

Q-21

As per code the destination and source vectors for each of 3 threads are.

Thread-1

 $v\_dest \rightarrow v_1$  $v\_src \rightarrow v_2$ 

Thread-3

 $v\_dest \rightarrow v_3$  $v\_src \rightarrow v_1$ 

Thread-2

 $v\_dest \rightarrow v_2$  $v\_src \rightarrow v_3$ 

Define execution as below.

$T_1$   $T_2$   $T_3$   
 $\text{pth\_mutex\_lock}(v\_dest \rightarrow \text{lock})$

Let execution order be like. for vector\_add() by all threads

- ①  $\text{pthread\_mutex\_lock}(v\_dest \rightarrow \text{lock})$  // by  $T_1$  so  $v_1$  lock.  
 $\downarrow$  interrupt & switch to  $T_2$
- ②  $\text{pthread\_mutex\_lock}(v\_dest \rightarrow \text{lock})$  // by  $T_2$  so  $v_2$  lock.  
 $\downarrow$  interrupt & switch to  $T_3$
- ③  $\text{pthread\_mutex\_lock}(v\_dest \rightarrow \text{lock})$  // by  $T_3$  so  $v_3$  lock.

If the above order of execution is done then all  $v_1, v_2, v_3$  as destination are locked by  $T_1, T_2, T_3$  threads.

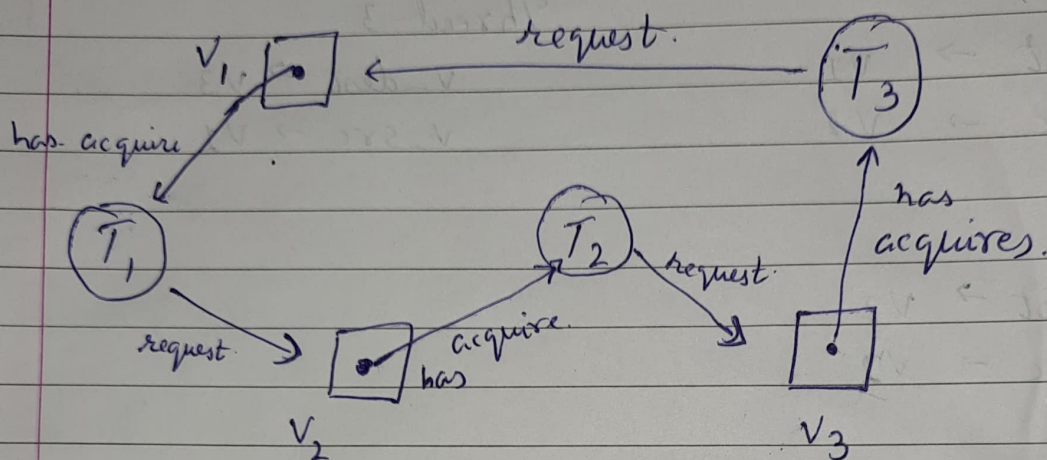
Now in order to lock any  $v_1, v_2, v_3$  no thread can acquire lock as its already taken by other thread. So this situation results in deadlock.



The Resource Allocation Graph is defined as.

$T_1, T_2, T_3 \Rightarrow$  threads.

$V_1, V_2, V_3 \Rightarrow$  Resources (with only one instance)



Here from above can be seen that there is the circular wait in resource allocation graph which shows existence of deadlock in the code.