

BUSINESS PROCESS MANAGEMENT

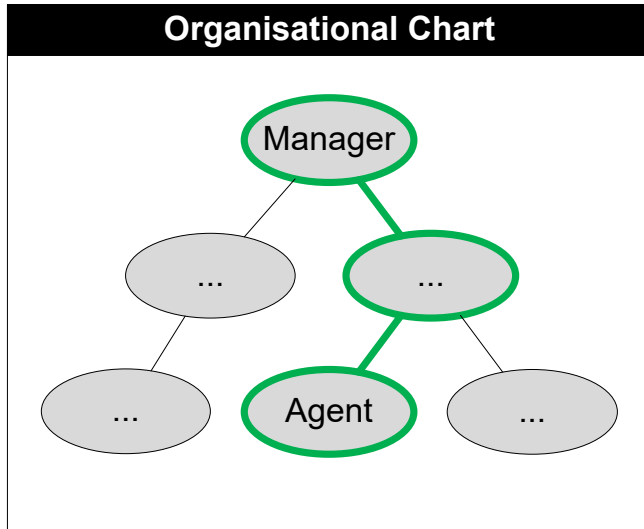
MODEL QUERY III:
THE DIAGRAMED MODEL QUERY LANGUAGE (GRAPH MATCHING)

- Motivation for a Diagrammed Model Query Language (DMQL)
- Requirements of DMQL
- The Diagrammed Model Query Language (DMQL) incl. Parallel Tool Presentation
- Further Information on the Topic Model Query

MOTIVATING EXAMPLE: BUSINESS PROCESS COMPLIANCE MANAGEMENT

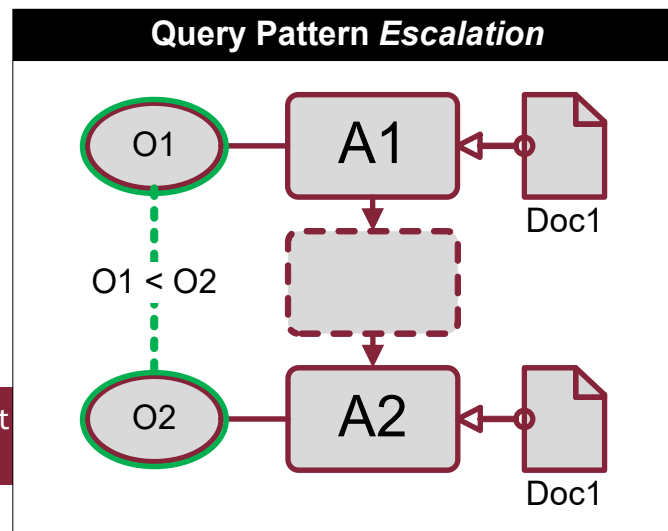


Organisational Chart

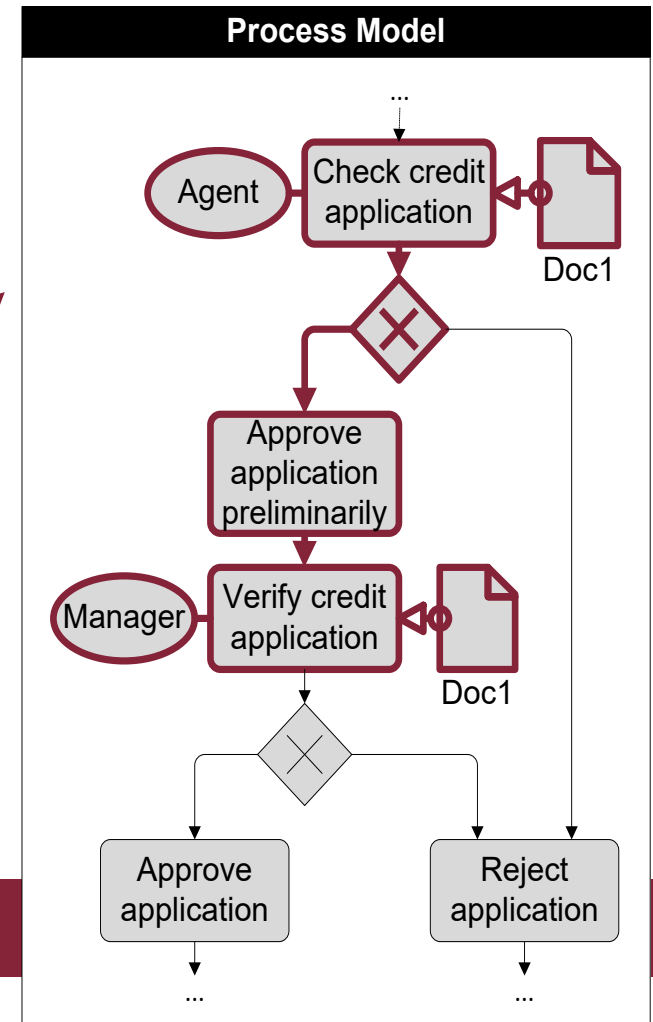


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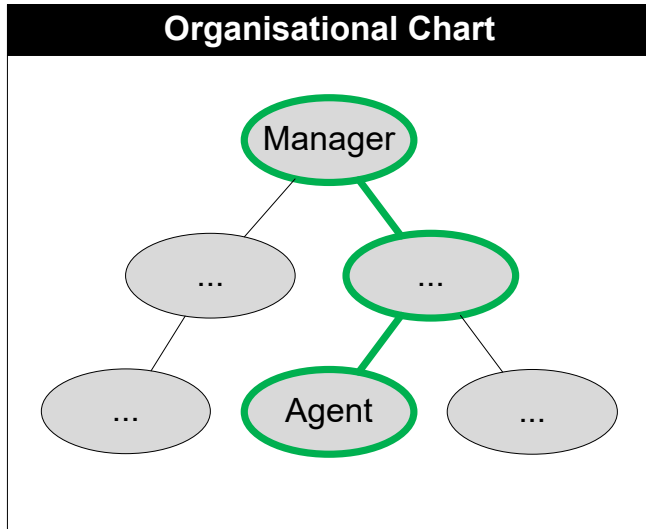
Query Pattern Escalation



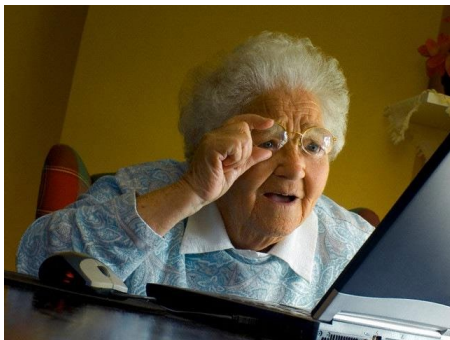
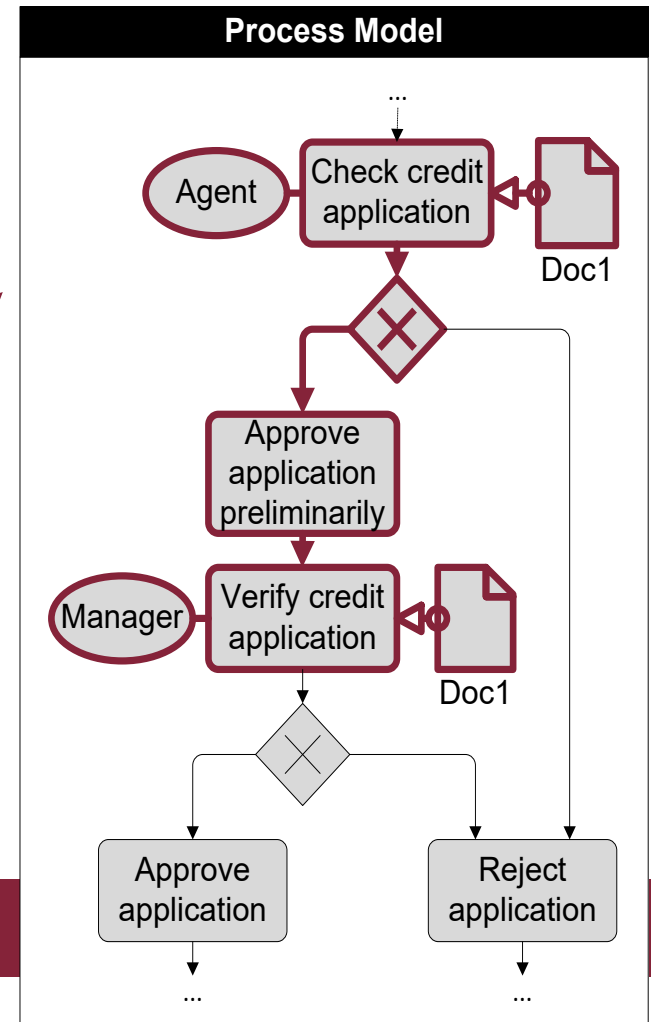
Process Model



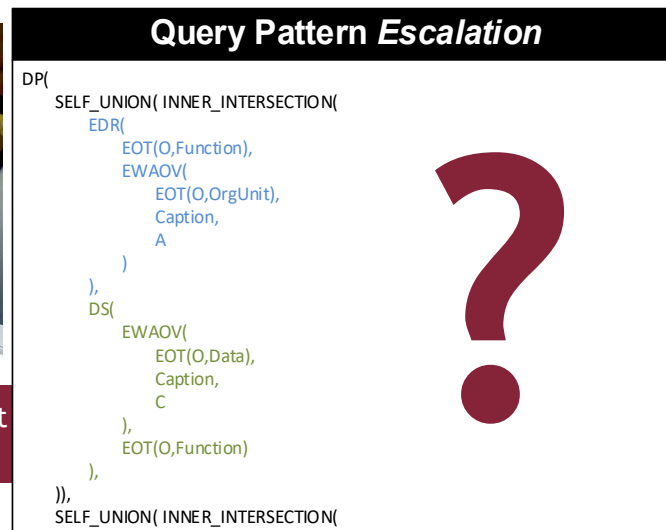
GMQL (AND RELATED GRAMMAR-BASED QUERY LANGUAGES)



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Business Process Management
Prof. Dr. Patrick Delfmann

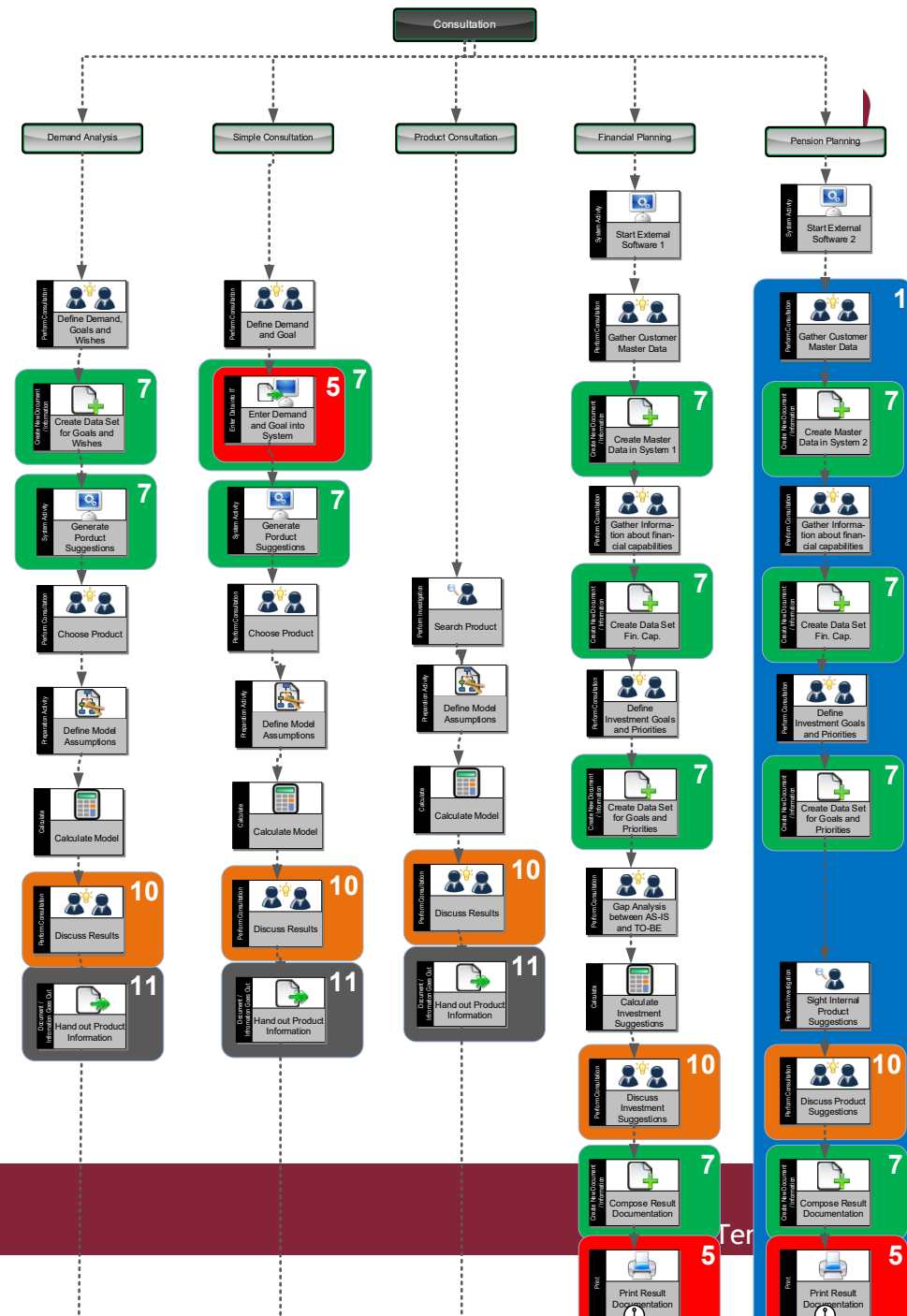


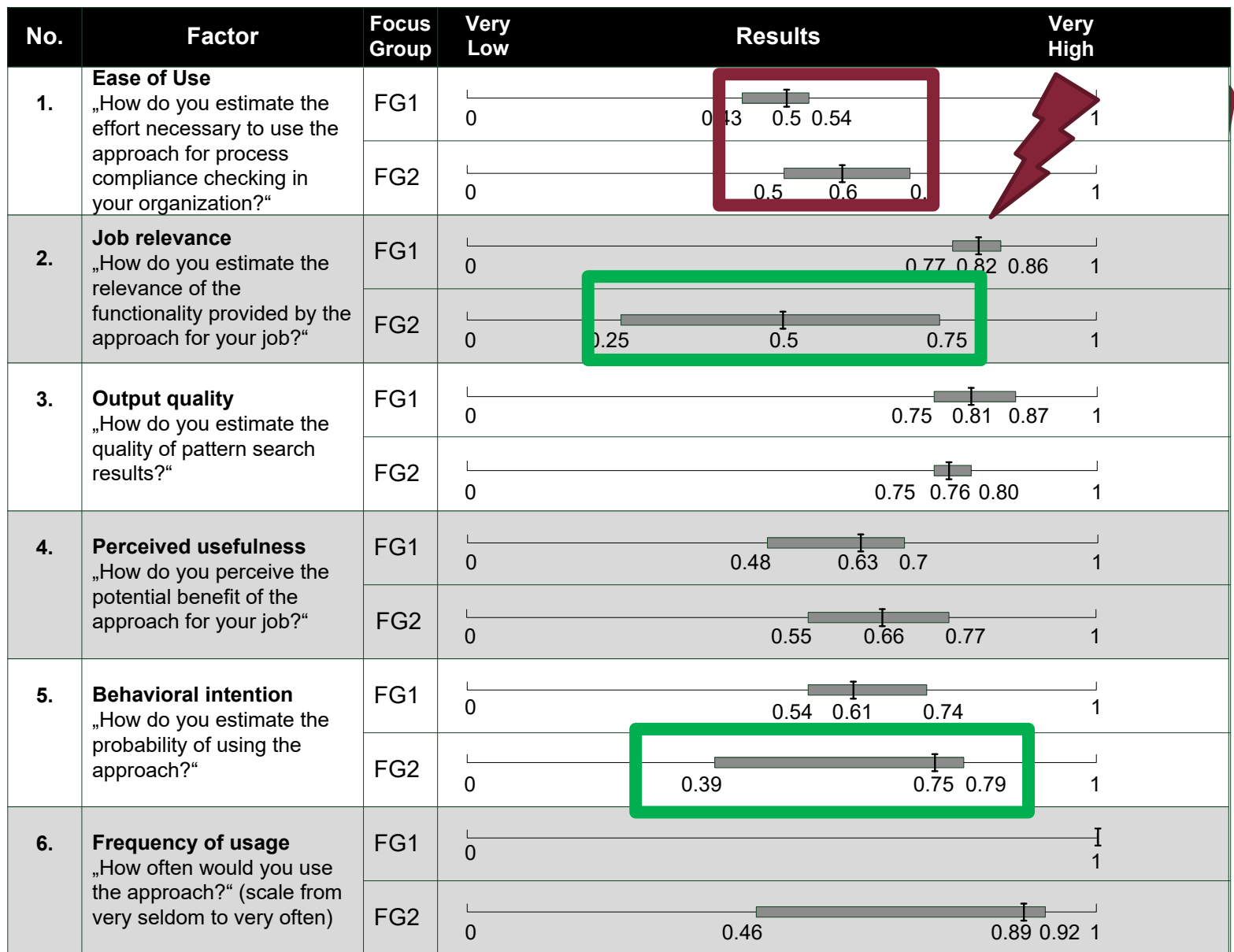
- Business processes of an IT service provider for banks
- Banking software
- PICTURE models
- 13 exemplary compliance rules from WpHG, InvG, WpDVerOV, KWG, GWG



EVALUATION

- Several matches
- Potential compliance violations
- Several matches were actual compliance violations
- **Impact:** adaptations of business software for several 1000s of banks





AGENDA

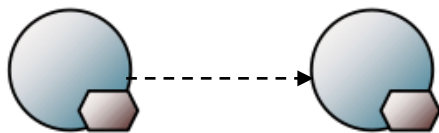


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- Further Information on the Topic Model Query

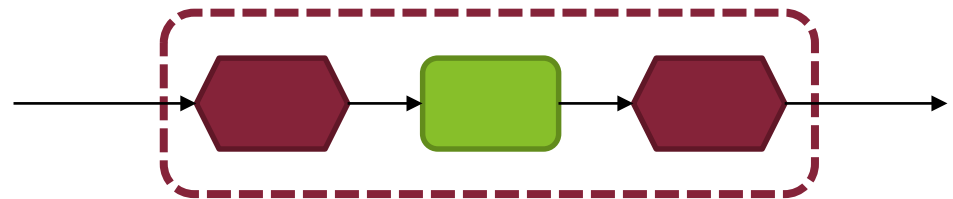
THE DIAGRAMED MODEL QUERY LANGUAGE (DMQL)



- Requirements: see GMQL
- Plus:
 - graphical specification of patterns
 - check different kinds of models simultaneously (cf. example)
- A formalization of the graphic pattern is used as input for a graph pattern matching algorithm



DMQL pattern



what's found in an EPC, e.g.

THE DIAGRAMED MODEL QUERY LANGUAGE (DMQL)



- Requirements: see GMQL
- Plus:
 - graphical specification of patterns
 - check different kinds of models simultaneously (cf. example)
- A formalization of the graphic pattern is used as input for a graph pattern matching algorithm
- Graph pattern matching incl.
 - **isomorphic** parts (edge-edge-mapping)
 - **homeomorphic** parts (edge-path-mapping)
 - vertex and edge **attributes**
 - logical **combination** of pattern elements

THE DIAGRAMED MODEL QUERY LANGUAGE (DMQL)



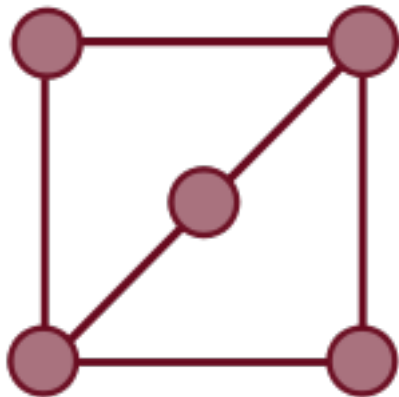
- **Homeomorphism**: A graph G and G' are homeomorphic, if there is a graph isomorphism from a subdivision of G , to a subdivision of G'



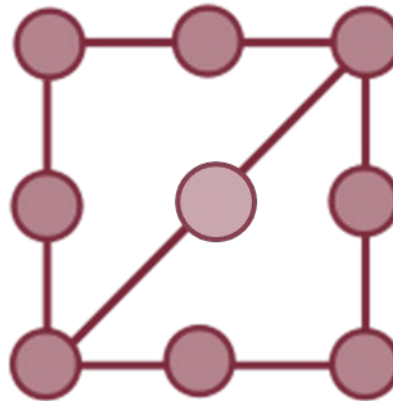
THE DIAGRAMED MODEL QUERY LANGUAGE (DMQL)



- **Homeomorphism**: A graph G and G' are homeomorphic, if there is a graph isomorphism from a subdivision of G , to a subdivision of G'



G



G'

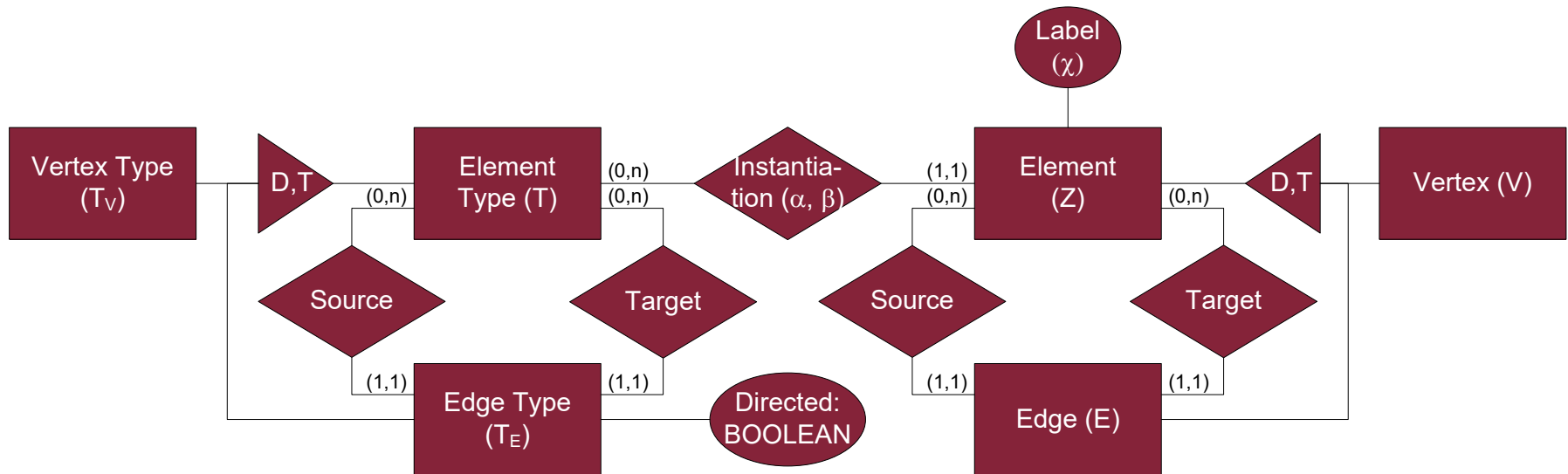
AGENDA



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BASIC SPECIFICATIONS

MULTIPLE MODELING LANGUAGES AND VIEWS



PRELIMINARIES: INTEGRATED CONCEPTUAL MODELS

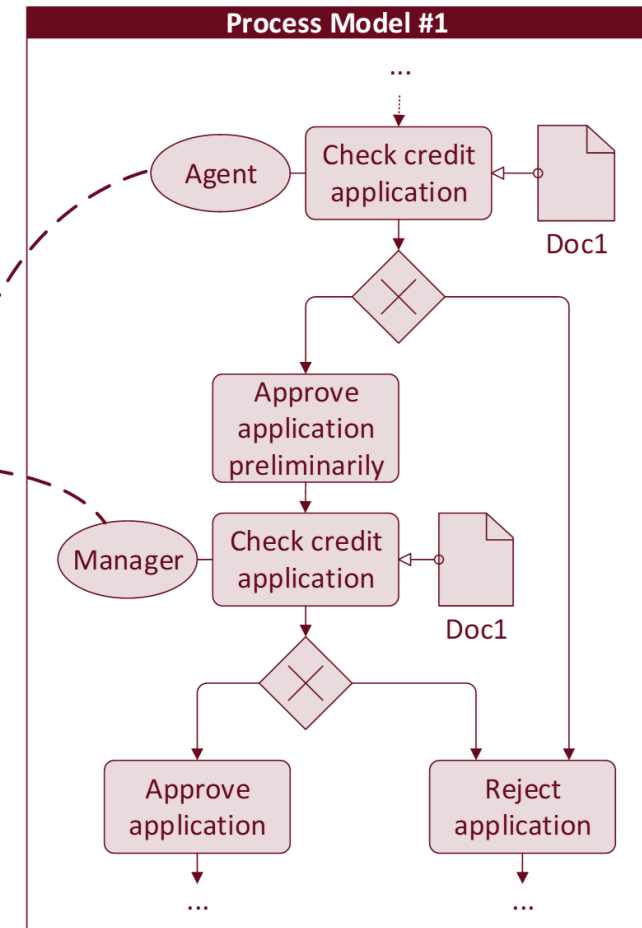
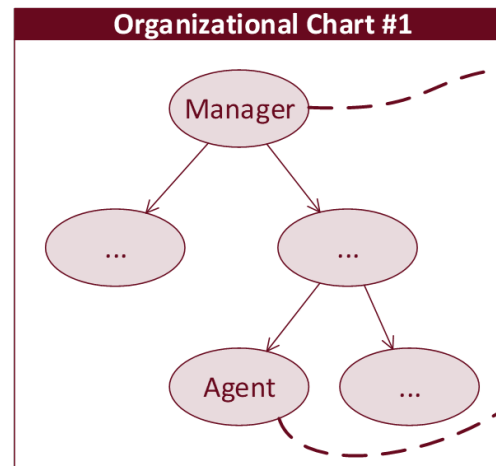


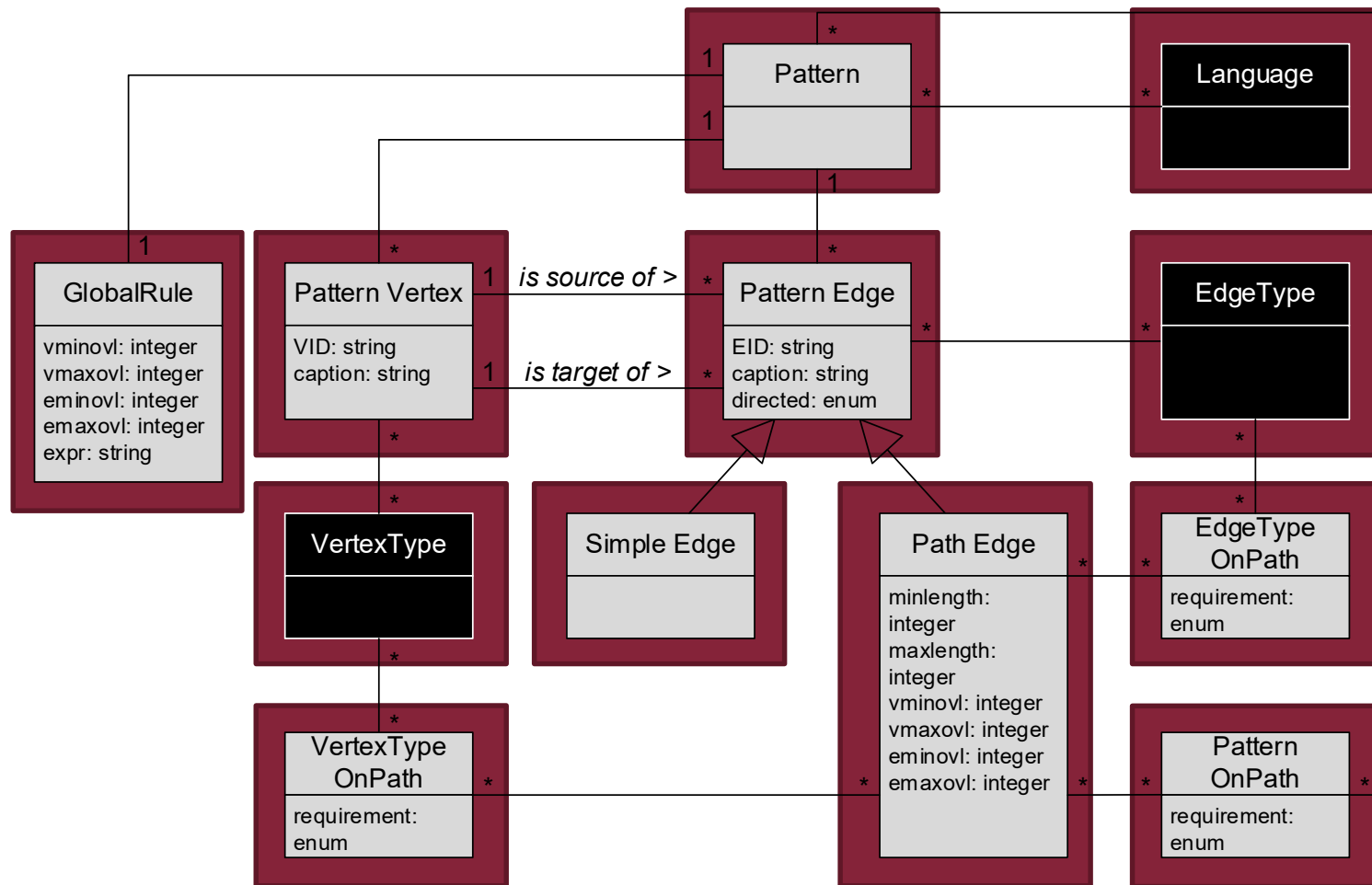
Vertex	Model
Agent	OrgChart
Manager	OrgChart
Agent	Process Model
Manager	Process Model

Vertex Type	Language
OrgUnit	OrgChart
OrgUnit	Process Model

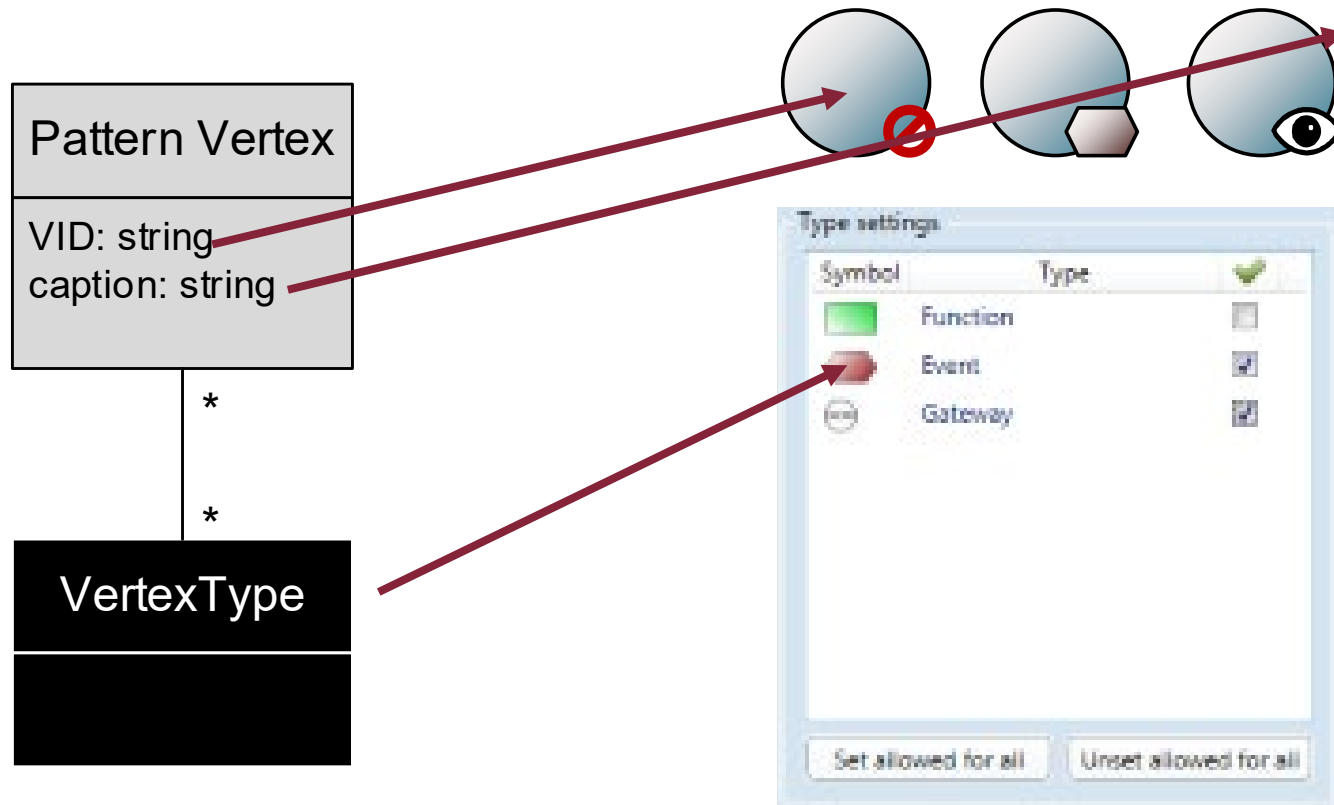
Vertex	Vertex Type
Agent	OrgUnit
Manager	OrgUnit
...	...

Model	Language
OrgChart #1	OrgChart
Process Model #1	Process Model

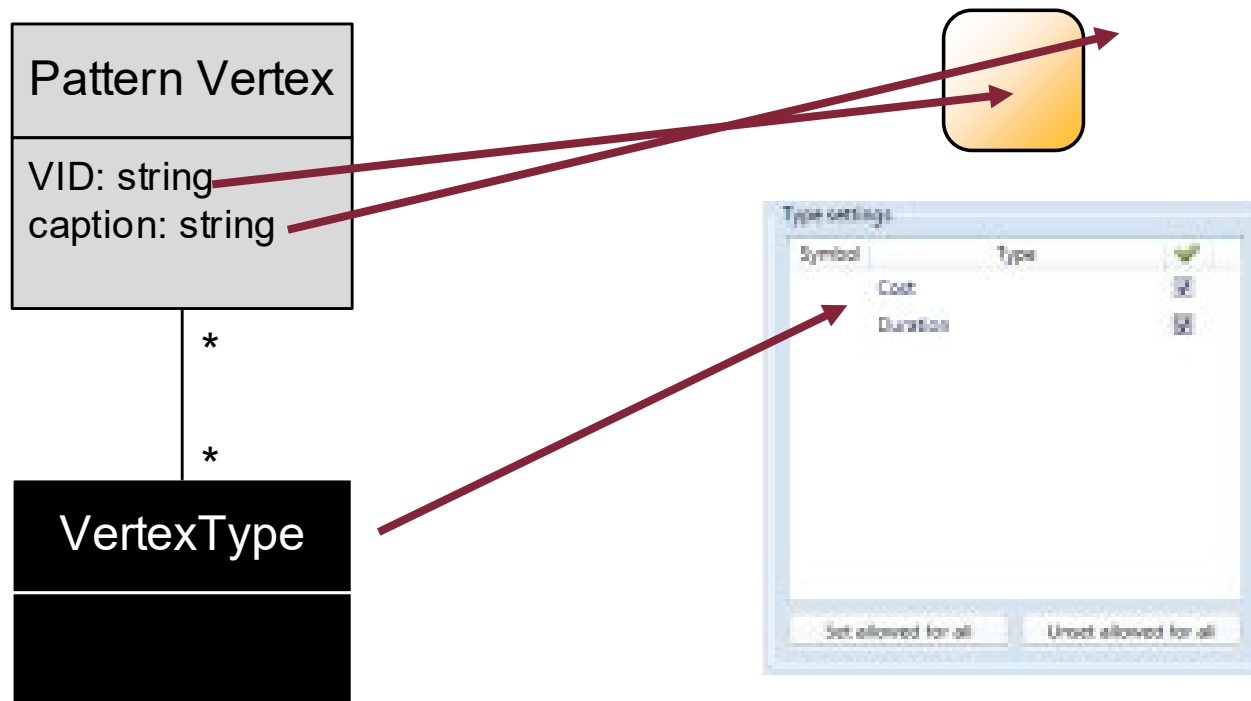




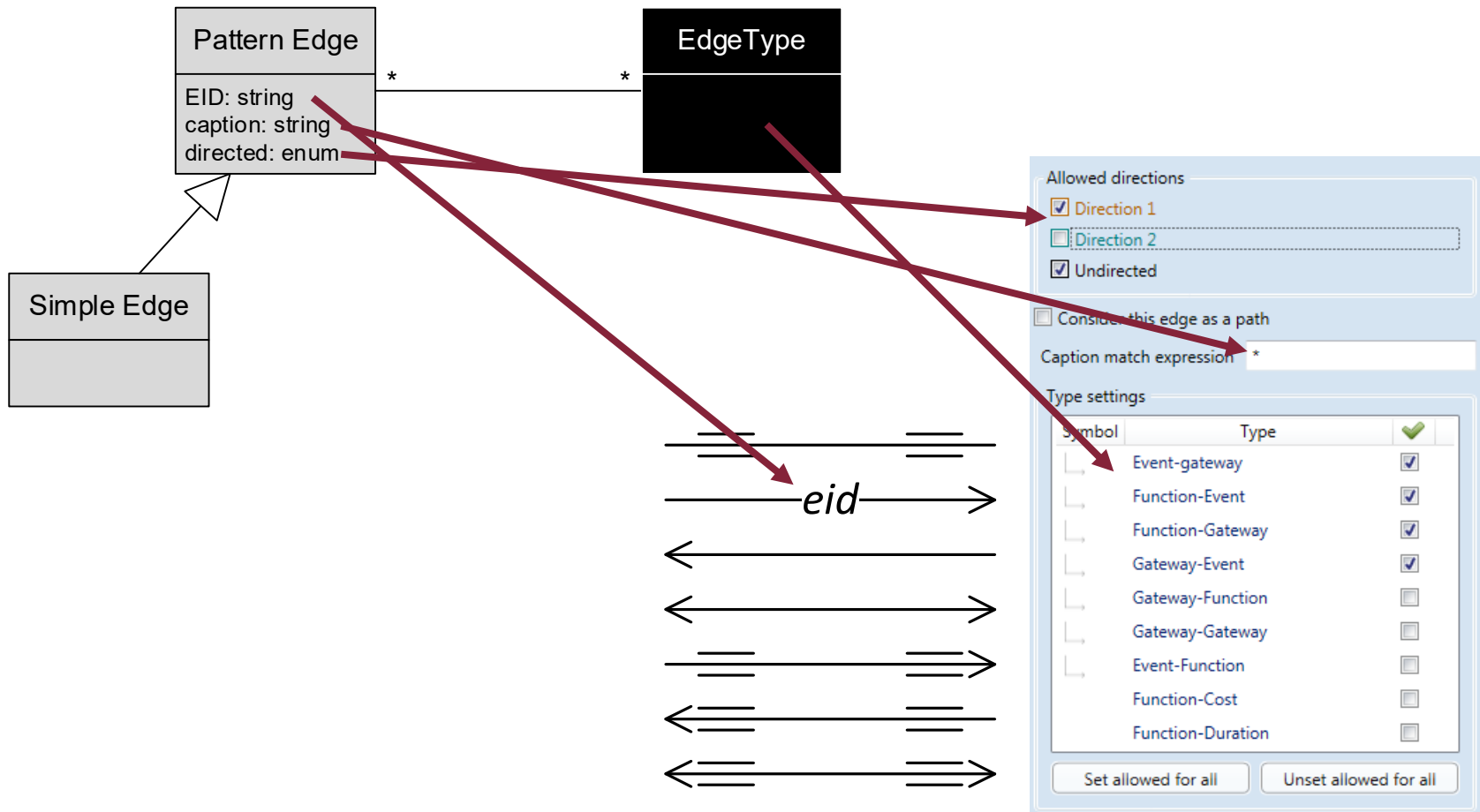
CONCRETE SYNTAX



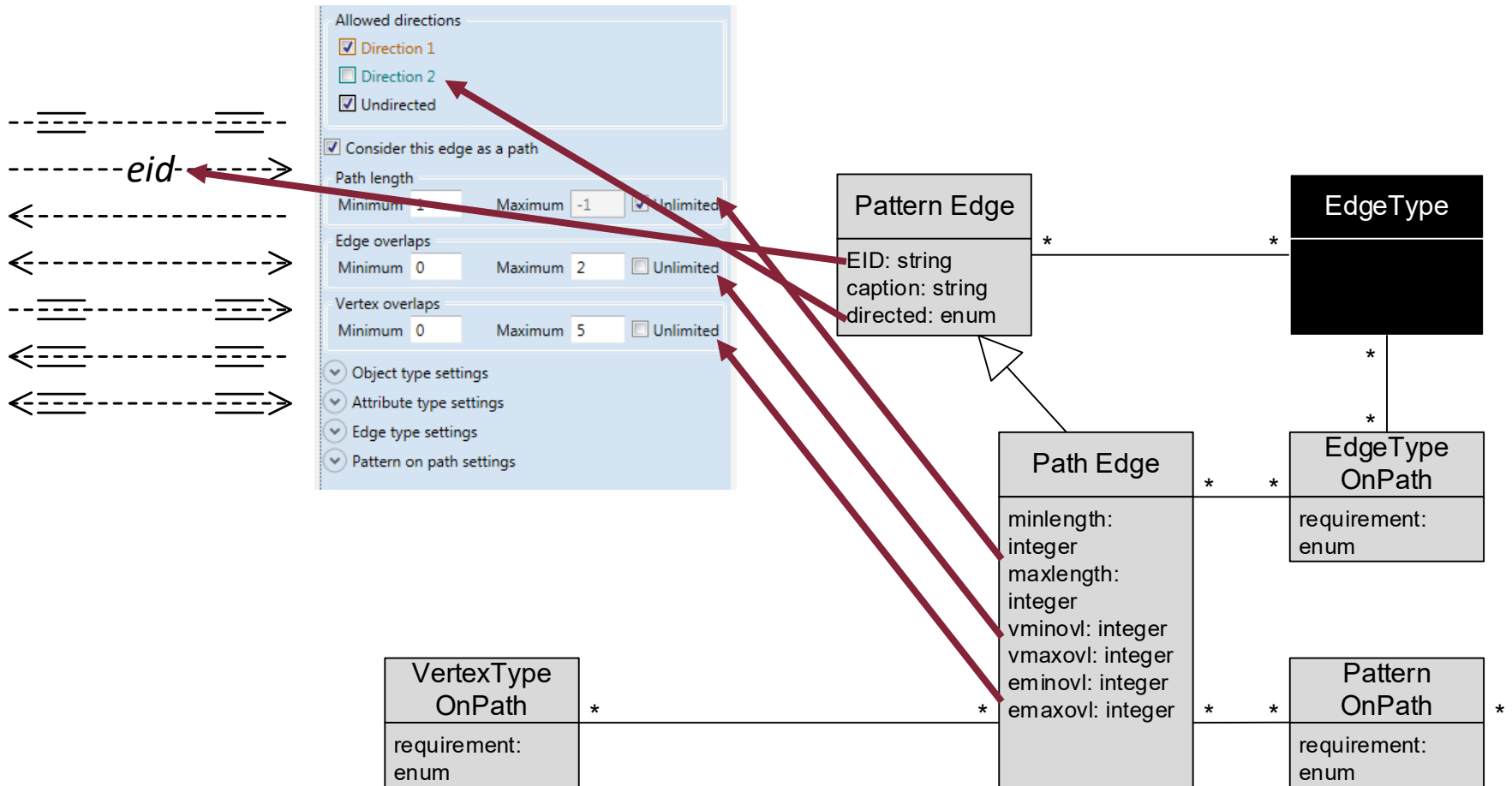
CONCRETE SYNTAX



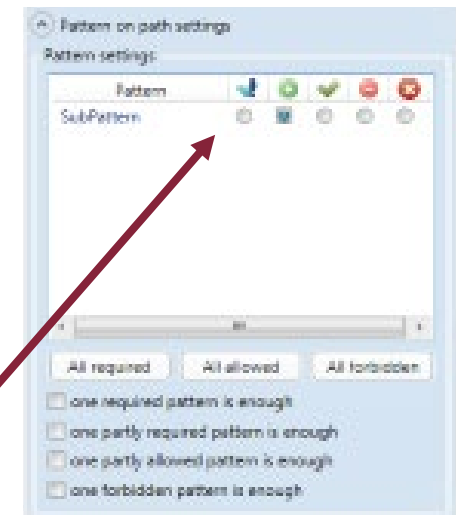
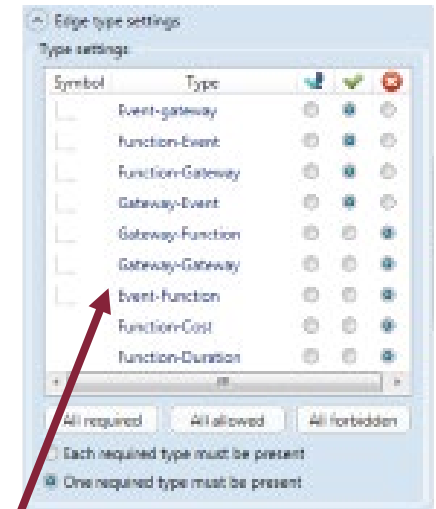
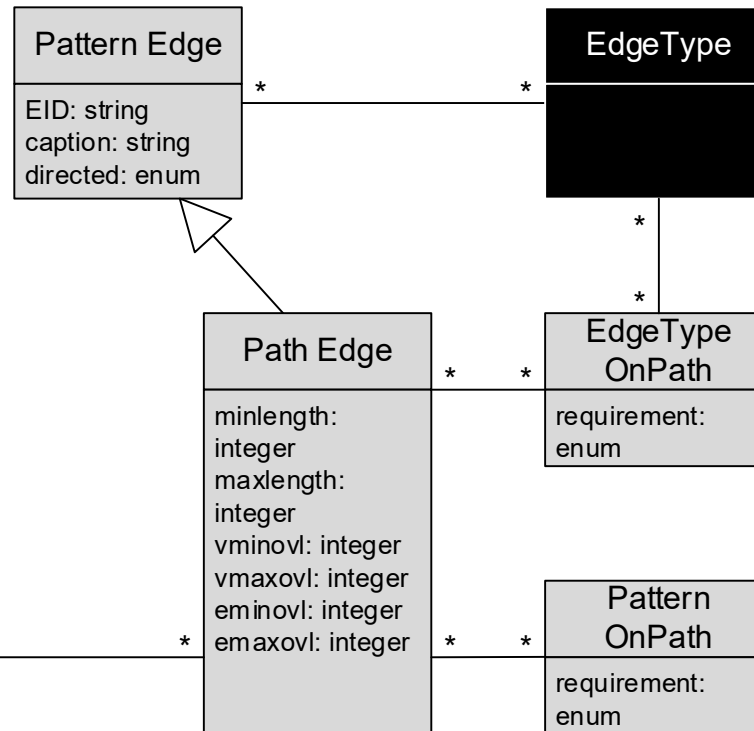
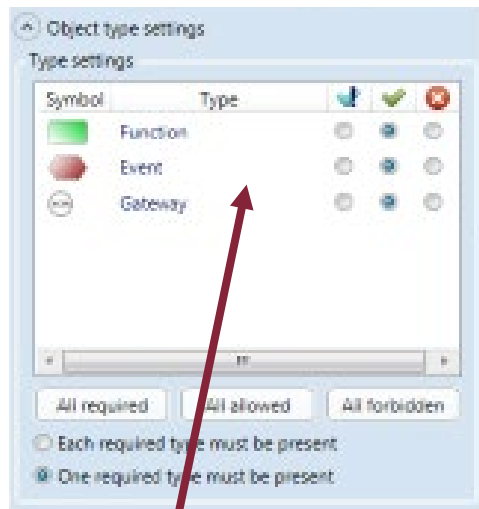
CONCRETE SYNTAX



CONCRETE SYNTAX



CONCRETE SYNTAX



- Global rules can be used to formulate constraints on patterns
- An exemplary Global Rule could look like this:
[a].[costs] < 100 OR [a].Caption LIKE “*service*”
Meaning: In order to be mapped to the pattern vertex named „a“, the attribute „cost“ of a vertex candidate must carry a value lower than 100, or its name should contain the word „service“.
- With rules, we can relate the properties of any pattern element to each other (e.g., “the name of the edge x must be different from the value of the attribute y of the vertex z”)

CONCRETE SYNTAX



GLOBAL RULES

- Global rules can be composed using the following concepts:

Syntax	Supported elements	Description
[a].Value	Instance Attributes	The value of an Instance Attribute
[a].[b].Value	Objects	The value of an element's Type Attribute
[a].Type	Objects, Instance Attributes, Edges	The type of the element as String
[a].[Caption].Value	Objects, Instance Attributes, Edges	The caption of the element

CONCRETE SYNTAX



GLOBAL RULES

- Concepts can be compared using the following operators

Data Type	Operator	Meaning
String	==	Two equal strings.
	!=	Two non-equal strings.
	LIKE	Like equal, but allows wildcards. [The wildcard '*' allows any number of arbitrary characters, '?' allows one arbitrary character.]
Number	==	Two equal numbers.
	!=	Two non-equal numbers.
	<	A number less than another.
	<=	A number less than or equal to another.
	>	A number greater than another.
	>=	A number greater than or equal to another.
Boolean	==	Two equal boolean values.
	!=	Two non-equal boolean values.

- Global rules also allow aggregate functions on the following expressions

Expression	Description
PREDECESSORS	Neighbors connected with a directed edge that points to the element.
SUCCESSORS	Neighbors connected with a directed edge that points to the neighbor.
UNDIRECTED_NEIGHBORS	Neighbors connected with an undirected edge.
NEIGHBORS	All neighbor nodes.
OUTGOING_EDGES	Directed edges pointing away from the element.
INCOMING_EDGES	Directed edges pointing to the element.
UNDIRECTED_EDGES	Edges that are connected to the element and have no direction.
EDGES	All edges connected to the element.
NODES	Requires that the element is a path, all nodes on that path.

CONCRETE SYNTAX

GLOBAL RULES



- This is how the aggregate functions are used

Function	Example Rule	Description
Count	<code>[a].UNDIRECTED_EDGES.Count()</code>	The number of elements specified. In the example, the number of undirected edges connected to a.
Max	<code>[a].SUCCESSORS.Max([duration])</code>	The maximum value of the attribute among the given elements. In the example, the value of the highest duration among the succeeding neighbors of a.
Min	<code>[a].NEIGHBORS.Min([service_time])</code>	The minimum value of the attribute among the given elements. In the example, the lowest value of service_time among all neighbors of a.
Avg	<code>[p].NODES.Avg([duration])</code>	The average of all attribute values of the specified elements. In the example, the average of all occurrences of the attribute duration on nodes on path p.
Sum	<code>[p].NODES.Sum([duration])</code>	The sum of all attribute values of the specified elements. In the example, the sum of all occurrences of the attribute duration on nodes on path p.

SEMANTICS (OUTLINE)



- Patterns are mapped against a set of models
- Every **pattern vertex** is mapped to exactly **one vertex** of a **pattern occurrence**
- Every **simple edge** of the pattern is mapped to exactly **one edge** of a **pattern occurrence**
- Every **path edge** of the pattern is mapped to exactly **one path** of a **pattern occurrence**
- Hence, the matching algorithm is based on:
 - **Subgraph isomorphism** considering vertex and edge properties (simple pattern edges)
 - **Non-disjunctive, configurable subgraph homomorphism** (path pattern edges) considering vertex and edge properties

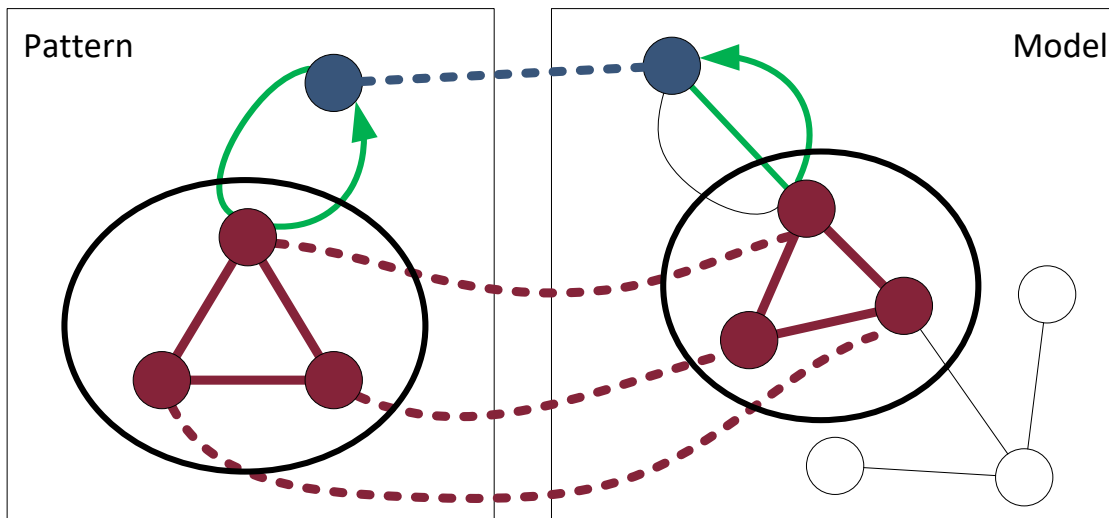
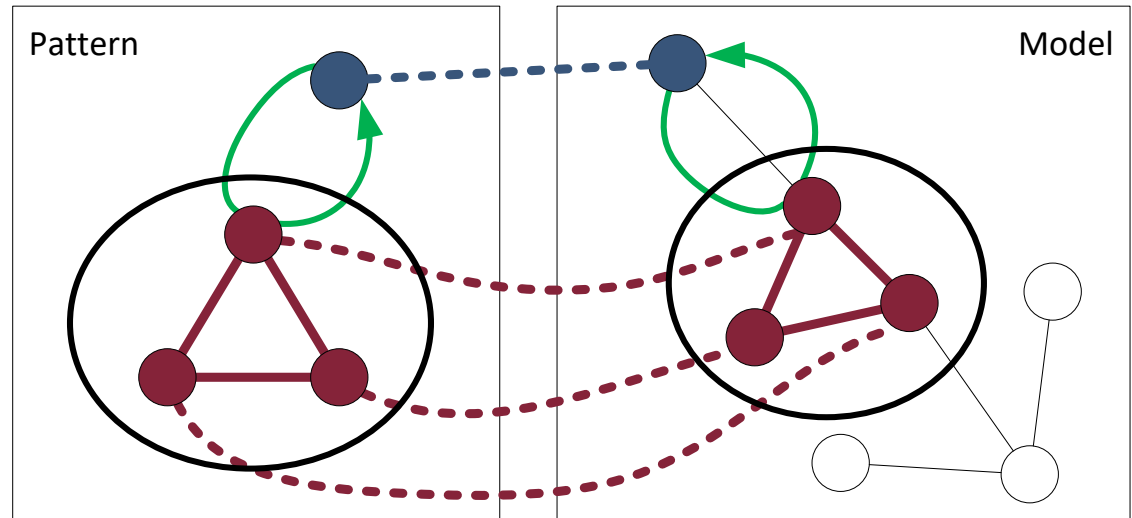
SEMANTICS (IN DETAIL)



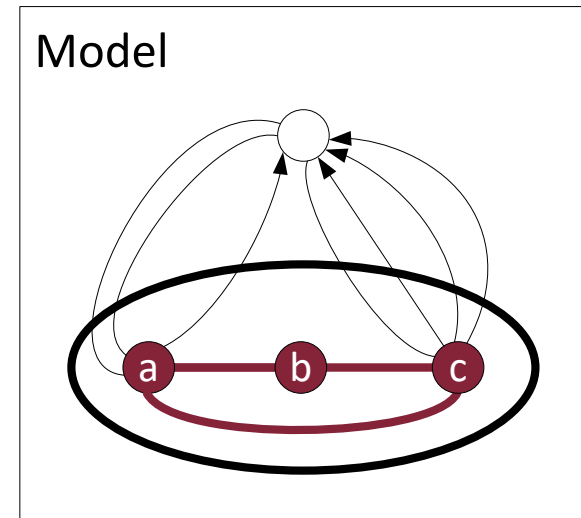
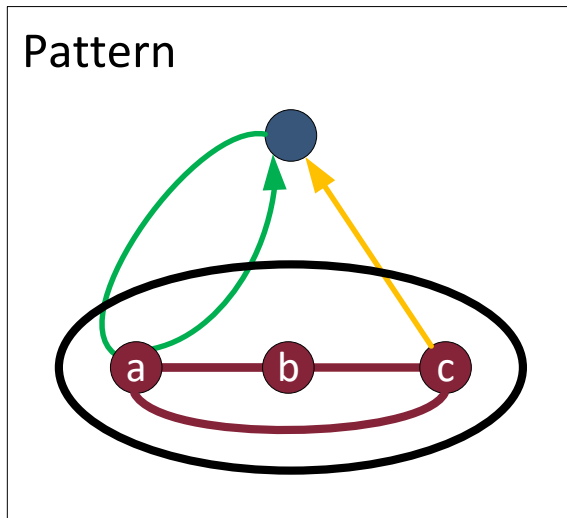
→ READING!

1. $v_{Qx}, v_{Qy} \in V_Q \wedge (v_{Qx}, v_{Qy}, n_Q) \in E_Q \wedge \varepsilon((v_{Qx}, v_{Qy}, n_Q)).minl = \varepsilon((v_{Qx}, v_{Qy}, n_Q)).maxl = 1 \wedge dir = \varepsilon((v_{Qx}, v_{Qy}, n_Q)).dir \Leftrightarrow$
2. $\exists^1 v_{ix}' \in V_i', \exists^1 v_{iy}' \in V_i': \alpha_i'(v_{ix}') \in \delta(v_{Qx}).vtypes, \alpha_i'(v_{iy}') \in \delta(v_{Qy}).vtypes, \delta(v_{Qx}).vcaption \sim \chi_i'(v_{ix}'),$
 $\delta(v_{Qy}).vcaption \sim \chi_i'(v_{iy}')$
3. $(\exists^1 (v_{ix}', v_{iy}', n_M) \in E_i': \varepsilon((v_{Qx}, v_{Qy}, n_Q)).ecaption \sim \chi_i'((v_{ix}', v_{iy}', n_M)), org \in \varepsilon((v_{Qx}, v_{Qy}, n_Q)).dir,$
 $\beta_i'((v_{ix}', v_{iy}', n_M)) \in \varepsilon((v_{Qx}, v_{Qy}, n_Q)).etypesr) \vee$
4. $(\exists^1 (v_{iy}', v_{ix}', n_M) \in E_i': \varepsilon((v_{Qx}, v_{Qy}, n_Q)).ecaption \sim \chi_i'((v_{iy}', v_{ix}', n_M)), opp \in \varepsilon((v_{Qx}, v_{Qy}, n_Q)).dir,$
 $\beta_i'((v_{iy}', v_{ix}', n_M)) \in \varepsilon((v_{Qx}, v_{Qy}, n_Q)).etypesr) \vee$
5. $(\exists^1 \{v_{ix}', v_{iy}', n_M\} \in E_i': \varepsilon((v_{Qx}, v_{Qy}, n_Q)).ecaption \sim \chi_i'(\{v_{ix}', v_{iy}', n_M\}), none \in \varepsilon((v_{Qx}, v_{Qy}, n_Q)).dir,$
 $\beta_i'(\{v_{ix}', v_{iy}', n_M\}) \in \varepsilon(\{v_{Qx}, v_{Qy}, n_Q\}).etypesr),$
6. $v_{Qx}, v_{Qy} \in V_Q \wedge (v_{Qx}, v_{Qy}, n_Q) \in E_Q \wedge \varepsilon((v_{Qx}, v_{Qy}, n_Q)).minl > 1, \varepsilon((v_{Qx}, v_{Qy}, n_Q)).maxl \neq 1 \wedge dir = \varepsilon((v_{Qx}, v_{Qy}, n_Q)).dir \Leftrightarrow$
7. $\exists^1 v_{ix}' \in V_i', \exists^1 v_{iy}' \in V_i': \alpha_i'(v_{ix}') \in \delta(v_{Qx}).vtypes, \alpha_i'(v_{iy}') \in \delta(v_{Qy}).vtypes, \delta(v_{Qx}).vcaption \sim \chi_i'(v_{ix}'),$
 $\delta(v_{Qy}).vcaption \sim \chi_i'(v_{iy}')$
8. $\exists^1 vertices(path_k(v_{xi}', v_{yi}')_{dir}) \in V_i', \exists^1 edges(path_k(v_{xi}', v_{yi}')_{dir}) \in E_i':$
 $\varepsilon((v_{Qx}, v_{Qy}, n_Q)).minl \leq length(path_k(v_{xi}', v_{yi}')_{dir}) \leq \varepsilon((v_{Qx}, v_{Qy}, n_Q)).maxl \wedge$
9. $\varepsilon((v_{Qx}, v_{Qy}, n_Q)).minvo \leq vovl(path_k(v_{xi}', v_{yi}')_{dir}) \leq \varepsilon((v_{Qx}, v_{Qy}, n_Q)).maxvo \wedge$
 $\varepsilon((v_{Qx}, v_{Qy}, n_Q)).mineo \leq eovl(path_k(v_{xi}', v_{yi}')_{dir}) \leq \varepsilon((v_{Qx}, v_{Qy}, n_Q)).maxeo \wedge$
10. $\forall t \in \varepsilon((v_{Qx}, v_{Qy}, n_Q)).vtypesr \exists v \in vertices(path_k(v_{xi}', v_{yi}')_{dir}) \mid \alpha_i'(v) = t \wedge$
 $\nexists v \in vertices(path_k(v_{xi}', v_{yi}')_{dir}) \mid \alpha_i'(v) \in \varepsilon((v_{Qx}, v_{Qy}, n_Q)).vtypesf \wedge$
11. $\forall t \in \varepsilon((v_{Qx}, v_{Qy}, n_Q)).etypesr \exists e \in edges(path_k(v_{xi}', v_{yi}')_{dir}) \mid \beta_i'(e) = t \wedge$
 $\nexists e \in edges(path_k(v_{xi}', v_{yi}')_{dir}) \mid \beta_i'(e) \in \varepsilon((v_{Qx}, v_{Qy}, n_Q)).etypesf \wedge$
12. $((vertices(path_k(v_{xi}', v_{yi}')_{dir}) \cup edges(path_k(v_{xi}', v_{yi}')_{dir})) \cap M_n'(\theta.Q) = M_p'(\theta.Q) \mid \theta.constraint = rer$
 $\wedge \theta.constraint = pre$

MAPPING EXAMPLES



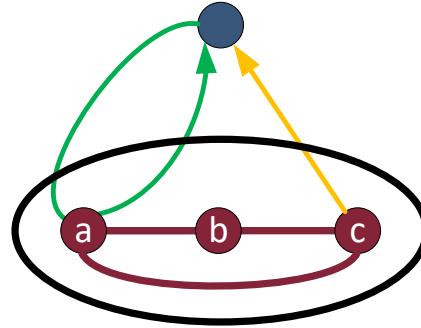
MAPPING EXAMPLES



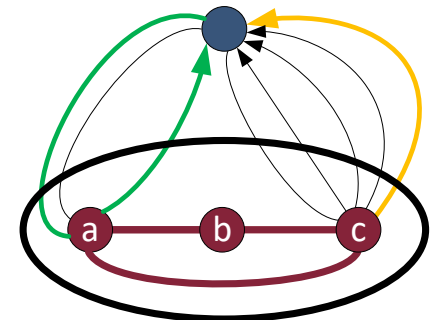
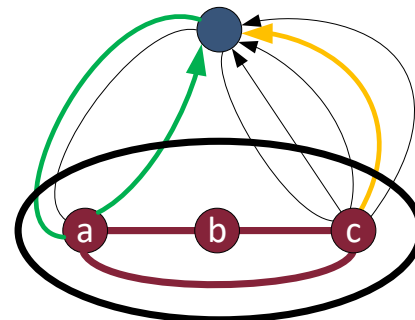
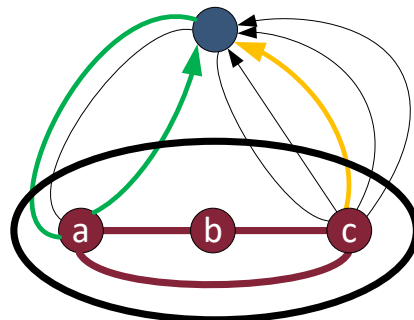
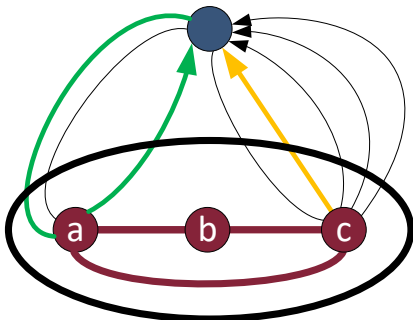
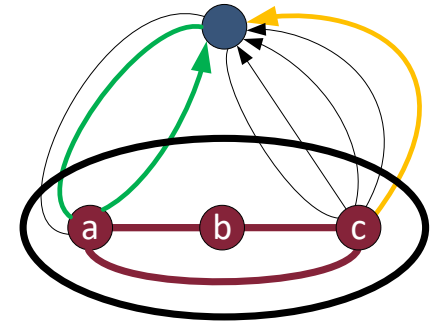
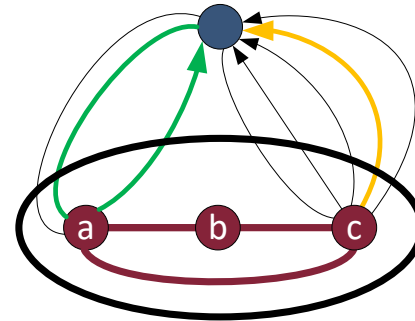
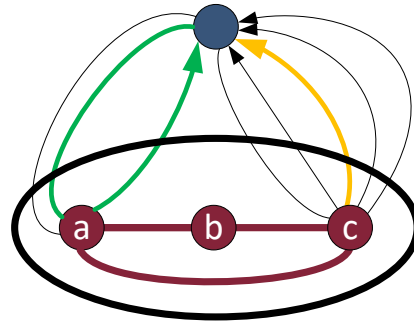
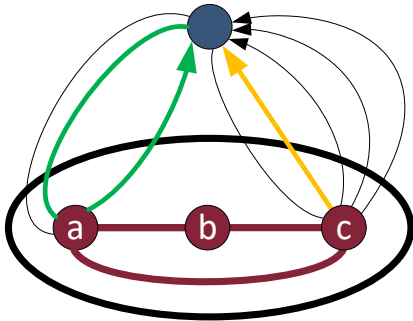
MAPPING EXAMPLES



Pattern



Model Mapping Possibilities



- A **new version** of DMQL also includes concepts to easier handle execution semantics of process models
 - Forbidden vertices
 - Forbidden edges
 - Forbidden paths
- For instance, a vertex **x** that is connected to a forbidden vertex **y** means that a mapping candidate for **x** is mapped only if it is **NOT** connected to another vertex that matches the properties of **y**

AGENDA



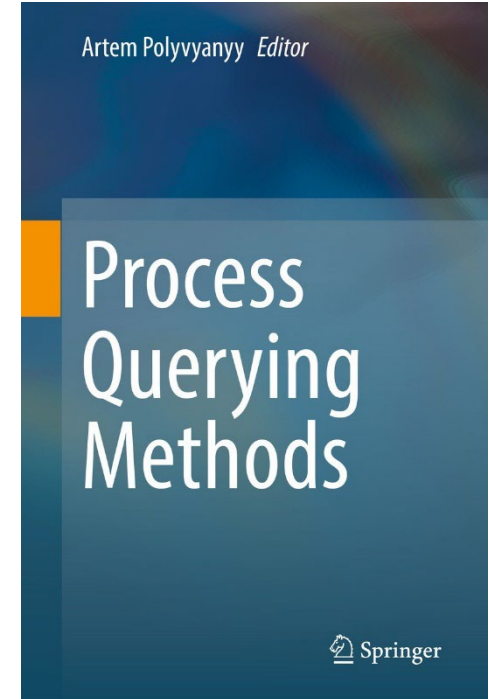
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- Delfmann, P.; Breuker, D.; Matzner, M.; Becker, J.: Supporting Information Systems Analysis Through Conceptual Model Query - The Diagramed Model Query Language (DMQL). *Communications of the Association for Information Systems (CAIS)* 37 (2015) 1, pp. 473-509.

OTHER MODEL QUERY LANGUAGES



- aFSA
 - APQL
 - BeehiveZ
 - BPMN VQL
 - BPMN-Q
 - BPQL
 - BP-QL
 - Cypher (Neo4j)
 - EMF-IncQuery
 - IPM-PQL
 - OCL
 - PPSL
 - SPARQL
 - Untanglings
 - VMQL
- Book on Process Query Languages



- Also: see Related Work Sections of GMQL and DMQL Readings

BUSINESS PROCESS MANAGEMENT

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INSTITUTE FOR **IS** RESEARCH

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