Gröbner Bases Native to 'pseudo'-Hodge Algebras, with Application to the Algebra of Minors

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Joint work with Joshua Grochow (Univ. of Colorado, Boulder, USA)

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- Gröbner bases are well-suited to both of the above!
- ► Gröbner bases give theoretical insight as well as are the key tool in effective methods

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Gröbner bases tend to obscure symmetry!

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Question

Develop a Gröbner basis theory which takes advantage if variety corresponding to ideal has large symmetry group, or is 'determinantal'

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- ► If A is an ASL, the product of two standard monomials can be straightened into a linear combination of 'smaller' standard monomials
- ► ASLs arise as coordinate rings of algebraic varieties, e.g. Grassmanians, determinantal varieties, flag varieties, Schubert varieties

Bideterminants (products of minors)

► Example of Hodge algebra - algebra of bideterminants

e.g.
$$A = \mathbb{F}[X_{1,1}, X_{1,2}, X_{2,1}, X_{2,2}, Y] / \langle X_{1,2} X_{2,1} - X_{1,1} X_{2,2} + Y \rangle$$

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 - \triangleright poly ring with one variable for each minor of $n \times m$ matrix
 - quotient by relations between minors
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 - ► standard monomials correspond to standard bitableaux
- ► Advantage smaller expressions for 'determinant-like' polynomials; bideterminants are reflect symmetries coming from the action (representation theory) of GL_n

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Question

Can we build a theory of Gröbner bases 'native' to p-ASLs, i.e. Gröbner theory without referencing the ideal J?

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- Product of standard monomials not necessarily standard, might require straightening
 - ► How do you define term order?
 - How would you define division of monomials?
 - ► What plays the role of monomial ideals?

Term Order & Division

- ightharpoonup A p-ASL term order on a p-ASL A is a total order \prec on standard monomials in A such that
 - ▶ 1 ≤ m
 - ▶ If $a \prec b$ and $c \leq d$, and ac, $bd \neq 0$, then

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- ► When does standard monomial m divide m':
 - ordinary division in the polynomial ring, or
 - ▶ m divides m' if there exists standard monomial f such that

$$LM(mf) = m'$$

Auxilliary Algebra of Leading Terms

► Given p-ASL A, algebra of leading terms w.r.t. A is another p-ASL A_{1t} on the same variables, and the same standard monomials such that for standard monomials m, m' no straightening

$$\pi_{lt}(\mathfrak{m})\cdot\pi_{lt}(\mathfrak{m}')=$$
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Proposition

Every p-ASL A admits two algebras of leading terms – A_{gen} where the product is never 0, and, A_{disc} where product is 0 unless mm' is also a standard monomial.

Definition of p-ASL Gröbner Basis

- ▶ Given p-ASL A, algebra of leading terms A_{lt} , and an ideal $I \subseteq A$, then $G \subseteq A$ is a p-ASL Gröbner basis if:
 - ▶ For all $f \in I$, there exists $g \in G$ such that $\pi_{lt}(LM(g))$ divides $\pi_{lt}(LM(f))$, or

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Our Main Result

Theorem (Grochow-N, 2025)

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Reduced	1	✓
Universal	1	√
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Corollary (Grochow-N, 2025)

The algebra of bideterminants has a p-ASL term order, thus we have a Gröbner basis theory (called bd-Gröbner bases).

Applications to Bideterminant Algebra

► Universal p-ASL Gröbner basis is a p-ASL Gröbner basis for any p-ASL term order and any algebra of leading terms

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Takeaway

- 1. Given all our machinery, the proof is one-line
- 2. In the ordinary case, universal Gröbner basis are known only for maximal minors and minors of size 2

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- ► Compute Weyl closure, b-functions, etc. using bd-Gröbner bases in the Weyl algebra
- ➤ See if we can develop a bipermanent Gröbner basis theory (codimension of singular locus of permanent hypersurface is unknown!)

References

- E. W. Mayr and A. R. Meyer. The complexity of the word problems for commutative semigroups and polynomial ideals. Advances in mathematics, 46(3):305-329, 1982.
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