

An Introduction to Programming through C++

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Lecture 7.1

Ch. 14: Arrays

Mark display variation

- Roll numbers are not in range 1 .. 100, but a larger range, e.g. 170010022.
- Marklist = 100 pairs of numbers: (rollno, marks),
- Teacher must enter roll number, marks into the computer.
- Later:
 - Students arrive and each student types in roll number r .
 - Program must print out marks if r is valid roll number.
 - If r is -1, then stop.

Program idea:

- Use an additional array for storing roll numbers.
- Store i th roll number in `rollno[i]` and i th mark into `marks[i]`.
- When students arrive: Examine each element of the roll number array and see if it equals r . If so print corresponding marks.

The program

```
int rollno[100]; double marks[100];  
for(int i=0; i<100; i++) cin >> rollno[i] >> marks[i];  
  
while(true){  
    int r; cin >> r;    // read in query roll number  
    if(r == -1) break;  
    bool found = false;  
    for(int i=0; i<100; i++){  
        if(rollno[i] == r){  
            cout << marks[i] << endl;  
            found = true;  
            break;  
        }  
    }  
    if(!found) cout << "Roll number not found.\n";  
}
```

Demo

- `generalRollNos.cpp`

Exercise

- Modify the program so that there are marks for two subjects.

What we discussed

- If the roll numbers are not in the range $0..N-1$ where N is the number of students, then we need to store the roll numbers also.
- Searching through the array to find data corresponding to a “key” is a common idiom.
 - Use of a “found” variable is a natural strategy

Next: Polynomial multiplication

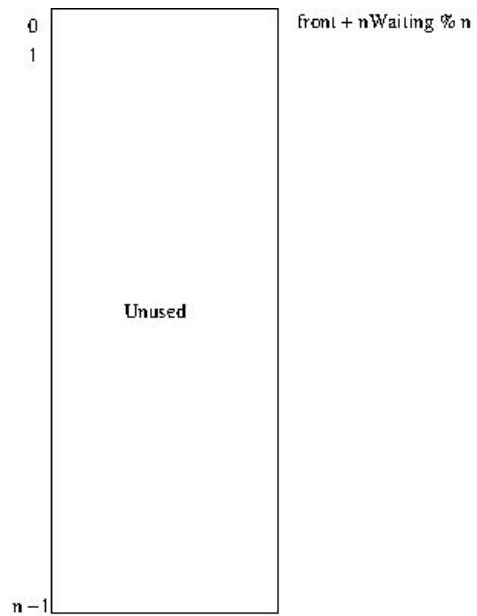


Taxi Dispatch: Idea 2

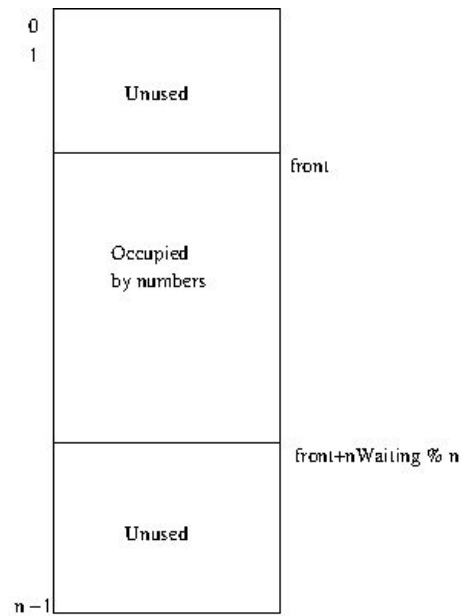
- Previous program seems to be copying data a lot...
 - Can we avoid the copying?
- Emulate what might happen without computers.
 - Names written on blackboard.
 - Arriving driver IDs written top to bottom.
 - When board bottom reached, begin from top if drivers have left.
 - Think of **driverID** as a circular array.
 - “Next” position after **driverID[n-1]** (bottom of board) is **driverID[0]** (top of board).

Invariants

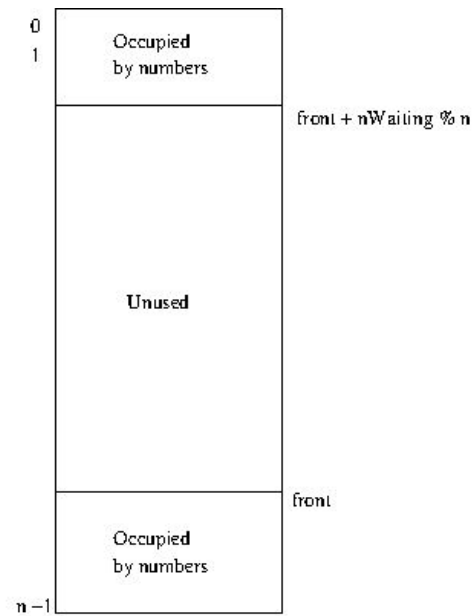
- **$nWaiting$** = number of waiting drivers. **$0 \leq nWaiting \leq n$**
- New variable **$front$** : position of earliest unassigned driver
= **0** initially.
 $0 \leq front < n$
- First valid driver ID is at **$driverID[front]$**
- Where is the next? **$driverID[front+1]$** ?
 - If **$front = n-1$** , then **$front+1 = n$** , which is not a valid index.
 - Next waiting driver is **$driverID[(front+1) \% n]$**
 - Last waiting driver id is: **$driverID[(front+nWaiting-1) \% n]$**
- Note that **$\%$** provides the effect of “wrapping around”.



(a) At the beginning



(b) After some time



(a) After more time

Processing driver arrival

```
if(nWaiting == n)  
  cout << "Queue full.\n";  
else{  
  int d; cin >> d;  
  driverID[(front+nWaiting) % n] = d;  
  nWaiting ++;  
}
```

```
// front + nWaiting % n : index of  
// empty position after end of queue.
```

Processing Customer Arrival

```
if(nWaiting == 0)  
    cout << "Try later.\n";  
else{  
    cout << "Assigning " <<  
        driverID[front] << endl;  
    front = (front + 1) % n;  
    nWaiting --;  
}
```

What we have discussed

- Different idea for solving taxi dispatch problem.
- New idea is better, copying of elements of ***driverID*** is avoided.
- Exercise: verify that the invariants indeed remain true after each customer or driver arrival.

Next: Disk intersection problem



Disk intersection problem

- Given a collection of n disks in the plane, decide if any two intersect.

Algorithm: check all possible pairs.

- Let i = index of smaller numbered disk.
- So i must range from 0 to $n-2$
- For each i , the other disk must range from $i+1$ to $n-1$.

Intersection check:

- Distance between centers must be less than sum of the radii.

Program

```
const int n=5;
double centerx[n], centery[n], radius[n];
for(int i=0; i<n; i++) cin >> centerx[i] >> centery[i] >> radius[i];

bool intersect = false;
for(int i=0; i<n-1; i++){
    for(int j=i+1; j<n; j++)
        if(pow(centerx[i]-centerx[j],2)+pow(centery[i]-centery[j],2) < pow(radius[i]+radius[j], 2)){
            intersect = true;
            break;
        }
}
cout << intersect << endl;
```

What we have discussed

- Given a set of objects, how do you go over all pairs.
- Exercise: Can you generalize the idea?
 - Write a program to check if any 3 of a given set of n points in the plane are collinear.

Next: arrays of graphics objects, conclusion of lecture sequence



Arrays of graphics objects

- We can create multiple graphics objects by using arrays.

Turtle t[3];

t[1].right(120);

A program

```
int main(){  
    initCanvas();  
    Turtle t[3];  
  
    t[1].left(120);  
    t[2].left(240);  
  
    repeat(8){  
        for(int i=0; i<3; i++) t[i].forward(100);  
        for(int i=0; i<3; i++) t[i].left(360.0/8);  
    }  
    wait(5);  
}
```

Demo

- 3poly.cpp

Arrays: concluding remarks

- Way to store many objects of the same type in memory, without having to separately define a variable for each.
- Index: used to choose an element.
 - Must be at least zero and at most array length $- 1$.
 - Can be an expression
- Index may sometimes play an active role:
 - When roll numbers are consecutive from 0 to $N-1$
 - Polynomial representation: i th coefficient stored in i th element
 - Taxi dispatch: indicates arrival order
- Index may not have significance:
 - When roll numbers took on arbitrary values
- Indexing into an array happens very fast, independent of how many elements are present in the array.

More remarks

Several idioms used with arrays:

- Scan through all elements to find a matching element (“linear search”)
- Do some calculation which determines the index where something is to be done (“histogram”)
- Use the array like a queue of elements which wait for something to happen, after which they leave the queue (“taxi dispatch”)

