Postgres-XC PG session #3

Michael PAQUIER Paris, 2012/02/02

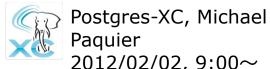


Agenda

- Self-introduction
- Highlights of Postgres-XC
- Core architecture overview
- Performance
- High-availability
- Release status

Self-introduction

- Michael Paquier, 27 years old.
- Based in Tokyo since 2009
- Working for NTT Data Intellilink
 - 500 employees
 - Website http://www.intellilink.co.jp
 - Company of NTT Data Group: 55,000 employees
- Working on DB system support mainly PostgreSQL
- PostgreSQL community member
 - pgbench shell-related features of 9.0
 - 2PC studies
- Core developer of Postgres-XC
- Other information:
 - Email: michael.paquier@gmail.com
 - Twitter: @michaelpq
 - Blog: http://michael.otacoo.com



Highlights

Highlights - Postgres-XC

- Cluster software focused on write-scalability
- Based on PostgreSQL
 - world's most advanced open source database
 - PostgreSQL license
- Same client APIs as PostgreSQL
 - Ease of application migration from existing PostgreSQL deployment
 - Same drivers, same front end, same SQL queries
- Licensing
 - PostgreSQL license (more or less BSD)
 - Free to use, modify and redistribute for commercial purposes

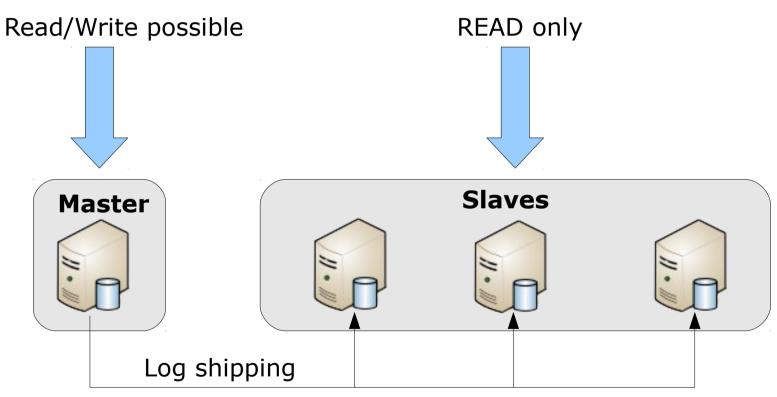
History

- Started through a collaboration between EnterpriseDB and NTT Open Source Software Center in January 2009
- Goal to build a PostgreSQL based clustering solution which can serve as an alternative to Oracle RAC
- Development is community-based, with resources gathered from NTT and EnterpriseDB
- Licensing terms changed from GPL to PostgreSQL license (same as Postgres) in 2011

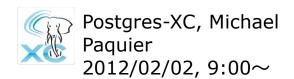
Core architecture

About PostgreSQL 9.1

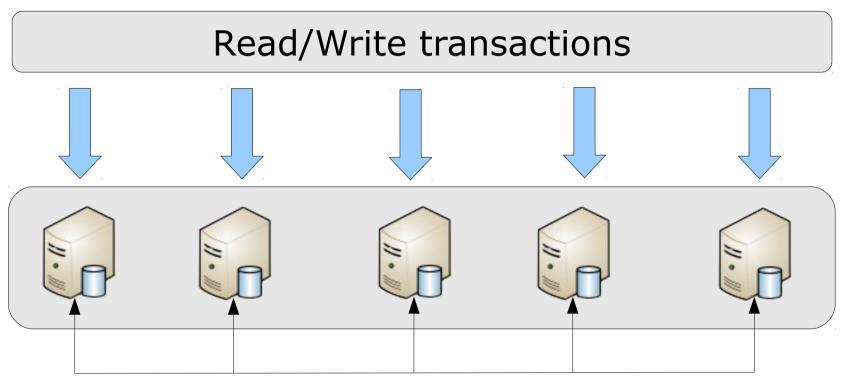
Streaming replication and HOT-Standby



- Asynchronous mode => timestamp view not consistent
- Synchronous mode => timestamp view consistent



And Postgres-XC itself?



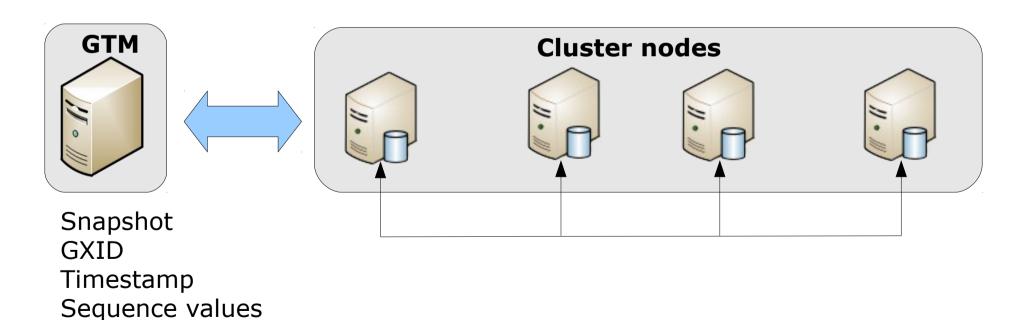
Distributed/replicated tables Same data consistency/transparency Same timestamp view

In short

- Symetric cluster of PostgreSQL
 - No Slave and no Master
 - Read and write scalability
 - Transparent Transaction Management
 - Shared-nothing structure
 - 3 types of nodes: GTM, Coordinator, Datanode

Node types – GTM (1)

- Designed for transparency
- Feeding of MVCC-related data: transaction ID, snapshot
- Cluster follows GTM timeline: timestamp, sequence



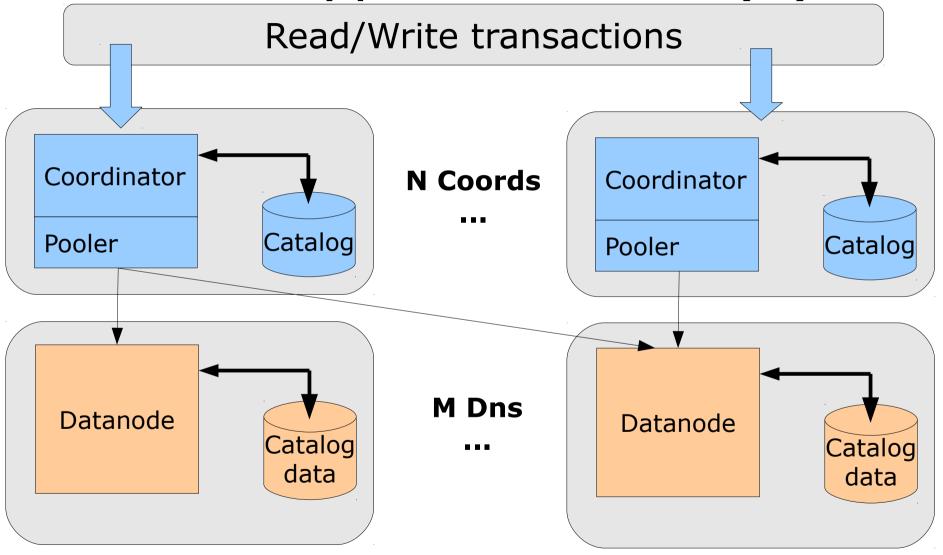
Node types - Coordinator (2)

- Point of contact for the application/client
- Management of remote node data
 - Parse and partially plan the statements
 - Determine the data to be fetched from remote nodes at planning or execution
 - Fetch the required data by issuing queries to dedicated Datanodes
 - Combine and process the data to evaluate the results of the query (if needed)
 - Pass the results to the applications
- Manages two-phase commit
- Stores catalog data: cluster-related information
- Needs and manages space for materializing results from remote nodes
- Binary based on the latest PostgreSQL release

Node types - Datanode (3)

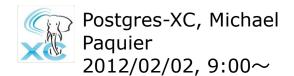
- More or less a PostgreSQL instance (remote node)
- Stores tables and catalogs
- Executes the queries from client Coordinator and returns results to it
- Data nodes can be made fault tolerant by Hot-Standby and Synchronous Replication technologies available of standard PostgreSQL
- Binary same as Coordinator, based on latest PostgreSQL release

Node types - Global (4)

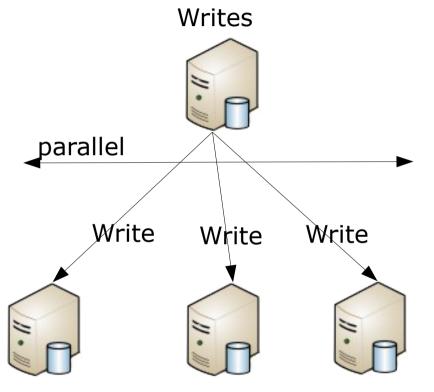


Data management

- Table types
 - Replicated table
 - Each row replicated to Datanodes
 - Statement based replication
 - Distributed table
 - Each row of the table is stored on one datanode, decided by one of following strategies
 - Hash
 - Round Robin
 - Modulo
- Managed by SQL extensions (CREATE TABLE)
- Possible to define subset of nodes



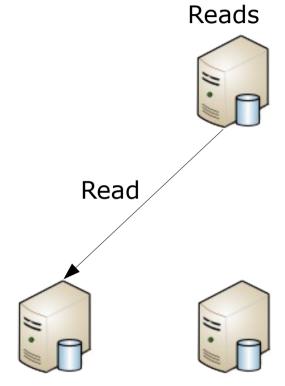
Replicated tables



col1	col2
1	45
2	23
3	34

col1	col2
1	45
2	23
3	34

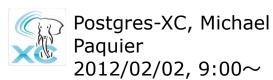
col1	col2
1	45
2	23
3	34



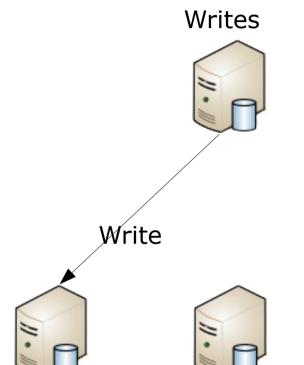
col1	col2
1	45
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3	34

col1	col2
1	45
2	23
3	34

col1	col2
1	45
2	23
3	34



Distributed tables



col1	col2
1	45
2	23
3	34



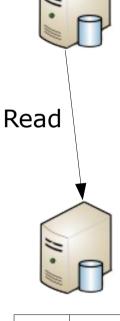
col1	col2
101	654
201	8
301	124



col1	col2
102	26
202	98
302	6



col1	col2
1	45
2	23
3	34



Reads

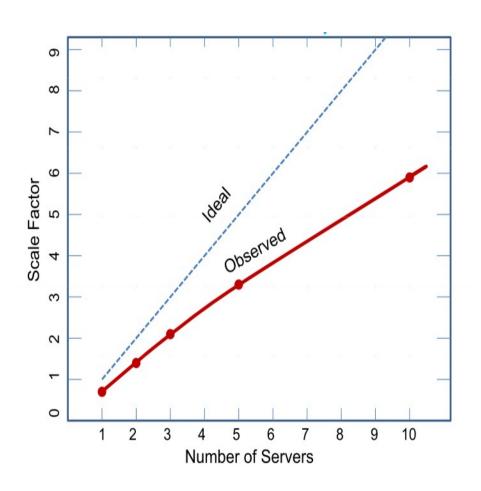
col1	col2
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Performance

Scalability measurements

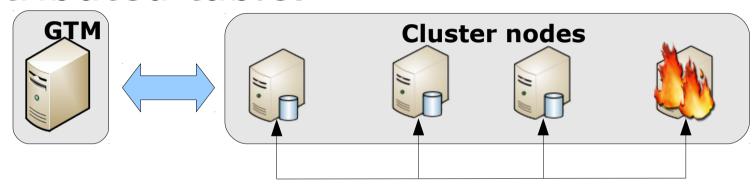


- Tests done with DBT-1 (TPC-W) benchmark with some minor modification to the schema
- 1 server = 1 coordinator + 1 datanode
- Coordinator is CPU bound
- Datanode is I/O bound
- CPU usage
 - Coordinator 30%
 - Datanode 70%

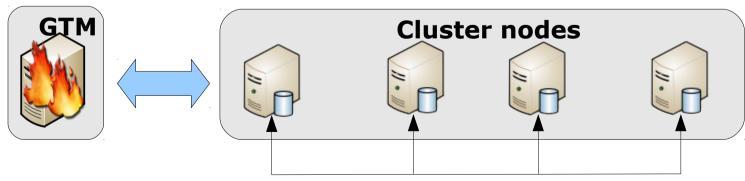
About high-availability

Cluster SPOF problem

 Datanode is a SPOF if it has a portion of distributed table.

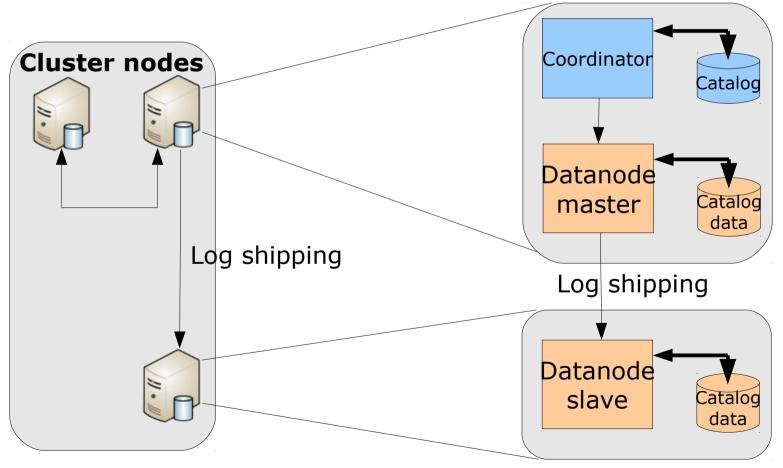


GTM case



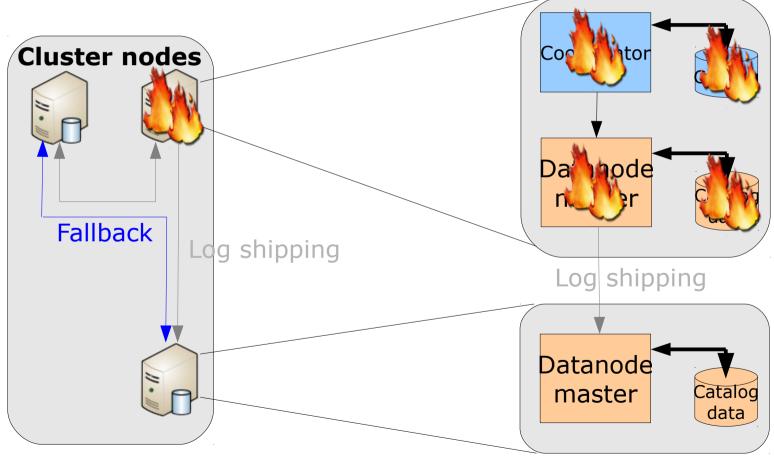
Datanode SPOF resolution (1)

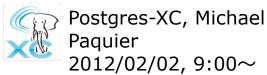
PostgreSQL 9.1 synchronous stream-rep



Datanode SPOF resolution (2)

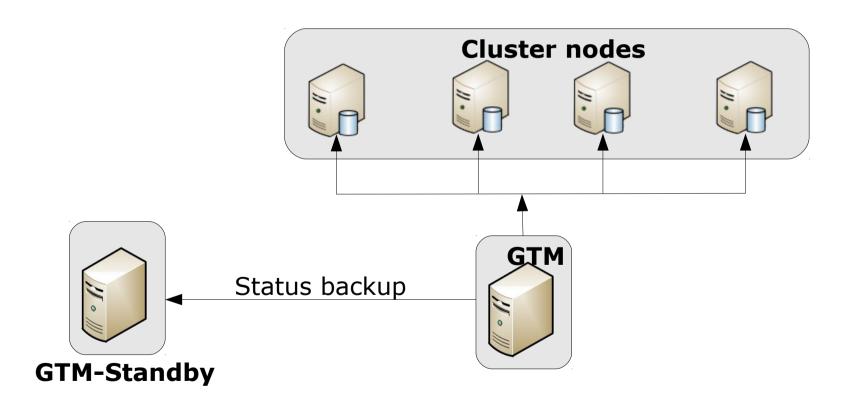
Fallback slave node





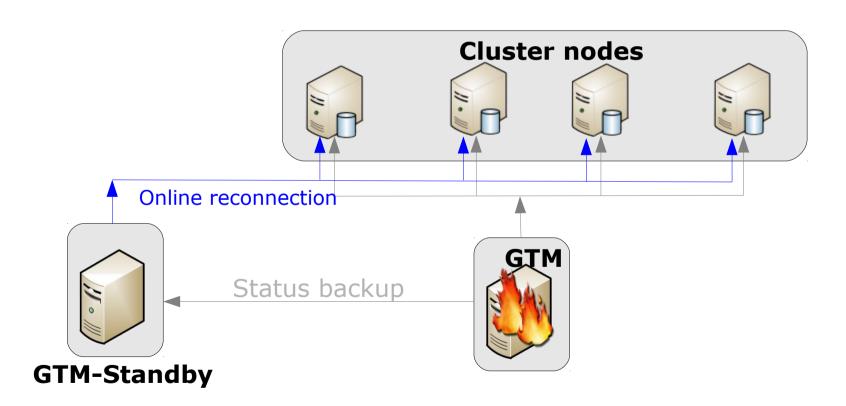
GTM SPOF resolution (1)

Use of a standby for GTM



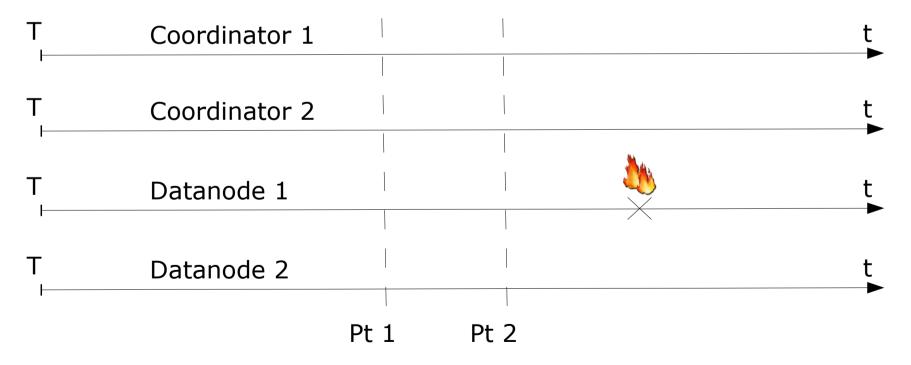
GTM SPOF resolution (2)

Fallback to standby and reconnect nodes



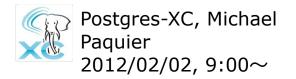
PITR – requirements (1)

- PITR, Point in-time recovery
 - Rollback the database to a given past state
 - Need consistent points to restore shared-nothing nodes



PITR – requirements (2)

- Transaction status has to be consistent in the cluster
- Each transaction must be either:
 - Committed/Prepared/Aborted/Running on all the involved nodes
 - We must avoid cases where transaction is prepared and committed partially, or prepared and rollbacked partially
- Write record in WALs of all the coordinators and datanodes at a moment when all the transaction statuses are consistents.
- External Application can provide such timing as with BARRIER
 - CREATE BARRIER barrier_id
- BARRIER:
 - Waits that partially committed or aborted transactions commit (2PC)
 - Blocks all transaction commit when running a barrier
 - Still needs a timeout functionality
- When running PITR, specify recovery_target_barrier in recovery.conf



Release status What now and next?

Current functionalities

- Up to 0.9.7 (current release of January 2012)
 - Based on PostgreSQL 9.1
 - 90%-92% of SQL
 - Major DDL/DML (TABLE, ROLE, VIEW, SELECT INTO, DEFAULT values)
 - General select support: support extension
 - HAVING, GROUP BY, ORDER BY, LIMIT, OFFSET, aggregate, window function, etc.
 - SQL-based cluster setting
 - Relation-size functions

About Postgres-XC 1.0

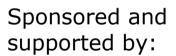
- Release on April 2012 (plan)
- Tablespace
- Triggers
- SERIAL
- Cluster bootstrap
- SELECT FOR UPDATE

After 1.0

- Node addition/removal
 - Move tuples from a node to another node
 - Ex: update of a distribution column
- Online server removal/addition
- Connection balancing between master and slave Datanodes for read transactions.
- SQL/MED, Foreign data wrapper (FDW) integration
- Installation, configuration, operation
- Global deadlock detection (global wait-for-graph)

Project ressources and contacts

- Project home
 - http://postgres-xc.sourceforge.net
- Developer mailing list
 - postgres-xc-developers@lists.sourceforge.net
 - postgres-xc-general@lists.sourceforge.net
- Contacts
 - michael.paquier@gmail.com
 - koichi.szk@gmail.com
- Twitter: @michaelpq
- Blog: http://michael.otacoo.com



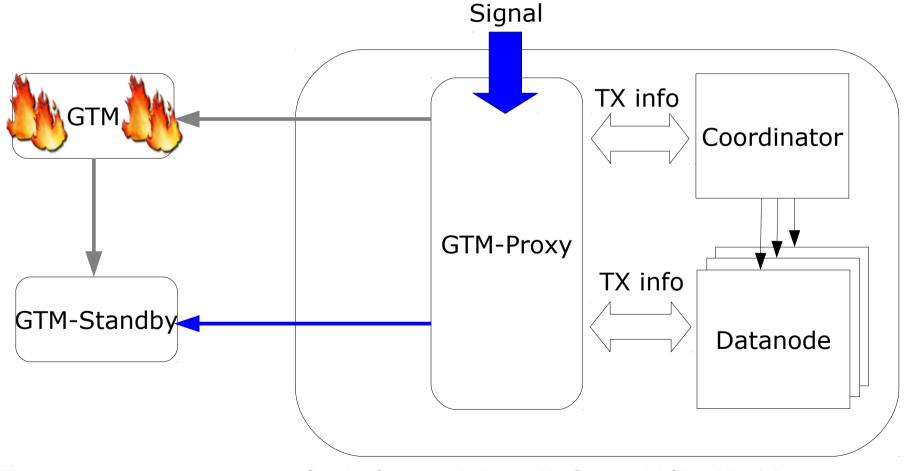


Thanks for your attention. Questions?



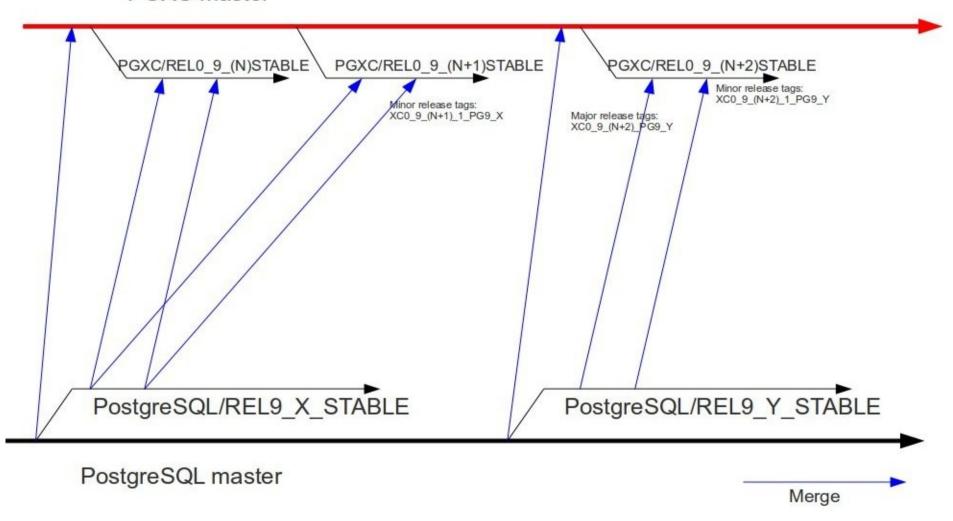
GTM Proxy reconnection

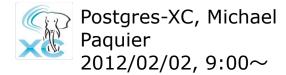
Signal GTM Proxy and reconnect nodes



Release policy

PGXC master

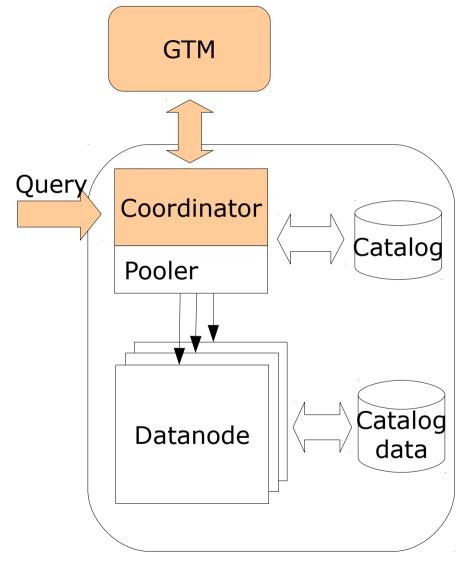




Key algorithm

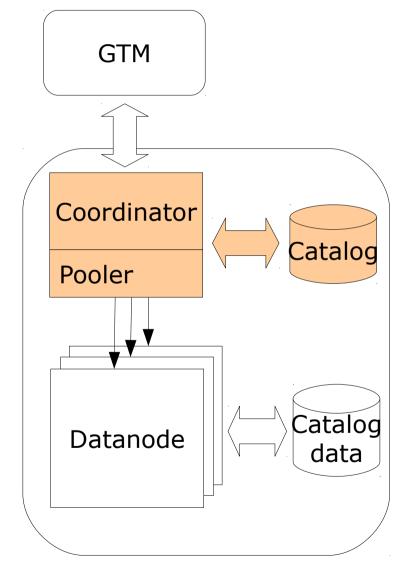
Query algorithm (1)

- Receive query from application
- Get snapshot,
 GXID and
 timestamp from
 GTM



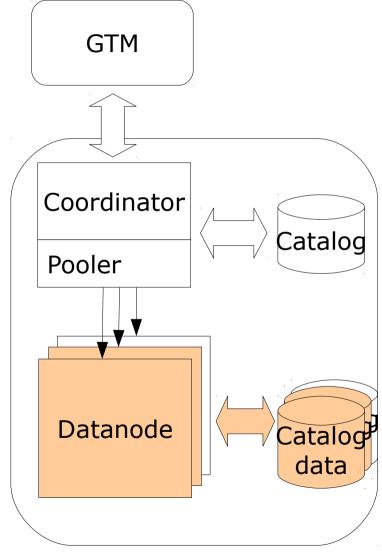
Query algorithm (2)

- Incoming statements: analyzer and rewriter
- Planning: analyze nodes to be involved. Build queries for local nodes (push down if necessary)



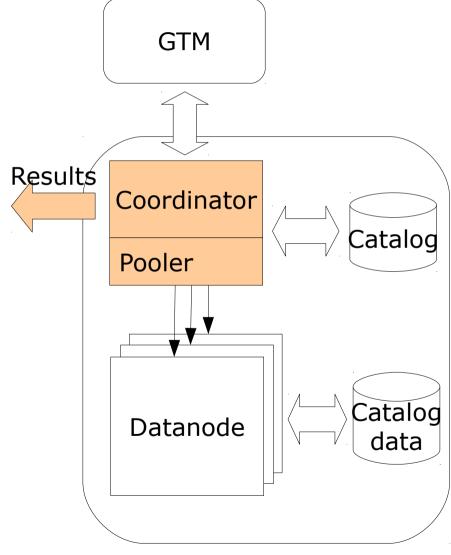
Query algorithm (3)

 Run queries on remote Datanodes and send back results to Coordinator

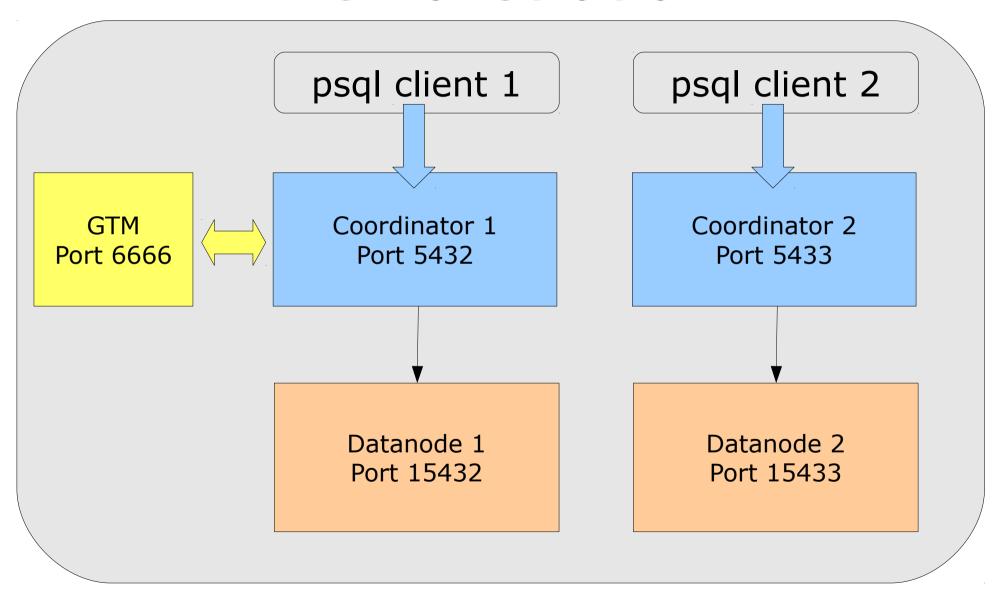


Query algorithm (4)

 Materialize results if necessary and send back to client



Demonstration



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