Introduction to Programming

Spring 2022

Algorithm Design and Recursion

- Searching
- •Recursive Problem Solving
- Sorting Algorithms
- Hard Problems

Sorting Algorithms

- •The basic sorting problem is to take a list and rearrange it so that the values are in
- -increasing (or ascending) order
- -decreasing (or descending) order

- •To start out, pretend you're the computer, and you're given a shuffled stack of index cards, each with a number.
- -How would you put the cards back in order?

- •One simple method is to look through the deck to find the smallest value and place that value at the front of the stack.
- •Then go through, find the next smallest number in the remaining cards, place it behind the smallest card at the front.
- •Repeat, until the stack is in sorted order!
- •Example

- •We already have an algorithm to find the smallest item in a list.
- -As you go through the list, keep track of the smallest one seen so far
- -updating it when you find a smaller one.
- •Now place that smallest number at the front of the list.
- •This sorting algorithm is known as a selection sort.

- •The algorithm has a loop, and each time through the loop the smallest remaining element is selected and moved into its proper position.
 - -For n elements, we find the smallest value and put it in the 0th position.
 - -Then we find the smallest remaining value from position 1 (n-1) and put it into position 1.
 - -The smallest value from position 2 (n-1) goes in position 2.
 - -Etc.

- •When we place a value into its proper position, we need to be sure we don't accidentally lose the value originally stored in that position.
 - -If the smallest item is in position 10, moving it into position 0 involves the assignment:

```
nums[0] = nums[10]
```

-This wipes out the original value in nums [0]!

•We can use simultaneous assignment to swap the values between nums [0] and nums [10]:

```
nums[0], nums[10] = nums[10], nums[0]
```

•Using these ideas, we can implement our algorithm, using variable bottom for the currently filled position, and mp is the location of the smallest remaining value.

Naive Sorting: Selection Sort

- •Rather than remembering the minimum value scanned so far, we store its position in the list in the variable mp.
- New values are tested by comparing the item in position i with the item in position mp.
- •bottom stops at the second to last item in the list. Why? Once all items up to the last are in order, the last item must be the largest!

•The selection sort is easy to write and works well for moderate-sized lists, but is not terribly efficient. We'll analyze this algorithm in a little bit.