

Complexity of Matrix Multiplication (Contd...)

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Lecture 3

Recall :-

$$M_{\langle k, m, n \rangle} : \mathbb{C}^{k \times m} \times \mathbb{C}^{m \times n} \rightarrow \mathbb{C}^{k \times n}$$

$(A, B) \mapsto A \times B$

is the matrix mult. map

$M_{\langle k, m, n \rangle}$ is a bilinear map, so can be thought of as a tensor in
 $(\mathbb{C}^{k \times m})^* \otimes (\mathbb{C}^{m \times n})^* \otimes (\mathbb{C}^{k \times n})$.

$$\omega = \inf \left\{ \gamma \in \mathbb{R} \mid \underset{\substack{\uparrow \\ \text{rank}}}{R}(M_{\langle n, n, n \rangle}) = O(n^\gamma) \right\}$$

We want to study ω

===== Machinery that gives you an u.b. on ω , given u.b. on
 $R(M_{\langle k, m, n \rangle})$ for fixed k, m, n =====

Recall Strassen showed $R(M_{\langle 2, 2, 2 \rangle}) \leq 7$.

Defn [Permutation of tensors] Let $t \in A \otimes B \otimes C$, where

$$t = \sum_{j=1}^k t_j, \text{ where } t_j = a_j^{(1)} \otimes a_j^{(2)} \otimes a_j^{(3)}$$

Let $\sigma \in S_3$

$$\sigma(t) = \sum_{j=1}^k \sigma(t_j), \text{ where } \sigma(t_j) = a_j^{(\sigma(1))} \otimes a_j^{(\sigma(2))} \otimes a_j^{(\sigma(3))}$$

Lemma $R(\sigma(t)) = R(t)$. Proof obvious

Defn [Restriction of tensors]. Let $t \in A \otimes B \otimes C$ be such that

$$t = \sum_{i=1}^k a_i \otimes b_i \otimes c_i.$$

Let $t' \in A' \otimes B' \otimes C'$ be such that

$$\rightarrow t' = \sum_{i=1}^k f_1(a_i) \otimes f_2(b_i) \otimes f_3(c_i), \text{ where } f_1: A \rightarrow A', f_2: B \rightarrow B',$$

$f_3: C \rightarrow C'$ are homomorphisms.

$\dots \dots \dots$ t' is called the restriction of t , and we denote $t' \leq t$.

$t_j : C \rightarrow C$ are isomorphisms.

We say t' is a restriction of t , and we denote $t' \leq t$.

Lemma If $t' \leq t$, $R(t') \leq R(t)$. We have equality when f_1, f_2, f_3 are isomorphisms.

Proof by defn \square

Lemma. Take $\alpha \in S_3$. $R(M_{\langle k, m, n \rangle}) = R(M_{\alpha \langle k, m, n \rangle})$

Proof

$$M_{\langle k, m, n \rangle} := \mathbb{C}^{K \times m} \times \mathbb{C}^{m \times n} \rightarrow \mathbb{C}^{k \times n}$$

$$(\mathbb{C}^{K \times m})^* \otimes (\mathbb{C}^{m \times n})^* \otimes \mathbb{C}^{k \times n}$$

\curvearrowleft Perm \curvearrowright

$$\begin{array}{ccc} m \times n & & K \times m \\ \downarrow \text{transpose} & & \downarrow \text{transpose} \\ n \times m & & m \times K \end{array}$$

\square

Defn [Direct sum of tensors] $t \in F^k \otimes F^m \otimes F^n$, $t' \in F^{k'} \otimes F^{m'} \otimes F^{n'}$.

$$t = \sum_{i=1}^n t_i, \text{ where } t_i = \underbrace{a_i}_{F^k} \otimes \underbrace{b_i}_{F^m} \otimes \underbrace{c_i}_{F^n}$$

a_i is a k -dimensional co-ordinate vector.

$$\rightarrow t' = \sum_{i=1}^{k'} t'_i \quad " \quad t'_i = \underbrace{a'_i}_{F^{k'}} \otimes \dots$$

a'_i are k' -dim co-ordinate vectors

$$\underline{a_i} = (a_i, -\frac{\overset{0's}{\cdots}}{k'}, \dots), \underline{b_i} \underset{F^{m+n'}}{\underset{\uparrow}{\cdots}}, \underline{c_i} \underset{F^{m+n'}}{\underset{\uparrow}{\cdots}}$$

$$\underline{a_i} \in F^{k+k'}$$

$$\underline{a'_i} = (-\frac{\overset{0's}{\cdots}}{k'}, a'_i), \underline{b'_i} = \dots, \underline{c'_i}$$

$$a' \quad b' \quad c'$$

$$t \oplus t' = \sum_{i=1}^k a_i \otimes b_i \otimes c_i + \sum_{i=1}^{k'} a'_i \otimes b'_i \otimes c'_i$$



Lemma $R(t \oplus t') \leq R(t) + R(t')$

Proof by defn \square

Conjecture This inequality is tight. (Strassen's additivity conjecture)

$$\begin{array}{l} \text{Defn [tensor product of tensors]} \quad t \in \mathbb{F}^k \otimes \mathbb{F}^m \otimes \mathbb{F}^n, \quad t' \in \mathbb{F}^{k'} \otimes \mathbb{F}^{m'} \otimes \mathbb{F}^{n'} \\ t = \sum_{i=1}^k t_i \quad t_i = a_i \otimes b_i \otimes c_i \\ \rightarrow a_i = (a^{(1)}, \dots, a^{(k)}) \quad | \quad t' = \sum_{i=1}^{k'} t'_i \quad t'_i = a'_i \otimes b'_i \otimes c'_i \\ \uparrow \mathbb{F}^k \quad \rightarrow a'_i = (\underline{a}^{(1)}, \dots, \underline{a}^{(k')}) \end{array}$$

$$(a_i \otimes a'_i)_{i,j} = a^{(i)} \times a^{(j)}$$

\uparrow $\mathbb{F}^{kk'}$ extend linearly to obtain $t \otimes t' \in \mathbb{F}^{kk'} \otimes \mathbb{F}^{mn'} \otimes \mathbb{F}^{nn'}$

Lemma $R(t \otimes t') \leq R(t)R(t')$

Proof by defn \square

$$\text{Observe } M_{\langle k, m, n \rangle} \otimes M_{\langle k', m', n' \rangle} = M_{\langle kk', mm', nn' \rangle}$$

Thm $R(M_{\langle k, m, n \rangle}) \leq r$. Then $w \leq 3 \log_{kmn} r$

$$\text{Proof } R(M_{\langle k, m, n \rangle} \otimes M_{\langle m, n, k \rangle} \otimes M_{\langle n, k, m \rangle}) \leq r^3$$

$$\Rightarrow R(M_{\langle kmn, kmn, kmn \rangle}) \leq r^3 = (kmn)^{3 \log_{kmn} r}$$

$$w \leq 3 \log_{kmn} r \quad \square$$

Strassen $R(M_{\langle 2,2,2 \rangle}) \leq 7 \Rightarrow \omega \leq 3 \log_8 7 \approx 2.8074$

Pan $R(M_{\langle 70,70,70 \rangle}) \leq 143640 \Rightarrow \omega \leq 2.796$

↑
There is an intuitive way.

Timeline of matrix multiplication exponent

Year	Bound on omega	Authors
1969	2.8074	Strassen ^[1]
1978	2.796	Pan ^[11]
1979	2.780	Bini, Capovani ^[it] , Romani ^[12]
1981	2.522	Schönhage ^[13]
1981	2.517	Romani ^[14]
1981	2.496	Coppersmith, Winograd ^[15]
1986	2.479	Strassen ^[16]
1990	2.3755	Coppersmith, Winograd ^[17]
2010	2.3737	Stothers ^[18]
2013	2.3729	Williams ^{[19][20]}
2014	2.3728639	Le Gall ^[21]
2020	2.3728596	Alman, Williams ^[3]
2022	2.37188	Duan, Wu, Zhou ^[2]

Border rank

Consider $M_{\langle 2,2,3 \rangle}$

$$\begin{array}{ccc}
 2 \times 2 & \times & 2 \times 3 \longrightarrow 2 \times 3 \\
 \begin{bmatrix} x_{11} & x_{12} \\ x_{21} & x_{22} \end{bmatrix} & \times & \begin{bmatrix} y_{11} & y_{12} & y_{13} \\ y_{21} & y_{22} & y_{23} \end{bmatrix} = \begin{bmatrix} z_{11} & z_{12} & z_{13} \\ z_{21} & z_{22} & z_{23} \end{bmatrix}
 \end{array}$$

$M_{\text{reduced } \langle 2,2,2 \rangle}$

$$\begin{bmatrix} x_{11} & x_{12} \\ x_{21} & x_{22} \end{bmatrix} \times \begin{bmatrix} y_{11} & y_{12} \\ y_{21} & y_{22} \end{bmatrix} = \begin{bmatrix} z_{11} & z_{12} \\ z_{21} & * \end{bmatrix}$$

Idea Do $M_{\langle 2,2,3 \rangle}$ using two $M_{\text{reduced } \langle 2,2,2 \rangle}$

Idea Do $M_{(2,2,3)}$ using two $M_{\text{reduced}} \langle 2,2,2 \rangle$

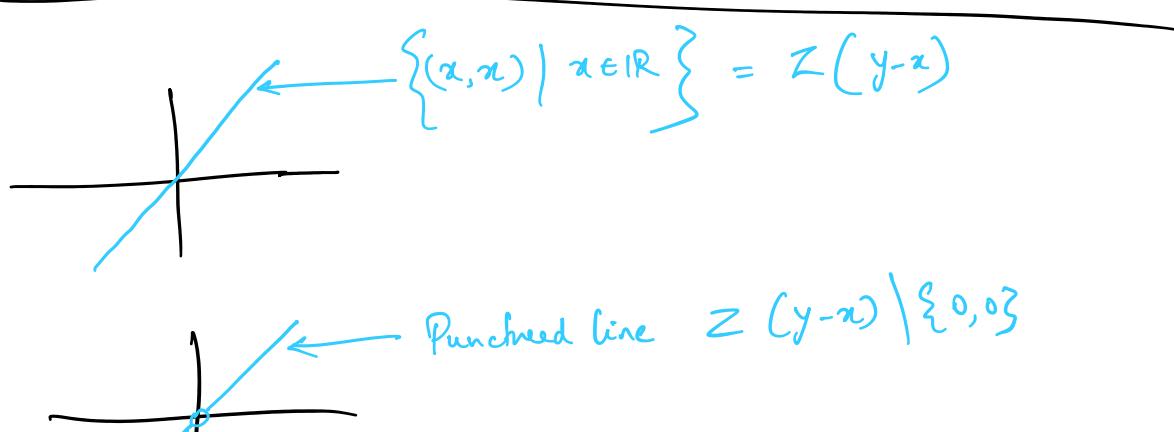
$M_{\text{reduced}} \langle 2,2,2 \rangle$ in 6 mults. trivially. Can you do in 5 mults?

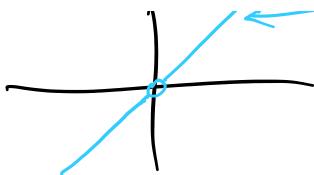
shows $R(M_{(2)}^{\text{red}}) \leq 6$. Bini et al. attempted to find a rank five expression for $M_{(2)}^{\text{red}}$. They searched for such an expression by computer. Their method was to minimize the norm of $M_{(2)}^{\text{red}}$ minus a rank five tensor that varied (see §4.6 for a description of the method), and their computer kept on producing rank five tensors with the norm of the difference getting smaller and smaller, but with larger and larger coefficients. Bini (personal communication) told me about how he lost sleep trying to understand what was wrong with his computer code. This went on for some time, when finally he realized there was nothing wrong.

$$M_{\text{reduced}} \langle 2,2,2 \rangle = \lim_{t \rightarrow 0} \frac{1}{t} \left[\begin{array}{l} (\alpha_{12} + t \alpha_{11}) \otimes (y_{12} + t y_{22}) \otimes z_{21} \\ 1 \\ + (\alpha_{21} + t \alpha_{11}) \otimes y_{11} \otimes (z_{11} + t z_{12}) - (\alpha_{12} \otimes y_{12} \otimes (z_{11} + z_{21} + t z_{22})) \\ 2 \quad 3 \\ - \alpha_{21} \otimes ((y_{11} + y_{12}) + t y_{21}) \otimes z_{11} + (\alpha_{12} + \alpha_{21}) \otimes (y_{12} + t y_{21}) \otimes (z_{11} + t z_{22}) \\ 4 \quad 5 \end{array} \right]$$

$M_{\text{reduced}} \langle 2,2,2 \rangle$ can be arbitrarily well approximated by a rank 5 tensor

- (*) $\underline{R}(M_{\text{reduced}} \langle 2,2,2 \rangle) \leq 5$. We'll use u.b. on border ranks to get an upper bound on w.
- (x) $\underline{R}(M_{(2,2,3)}) \leq 10$.





Defn [Zariski Closure] Zariski closure of the punctured line is
Zero set of the polynomial $y - z$.

If a set is Zariski closed, there is a finite set of polynomials whose common zeros vanish exactly on the set.

Defn [Segre Variety] (Parameterizes Rank 1 tensors)

$$\text{Seg}: \mathbb{P}(A_1) \times \dots \times \mathbb{P}(A_n) \longrightarrow \mathbb{P}(A_1 \otimes \dots \otimes A_n)$$

$$(a_1, \dots, a_n) \longmapsto a_1 \otimes \dots \otimes a_n$$

a_i have $\dim(A_i)$ homogeneous coordinates

$$S^1 = \text{Im}(\text{Seg}) \quad \leftarrow \text{Segre Variety.}$$

Defn [Secant line] Let V be a projective variety. A line L is called a secant line to V if L meets V in two or more points.

Defn Let $X \subseteq \mathbb{P}(V)$ be a projective variety. Define

$$S_n(X)^0 = \left\{ (x_1, \dots, x_n, z) \in X^{x_n} \times \mathbb{P}(V) \mid z \in \text{Span}(x_1, \dots, x_n) \right\}$$

$$\subseteq \text{Seg}(X^{x_n} \times \mathbb{P}(V)) \subseteq \mathbb{P}(V^{\otimes n+1}) \quad \leftarrow$$

Let $S_n(X) = \overline{S_n(X)^0}$. \curvearrowright Zariski closure

$S_n(X)$ is called the n^{th} abstract secant variety of X .

$$\pi^0: S_n(X)^0 \rightarrow \mathbb{P}(V) \quad \text{Likewise } \pi: S_n(X) \rightarrow \mathbb{P}(V)$$

$$(x_1, \dots, x_n, z) \mapsto z$$

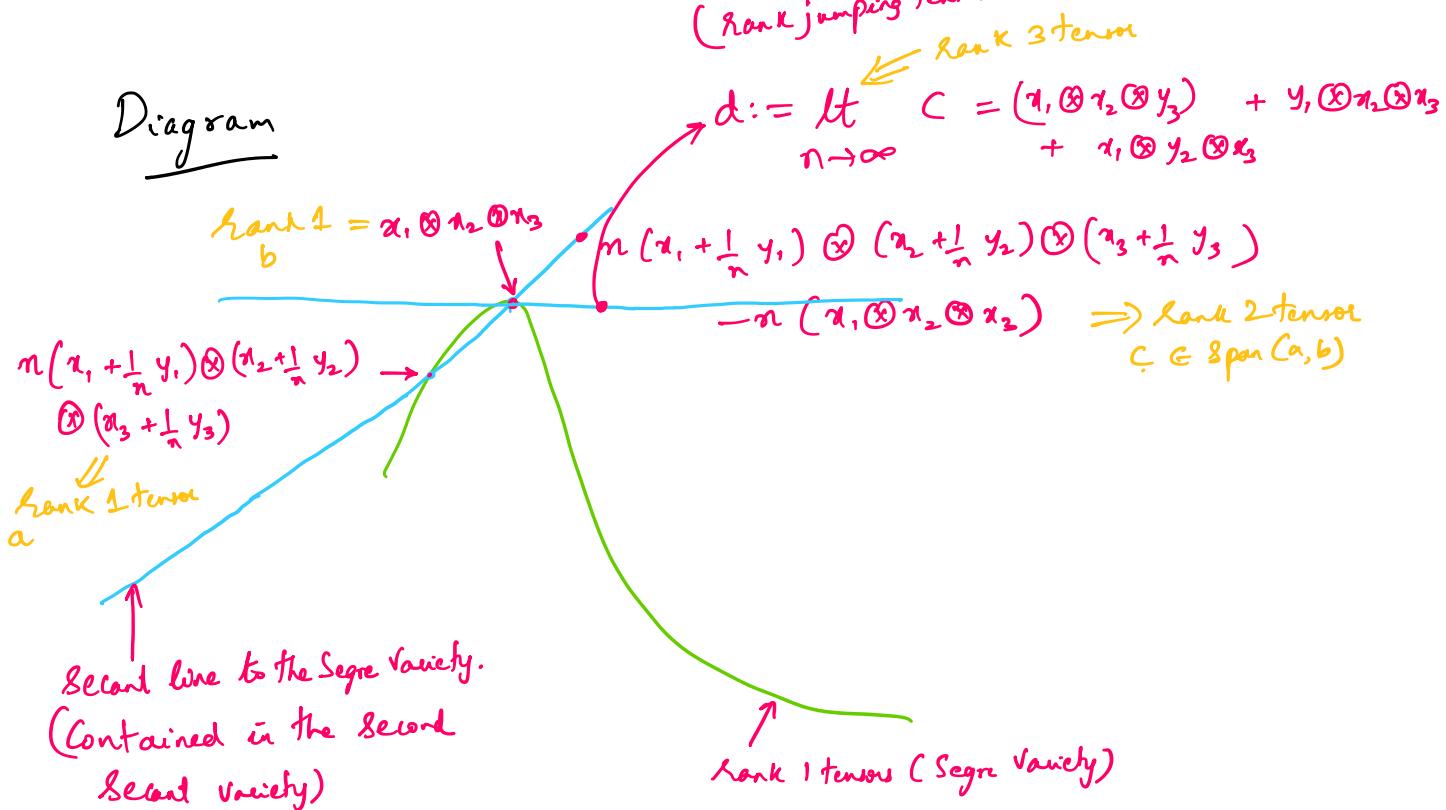
- π and π^0 is denoted $\sigma_n^0(X)$

- Image of Π^* is denoted $\sigma_n^*(x)$
- $\sigma_n(x) = \text{Im}(\Pi)$ (Called n th Selant variety of X)
 $(\sigma_1(x) = x)$

When $X = \text{Seg}(\mathbb{P}(A_1) \times \dots \times \mathbb{P}(A_n))$

$\sigma_n(x) \rightarrow$ Set of tensors of border rank at most n
(Rank jumping tensor)

Diagram



⑦ $M_{\text{reduced } \langle 2,2,2 \rangle}$ is a rank jumping tensor

Idea:- Interpolate between rank and border rank

Let ε be an indeterminate

Defn $[h\text{-rank}, \text{border rank}]$ $t \in \mathbb{F}^k \otimes \mathbb{F}^m \otimes \mathbb{F}^n$

$$\textcircled{1} R_h(t) = \min \left\{ h \mid \exists u_i \in \mathbb{F}[\varepsilon]^k, v_i \in \mathbb{F}[\varepsilon]^m, w_i \in \mathbb{F}[\varepsilon]^n : \sum_{i=1}^h u_i \otimes v_i \otimes w_i = \varepsilon^h t + O(\varepsilon^{h+1}) \right\}$$

$\uparrow \quad \downarrow \quad \uparrow \quad \downarrow \quad \uparrow \quad \downarrow \quad \uparrow \quad \downarrow \quad \uparrow \quad \downarrow$

$$1 \quad 1 \cdots \quad 1 \quad \int \sum_{i=1}^h u_i \cdot (v_i \otimes w_i)$$

$$t = \lim_{\epsilon \rightarrow 0} \frac{1}{\epsilon^h} \left(\sum_{i=1}^n u_i \otimes v_i \otimes w_i \right)$$

$$(2) \underline{R}(t) = \min_{h \in \mathbb{N}} R_h(t)$$

Lemma (1) $R_0(t) = \underline{R}(t)$

(2) $R_0(t) \geq R_1(t) \geq \dots = \underline{R}(t)$ \leftarrow

Proof By defn \square

Lemma (1) $\pi \in S_3 : R_n(\pi(t)) = R_n(t)$

(2) $R_{\max\{h,h'\}}(t \oplus t') \leq R_h(t) + R_{h'}(t')$

(3) $R_{h+h'}(t \oplus t') \leq R_h(t) R_{h'}(t')$

Proof Same as earlier lemmas .. use language of R_n \square

Lemma [turn approx. computations into real ones] There is a constant $C_h \leq \binom{h+2}{2}$

s.t $\# t$

$$\underline{R}(t) \leq C_h R_h(t)$$

Proof Counting \square

Thm $\underline{R}(M_{\langle k,m,n \rangle}) \leq \omega \Rightarrow \omega \leq 3 \log_{kmn} 1$

Proof Use R_h ... and same logic as previous thm that gave you u.b. on ω using an u.b. on $R(M_{\langle \# \rangle})$... make h small .. \square

(*) $\underline{R}(M_{\text{reduced } \langle 2,2,2 \rangle}) \leq 5 \Rightarrow \underline{R}(M_{\langle 2,2,3 \rangle}) \leq 10$
 $\Rightarrow \omega \leq 3 \log_{12} 10 \approx 2.78$

Schonhage's γ theorem

- \otimes Strassen & Pan turned u.b. of $R(M_{\langle *, * \rangle})$ into u.b. on w
- \otimes Bini et al. turned u.b. on $\underline{R}(M_{\langle *, * \rangle})$ into u.b. on w
- \otimes Schonhage turns u.b. on \underline{R} (not a mat mult.-tensor but it is a direct sum of matrix mult. tensor) into u.b. on w

Lemma ① $\underline{R}\left(M_{\langle k, 1, n \rangle} \oplus M_{\langle 1, m, 1 \rangle}\right) = kn + m$

② $\underline{R}(M_{\langle k, 1, n \rangle}) = kn$ & $\underline{R}(M_{\langle 1, m, 1 \rangle}) = m$

③ $\underline{R}\left(M_{\langle k, 1, n \rangle} \oplus M_{\langle 1, N, 1 \rangle}\right) \leq kn + 1$ when
 $N = (m-1)(k-1)$

Defn $I_{\langle n \rangle} \in \mathbb{F}^n \otimes \mathbb{F}^n \otimes \mathbb{F}^n$ is s.t

$$(I_{\langle n \rangle})_{i,i,i} = 1 \quad \forall i \in [n] \geq 0 \text{ elsewhere.}$$

$R(I_{\langle n \rangle}) = n$

Lemma $R(t) \leq r \iff t \leq I_{\langle r \rangle}$

Proof Skipped \boxtimes

Thm [γ -theorem] If $\underline{R}\left(\bigoplus_{i=1}^P M_{\langle k_i, m_i, n_i \rangle}\right) \leq r$, and $n > P$, then

$w \leq 3\gamma$, where γ satisfies

$$\sum_{i=1}^P (k_i m_i n_i)^\gamma = r \leftarrow$$

Proof Skipped \boxtimes

$\otimes \underline{R}(M_{\langle 4, 1, 3 \rangle} \oplus M_{\langle 1, 6, 1 \rangle}) \leq 13$

Using the theorem $\omega \leq 2.55$

Coppersmith-Winograd

Defn [easy C.W. tensor]

$$T_{q,CW} = \sum_{j=1}^q a_0 \otimes b_j \otimes c_j + a_j \otimes b_0 \otimes c_j + a_j \otimes b_j \otimes c_0$$

↑

$$\mathbb{C}^{q+1} \otimes \mathbb{C}^{q+1} \otimes \mathbb{C}^{q+1}$$

↑
expression has $3q$ terms. Turns out

$$\underline{R}(T_{q,CW}) = q+2.$$

dense sets of integers with no
3-term arithmetic progressions

$$\text{Then } \omega \leq \underbrace{\log\left(\frac{4}{27} \underline{R}(T_{q,CW})^3\right)}_{\log q}.$$

$$\text{Set } q = 8, \Rightarrow \omega \leq 2.41$$

Defn [big C.W. tensor]

$$T_{q,CW}^+ = T_{q,CW} + a_0 \otimes b_0 \otimes c_{q+1} + a_0 \otimes b_{q+1} \otimes c_0 + a_{q+1} \otimes b_0 \otimes c_0$$

$$\mathbb{C}^{q+2} \otimes \mathbb{C}^{q+2} \otimes \mathbb{C}^{q+2}$$

$$\text{Then } \underline{R}(T_{q,CW}^+) = q+2, \text{ and}$$

$$\omega \leq \underbrace{\log\left(\frac{4}{27} \underline{R}(T_{q,CW}^+)^{3/k}\right)}_{\log q}$$

C.W. analyzed $T_{q,CW}^+ \otimes T_{q,CW}^+$

Latest work analyzes $(T_{q,CW}^+)^{\otimes 32} \Rightarrow \omega < 2.3728639$

Latest work analyses $(T_{q,CW}^+)^{\wedge}$ $\Rightarrow w < 2.3$ \uparrow

People have shown that taking higher & higher powers of $T_{q,CW}^+$
won't improve beyond 2.3

Conjecture [Asymptotic Rank Conjecture]

$$\lim_{n \rightarrow \infty} R(T_{2,CW}^{\otimes n})^{1/n} = 3 \quad (\text{Theorem} \Rightarrow w=2)$$

Conjecture [No 3 disjoint equivoluminous subsets Conjecture] Let H be an abelian group.

Let $m_1, \dots, m_n \in H$. This satisfies N3DES property if

$\# S, T, U \subseteq [n], S, T, U$ disjoint

$$\sum_{i \in S} m_i \neq \sum_{i \in U} m_i \neq \sum_{i \in T} m_i$$

The actual conjecture is that there is an H of size

$$|H| \leq 2^{o(n)} \quad [\text{Theorem Conjecture} \Rightarrow \text{asymptotic rank conjecture}]$$

e.g. $H = \mathbb{Z}_2^n$ $m_i = e_i$ satisfies
N3DES property, but $|H| = 2^n$

* truth of the above conjecture \Rightarrow one of the sunflower conjectures is false..