Semantic Web - Tutorial #3

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Correction Regarding Assignment 2A, Task 2.1

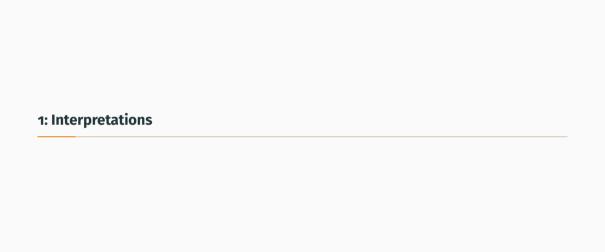
Correction Regarding Assignment 2A, Task 2.1

Consider the interpretation $I=(\Delta^I,\cdot^I)$, where

$$\begin{split} \Delta^I &= \{n_1, n_2, p_1, j_1, i_1\} \\ \mathsf{Nurse}^I &= \{n_1, n_2\} \\ \mathsf{Patient}^I &= \{p_1\} \\ \mathsf{Janitor}^I &= \{j_1, n_2, p_1, i_1\} \\ \mathsf{Illness}^I &= \{i_1\} \\ \mathsf{Physician}^I &= \{n_1, n_2, p_1, i_1\} \\ \mathsf{takesCareOf}^I &= \{(n_1, p_1), (j_1, p_1)\} \\ \mathsf{canTreat}^I &= \{(j_1, i_1), (j_1, p_1)\} \end{split}$$

I satisfies \mathcal{T}_1 but does not satisfy neither of the statements (4)–(6). Both nurses n_1 and n_2 only take care of patients (n_2 satisfies it vacuously, since it takes care of no patient), but it violates (4) since n_2 takes care of no patient. Moreover, (5) is violated, because j_1 , a janitor, takes care of the patient n_1 . Statement (6) is violated, because n_2 treats the illness n_3 .





1: Interpretations

Consider the following knowledge base $\mathcal{K}_1=(\mathcal{T}_1,\mathcal{A}_1)$ in \mathcal{ALC} :

$$\mathcal{T}_1 = \{ \exists R.A \sqsubseteq B,$$

$$A \sqcap D \equiv E,$$

$$A \sqsubseteq \neg \exists Q.D \}$$

$$\mathcal{A}_1 = \{ y : A, (x, y) : R, (y, z) : Q \}$$

Task: Formally define an interpretation I that satisfies \mathcal{K}_1 .

$$\begin{split} I &= (\Delta^I, \cdot^I) \text{ with } \\ \Delta^I &= \{x_1, y_1, z_1\} \\ x^I &= \{x_1\} \\ y^I &= \{y_1\} \\ z^I &= \{z_1\} \\ A^I &= \{x_1, y_1, z_1\} \\ B^I &= \{x_1\} \\ D^I &= \emptyset \\ E^I &= \emptyset \\ R^I &= \{(x_1, y_1)\} \\ Q^I &= \{(y_1, z_1)\} \end{split}$$

1: Interpretations

Task: Demonstrate that I satisfies \mathcal{K}_1 .

- \blacktriangleright We have to show that I is a model of both \mathcal{A}_1 and \mathcal{T}_1 .
- ▶ For A_1 , show that:
 - ▶ for each $a:C,a^I\in C^I$ and
 - for each $(a,b):r,(a^I,b^I)\in r^I)$ (where C is a concept name, and r a role name)
- 1) Only one concept assertion: y:A. From \cdot^I , we have $y^I=y_1$, and $y_1\in A^I.$ Thus, $y^I\in A^I.$
- 2.1) (x,y):R. We have $x^I=x_1$ and $y^I=y_1$, and $(x_1,y_1)\in R^I.$ Thus, $(x^I,y^I)\in R^I.$
- 2.2) (y,z):Q. We have $y^I=y_1$ and $z^I=z_1$, and $(y_1,z_1)\in Q^I.$ Thus, $(y^I,z^I)\in Q^I.$

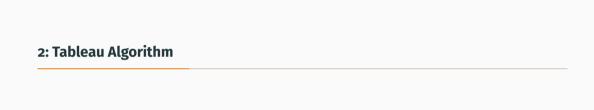


1: Interpretations

For \mathcal{T}_1 , we have to show that $C^I \subseteq D^I$, for each $C \subseteq D$. (Note that $C \equiv D$ is equivalent to $C \subseteq D$ and $D \subseteq C$.)

- ▶ $\exists R.A \sqsubseteq B$. We have $(\exists R.A)^I = \{x_1\}$. Also, $B^I = \{x_1\}$. It follows that $\{x_1\} \subseteq \{x_1\}$. Thus, $(\exists R.A)^I \subseteq B^I$.
- $lacksquare A \cap D \equiv E.$ In this case, it suffices to show that $A^I \cap D^I = E^I.$ We have $A^I = \{x_1, y_1, z_1\}$ and $D^I = \emptyset.$ Thus, $A^I \cap D^I = \emptyset = E^I.$
- ▶ $A \sqsubseteq \neg \exists Q.D$ }. We have $(\exists Q.D)^I = \emptyset$. Thus, $(\neg \exists Q.D)^I = \Delta^I$. Also, $A^I = \{x_1, y_1, z_1\} \subseteq \Delta^I$. It follows that $A^I \subseteq (\neg \exists Q.D)^I$.





2: Tableau Algorithm

Consider the following Knowledge base $\mathcal{K}_2 = (\mathcal{T}_2, \mathcal{A}_2)$ with

$$\mathcal{T}_2 = \{ A \sqsubseteq D, \\ \neg (B \sqcup \neg D) \sqsubseteq E, \\ D \sqsubseteq \neg \exists R.E, \\ B \sqcap E \sqsubseteq \bot, \\ \neg \neg D \sqsubseteq E \}$$

$$\mathcal{A}_2 = \{ x : \neg (\neg D \sqcap \neg E), y : A \sqcap B, (y, x) : R \}$$

Task: Transfer both \mathcal{T}_2 and \mathcal{A}_2 to NNF.

$$\neg (B \sqcup \neg D) \sqsubseteq E$$

$$\blacktriangleright \ \neg B \sqcap D \sqsubseteq E$$

$$D \sqsubseteq \neg \exists R.E$$

$$ightharpoonup D \sqsubseteq \forall R. \neg E$$

$$\neg \neg D \sqsubseteq E$$

$$ightharpoonup D \sqsubseteq E$$

$$x: \neg(\neg D \sqcap \neg E)$$

$$ightharpoonup x:D\sqcup E$$



2: Tableau Algorithm

Is \mathcal{K}_2 consistent or not? In either case, apply the tableau algorithm for your answer. Document and explain each step you take when applying the tableau algorithm.



Solution 1

TBox $\mathcal{T}_2 = \{A \sqsubseteq D,$ $\neg B \sqcap D \sqsubseteq E$, $D \sqsubseteq \forall R. \neg E$, $B \sqcap E \sqsubseteq \bot$,

 $D \sqsubseteq E$

ABox

$$\mathcal{A}_2 = \{x : D \sqcup E, \\ y : A \sqcap B\}$$

$$\{x: D \sqcup E, \\ y: A \sqcap B, \}$$

$$A \sqcap B$$
,

2. 3. 4.

5. 7.

8.

9.

10.

1.

11.

12.

13.

14. 15.

16.

17.

 $x : \neg B$ \otimes

18. 13, 17

19.

20.

 $x : \neg D$ x:E 17, 19 \otimes \otimes

19, 20

 $x: \neg B \sqcup \neg E$

 $x : D \sqcup E$

 $y:A\sqcap B$

(y,x):R

y:A

u : B

 $y: \neg A \sqcup D$

 $y : \neg D$

 \otimes

 $x: B \sqcup \neg D$

 $x : \bot$

x : E

 \otimes

x : D

13, 14

y:D

 $y : \neg A$

 \otimes

4,7

x : B

 $x: (\neg B \sqcup \neg E) \sqcup \bot$

 $x : \neg E$

 $x : \neg D \sqcup E$

17, 20

 $y: \neg D \sqcup (\forall R. \neg E)$

 $y : \forall R. \neg E$

 $x : \neg E$ $x: \neg(\neg B \sqcap D) \sqcup E$

> x : E \otimes

 $x : \neg D$ 10, 12 x : E

10, 14

(⊔, 16) \sqsubseteq , $D \sqsubseteq E$ (⊔, 18)

(山. 15)

ABox

ABox

ABox

(□, 2)

(□, 2)

(⊔, 6)

(⊔, 8)

(∀, 9)

(山, 11)

(⊔, 12)

 $(\Box, A \Box D)$

 $(\Box, D \Box \forall R. \neg E)$

 $(\Box, \neg B \sqcap D \sqsubseteq E)$

 $(\Box, B \sqcap E \Box \bot)$

Solution 2

TBox

$$\mathcal{T}_2 = \{ A \sqsubseteq D, \\ \neg B \sqcap D \sqsubseteq E, \\ D \sqsubseteq \forall R. \neg E, \\ B \sqcap E \sqsubseteq \bot, \\ D \sqsubseteq E \}$$

ABox

$$\mathcal{A}_2 = \{x : D \sqcup E, \\ y : A \sqcap B, \\ (y, x) : R\}$$

1.
$$x:D\sqcup E$$

2.
$$y: A \sqcap B$$

3. $(y, x): R$

$$y: \neg A \sqcup D$$

7.
$$y: \neg A$$
 $y: D$
8. \otimes $y: \neg D \sqcup E$

$$\otimes$$
 $y: \neg D \sqcup E$

$$y: \neg D \qquad y$$

9.
$$y:\neg D \qquad y:E$$
 10.
$$\otimes \qquad y:\neg (B\sqcap E)\sqcup\bot$$

11.

12.
$$y: \neg B \quad y: \neg E$$

5, 12

 $y: \neg B \sqcup \neg E$

9, 12

 $y: \bot$

$$(\sqcap, 2)$$

 $(\sqcap, 2)$
 $(\sqsubseteq, A \sqsubseteq D)$

ABox

$$\subseteq$$
, $A \sqsubseteq D$)

$$(\sqcup, 6)$$

 $(\sqsubseteq, D \sqsubseteq E)$

$$(\sqsubseteq, B \sqcap E \sqsubseteq \bot)$$



Let
$$\mathcal{S}_1=(N_C,N_R,N_O)$$
 be a signature where
$$N_C=\{A,B,D,E,F,G\}$$

$$N_R=\{P,Q\}$$

$$N_O=\{x,y\}$$

Consider the following *Tbox* and *Aboxes*:

$$\mathcal{T}_{3} = \{ A = \neg (E \sqcup F), \\ B \sqcup D = A, \\ E = \exists P.G, \\ A \sqcap \exists P.G = \bot, \\ E \sqcup F = \neg \exists Q.G \}$$

$$\mathcal{A}_{3} = \{ x : G, \\ (x,y) : Q \}$$

$$\mathcal{A}_{4} = \{ x : G, \\ y : B, \\ (x,y) : Q \}$$



Task: Is the knowledge base $K_3 = (T_3, A_3)$ consistent? If your answer is positive, then define an interpretation I that satisfies K_3 . Otherwise, prove that K_3 is inconsistent.

$$\mathcal{K}_{3} = (\mathcal{T}_{3}, \mathcal{A}_{3})$$

$$\mathcal{T}_{3} = \{ A = \neg(E \sqcup F), \\
B \sqcup D = A, \\
E = \exists P.G, \\
A \sqcap \exists P.G = \bot, \\
E \sqcup F = \neg \exists Q.G \}$$

$$\mathcal{A}_{3} = \{ x : G, \\
(y,x) : Q \}$$

Yes, \mathcal{K}_3 is consistent. We define $I = (\Delta^I, \cdot^I)$ with

$$A^{I} = \{x_{1}, y_{1}\}$$

$$B^{I} = \emptyset$$

$$D^{I} = \{y_{1}\}$$

$$E^{I} = \emptyset$$

$$F^{I} = \emptyset$$

$$G^{I} = \{x_{1}\}$$

$$Q^{I} = \{(y_{1}, x_{1})\}$$

$$Q^{I} = \{y_{1}\}$$

$$Q^{I} = \{y_{1}\}$$

Task: Is the knowledge base $K_4 = (\mathcal{T}_3, \mathcal{A}_4)$ consistent? If your answer is positive, then define an interpretation I that satisfies K_4 . Otherwise, prove that K_4 is inconsistent.

```
\mathcal{K}_{4} = (\mathcal{T}_{3}, \mathcal{A}_{4})

\mathcal{T}_{3} = \{ A = \neg (E \sqcup F), \\
B \sqcup D = A, \\
E = \exists P.G, \\
A \sqcap \exists P.G = \bot, \\
E \sqcup F = \neg \exists Q.G \}

\mathcal{A}_{4} = \{ x : G, \\
y : B, \\
(y, x) : Q \}
```

▶ The knowledge base K_4 is inconsistent. Let us assume, for contradiction, that K_4 is consistent. Then, there is an interpretation I that satisfies K_4 .



1.
$$\xrightarrow{(A_4)} x : G \text{ and } y : B \text{ and } (y, x) : P$$

$$2. \stackrel{(1.)}{\Longrightarrow} x^I \in G^I \text{ and } y^I \in B^I \text{ and } (y^I, x^I) \in P^I$$

$$3. \stackrel{(\mathcal{T}_3)}{\Longrightarrow} B \sqcup D \sqsubseteq A$$

$$4. \stackrel{(3.)}{\Longrightarrow} B^I \cup D^I \subseteq A^I$$

$$5. \stackrel{(2.,4.)}{\Longrightarrow} y^I \in A^I$$

$$6. \stackrel{(\mathcal{T}_3)}{\Longrightarrow} A \sqcap \exists P.G \sqsubseteq \bot$$

$$7. \stackrel{(6.)}{\Longrightarrow} A^I \cap (\exists P.G)^I \subseteq \emptyset$$

$$8. \xrightarrow{(2.,5.)} y^I \in A^I \text{ and } x^I \in G^I \text{ and } (y^I,x^I) \in P^I$$

9.
$$\xrightarrow{(7.,8.)} \{y^I\} \cap \{y^I\} \subseteq \emptyset$$
 contradiction



Task: Prove that $(\mathcal{T}_3, \mathcal{A}_3) \models y : A$.

▶ We need to show that for every interpretation I, if $I \models (\mathcal{T}_3, \mathcal{A}_3)$, then $I \models y : A$. Let us assume $I \models (\mathcal{T}_3, \mathcal{A}_3)$. We will show that $y^I \in A^I$. As $I \models (\mathcal{T}_3, \mathcal{A}_3)$, we get that $I \models \mathcal{T}_3$ and $I \models \mathcal{A}_3$.



$$1. \stackrel{(\mathcal{A}_3)}{\Longrightarrow} x: G \text{ and } (y,x): Q$$

$$2. \stackrel{(1.)}{\Longrightarrow} x^I \in G^I \text{ and } (y^I, x^I) \in Q^I$$

$$3. \stackrel{(\mathcal{T}_3)}{\Longrightarrow} E \sqcup F \sqsubseteq \neg \exists Q.G$$

$$5. \stackrel{(4.)}{\Longrightarrow} E^I \cup F^I \subseteq \Delta^I \setminus (\exists Q.G)^I$$

$$6. \stackrel{(5.)}{\Longrightarrow} E^I \cup F^I \subseteq \Delta^I \setminus \{y^I\}$$

$$7. \stackrel{(\mathcal{T}_3)}{\Longrightarrow} A \equiv \neg (E \sqcup F)$$

$$8. \stackrel{(7.)}{\Longrightarrow} A^I = \neg (E^I \cup F^I)$$

$$9. \stackrel{(6.)}{\Longrightarrow} y^I \not\in (E^I \cup F^I)$$

$$10. \xrightarrow{(6.,9.)} y^I \in \neg (E^I \cup F^I)$$

11.
$$\stackrel{(7.,10.)}{\Longrightarrow} y^I \in A^I$$

