
               i = parse_if(i, data)[1]
           elif parse_keyword(i, data) == "While":
59
               i += 5
               i = parse_parenthesis(i, data)[1]
           else:
62
               str += data[i]
63
           i += 1
64
      return set(parse_variables(str))
65
66
67 def parseTestVariables (data):
```

```
test_index = [ m. start() for m in re.finditer('test=', data) ]
68
       11 = []
69
       for it in test_index:
70
           11 += parse_variables(parse_next_parenthesis(it, str)[0])
71
       return 11
72
  def AccessedGlobal(data):
      #(i)right-hand side of assignment
75
      #(ii) condition of branching/iteration
      #(iii) return
       11 = []
       length = len(data)
       i = 0
80
       while i < length - 1:
81
           # checking for keyword
82
           if parse_keyword(i, data) == "AugAssign":
               i += 9
               tmp = parse_parenthesis(i, data)
               11 += parse_variables(tmp[0])
86
               i = tmp[1]
           elif parse_keyword(i, data) == "Assign":
88
               i += 6
               tmp = parse_parenthesis(i, data)
90
               ryt = tmp[0]. split("value=")[1]
               11 += parse_variables(ryt)
92
               i = tmp[1]
93
           elif parse_keyword(i, data) == "If":
94
               i += 2
95
               tmp = parse_parenthesis(i, data)
               11 += parseTestVariables(tmp[0])
               i = tmp[1]
           elif parse_keyword(i, data) == "While":
99
               i += 5
               tmp = parse_parenthesis(i, data)
               11 += parseTestVariables(tmp[0])
               i = tmp[1]
103
           elif parse_keyword(i,data) == "Return":
               i += 6
105
               tmp = parse_parenthesis(i, data)
106
               11 += parse_variables(tmp[0])
107
               i = tmp[1]
108
           i += 1
109
```

```
return set(11)
110
111
def ModifiedGlobal(data):
       ss = target_of_assignment(data)
113
       modifiedVarList = parse_variables(ss)
114
       return Global(data) & set(modifiedVarList)
115
  def make_lub_string(llist): # assumption list containing string elemnts
118
       if len(1list) == 0:
119
           return "Low"
120
       if len(1list) == 1:
           return llist[0]
122
       tmp = set(1list)
123
       uniq_list = list(tmp)
124
       if len(uniq_list) == 1:
125
           return str(uniq_list[0])
126
       ret = ""
127
       ret += uniq_list[0]
128
       i = 1
       while i < len(uniq_list):</pre>
130
           ret += " " + const.oplus + " "
           ret += uniq_list[i]
           i += 1
133
       return ret
134
135
136
  def make_glb_string(llist): # assumption list containing string elemnts
137
       # type: (list) -> string
138
       if len(11ist) == 0:
139
           return "High"
140
       if len(1list) == 1:
141
           return llist[0]
142
       uniq_list = list(set(llist))
143
       if len(uniq_list) == 1:
           return uniq_list[0]
145
       ret = ""
       ret += uniq_list[0]
147
       i = 1
148
       while i < len(uniq_list):</pre>
149
           ret += " " + const.otime + " "
150
           ret += uniq_list[i]
151
```

```
i += 1
152
       return ret
153
154
155
  def split_through_orelse(if_str):
156
       # find first body word
157
       i = if_str.find("body=[")
       i = parse\_square\_br(i + 5, if\_str)[1] + 1
159
       return ["{" + if_str[1:i] + "}", if_str[i + 7:]]
161
  def parse_keyword(i, data):
163
       # checking for
164
       funLen = len("Expr(value=Call(func=Name(id='")
165
       attrLen = len("Expr(value=Call(func=Attribute(value=Name(id='")
166
167
       if i + 6 < len(data) - 1 and data[i:i + 6] == 'Assign':
168
           return "Assign"
169
       if i + 9 < len(data) - 1 and data[i:i + 9] == 'AugAssign':
           return "AugAssign"
       if i + 2 < len(data) - 1 and data[i:i + 2] == 'If':
172
           return "If"
173
       if i + 5 < len(data) - 1 and data[i:i + 5] == 'While':
174
           return "While"
       if i + 11 < len(data) - 1 and data[i:i + 11] == 'FunctionDef':
176
           return "FunctionDef"
       if i + 6 < len(data) - 1 and data[i:i+6] == "Return":
178
           return "Return"
179
       if i + funLen < len(data) - 1 and data[i:i + funLen] == "Expr(value=
180
      Call(func=Name(id='":
     return "fun_call"
181
       if i + attrLen < len(data) - 1 and data[i:i + attrLen] == "Expr(value
182
      =Call(func=Attribute(value=Name(id='":
           if extract_variavle_name(i+attrLen, data) == 'thread':
183
               return "thread_fun_call"
           else:
185
         return "set_clear_wait"
       return "none"
187
188
  def parse_square_br(i, data):
189
       if data[i] != '[':
190
           print "Error: [ is missing"
191
```

```
return []
192
       ret = "["]
193
       count = 1
194
       i += 1
195
       while count > 0 and i < len(data) - 1:
196
            if data[i] == '[':
197
                count += 1
198
            if data[i] == ']':
199
                count -= 1
            ret += data[i]
201
            i += 1
202
       return [ret, i]
203
204
205
   def parse_parenthesis(i, data):
206
       # type: (int , string) -> string
207
       if data[i] != '(':
208
            print data[i-4:i+4], data[i]
209
            print "Error: ( is missing"
210
            return []
211
       ret = "("
       count = 1
213
       i += 1
214
       while count > 0 and i < len(data) - 1:
215
            if data[i] == '(':
                count += 1
217
            if data[i] == ')':
218
                count -= 1
219
            ret += data[i]
220
            i += 1
       return [ret, i]
222
223
  def parse_next_parenthesis(i, data):
224
       while i < len(data)-1 and data[i] != '(':
225
            i += 1
226
       if i == len(data)-1:
            print "No parenthesis in string"
228
            return ["",i]
229
       ret = "("
230
       count = 1
       i += 1
232
       while count > 0 and i < len(data) - 1:
233
```

```
if data[i] == '(':
234
                count += 1
235
           if data[i] == ')':
236
                count -= 1
            ret += data[i]
238
            i += 1
239
       return [ret, i]
241
  def extract_variavle_name(startpos, line):
       # string -> string
243
       var = ""
       while line[startpos] != "'":
245
           var += line[startpos]
246
            startpos += 1
247
       return var
248
249
  def target_of_assignment(str): # find all targets
250
       # string -> list
251
       targets\_ptrn = r"targets = \setminus [.*?\setminus]"
252
       ctargets_ptrn = re.compile(targets_ptrn)
       temp_list = ctargets_ptrn.findall(str)
254
       ret = ''.join(temp_list) # converting to string
255
       return ret
256
258
  def parse_variables(line):
       # type: (string) -> list
260
       # type: (str) -> object
261
       id_index = [m. start() for m in re.finditer('id=', line)]
262
       var list = []
263
       for it in id_index:
           vname = extract_variavle_name(it + 4, line)
265
           if vname == "False" or vname == "True":
                continue
267
            var_list.append(vname)
       return var list
269
270
271
       multiple_assign(assign_str, target_id_index, PC):
272
       global line
273
       tmp = assign_str.split("value", 1)
274
       rvalue = parse_variables(tmp[0])
275
```

```
lvalue = parse_variables(tmp[1])
276
       pc_update(PC, lvalue)
277
278
       # printing denning's rule
279
       for it in rvalue:
280
           # left = make_lub_string(dict[key])
281
            if len(lvalue) == 0:
                print "low " + const.lt + " " + it
283
                line += 1
284
            else:
285
                print make_lub_string(lvalue), const.lt, it
286
                line += 1
287
       return PC[:]
288
289
   def pc_update(PC, list):
290
       #print "pre update", PC
291
       for pc in PC:
292
            pc += list
293
       #print "post update",PC
294
  def label_to_lub(llist):
       if len(1list) == 1:
296
            return llist[0]
       ret = ""
298
       ret += 11ist[0]
299
       i = 1
300
       while i < len(llist):</pre>
301
            ret += " "+const.oplus+" "+ 11ist[i]
302
            i +=1
303
       return ret
304
305
306
  def make_lub_string_label(llist):
307
       if len(1list) == 0:
            return "Low"
309
       if len(1list) == 1:
310
            return label_to_lub(label[llist[0]])
       tmp = set(llist)
312
       uniq_list = list(tmp)
313
       if len(uniq_list) == 1:
314
            return label_to_lub(label[uniq_list[0]])
315
       ret = ""
316
       ret += label_to_lub(label[uniq_list[0]])
317
```

```
i = 1
318
       while i < len(uniq_list):</pre>
319
            ret += " " + const.oplus + " "
320
            ret += label_to_lub(label[uniq_list[i]])
            i += 1
322
       return ret
  def convertToLabelList(11):
325
       11ist = list(set(11))
326
       global label
327
       if len(1list) == 0:
328
           return llist
329
       ret = []
330
       for it in llist:
            if it not in label:
                ret += [it]
333
                continue
334
            ret += label[it]
335
       return list(set(ret))
336
  def make_lubPC(PC):
338
       11ist = []
       for 1t in PC:
340
            11ist += 1t
       return list(set(llist))
342
343
344 def assign_denning(G, assign_str, PC): # applying dennig's model on
      assignments
       #pdb.set_trace()
345
       global line
346
       global label
347
       global output
348
       ss = assign_str.split("value")
       target_id_index = [m. start() for m in re.finditer('id=', ss[0])]
350
       if len(target_id_index) > 1:
352
            return multiple_assign(assign_str, target_id_index, PC)
354
       if "id='" in ss[0]:
355
            left = extract_variavle_name(0, ss[0].split("id='")[1])
356
       else:
357
            left = ['const']
358
```

```
id_index = [m. start() for m in
359
                re.finditer('id=', ss[1])] # list of starting index of
360
      variables in right part of string
       rvalue = []
361
       if len(id_index) == 0:
362
           # rvalue.append("low")
363
           pass
       else:
365
           for it in id_index:
                startpos = it + 4
367
                vname = extract_variavle_name(startpos, ss[1])
                """Exclusion of False keyword"""
369
                if vname == "False" or vname == "True":
370
                    continue
371
                rvalue.append(vname)
372
       #pc_update(PC, rvalue)
373
374
       ret = ""
375
       1 = convertToLabelList(rvalue)
376
       #11 = 1 + lambda(pc)
       #updating PC
378
       if left in G:
           for pc in PC:
380
                output.append(make_lub_string(convertToLabelList(pc)+1)+ " " +
381
      const.lt + " " + make_lub_string(label[left]))
       else:
382
           # update dynamic label, for all x (MNG(BR)) union \{pc\}) [lambda(x)=1]
383
      +lambda(x)+lambda(pc)]
           if left not in label:
384
                label[left] = []
385
           label[left] = list(set(label[left] + make_lubPC(PC[:]+[1])))
387
           if left in label[left]:
                label[left].remove(left)
389
           #print label[left], "changed", left
391
       line += 1
       return PC[:]
393
394
395
  def augAssign_denning(called_by_fun, fun_global, augAssign_str, PC):
       \# i = augAssign_str.find("id=")
397
```

```
# Su = [extract_variavle_name(i + 4, augAssign_str)]
398
       global output
399
       Sr = parse_variables(augAssign_str)
400
       for pc in PC:
401
           output.append( make_lub_string(pc) + " " + const.lt + " " + Sr[0])
402
       return PC[:]
403
405 \text{ PCA} = []
  def if_denning(G, called_by_fun, fun_global, if_str, rest, PC):
407
       # type: (list, list, string, dict, string) -> print rules
       #pdb.set_trace()
409
       #print "begin IF", PC
410
       if "orelse=" not in if_str:
411
           # print "termination", if_str
412
           if if_str[0:2] == "[]": # absence of else part
413
                return []
414
           else: # handling else part
415
                else\_str = if\_str
416
                continuous\_parse(G,[], ccalled\_by\_fun, fun\_global, else\_str, PC
      [:])
                return []
419
       tmp = split_through_orelse(if_str)
       if_half = tmp[0]
421
       ladder = tmp[1]
422
423
       if if_str[1:5] != "test":
424
           print "Error test not found in if"
425
426
       """ extract test = ...() from if half """
427
       i = if_half.find("(")
428
       tmp = parse_parenthesis(i, if_half)
       test_str = tmp[0]
430
       #parent_list += parse_variables(test_str)
       i = tmp[1]
432
       pc_update(PC, parse_variables(test_str))
434
       """then extract body part and process like normal AST text """
435
       # body processing
436
       body_onward_str = if_half[i:] ### Asumption : Compare string always
437
      followed by body = [...] imediatly
```

```
# setting i to location of [ in body_str: ,body = [...
438
       i = body_onward_str.find("[")
439
       body_str = parse_square_br(i, body_onward_str)[0]
440
       memo = \{ \}
441
       memo2 = {}
442
       PC1 = copy._deepcopy_list(PC, memo)
443
       PC2 = copy._deepcopy_list(PC, memo2)
      PCA = list(continuous_parse(G,[], called_by_fun, fun_global, body_str,
445
       PC1[:]))
      PCB = list(continuous_parse(G,[], called_by_fun, fun_global, ladder,
446
      PC2[:]))
       #debugPrint (PCA)
447
       #debugPrint (PCB)
448
       #print "end IF",PCB + PCA
449
       return PCB + PCA
450
  def while_denning(iteration, G, called_by_fun, fun_global, while_str, PC):
452
       # debug print "printing while_str", while_str[6:10]
453
       compare = "()"
454
       if while_str [6:10] == "Name":
           tmp = parse_parenthesis(10, while_str)
456
           compare = tmp[0]
           i = tmp[1]
458
       elif while_str[6:10] == "Comp":
           tmp = parse_parenthesis(13, while_str)
460
           compare = tmp[0]
461
           i = tmp[1]
462
       pc_update(PC, parse_variables(compare))
463
       # body processing
464
       body_onward_str = while_str[i:] ### Asumption : Compare string always
465
      followed by body = [...] imediatly
       # setting i to location of [ in body_str: ,body = [...
466
       i = body_onward_str.find("[")
       body_str = parse_square_br(i, body_onward_str)[0]
468
       memo = \{ \}
       memo2 = \{\}
470
       PC1 = copy._deepcopy_list(PC, memo)
       PC2 = copy._deepcopy_list(PC, memo2)
472
473
      PCB = list(continuous_parse(G, [], called_by_fun, fun_global, body_str,
474
       PC2[:]))
       return PC1 + PCB
475
```

```
476
  def set_clear_denning(G, called_by_fun, fun_global, expr_str, PC):
477
       #ASSUMPTION SEMAPHORE VAR IS ALWAYS GLOBAL
478
       global output
479
       i = expr_str.find("Call(")
480
       i += 4
481
       call_str = parse_parenthesis(i, expr_str)[0]
       i = call_str.find("Attribute(")
483
       attribute_str = parse_parenthesis(i, call_str)[0]
485
       i = attribute_str.find("Name(")
       i += 4
487
       # name_str = parse_parenthesis(i, attribute_str)[0]
488
       i = attribute_str.find("id=")
489
       var_name = extract_variavle_name(i + 4, attribute_str)
490
       i = attribute_str.find("attr=")
       attr = extract variavle name(i + 6, attribute str)
492
       if attr == "set" or attr == "clear":
493
           # treat it like AugAssign s0 += 1
494
           # treat it like AugAssign s0 -= 1
           #pc_update(PC, convertToLabelList([var_name]))
496
           #print "set clear -> ",PC
           for pc in PC:
498
               output.append( make_lub_string(pc+[var_name])+ " " + const.lt +
       " " + make_lub_string(label[var_name])) #label[left]
       elif attr == "wait":
500
           # global_while_list.append(var_name)
501
           #pc_update(PC, convertToLabelList([var_name]))
502
503
           #print "wait -> ",PC
504
505
  def extract_Globals(fun_str):
506
       global_index = [m.start() for m in re.finditer("Global\(", fun_str)]
       globals = \{\}
508
       for it in global_index:
           global_str = parse_parenthesis(it + 6, fun_str)[0]
510
           sq_str = parse_square_br(global_str.find("["), global_str)[0]
           ss = sq_str.strip("["].strip("]")
512
           sslist = ss.split(",")
513
           for it in sslist:
514
               if it == '':
515
                    continue
516
```

```
globals [(it.strip("'"))] = 1
517
       return globals
518
519
520
  def fun_denning(fun_str,PC):
521
       fun_name = extract_variavle_name(fun_str.find("name=") + 6, fun_str)
522
       fun_globals = extract_Globals(fun_str)
       funPC = []
524
       PC = continuous_parse(G, funPC, fun_name, fun_globals, fun_str, PC)
       return PC + funPC
526
528 # global var for counting
529 \text{ ww} = \text{ww}1 = \text{ww}2 = 1
530
  # global while list
532 global_while_list = []
533
534
  def parse_if(i, data):
535
       tmp = parse_parenthesis(i, data)
       i = tmp[1]
537
       if_str = tmp[0]
       rest = data[i:]
539
       return [if_str, i, rest]
541
  def uniq(1):
542
       11 = []
543
       for it in 1:
544
            11.append(list(set(it)))
545
       return 11
546
547
  fun_hash = \{\}
548
  def duplicateRemoval(PC):
550
       tmpPC = []
551
       for pc in PC:
552
            tmppc = list(set(pc))
            tmppc.sort()
554
            tmpPC.append(tmppc)
555
       tmpPC.sort()
556
       return list(tmpPC for tmpPC, _ in itertools.groupby(tmpPC))
557
558
```

```
559 recCount =0
  def continuous_parse(G, funPC, called_by_fun, fun_global, data, PC):
       #pdb.set trace()
561
       # type: (object, object) -> object
562
       global recCount
563
       length = len(data)
564
       i = 0
       recCount += 1
566
       debugPrint (recCount)
       debugPrint("Continuous_parse before loop:")
568
       while i < length - 1:
           #checking for keyword
570
           if parse_keyword(i,data) == "FunctionDef": #skipping all function
571
      definition
               tmp = parse parenthesis(i+11, data)
572
               i = tmp[1]
573
           if parse_keyword(i,data) == "thread_fun_call": #parsing function
574
      call used in threads
               tmp = parse_parenthesis(i+4,data)
575
               i = tmp[1]
               if 'start_new_thread' in tmp[0]:
577
                    #print "got thread call"
                    fi = tmp[0].find("args=[")
579
                    funName = extract_variavle_name(fi+15,tmp[0]) #len(args=[
      Name (id = ') = 15
                    if funName in fun_hash:
581
                        findex = fun_hash[funName]
582
                    else:
583
                        print "Function not found but prgram called function"
584
                    findex += 11
585
                    tmp = parse_parenthesis(findex, data)
                    fun_str = tmp[0]
587
                    tmp = fun_denning(fun_str,[[]])
                   memo = \{ \}
589
                    ls = duplicateRemoval(tmp)
                   PC = copy._deepcopy_list(ls, memo)
591
           if parse_keyword(i,data) == "fun_call": #parsing functioncalls
593
               tmp = parse_parenthesis(i+4,data)
594
               #print "got fun call"
595
               funName = extract_variavle_name(i + len("Expr(value=Call(func=
596
      Name(id='"), data)
```

```
i = tmp[1]
597
                if funName in fun_hash:
598
                    findex = fun_hash[funName]
599
                else:
600
                    print "Function not found but prgram called function"
601
                findex += 11
602
                tmp = parse_parenthesis(findex, data)
                fun_str = tmp[0]
604
                tmp = fun_denning(fun_str,PC)
               memo = \{ \}
606
                1s = duplicateRemoval(tmp)
               PC = copy._deepcopy_list(ls, memo)
608
609
           if parse_keyword(i,data) == "Return":
610
                debugPrint("Return stmt:")
611
                funPC += PC
612
           if parse_keyword(i, data) == "set_clear_wait":
613
                i += 4
614
                tmp = parse_parenthesis(i, data)
615
                expr_str = tmp[0]
                if expr_str.find("'set'") != -1 or expr_str.find("'clear'") !=
617
      -1 or expr_str.find("'wait'") != -1:
                    i = tmp[1]
618
                    set_clear_denning(G, called_by_fun, fun_global, expr_str,
619
      PC)
           if parse_keyword(i, data) == "AugAssign":
620
                i += 9
621
                tmp = parse_parenthesis(i, data)
622
                augAssign_str = tmp[0]
623
                i = tmp[1]
624
                tmp = augAssign_denning(parent_list[:], global_while_list,
625
      called_by_fun , fun_global , augAssign_str , PC[:])
           if parse_keyword(i, data) == "Assign":
627
                global ww
               ww += 1
629
                i += 6
                tmp = parse_parenthesis(i, data)
631
                assign_str = tmp[0]
632
                i = tmp[1]
633
                if "value=Name(id='threading'" in assign_str:
634
                    continue
635
```

```
assign_denning(G, assign_str, PC[:])
636
            elif parse_keyword(i, data) == "If":
637
                debugPrint("If stmt:")
638
                1s = duplicateRemoval(PC)
639
                memo = \{ \}
                PC_reset = copy._deepcopy_list(ls, memo)
641
                global ww1
                ww1 += 1
643
                i += 2
                tmp = parse_if(i, data)
645
                if_str = tmp[0]
                i = tmp[1]
647
                rest = tmp[2]
648
                if_denning(G, called_by_fun, fun_global, if_str, rest, PC[:])
649
                memo = \{ \}
650
                PC = copy._deepcopy_list(PC_reset, memo)
            elif parse_keyword(i, data) == "While":
652
                debugPrint("while stmt:")
653
                global ww2
654
                ww2 += 1
                i += 5
656
                tmp = parse_parenthesis(i, data)
                while_str = tmp[0]
658
                i = tmp[1]
                it = 1
660
                1s = duplicateRemoval(PC)
661
                memo = \{ \}
662
                PC_reset = copy._deepcopy_list(ls, memo)
663
                while( it <= iteration ):</pre>
                    memo = \{ \}
665
                    lastPC = copy._deepcopy_list(PC, memo)
                    tmp = duplicateRemoval(PC)
667
                    memo = \{ \}
                    PC = copy._deepcopy_list(tmp, memo)
669
                    #print "### While Iteration:", it
                    tmp = while_denning(it, G, called_by_fun, fun_global,
671
      while_str, PC[:])
                    1s = duplicateRemoval(tmp)
672
                    memo = \{\}
673
                    PC = copy._deepcopy_list(ls, memo)
674
                    #print lastPC,"|--|", PC
675
                    #print "-> PC:", PC
676
```

```
if lastPC == PC:
677
                       #print "Saturation point of loop!", PC, lastPC
678
679
                   it += 1
680
               memo = \{ \}
681
               PC = copy._deepcopy_list(PC_reset, memo)
682
           i += 1
       debugPrint(recCount)
684
       recCount -= 1;
       debugPrint("END continuous parse:")
686
       return PC[:]
687
688
  def process_label_fille(filename):
689
      \# a = x, y, z
690
       label = \{\}
691
      with open (filename, "r") as inputfile:
           for line in inputfile:
693
               11 = "".join(line.split())
694
               tmp = 11.split("=")
695
               var = tmp[0]
               lineList = tmp[1].split(",")
697
               label[var] = lineList
       return label
699
  ############# main
701
      with open(sys.argv[1], "r") as inputfile:
      # data = inputfile.read().replace('\n', '').replace(' ','')
703
      data = "".join(inputfile.read().split())
704
705 \ 11ist = []
706 dummy = []
su_sr_list = []
1ine = 0
709 PC = [[]]
710 \ label = \{ \}
11  label = process_label_fille(sys.argv[2])
712 \text{ gg} = []
output = []
714 for key in label:
      gg.append(key)
716 G = set(gg)
fun_call_index = [m. start() for m in re.finditer("FunctionDef", data)]
```

```
for index in fun_call_index:
    varName = extract_variavle_name(index+len("FunctionDef(name='"),data)
    fun_hash[varName] = index

721
722 #G = Global(data)
723 continuous_parse(G,"", [], dummy, data, PC)

724
725 for it in list(unique_everseen(output)):
    print it
```

Appendix D

Python Script category 4: Constraint Generator

```
1 #INPUT P - the set of principals that have a stake in the computation.
       p - computing athority
        S - set of all principals in system.
        label file
from more_itertools import unique_everseen
6 import sys
7 import re, pdb
8 import copy
9 import itertools
iteration = 20
debug = 0
def debugPrint(x):
      if debug == 1:
          print x
17 class const:
      otime = "*" # u"\u2295"
      oplus = "+" # u" \setminus u2297"
      1t = "<=" # u" \ u 2264"
 def SemanticsOfProgram(P,p,c,PC,S):
      if p not in P:
          print "MISSUSE";
25
      for x in AccessedGlobal(c):
```

```
if p not in R(lamda(x)):
               print "MISSUSE";
28
      #intialization
29
      for x in Global(c):
30
          M[x] = Md[x]
31
          lamda[x] = lamdad[x]
      for x in ((VA(c) - Global(c)) | set(PC)):
          M[x] = 0
34
          lamda[x] = (p, S, set([p]))
36
  def VA(data):
      return set(parse_variables(data))
38
39
40 def Global(data): #discard all whiles, ifs and functions -> then remaining
     code will have only globals.
      str = ""
41
      length = len(data)
42
      i = 0
43
      while i < length - 1:
44
          # checking for keyword
          if parse_keyword(i, data) == "FunctionDef":
46
               i += 11
               i = parse_parenthesis(i, data)[1]
48
          elif parse_keyword(i, data) == "Expr(":
               i += 4
50
               i = parse_parenthesis(i, data)[1]
51
          elif parse_keyword(i, data) == "AugAssign":
52
               i += 9
53
               i = parse_parenthesis(i, data)[1]
          elif parse_keyword(i, data) == "If":
               i += 2
               i = parse_if(i, data)[1]
          elif parse_keyword(i, data) == "While":
59
               i += 5
               i = parse_parenthesis(i, data)[1]
          else:
               str += data[i]
63
          i += 1
64
      return set(parse_variables(str))
65
66
67 def parseTestVariables (data):
```

```
test_index = [ m. start() for m in re.finditer('test=', data) ]
68
       11 = []
69
       for it in test_index:
70
           11 += parse_variables(parse_next_parenthesis(it, str)[0])
71
       return 11
72
  def AccessedGlobal(data):
       #(i)right-hand side of assignment
75
       #(ii) condition of branching/iteration
       #(iii) return
       11 = []
       length = len(data)
       i = 0
80
       while i < length - 1:
81
           # checking for keyword
82
           if parse_keyword(i, data) == "AugAssign":
               i += 9
84
               tmp = parse_parenthesis(i, data)
85
               11 += parse_variables(tmp[0])
86
               i = tmp[1]
           elif parse_keyword(i, data) == "Assign":
88
               i += 6
               tmp = parse_parenthesis(i, data)
90
               ryt = tmp[0]. split("value=")[1]
               11 += parse_variables(ryt)
92
               i = tmp[1]
93
           elif parse_keyword(i, data) == "If":
94
               i += 2
95
               tmp = parse_parenthesis(i, data)
               11 += parseTestVariables(tmp[0])
               i = tmp[1]
98
           elif parse_keyword(i, data) == "While":
99
               i += 5
               tmp = parse_parenthesis(i, data)
101
               11 += parseTestVariables(tmp[0])
102
               i = tmp[1]
103
           elif parse_keyword(i,data) == "Return":
               i += 6
105
               tmp = parse_parenthesis(i, data)
106
               11 += parse_variables(tmp[0])
107
               i = tmp[1]
108
           i += 1
109
```

```
return set(11)
110
111
def ModifiedGlobal(data):
       ss = target_of_assignment(data)
113
       modifiedVarList = parse_variables(ss)
114
       return Global(data) & set(modifiedVarList)
  def make_lub_string(llist): # assumption list containing string elemnts
       if len(1list) == 0:
119
           return "Low"
120
       if len(1list) == 1:
121
           return llist[0]
122
       tmp = set(11ist)
123
       uniq_list = list(tmp)
124
       if len(uniq_list) == 1:
125
           return str(uniq_list[0])
126
       ret = ""
127
       ret += uniq_list[0]
128
       i = 1
       while i < len(uniq_list):</pre>
130
           ret += " " + const.oplus + " "
           ret += uniq_list[i]
           i += 1
133
       return ret
134
135
136
  def make_glb_string(llist): # assumption list containing string elemnts
137
       # type: (list) -> string
138
       if len(11ist) == 0:
139
           return "High"
140
       if len(1list) == 1:
141
           return llist[0]
       uniq_list = list(set(llist))
143
       if len(uniq_list) == 1:
           return uniq_list[0]
145
       ret = ""
       ret += uniq_list[0]
147
       i = 1
148
       while i < len(uniq_list):</pre>
149
           ret += " " + const.otime + " "
150
           ret += uniq_list[i]
151
```

```
i += 1
152
       return ret
153
154
155
  def split_through_orelse(if_str):
156
       # find first body word
157
       i = if_str.find("body=[")
158
       i = parse\_square\_br(i + 5, if\_str)[1] + 1
159
       return ["{" + if_str[1:i] + "}", if_str[i + 7:]]
161
  def parse_keyword(i, data):
163
       # checking for
164
       funLen = len("Expr(value=Call(func=Name(id='")
165
       attrLen = len("Expr(value=Call(func=Attribute(value=Name(id='")
166
167
       if i + 6 < len(data) - 1 and data[i:i + 6] == 'Assign':
168
           return "Assign"
169
       if i + 9 < len(data) - 1 and data[i:i + 9] == 'AugAssign':
170
           return "AugAssign"
171
       if i + 2 < len(data) - 1 and data[i:i + 2] == 'If':
172
           return "If"
173
       if i + 5 < len(data) - 1 and data[i:i + 5] == 'While':
174
           return "While"
       if i + 11 < len(data) - 1 and data[i:i + 11] == 'FunctionDef':
176
           return "FunctionDef"
       if i + 6 < len(data) - 1 and data[i:i+6] == "Return":
178
           return "Return"
179
       if i + funLen < len(data) - 1 and data[i:i + funLen] == "Expr(value=
180
      Call(func=Name(id='":
     return "fun_call"
181
       if i + attrLen < len(data) - 1 and data[i:i + attrLen] == "Expr(value
182
      =Call(func=Attribute(value=Name(id='":
           if extract_variavle_name(i+attrLen, data) == 'thread':
183
               return "thread_fun_call"
           else:
185
         return "set_clear_wait"
       return "none"
187
188
  def parse_square_br(i, data):
189
       if data[i] != '[':
190
           print "Error: [ is missing"
191
```

```
return []
192
       ret = "["]
193
       count = 1
194
       i += 1
195
       while count > 0 and i < len(data) - 1:
196
            if data[i] == '[':
197
                count += 1
            if data[i] == ']':
199
                count -= 1
            ret += data[i]
201
            i += 1
202
       return [ret, i]
203
204
205
   def parse_parenthesis(i, data):
206
       # type: (int , string) -> string
207
       if data[i] != '(':
208
            print data[i-4:i+4], data[i]
209
            print "Error: ( is missing"
            return []
       ret = "("
       count = 1
       i += 1
214
       while count > 0 and i < len(data) - 1:
215
            if data[i] == '(':
216
                count += 1
217
            if data[i] == ')':
218
                count -= 1
219
            ret += data[i]
220
            i += 1
       return [ret, i]
222
  def parse_next_parenthesis(i, data):
224
       while i < len(data)-1 and data[i] != '(':
225
            i += 1
226
       if i == len(data)-1:
            print "No parenthesis in string"
228
            return ["",i]
229
       ret = "("
230
       count = 1
       i += 1
232
       while count > 0 and i < len(data) - 1:
233
```

```
if data[i] == '(':
234
                count += 1
235
            if data[i] == ')':
236
                count -= 1
            ret += data[i]
238
            i += 1
239
       return [ret, i]
241
  def extract_variavle_name(startpos, line):
242
       # string -> string
243
       var = ""
       while line[startpos] != "'":
245
            var += line[startpos]
246
            startpos += 1
247
       return var
248
249
  def target_of_assignment(str): # find all targets
250
       # string -> list
251
       targets\_ptrn = r"targets = \setminus [.*?\setminus]"
252
       ctargets_ptrn = re.compile(targets_ptrn)
253
       temp_list = ctargets_ptrn.findall(str)
254
       ret = ''.join(temp_list) # converting to string
255
       return ret
256
258
  def parse_variables(line):
       # type: (string) -> list
260
       # type: (str) -> object
261
       id_index = [m. start() for m in re.finditer('id=', line)]
262
       var list = []
263
       for it in id_index:
264
            vname = extract_variavle_name(it + 4, line)
265
            if vname == "False" or vname == "True":
                continue
267
            var_list.append(vname)
       return var list
269
270
271
       multiple_assign(assign_str, target_id_index, PC):
272
       global line
273
       tmp = assign_str.split("value", 1)
274
       rvalue = parse_variables(tmp[0])
275
```

```
lvalue = parse_variables(tmp[1])
276
       pc_update (PC, lvalue)
277
278
       # printing denning's rule
279
       for it in rvalue:
280
           # left = make_lub_string(dict[key])
281
            if len(lvalue) == 0:
                print "low " + const.lt + " " + it
283
                line += 1
            else:
285
                print make_lub_string(lvalue), const.lt, it
286
                line += 1
287
       return PC[:]
288
289
   def pc_update(PC, list):
290
       #print "pre update", PC
291
       for pc in PC:
292
            pc += 1ist
293
       #print "post update",PC
294
   def label_to_lub(llist):
       if len(1list) == 1:
296
            return llist[0]
       ret = ""
298
       ret += 11ist[0]
299
       i = 1
300
       while i < len(llist):</pre>
301
            ret += " "+const.oplus+" "+ 11ist[i]
302
            i +=1
303
       return ret
304
305
306
  def make_lub_string_label(llist):
307
       if len(1list) == 0:
            return "Low"
309
       if len(1list) == 1:
310
            return label_to_lub(label[llist[0]])
       tmp = set(llist)
       uniq_list = list(tmp)
313
       if len(uniq_list) == 1:
314
            return label_to_lub(label[uniq_list[0]])
315
       ret = ""
316
       ret += label_to_lub(label[uniq_list[0]])
317
```

```
i = 1
318
       while i < len(uniq_list):</pre>
319
           ret += " " + const.oplus + " "
320
           ret += label_to_lub(label[uniq_list[i]])
           i += 1
322
       return ret
  def convertToLabelList(llist):
325
       global label
326
       if len(1list) == 0:
327
           return llist
       ret = []
329
       for it in llist:
330
           if it not in label:
                ret += [it]
                continue
333
           ret += label[it]
334
       return list(set(ret))
336
  def assign_denning(G, assign_str, PC): # applying dennig's model on
      assignments
338
       #pdb.set_trace()
       global line
339
       global label
       global output
341
       ss = assign_str.split("value")
342
       target_id_index = [m. start() for m in re.finditer('id=', ss[0])]
343
344
       if len(target_id_index) > 1:
345
           return multiple_assign(assign_str, target_id_index, PC)
346
347
       if "id='" in ss[0]:
348
           left = extract_variavle_name(0, ss[0].split("id='")[1])
       else:
350
           left = ['const']
       id_index = [m. start() for m in
352
                re.finditer('id=', ss[1])] # list of starting index of
      variables in right part of string
       rvalue = []
354
       if len(id_index) == 0:
355
           # rvalue.append("low")
356
           pass
357
```

```
else:
358
           for it in id_index:
359
                startpos = it + 4
360
                vname = extract_variavle_name(startpos, ss[1])
361
                """Exclusion of False keyword"""
362
                if vname == "False" or vname == "True":
363
                    continue
                rvalue.append(vname)
365
       #pc_update(PC, rvalue)
367
       ret = ""
       1 = convertToLabelList(rvalue)
369
       #11 = 1 + lambda(pc)
370
       #updating PC
371
       for pc in PC:
372
           pc += 1
373
       if left in G:
374
           for pc in PC:
375
                output.append(make_lub_string(pc)+ " " + const.lt + " " +
376
      make_glb_string(label[left]))
       else:
377
           # update dynamic label, for all x (MNG(BR) \text{ union } \{pc\}) [lambda(x)=1
      +lambda(x)+lambda(pc)
           if left not in label:
                label[left] = []
380
           label[left] += 1
381
           #print label[left], "changed", left
382
       line += 1
383
       return PC[:]
384
385
386
  def augAssign_denning(called_by_fun, fun_global, augAssign_str, PC):
387
       \# i = augAssign_str.find("id=")
       # Su = [extract_variavle_name(i + 4, augAssign_str)]
389
       global output
       Sr = parse_variables(augAssign_str)
391
       pc_update(PC, Sr)
       for pc in PC:
393
           output.append( make_lub_string(pc) + " " + const.lt + " " + Sr[0])
394
       return PC[:]
395
^{397} PCA = []
```

```
398
  def if_denning(G, called_by_fun, fun_global, if_str, rest, PC):
399
       # type: (list, list, string, dict, string) -> print rules
400
       #pdb.set trace()
401
       #print "begin IF",PC
402
       if "orelse=" not in if_str:
403
           # print "termination", if_str
           if if_str[0:2] == "[]": # absence of else part
405
                return []
           else: # handling else part
407
                else\_str = if\_str
408
                continuous\_parse(G,[], ccalled\_by\_fun, fun\_global, else\_str, PC
409
      [:])
                return []
410
411
       tmp = split_through_orelse(if_str)
412
       if_half = tmp[0]
413
       ladder = tmp[1]
414
415
       if if_str[1:5] != "test":
           print "Error test not found in if"
417
       """ extract test = ...() from if half """
419
       i = if_half.find("(")
       tmp = parse_parenthesis(i, if_half)
421
       test_str = tmp[0]
422
       #parent_list += parse_variables(test_str)
423
       i = tmp[1]
424
       pc_update(PC, convertToLabelList(parse_variables(test_str)))
425
426
       """then extract body part and process like normal AST text """
427
       # body processing
428
       body_onward_str = if_half[i:] ### Asumption : Compare string always
      followed by body = [...] imediatly
       # setting i to location of [ in body_str: ,body = [...
       i = body_onward_str.find("[")
431
       body_str = parse_square_br(i, body_onward_str)[0]
      memo = \{ \}
433
       memo2 = \{\}
434
      PC1 = copy._deepcopy_list(PC, memo)
435
       PC2 = copy._deepcopy_list(PC, memo2)
436
      PCA = list(continuous_parse(G,[], called_by_fun, fun_global, body_str,
437
```

```
PC1[:]))
      PCB = list(continuous_parse(G,[], called_by_fun, fun_global, ladder,
438
      PC2[:]))
      #debugPrint (PCA)
439
      #debugPrint (PCB)
440
      #print "end IF", PCB + PCA
441
       return PCB + PCA
443
  def while_denning(iteration, G, called_by_fun, fun_global, while_str, PC):
       # debug print "printing while_str", while_str[6:10]
445
       compare = "()"
       if while_str[6:10] == "Name":
447
           tmp = parse_parenthesis(10, while_str)
448
           compare = tmp[0]
449
           i = tmp[1]
450
       elif while_str[6:10] == "Comp":
451
           tmp = parse_parenthesis(13, while_str)
452
           compare = tmp[0]
453
           i = tmp[1]
454
       pc_update(PC, convertToLabelList(parse_variables(compare)))
       # body processing
456
       body_onward_str = while_str[i:] ### Asumption : Compare string always
      followed by body = [...] imediatly
      # setting i to location of [ in body_str: ,body = [...
       i = body_onward_str.find("[")
459
       body_str = parse_square_br(i, body_onward_str)[0]
460
      memo = \{ \}
461
      memo2 = \{\}
462
      PC1 = copy._deepcopy_list(PC, memo)
463
      PC2 = copy._deepcopy_list(PC, memo2)
464
      PCB = list(continuous_parse(G, [], called_by_fun, fun_global, body_str,
466
       PC2[:]))
       return PC1 + PCB
467
  def set_clear_denning(G, called_by_fun, fun_global, expr_str, PC):
469
      #ASSUMPTION SEMAPHORE VAR IS ALWAYS GLOBAL
       global output
471
       i = expr_str.find("Call(")
472
       i += 4
473
       call_str = parse_parenthesis(i, expr_str)[0]
474
       i = call_str.find("Attribute(")
475
```

```
i += 9
476
       attribute_str = parse_parenthesis(i, call_str)[0]
477
       i = attribute str.find("Name(")
478
       i += 4
479
       # name_str = parse_parenthesis(i, attribute_str)[0]
480
       i = attribute_str.find("id=")
481
       var_name = extract_variavle_name(i + 4, attribute_str)
       i = attribute str.find("attr=")
483
       attr = extract_variavle_name(i + 6, attribute_str)
       if attr == "set" or attr == "clear":
485
           # treat it like AugAssign s0 += 1
           # treat it like AugAssign s0 -= 1
487
           pc_update(PC, convertToLabelList([var_name]))
488
           #print "set clear -> ",PC
489
           for pc in PC:
490
               output.append( make_lub_string(pc)+ " " + const.lt + " " +
491
      make_glb_string(label[var_name])) #label[left]
       elif attr == "wait":
492
           #global_while_list.append(var_name)
493
           pc_update(PC, convertToLabelList([var_name]))
           #print "wait -> ",PC
495
  def extract_Globals(fun_str):
497
       global_index = [m. start() for m in re.finditer("Global\(", fun_str))]
498
       globals = \{\}
499
       for it in global_index:
500
           global_str = parse_parenthesis(it + 6, fun_str)[0]
501
           sq_str = parse_square_br(global_str.find("["), global_str)[0]
502
           ss = sq_str.strip("["].strip("]")
503
           sslist = ss.split(",")
504
           for it in sslist:
505
               if it == '':
506
                    continue
               globals [(it.strip("'"))] = 1
508
       return globals
509
510
511
  def fun_denning(fun_str,PC):
512
       fun_name = extract_variavle_name(fun_str.find("name=") + 6, fun_str)
513
       fun_globals = extract_Globals(fun_str)
514
       funPC = []
515
      PC = continuous_parse(G, funPC, fun_name, fun_globals, fun_str, PC)
516
```

```
return PC + funPC
517
518
519 # global var for counting
520 \text{ ww} = \text{ww}1 = \text{ww}2 = 1
522 # global while list
  global_while_list = []
524
  def parse_if(i, data):
526
       tmp = parse_parenthesis(i, data)
       i = tmp[1]
528
       if_str = tmp[0]
529
       rest = data[i:]
530
       return [if_str , i , rest]
531
532
  def uniq(1):
533
       11 = []
534
       for it in 1:
535
            11 . append(list(set(it)))
       return 11
537
  fun_hash = \{\}
539
  def duplicateRemoval(PC):
541
       tmpPC = []
542
       for pc in PC:
543
            tmppc = list(set(pc))
544
            tmppc.sort()
545
            tmpPC . append ( tmppc )
546
       tmpPC.sort()
547
       return list(tmpPC for tmpPC, _ in itertools.groupby(tmpPC))
548
550 recCount =0
   def continuous_parse(G, funPC, called_by_fun, fun_global, data, PC):
       #pdb.set_trace()
552
       # type: (object, object) -> object
       global recCount
554
       length = len(data)
555
       i = 0
556
       recCount += 1
557
       debugPrint (recCount)
558
```

```
debugPrint("Continuous_parse before loop:")
559
       while i < length - 1:
560
           #checking for keyword
561
           if parse_keyword(i,data) == "FunctionDef": #skipping all function
562
      definition
               tmp = parse_parenthesis(i+11, data)
563
               i = tmp[1]
           if parse_keyword(i,data) == "thread_fun_call": #parsing function
565
      call used in threads
               tmp = parse_parenthesis(i+4,data)
566
               i = tmp[1]
               if 'start_new_thread' in tmp[0]:
568
                    #print "got thread call"
569
                    fi = tmp[0]. find("args=[")
570
                    funName = extract_variavle_name(fi+15,tmp[0]) #len(args=[
571
      Name (id = ') = 15
                    if funName in fun_hash:
572
                        findex = fun_hash[funName]
573
574
                        print "Function not found but prgram called function"
                    findex += 11
576
                    tmp = parse_parenthesis(findex, data)
                    fun_str = tmp[0]
578
                    tmp = fun_denning(fun_str,[[]])
                   memo = \{ \}
580
                    1s = duplicateRemoval(tmp)
581
                   PC = copy._deepcopy_list(ls, memo)
582
583
           if parse_keyword(i,data) == "fun_call": #parsing functioncalls
584
               tmp = parse_parenthesis(i+4,data)
585
               #print "got fun call"
586
               funName = extract_variavle_name(i + len("Expr(value=Call(func=
587
      Name(id='"),data)
               i = tmp[1]
588
               if funName in fun_hash:
                    findex = fun hash[funName]
590
               else:
                    print "Function not found but prgram called function"
592
               findex += 11
593
               tmp = parse_parenthesis(findex, data)
594
               fun str = tmp[0]
595
               tmp = fun_denning(fun_str,PC)
596
```

```
memo = \{ \}
597
                1s = duplicateRemoval(tmp)
598
               PC = copy._deepcopy_list(ls, memo)
599
600
           if parse_keyword(i,data) == "Return":
601
                debugPrint("Return stmt:")
602
                funPC += PC
           if parse_keyword(i, data) == "set_clear_wait":
604
               tmp = parse_parenthesis(i, data)
606
                expr_str = tmp[0]
                if expr_str.find("'set'") != -1 or expr_str.find("'clear'") !=
608
      -1 or expr_str.find("'wait'") != -1:
                    i = tmp[1]
609
                    set_clear_denning(G, called_by_fun, fun_global, expr_str,
610
      PC)
           if parse_keyword(i, data) == "AugAssign":
611
                i += 9
612
                tmp = parse_parenthesis(i, data)
613
                augAssign_str = tmp[0]
                i = tmp[1]
615
               tmp = augAssign_denning(parent_list[:], global_while_list,
      called_by_fun , fun_global , augAssign_str , PC[:])
               memo = \{ \}
                1s = duplicateRemoval(tmp)
618
               PC = copy._deepcopy_list(ls, memo)
619
620
           if parse_keyword(i, data) == "Assign":
621
                global ww
622
               ww += 1
623
                i += 6
624
                tmp = parse_parenthesis(i, data)
625
                assign_str = tmp[0]
                i = tmp[1]
627
                if "value=Name(id='threading'" in assign_str:
629
                tmp = assign_denning(G, assign_str, PC[:])
               memo = \{ \}
631
                1s = duplicateRemoval(tmp)
632
               PC = copy._deepcopy_list(ls,memo)
633
           elif parse keyword(i, data) == "If":
634
                debugPrint("If stmt:")
635
```

```
global ww1
636
                ww1 += 1
637
                i += 2
638
                tmp = parse_if(i, data)
639
                if_str = tmp[0]
640
                i = tmp[1]
641
                rest = tmp[2]
                tmp = if_denning(G, called_by_fun, fun_global, if_str, rest, PC
643
       [:])
                1s = duplicateRemoval(tmp)
644
                memo = \{ \}
                PC = copy._deepcopy_list(1s, memo)
646
            elif parse_keyword(i, data) == "While":
647
                debugPrint("while stmt:")
648
                global ww2
649
                ww2 += 1
                i += 5
651
                tmp = parse_parenthesis(i, data)
652
                while_str = tmp[0]
653
                i = tmp[1]
                it = 1
655
                while( it <= iteration ):</pre>
                     memo = \{ \}
657
                     lastPC = copy._deepcopy_list(PC, memo)
658
                     tmp = duplicateRemoval(PC)
659
                     memo = \{ \}
660
                     PC = copy._deepcopy_list(tmp, memo)
661
                     #print "### While Iteration:", it
662
                     tmp = while_denning(it, G, called_by_fun, fun_global,
663
       while_str, PC[:])
                     1s = duplicateRemoval(tmp)
                     memo = \{ \}
665
                     PC = copy._deepcopy_list(ls, memo)
                     #print lastPC," | -- | ", PC
667
                     #print "-> PC:", PC
                     if lastPC == PC:
669
                         #print "Saturation point of loop!"
670
                         break
671
                     it += 1
672
            i += 1
673
       debugPrint (recCount)
674
       recCount -= 1;
675
```

```
debugPrint("END continuous parse:")
676
      return PC[:]
677
678
  def process_label_fille(filename):
679
      \# a = x, y, z
680
      label = \{\}
681
      with open(filename, "r") as inputfile:
          for line in inputfile:
683
              11 = "".join(line.split())
              tmp = 11. split("=")
685
              var = tmp[0]
686
              lineList = tmp[1].split(",")
687
              label[var] = lineList
688
      return label
689
  ############ main
      with open(sys.argv[1], "r") as inputfile:
      # data = inputfile.read().replace('\n', '').replace(' ','')
      data = "".join(inputfile.read().split())
695 11ist = []
696 dummy = []
su_sr_list = []
698 \quad 1ine = 0
699 PC = [[]]
700 \, label = \{\}
701 label = process_label_fille(sys.argv[2])
gg = []
output = []
for key in label:
      gg.append(key)
706 G = set(gg)
fun_call_index = [m. start() for m in re.finditer("FunctionDef", data)]
  for index in fun_call_index:
      varName = extract_variavle_name(index+len("FunctionDef(name='"), data)
      fun_hash[varName] = index
_{712} #G = Global (data)
713 continuous_parse(G, "", [], dummy, data, PC)
714
for it in list (unique_everseen (output)):
  print it
```

Appendix E

Python Script 2: Constraint Checker

```
"""It takes two files as input First: constraint file Second: label file
2 """Format of constraint file: no of constraint
                                   x + y \ll z
                                   a \le b
 """Format of label file: u = [['s1'], [x,y,...], [a,b,...]] single qute'
      are optional
                              v = [['s2'], [x,y,...], [a,b,...]]
      0.00
9
11 import sys
12
if len(sys.argv) != 3:
      print "Error: wrong parameter (constraint file labelfile)"
      exit(0)
def process_constraint_file():
      cons = []
18
      with open(sys.argv[1], "r") as inputfile:
20
          constraints = -1
21
          for line in inputfile:
22
              line = "".join(line.split())
              if line == "":
                  break
25
              if i == 1:
```

```
constraints = int(line)
                    i += 1
28
                    continue
29
               tmp = line.split("<=")
30
               left = tmp[0]
31
               right = tmp[1]
               left_list = left.split("+")
               cons.append([left_list , right])
34
      return cons
36
  def process_label_file():
      labels = \{\}
38
      with open(sys.argv[2], "r") as inputfile:
39
           for line in inputfile:
40
               line = "".join(line.split())
41
               line = "".join(line.split("'"))
               if line == "":
43
                   break
               tmp = line.split("=")
45
               left = tmp[0]
               right = tmp[1]
47
               tmp_s = right.strip("[").strip("]").split("],[")
               owner = tmp_s[0]
49
               first_list = tmp_s[1].split(",")
               second_list = tmp_s[2].split(",")
51
               labels[left] = [set(first_list), set(second_list), owner]
52
      return labels
53
54
  def join(label1, label2):
56
      R = label1[0].intersection(label2[0])
57
      W = label1[1] | label2[0]
58
      return [R,W]
60
  def can_flow(label1, label2):
      return label1 [0]. issuperset (label2 [0]) and label1 [1]. issubset (label2
62
      [1])
63
  def sat (cons, labels):
64
      for constraint in cons:
65
           tmp = set(constraint[0])
66
           left_list = list(tmp)
```

```
right = constraint[1]
68
           for it in left_list:
69
               if not can_flow(labels[it], labels[right]):
70
                   return False
71
       return True
72
  def can_perform(subject, constraint, labels):
       if subject not in labels [constraint [1]][1]:
75
           return False
      tmp = set(constraint[0])
       left_list = list(tmp)
       for it in left_list:
           if subject not in labels [it][0]:
80
               return False
81
       return True
82
  def can_perform_set(cons, labels):
84
      # type: (list, dict{set,set,"str"}) -> set
85
       s = labels [cons [0][1]][1]
86
       for constraint in cons:
          tmp = set(constraint[0])
88
           left_list = list(tmp)
           right = constraint[1]
90
           for it in left_list:
               s = s.intersection(labels[it][0])
92
               if len(s) == 0:
93
                   return s
94
           s = s.intersection(labels[right][1])
95
       return s
96
97
  def can_perform_all(subject, cons, labels):
98
       for constraint in cons:
99
           if can_perform(subject, constraint, labels) == False:
               return False
101
       return True
102
103
  print """ Enter Choice
105
            1. Satisfied or not
106
            2. Can given subject perform all constraints
107
            3. Checking existance of subjects who follows constraints """
108
choice = raw_input()
```

```
cons = process_constraint_file()
labels = process_label_file()

if choice == '1':
    print sat(cons, labels)

elif choice == '2':
    print "Enter subject"

subject = raw_input()
    print can_perform_all(subject, cons, labels)

elif choice == '3':
    print can_perform_set(cons, labels)
```

Appendix F

Copy Programs

```
2 #Procedure copy1
y \#z=0
4 \# y = 0
_{5} if x == 0:
  z=1
7 if z = 0:
 y=1
# Procedure copy2
12 z = 1
y = -1
while z == 1:
y = y + 1
 if y == 0:
z = x
 else:
  z = 0
22 #copy3 synchronization flow
23 import thread
24 import time
25 import threading
27 #x=7
y = 6
```

116 Copy Programs

```
29 #def copy3(x,y): # copy x to y
s0 = threading.Event()
s1 = threading.Event()
 def thread1():
33
      global x
34
      if x==0:
          s0.set()
36
      else:
          s1. set()
38
  def thread2():
40
      global y
41
      s0.wait()
42
      s0.clear()
43
      y=1
44
      s1. set()
45
46
  def thread3():
47
      global y
      s1. wait()
49
      s1.clear()
      y=0
51
      s0.set()
53
 try:
      thread.start_new_thread(thread1,())
55
      thread.start_new_thread(thread2,())
      thread.start_new_thread(thread3,())
57
  except:
      print "Error: unable to start thread"
59
62 #copy4 Global flow in concurrent programs
63 import thread
64 import time
 import threading
66
67
def thread1():
      global x
69
      global e0
```

```
global e1
71
      if x==0:
72
          e0 = False
      else:
74
          e1 = False
75
76
  def thread2():
      global e0
78
      global e1
      global y
80
      while e0:
81
          pass
82
83
      y = 1
84
      e1 = False
85
86
  def thread3():
87
      global e1
88
      global e0
89
      global y
      while e1:
91
92
          pass
      y = 0
93
      e0 = False
95
96 try:
      thread.start_new_thread(thread1,())
97
      thread.start_new_thread(thread2,())
98
      thread.start_new_thread(thread3,())
99
  except:
100
       print "Error: unable to start thread"
101
102
103
104 #copy5
y = 0
while x==0:
   pass
y = 1
111 #copy 6
112 z = 0
```

118 Copy Programs

References

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