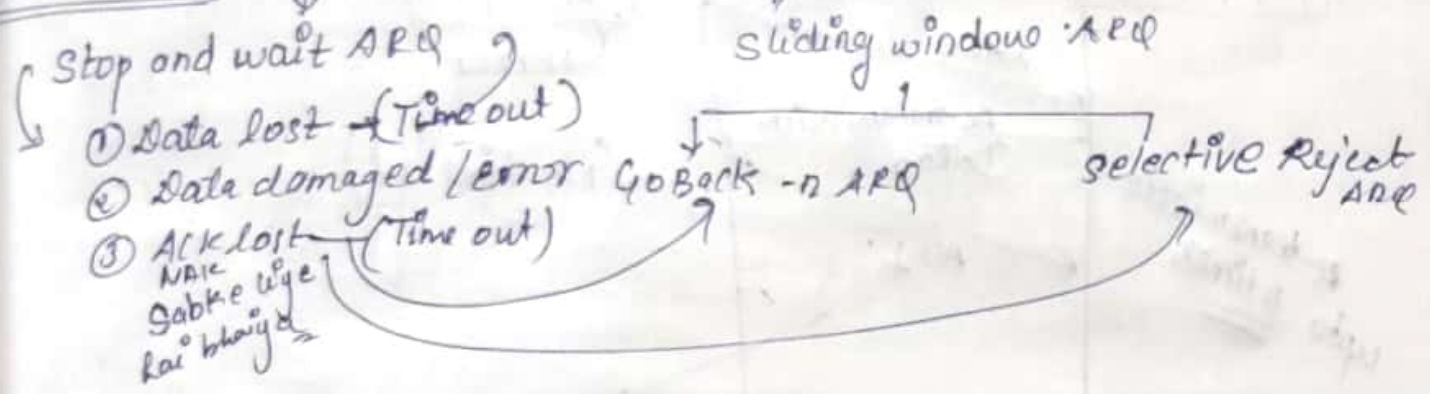
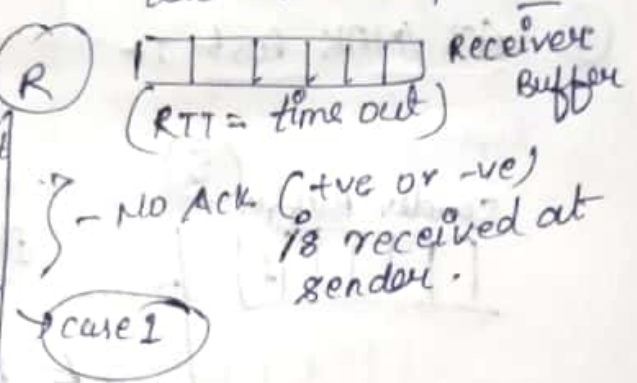
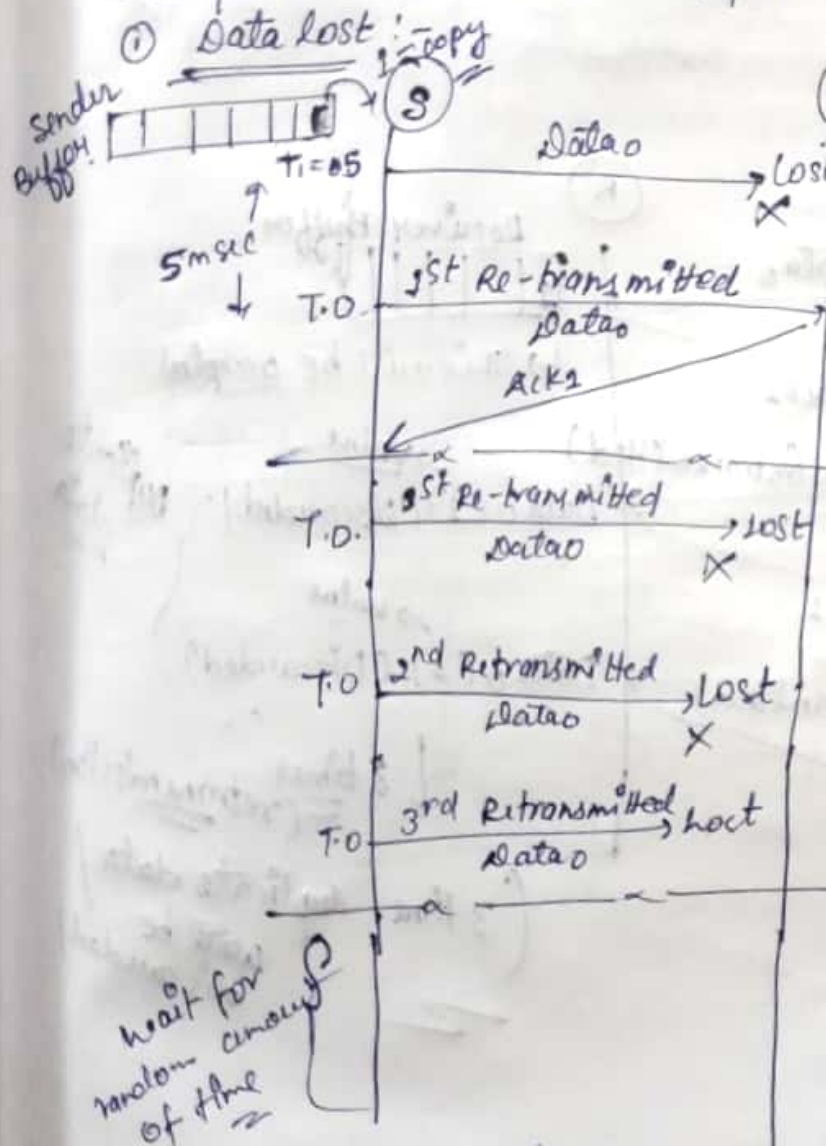


Error Control :-



Stop and wait ARQ :-

After establishment of connection data is transferred.



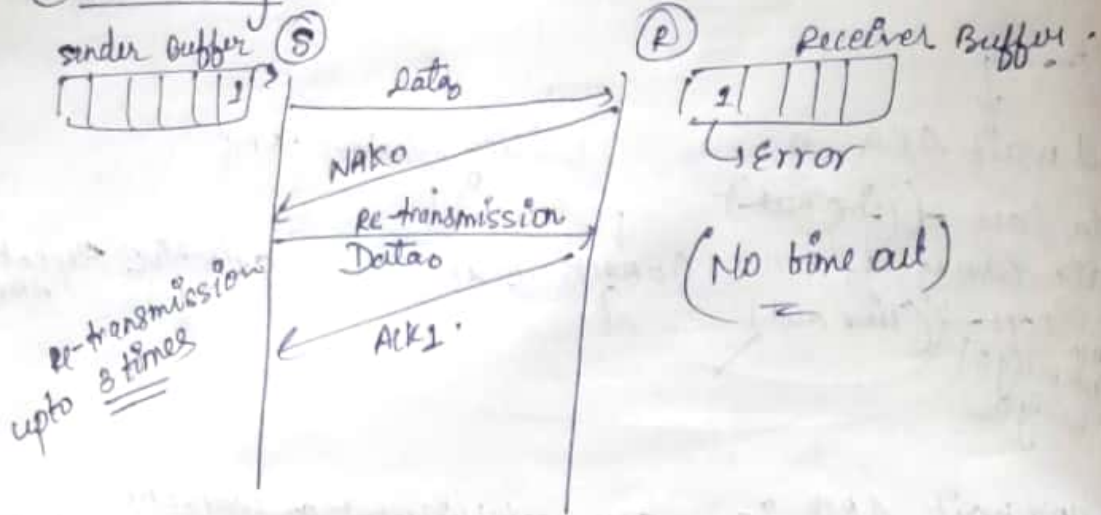
(Max. retransmission = 3) allowed.

After 3rd time out the sender will assume and Receiver is not connected

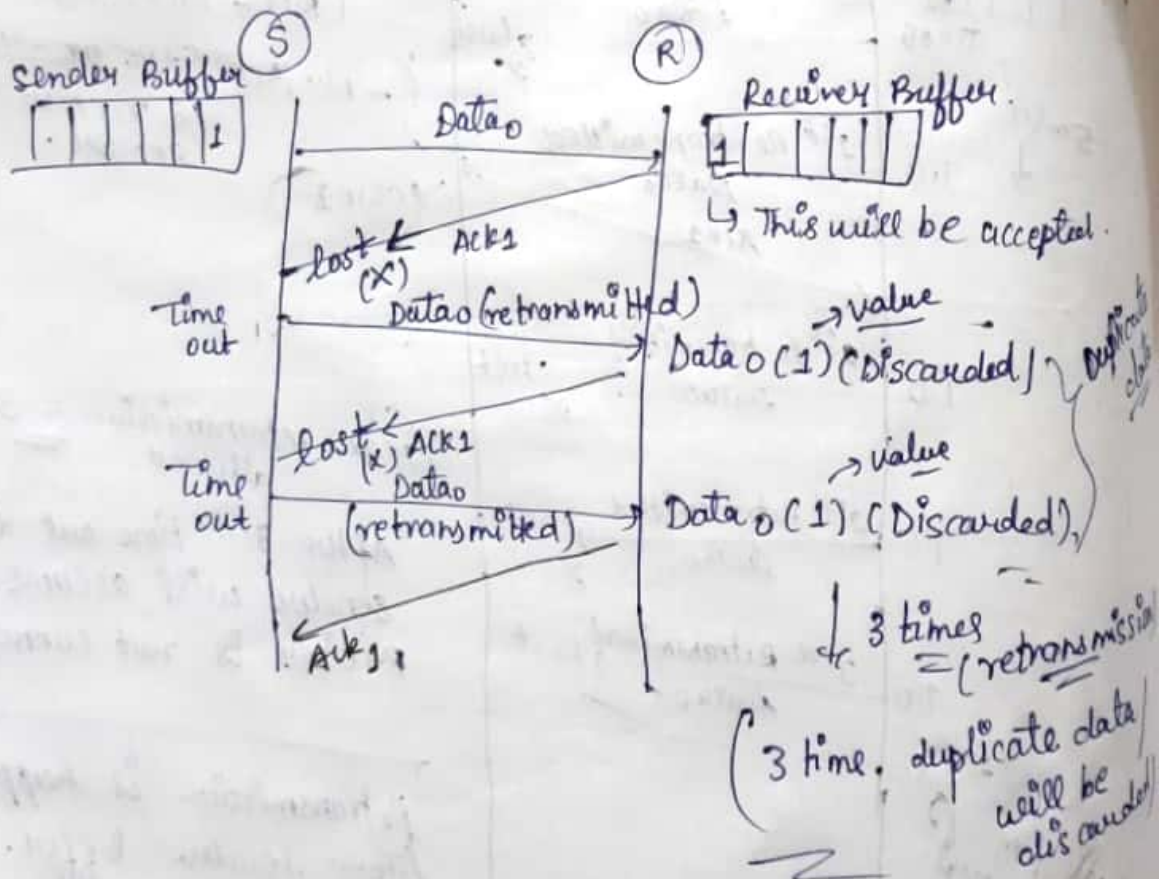
Retransmission is happened from sender buffer. Once the Ack is received then only 1 data is discarded.

RTT = Time out = TTL
 ↓
time to limit

② Data Damaged :-



③ Ack / NAK lost :-



30/march/2022

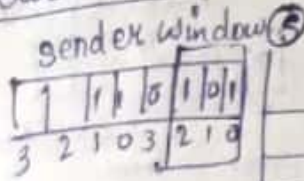
Sliding window ARQ

→ Data lost
→ Data damaged
→ ACK lost

0 → 1 → 2 → 3
0 → 2 → 3
Sequence is not allowed

Go-back-n ARQ :-

① Data lost :-



modulo(n)
n = 4 (0-3)
size = n-1
= 3

Retransmission

Data 0

Data 1 lost X

Data 2

Ack 1

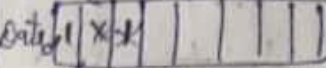
Data 1

Ack 2

Data 2

Data 3

Receiver window

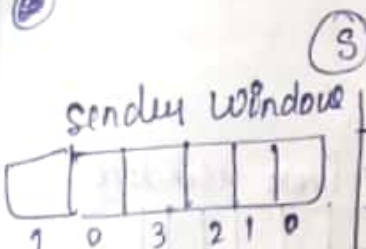


Data 1 Discarded
Data 2 Discarded

Data 1

TCP = Transmission control protocol

selective / Reject ARQ



Data 0

Data 1 lost X

Data 2

Data 3

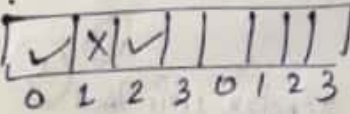
NAK 1

Data 1

Retransmitted

Ack 4

Receiver window

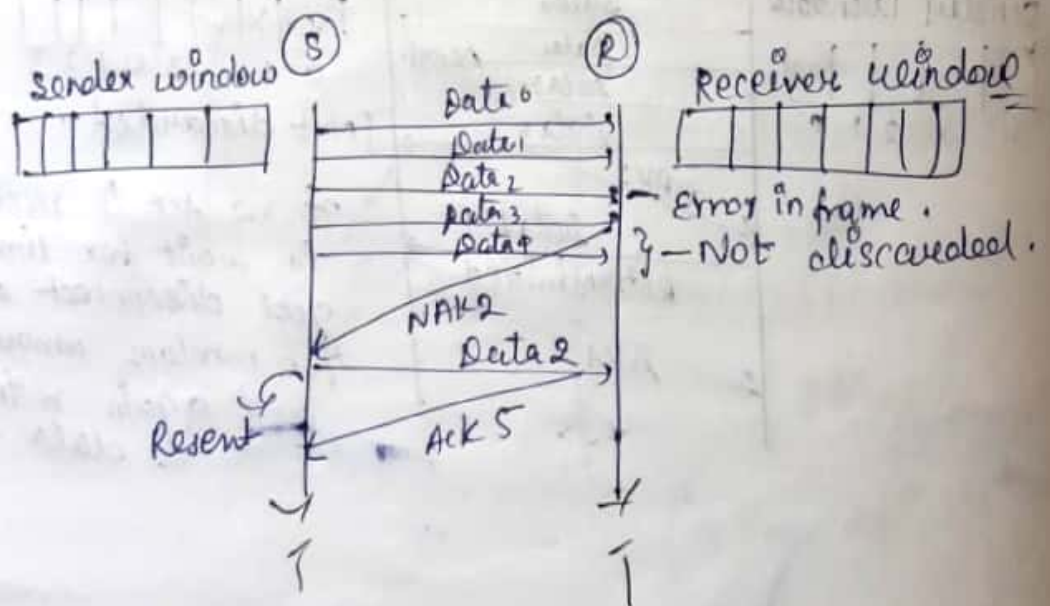
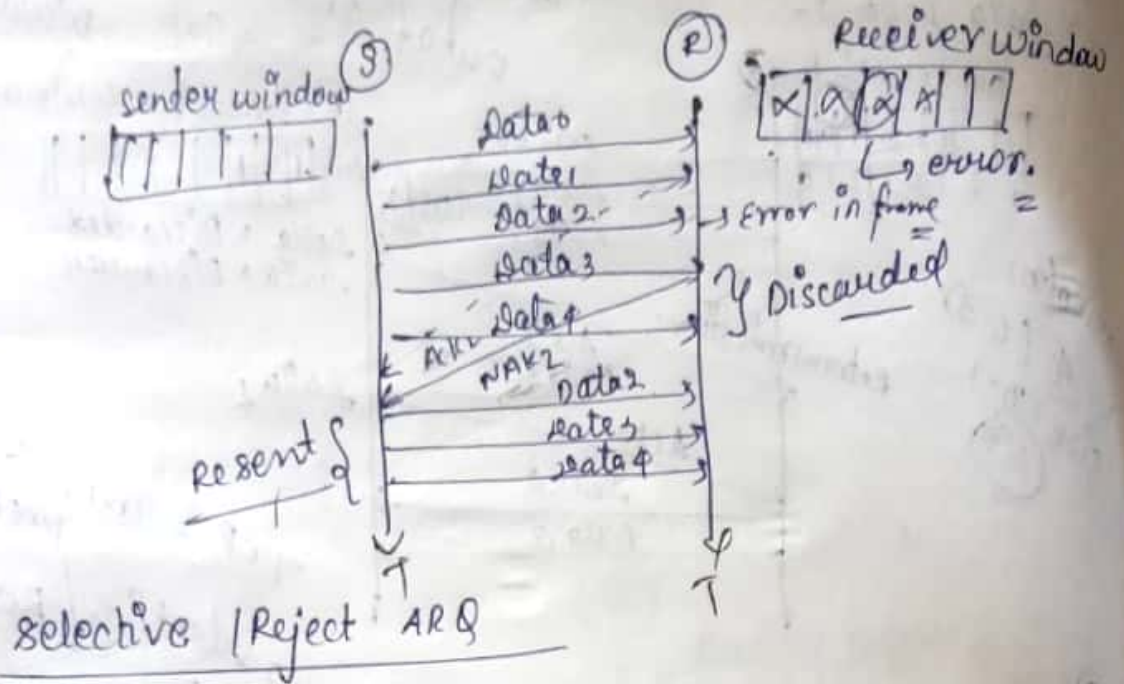


Not discarded

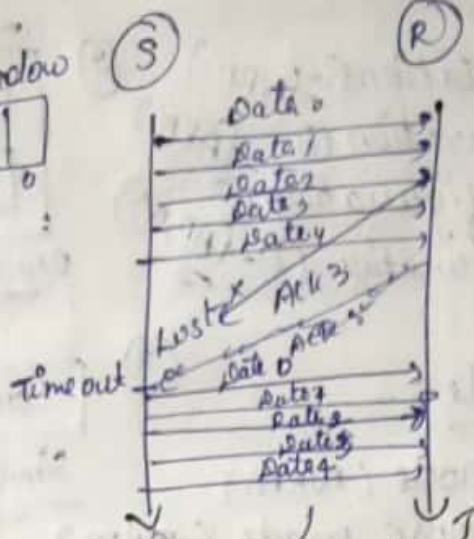
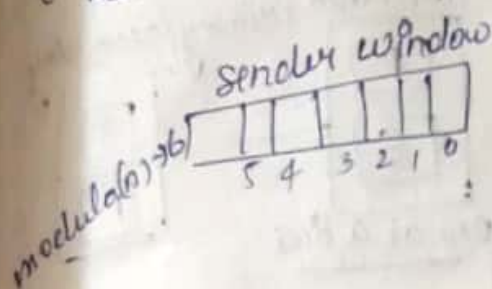
If no ACK is received then it wait for time out or gets disconnect and wait for random amount of time and again retransmit data.

③ Data damaged :-

Go-back-n ARQ



⑤ ACK lost
(Time-out) Go-back-n ARQ :-



If ACKS is received within the Time out then the Ack3 - which is lost is meaningless

Selective / Reject ARQ = Same

Data link layers protocols :-

Synchronous protocols

Asynchronous protocols

character oriented protocols
Bit oriented protocols

X-modem

Z

Start

BSC

SDLC
HDLC

LAPs

LANs

High level data link control protocol

HDLC :- DHCD is dynamic host control protocol.

Station type

Primary station
combined station

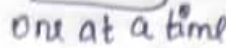
Primary as well as secondary

⑤ Configurations of

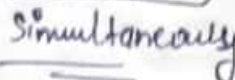
HDLC

① station type :-
↳ Primary station
↳ combined station

- ② Configuration of stations:- (NRM)
- ↳ Unbalanced configuration ① (ARM)
 - ↳ Symmetrical configuration ② (ARM)
 - ↳ Balanced configuration ③ (ARM)



- Normal Response mode (NRM)
- Asynchronous response mode (~~N~~ARM)
- Asynchronous Balanced mode (ABM)



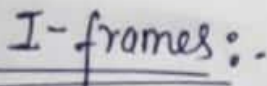
Remembrance

④ frames:-

- Supervisory

→ Information for management purpose

I-frames:-



problem:- user info. user want to send this data.
 0111110 flag 0111110 flag

To overcome this → Bit stuffing / char stuffing:-

character stuffing:-

Ex. Hello flag

Flag Hello ESC flag flag
 (S) (character stuffing) (E)

Ex. Flag Hello flag flag information

flag ESC flag Hello ESC flag ESC flag flag
 (S) (E) (E)

Bit stuffing:-

S-frame:-

flag | Address | control | fcs | flag

U-frame:-

flag | Address | control | information | fcs | flag

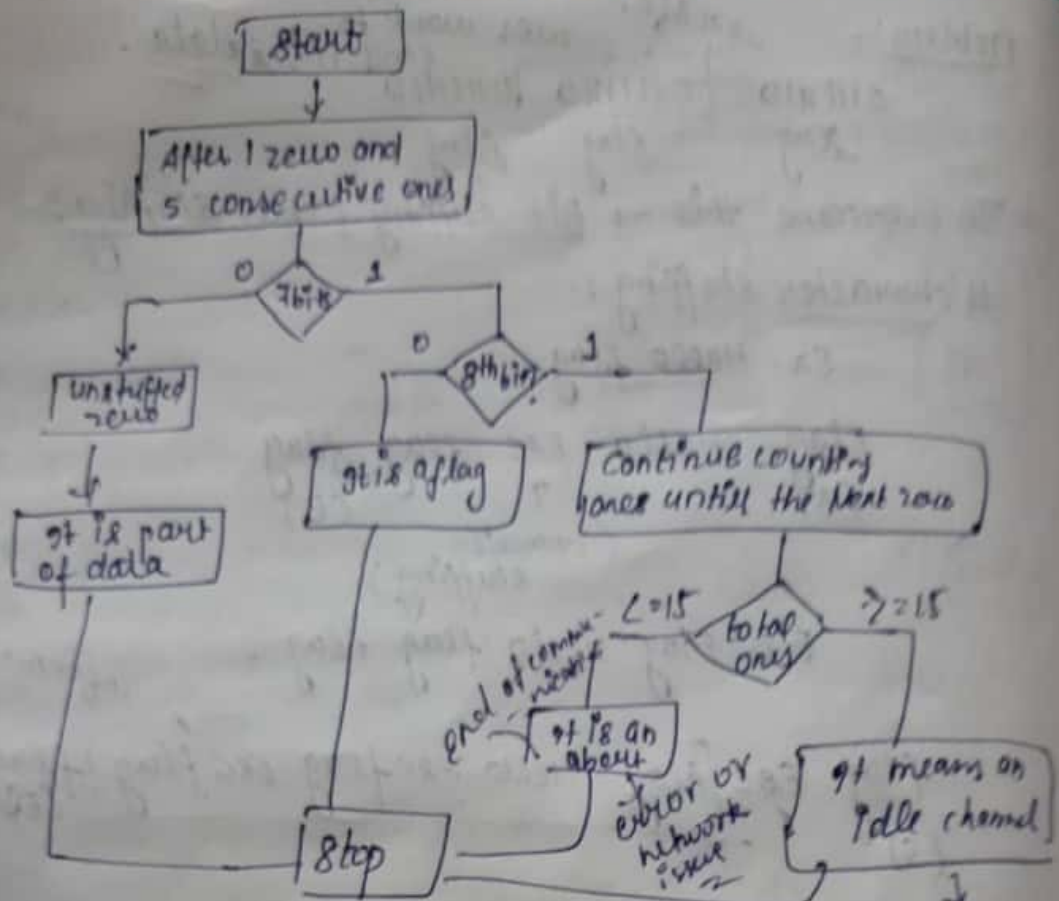
Not data / user information
 management related info.

Bit stuffing:-

flag → 0111110

ESC

next page



(01111110) - if you want to send this data. (free for transmission)
we can write it as, -

01111110 011111010 01111110
flag information flag

01111110
01111110

if you want to send 011111010
so 011111010

LAN (Local Area Network)

802.1, 802.2

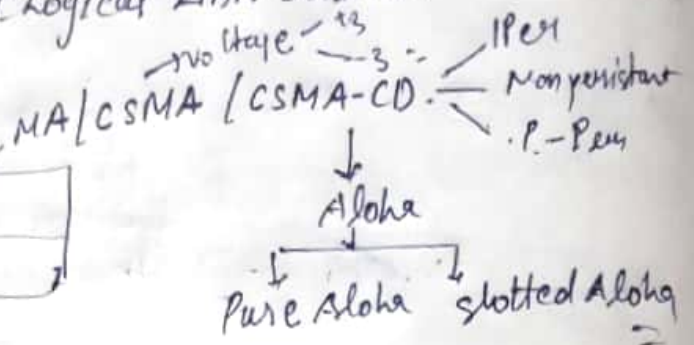
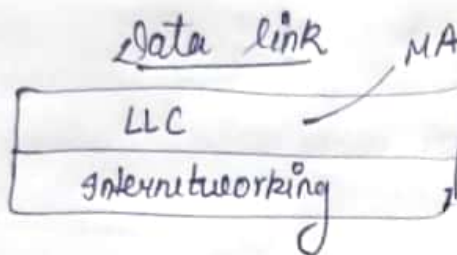
IEEE Project
↓ Protocols for sending data

IEEE 802.3, IEEE 802.4, IEEE 802.5
ethernet Token bus token bus token ring

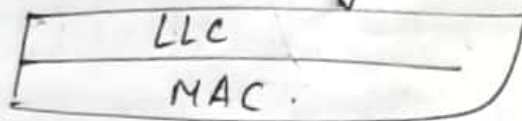
token ring

IEEE → 802.1 → Internetworking. (MAC)

IEEE → 802.2 → LLC (Logical Link Control).

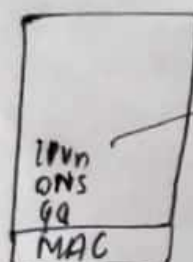
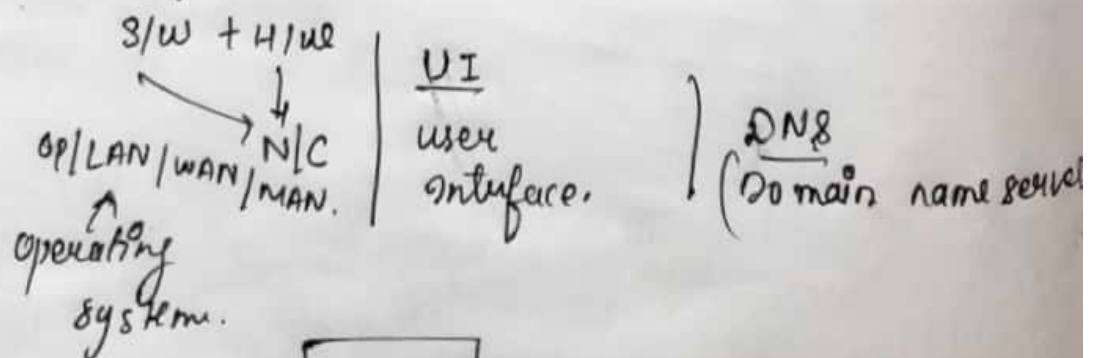


Data link layer



* 802.3 frame format (ethernet):

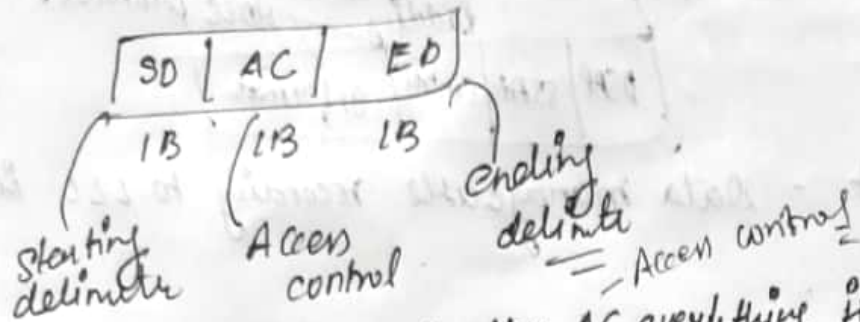
Combination of software and hardware.



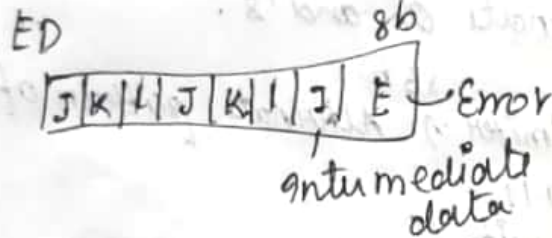
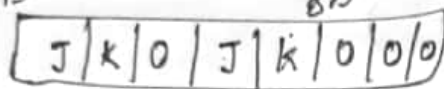
software attribute
in hardware

Ring can be configured clockwise or anticlockwise.
Token will rotate in direction of ring.

Token field :-



SD → indicates that after AC everything is part of data JK flipflop.



$E = 0$ (No error)

$E = 1$ (Error in token field.)

duplicate lost

provided to wrong system

$I = 1$ (some intermediate part of data is in channel, don't close)

$I = 0$ (last frame is received (use full close))

AC 8b.



destination is found

C - data is copied.

A

0 - receiver is down/crashed.

0

0 - data lost

1

1 - Not possible.

0

1 - Success

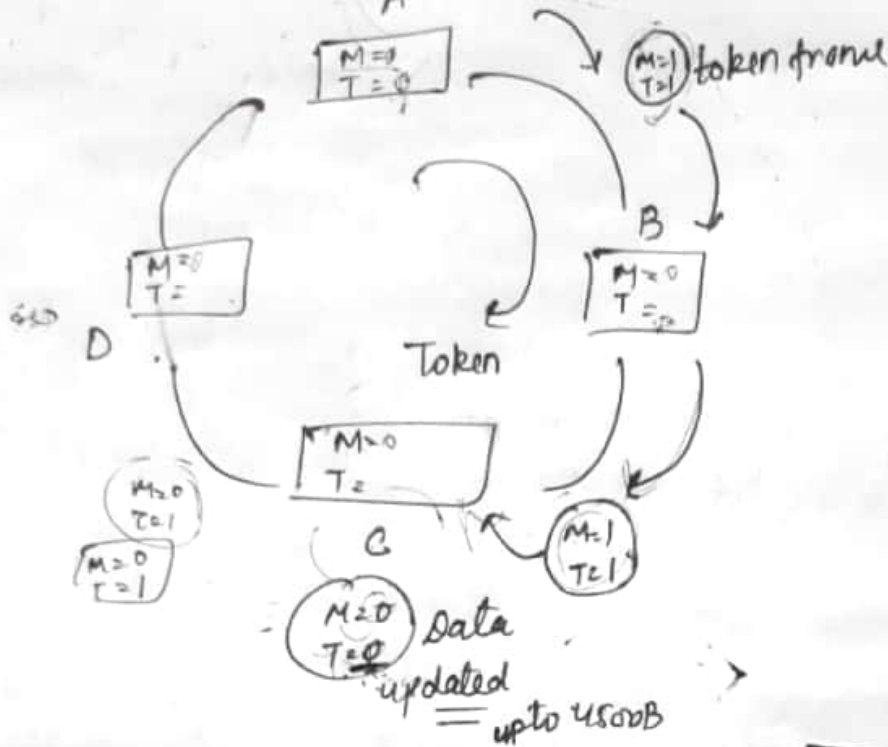
1

8/04/2022

802.5

token ring
A*

M=1 (monitor)
M=0 (Normal)
T=1 (Token)
T=0 (Data)



Abort

SD ED
1B 1B

SD	AC	FC	DA	SA	Data & padding	CRC	ED	FS
1B	1B	1B	2-6B	2-6B	4B	4B	1B	1B

Size of PDU

PDU: DSAP | SSAP | info | control

HDLC (sema)
mini size = PDU size

SD AC ED
1B 1B 1B
Token

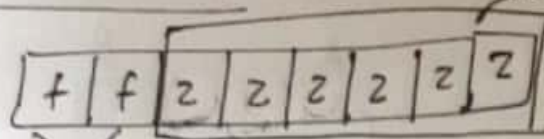
9D

J K 0 J K 0 0 0

ED

J K 1 J K 1 I E

FC (frame control) :-



either data
or control.

00 - Data

01 - Control

10 - Data

11 - Control.

control
signals

flag

z -> AMP (Active monitor present).

z -> claim (To become new monitor)

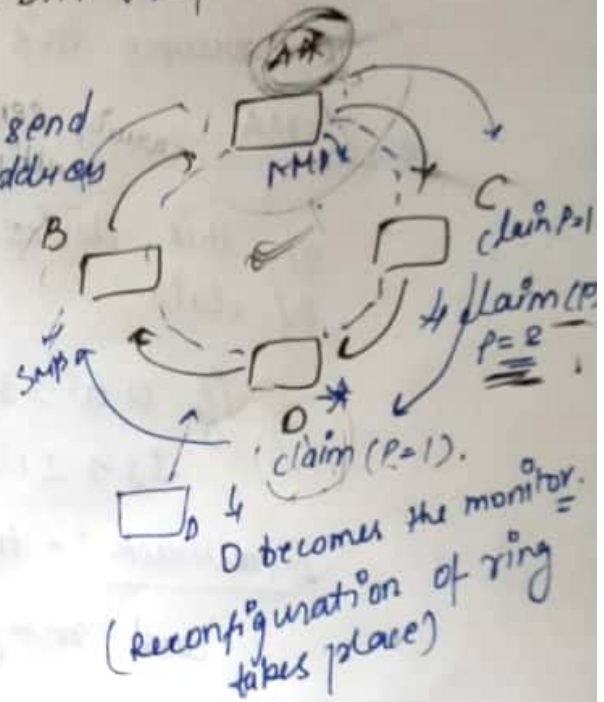
z -> Beacon (stop transmission)

z -> purge (clean the ring) (start fresh comm)

z -> SMP (secondary monitor present).

z -> DAT (Duplicate Address test)

DAT: If external station wants to send data monitor will check MAC address of both devices. If same then both are discarded.



802.4 (802.3 & 802.5) :-

Ethernet

CSMA/CD

MA : multiple access

whenever shared channel is used.

To overcome this

CSMA carries sense multiple access.

→ on this device will check voltage before transmission of data.

if 0 to $\pm 3V$ send data (channel is free).

± 3 to ± 15 → channel is ^{not} free.

Drawback : - two device check voltage simultaneously and may send data together.

↓ To overcome this.

CSMA/CD → CSMA with collision detection

Continuously monitor voltage in the channel. The device which started the transmission of data if more than $\pm 15V$ then stop transmission and wait for random amount of time and also check continuously the voltage in the channel.

- 1-persistent CSMA : Entire data will be inserted to channel (100%).
- Non-persistent CSMA → May or may not (Not for first time)
→ 2 times check.
- P-persistent CSMA :
 - $P = 50\%$
 - $P = 60\%$

$P =$ probability.
With P prob. data is inserted.

