

Artificial Intelligence 1

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Definition

A list L is a sequence of terms. A list is itself a term.

Example:

$[1, 2, 3, 4, 5]$

$[p(X, Y), 3, [X, Y, Z], q([a, b, c], 5), 7]$

Alternative definition:

Definition

A list L is either

1. the empty list $[]$ or
2. a term T followed by a list L' : $[T|L']$

For $L = [H|T]$ we call H the head and T the tail of the list L .

Lists 2/8

```
?- [1,2,3,4,5] = [H|T].  
H = 1,  
T = [2, 3, 4, 5].
```

We could define:

```
?- head([1,2,3,4,5],X).  
X = 1.  
  
?- tail([1,2,3,4,5],X).  
X = [2, 3, 4, 5].
```

How?

```
head([H|_],H).  
tail([_|T],T).
```

Lists 3/8

`member(E,L)`: true iff E is element of the list L.

Definition:

```
member(X, [X|_]) .  
member(X, [_|Y]) :- member(X,Y) .
```

```
?- member(2, [1,2,3]).  
true.  
  
?- member(4, [1,2,3]).  
false.  
  
?- member(X, [1,2,3]).  
X=1 ;  
X=2 ;  
X=3 .  
  
?- member(X, Y).  
Y = [X|_G280] ;  
Y = [_G16, X|_G20] ;  
Y = [_G16, _G19, X|_G23] ; ...
```

`append(L1,L2,L)`: true iff L is the concatenation of L1 and L2.

Definition:

```
append([],L,L).  
append([X|L1], L2, [X|L3]) :- append(L1, L2, L3).
```

```
?- append([1,2], [3,4], [1,2,3,4]).  
true.
```

```
?- append([1,2], [X,Y], [1,2,3,4]).  
X = 3,  
Y = 4.
```

```
?- append([1,2], [3,4], L).  
L = [1,2,3,4].
```

`reverse(L1,L2)`: true iff L1 is the reverse list of L2.

```
?- reverse([1,2,3],[3,2,1]).  
true.
```

```
?- reverse([1,2,3], L).  
L = [3,2,1].
```

Definition:

```
reverse([], []).  
reverse([X|Y],L) :- reverse(Y,Z), append(Z,[X],L).
```

`length(L,X)`: true iff X is the number of elements in L.

```
?- length([1,2,3,4],4).
```

```
true.
```

```
?- length([1,2,3,4],X).
```

```
X = 4.
```

```
?- length(X,4).
```

```
X = [_G2369, _G2372, _G2375, _G2378].
```

Definition:

```
length([],0).
```

```
length([_|X], N) :- length(X,M), N is (M+1).
```


Insertion sort in Prolog:

```
sorted([]).  
sorted([X]).  
sorted([X|[Y|Z]]) :- X =< Y, sorted([Y|Z]).  
  
insert(X, [], [X]).  
insert(X, [Y|Z], [X|[Y|Z]]) :- X =< Y.  
insert(X, [Y|Z], [Y|U]) :- Y < X, insert(X, Z, U).  
  
sort(X, X) :- sorted(X).  
sort([X|Y], Z) :- sort(Y, U), insert(X, U, Z).
```

Examples:

```
?- sorted([1,2,3]).
```

```
true.
```

```
?- sorted([2,3,1]).
```

```
false.
```

```
?- sort([5,1,4,2,3],X).
```

```
X = [1, 2, 3, 4, 5].
```

```
?
```

The cut operator 1/4

Prolog always searches for all possible solutions for a query (via backtracking).

```
a(1). a(2). b(3).  
p(X) :- a(X).  
p(X) :- b(X).
```

```
?- p(X).  
X = 1;  
X = 2;  
X = 3.
```

The cut operator 2/4

However, it can be reasonable to abort backtracking at some time

- ▶ for efficiency reasons: sometimes returning a single solution is sufficient
- ▶ for programmatic reasons: maybe the use of another clause should be avoided if a usable clause has already been found

The cut operator ! can be used to abort backtracking

```
a(1). a(2). b(3).  
p(X) :- !, a(X).  
p(X) :- b(X).
```

```
?- p(X).  
X = 1;  
X = 2.
```

The cut operator 3/4

- ▶ The cut operator ! is an atom and can be used as such in the body of any Prolog rule
- ▶ It has no influence on derivations as the atom ! is always satisfied.
- ▶ Side effect: the derivation of the atom in the head of the current rule is “frozen”
 - ▶ All variable assignments are fixed
 - ▶ No further clauses with the same head are tried

Compare:

```
a(1). a(2). b(3).  
p(X) :- !, a(X).  
p(X) :- b(X).
```

```
?- p(X).  
X = 1;  
X = 2.
```

```
a(1). a(2). b(3).  
p(X) :- a(X), !.  
p(X) :- b(X).
```

```
?- p(X).  
X = 1.
```

The cut operator 4/4

Another example:

```
p :- a, b.
```

```
p :- c.
```

$$p \Leftarrow (a \wedge b) \vee c.$$

```
p :- a, !, b.
```

```
p :- c.
```

$$p \Leftarrow (a \wedge b) \vee (\neg a \wedge c).$$

Note that for data basis $\{a., c.\}$ we can derive p in the left program but not in the right program.

Text output

- ▶ Prolog allows for text output during execution with the special predicate `write` (and `write_ln`)
- ▶ The use of text output has no effect on the derivation, atom with `write` are always satisfied.
- ▶ Side effect is text output

Beispiele:

```
?- write(john).  
john  
true.  
  
?- write('A whole sentence').  
A whole sentence  
true.  
  
?- write(X).  
_G309  
true.
```

Example: Path finding in directed graphs

Problem:

- ▶ Given a directed graph with nodes k_1, \dots, k_n
- ▶ Is there a path between any k_i and k_j (for arbitrary i, j)?

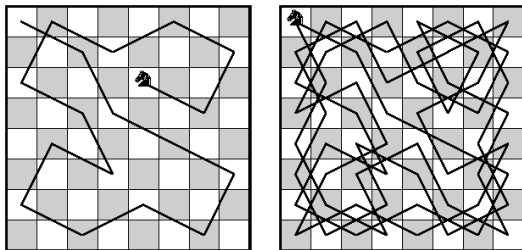
paths.pl:

```
edge(k1,k3).  
edge(k2,k5).  
edge(k4,k2).  
edge(k5,k1).  
path(X,Y) :- edge(X,Y).  
path(X,Y) :- edge(X,Z), path(Z,Y).
```

```
?- path(k1,k2).  
false.
```

```
?- path(k2,k1).  
true.
```


Example: knight's tour 1/2



Question: is there a sequence of moves of a knight on a chessboard such that the knight visits every square exactly once (i. e., is there a Hamilton tour for a knight?).

Simplified question here: is there a path from a square (x_1, y_1) to some other square (x_2, y_2) ?

Example: knight's tour 2/2

knight.pl:

```
legal(1). ... legal(8).
move(X1,Y1,X2,Y2) :- legal(X1), legal(X2), legal(Y1), legal(Y2),
    X2 is X1+1, Y2 is Y1+2.
move(X1,Y1,X2,Y2) :- legal(X1), legal(X2), legal(Y1), legal(Y2),
    X2 is X1+1, Y2 is Y1-2.
move(X1,Y1,X2,Y2) :- legal(X1), legal(X2), legal(Y1), legal(Y2),
    X2 is X1-1, Y2 is Y1+2.
move(X1,Y1,X2,Y2) :- legal(X1), legal(X2), legal(Y1), legal(Y2),
    X2 is X1-1, Y2 is Y1-2.
move(X1,Y1,X2,Y2) :- legal(X1), legal(X2), legal(Y1), legal(Y2),
    X2 is X1+2, Y2 is Y1+1.
move(X1,Y1,X2,Y2) :- legal(X1), legal(X2), legal(Y1), legal(Y2),
    X2 is X1+2, Y2 is Y1-1.
move(X1,Y1,X2,Y2) :- legal(X1), legal(X2), legal(Y1), legal(Y2),
    X2 is X1-2, Y2 is Y1+1.
move(X1,Y1,X2,Y2) :- legal(X1), legal(X2), legal(Y1), legal(Y2),
    X2 is X1-2, Y2 is Y1-1.

path(X1,Y1,X1,Y1,L).
path(X1,Y1,X2,Y2,L) :- legal(X3), legal(Y3), move(X1,Y1,X3,Y3),
    not(member([X3,Y3],L)), path(X3,Y3, X2,Y2, [[X3,Y3]|L]).
path(X1,Y1,X2,Y2) :- path(X1,Y1,X2,Y2,[[X1,Y1]]).
```

Chapter 2.2: Prolog

Summary

Chapter 2.2: Summary

- ▶ First-order logic and Horn logic
- ▶ Prolog programs
 - ▶ data base D
 - ▶ rule base R
- ▶ Prolog interpreter SWI-Prolog
- ▶ Arithmetic operators $+$, $-$, $*$, $/$, comparison operators $=$, $=:=$, $=\backslash=$, $<$, $>$, $>=$, $=<$, assignment is
- ▶ Lists and list operations
- ▶ Cut operator, text output