

# **CS305**

# **Computer Architecture**

## **Introduction to Pipelining**

Bhaskaran Raman  
Room 406, KR Building  
Department of CSE, IIT Bombay

<http://www.cse.iitb.ac.in/~br>

# Some Fun Videos to Watch

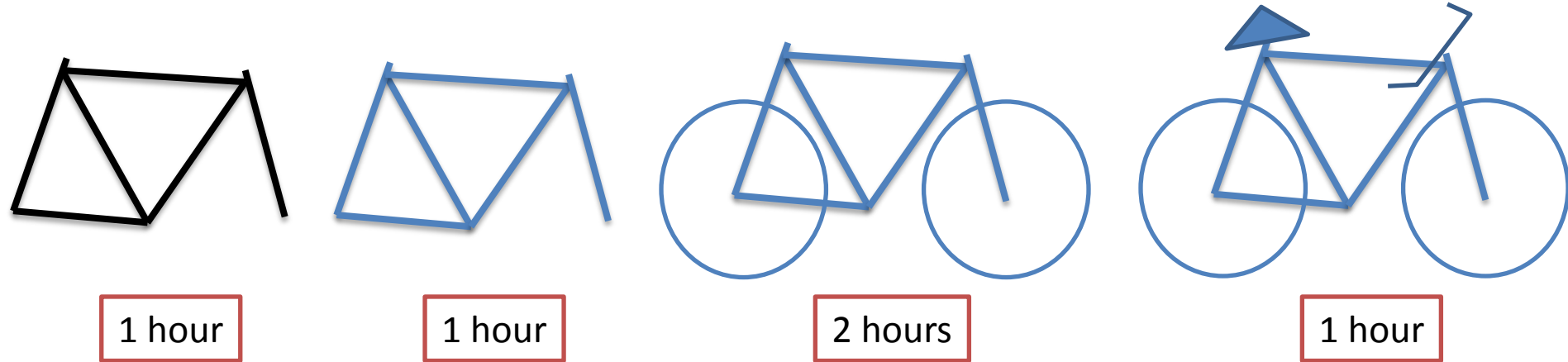
- Charlie Chaplin's critique of industrialization
  - [https://www.youtube.com/watch?v=a-FbVF1x1\\_U](https://www.youtube.com/watch?v=a-FbVF1x1_U)
- Modern Times, Assembly Line with Charlie Chaplin
  - <https://www.youtube.com/watch?v=QdwH84AT5fU>
- I Love Lucy – The Candy Wrapping Job
  - <https://www.youtube.com/watch?v=Clr6Zyjj7iI>

# Lessons Learnt

- Pipelining is natural (at least for machines)
- Overall speed only as good as the slowest unit
- Non-uniformity is bad for pipeline
- Exceptions are really bad

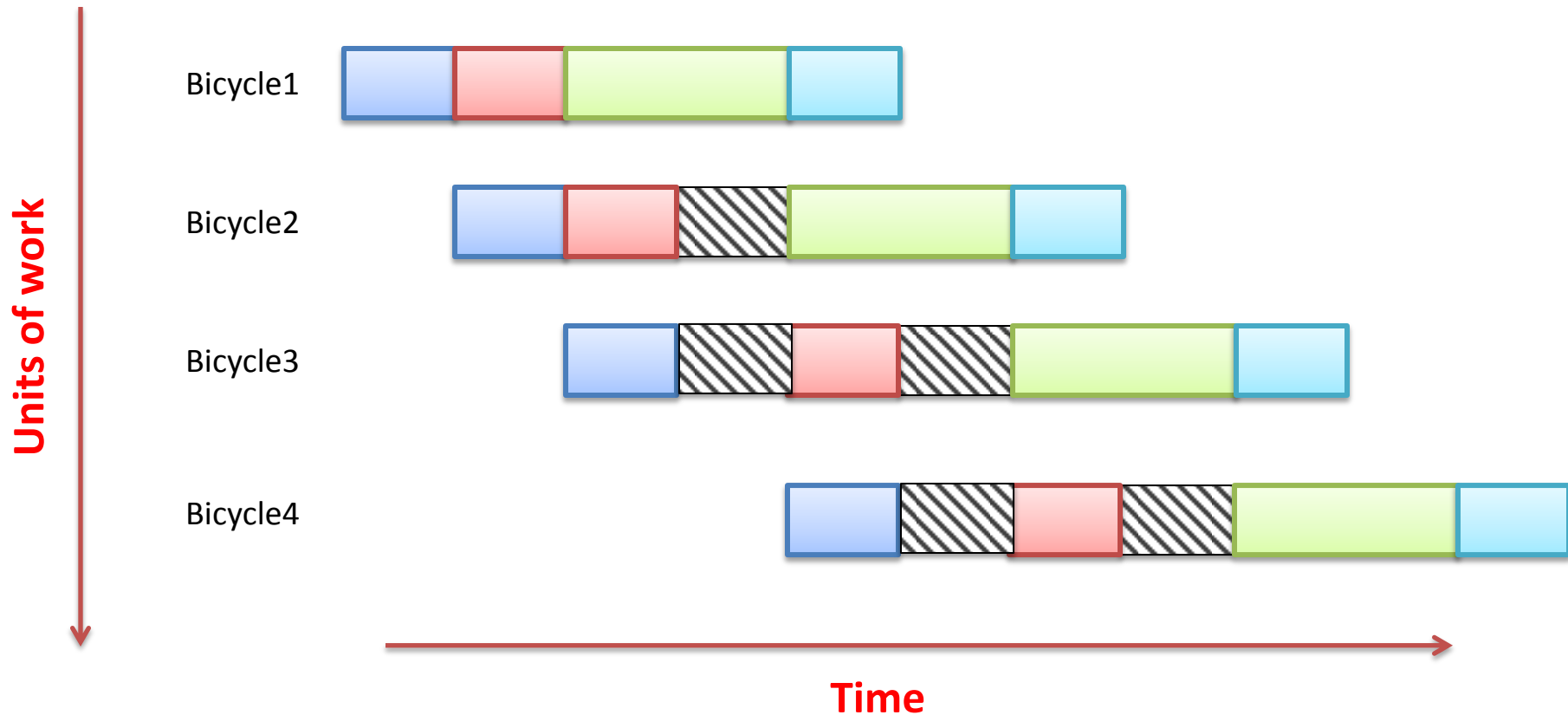
# Example Task: Bicycle Manufacture

Make frame → Paint frame → Fit wheels → Fit accessories

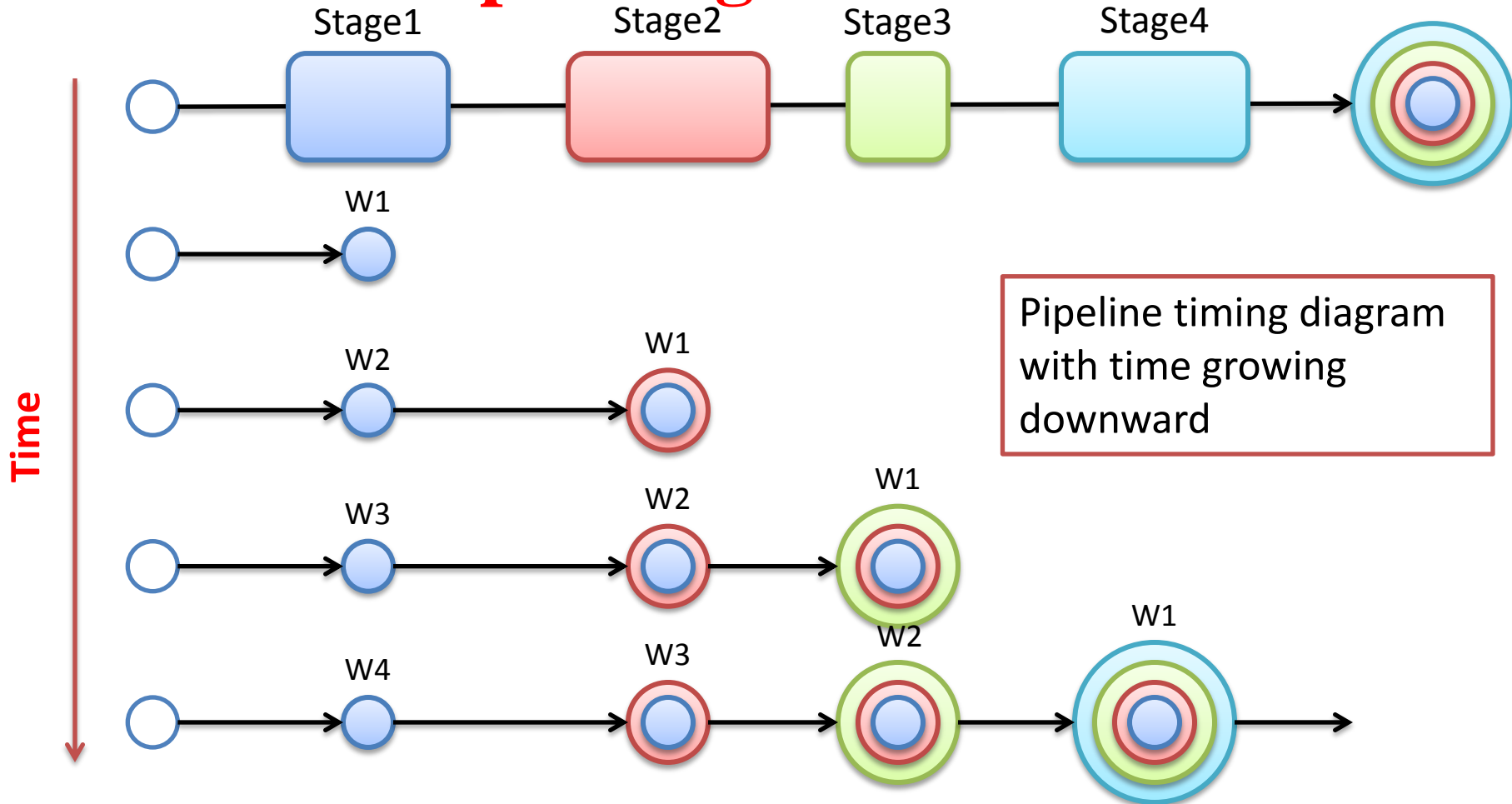


- Time to make 1 bicycle?
- Time to make 10 bicycles?
  - With pipelining? Without pipelining?
  - Pipelining:
    - “Make frame” for 2<sup>nd</sup> bicycle, when “Paint frame” for 1<sup>st</sup> bicycle
    - “Make frame” for 3<sup>rd</sup> bicycle, “Paint frame” for 2<sup>nd</sup> bicycle, when “Fit wheels” for 1<sup>st</sup> cycle
    - And so on...

# Pipeline Timing Diagram

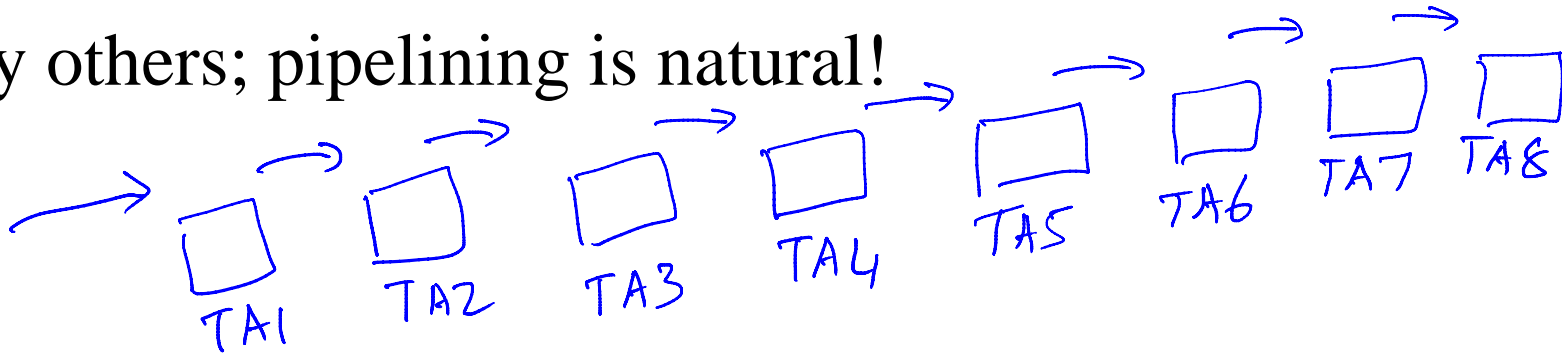
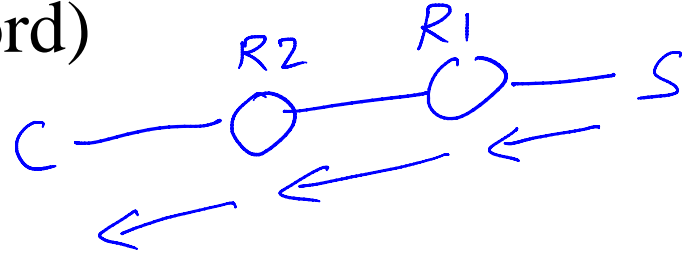


# Pipelining in Abstract



# Other Examples of Pipelining

- Water pipeline (origin of word)
- Network packets
- Course project evaluation
- Many others; pipelining is natural!



# Pipeline Speedup

- Time taken to manufacture  $N$  bicycles
  - Without pipelining:  $5N$
  - With pipelining:  $5 + (N-1) \times 2$
  - Speedup =  $5/2$  for large  $N$
- Suppose “Fit wheels” also takes 1 hour?
  - Speedup =  $4N/(4+N-1) = 4$  for large  $N$
- Suppose “Fit wheels” takes 2 hours, but is broken into “Fit front wheel”, “Fit rear wheel” each of 1 hour?
  - Speedup =  $5N/(5+N-1) = 5$  for large  $N$



# Pipeline Speedup (continued)

- In general,  $k$  stages, each of time  $t$
- Time for  $N$  units:
  - Without pipeline:  $N \times kt$
  - With pipelining:  $kt + (N-1)t$ 
    - Pipeline setup time = time to fill pipeline = time to first output
    - Is a significant component, for small  $N$
  - Speedup =  $k$  for large  $N$
  - This is the *ideal* speedup

# Ideal Pipeline Speedup

- Ideal pipeline speedup = number of pipeline stages
- Necessary conditions:
  - All stages are of equal length in time
  - Enough hardware/hands: e.g. same spanner/person/machine cannot be used for “fit wheels” and “fit accessories”
  - Many, many more conditions (stated later)

# Latency versus Throughput

- Latency: from perspective of single unit of work
- Throughput: rate at which work completes

# Pipelining in MIPS

Stage-1: Instruction fetch,  $PC=PC+4$  (all instructions)

IF

Stage-2: Reg read, branch target computation (all instructions)

ID

Stage-3: ALU operation (reg-reg)  
Memory address computation (lw, sw)  
Branch condition computation (beq)

EX

Stage-4: Memory access (lw, sw)  
Reg write back (reg-reg)

MEM

Stage-5: Reg write back (lw)

WB



# Pipelining for Sequence of lw

lw \$t0, 0(\$sp)



lw \$t1, 4(\$sp)



lw \$t2, 8(\$sp)



Such pipelined execution of MIPS instructions is *potentially* possible

# Pipelining for Sequence with Other Instructions

`add $t0, $t1, $t2`



`ori $t1, $s0, 15`



`lw $t2, 8($sp)`



Effect of moving WB to stage-5 of reg-reg ?

Latency of those instructions increases

Instruction throughput not affected!

# Summary

- Pipelining a natural idea for speed-up
- Latency versus throughput
- MIPS: uniformity of instructions allows pipelining
- Issues with pipelining: hazards
  - Charlie Chaplin wants to scratch under his arm, take a break
  - Gets slowed down by buzzing bee
  - Pipeline stops → startup delays