

SYSTEMS ENGINEERING

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SE Agenda

- Motivation / Scalability
- Consensus
- Paxos
- Raft
- CAPTheorem: DynamoDB vs Spanner
- Chubby and K8s
- Memcached / REDIS
- Object Store: S3 and Ceph
- Virtual Time / Vector Time / Real Time
- DynamoDB (eventual consistency)
- Spanner

Administrative Stuff

- Slide decks, calendar, etc on Opal
- Exams:
 - written or oral depending on program
- Exercises:
 - on Fridays



A SCALABLE ARCHITECTURE

SEARCH QUERIES

- ➤ Search queries in 2011:
 - ➤ Google: about 34,000 queries per second (1-3 billions per day)
 - ➤ Yahoo: about 3,200 queries per second
 - ➤ Bing: about 930 queries per second
- ➤ Search queries in Sept 2019:
 - ➤ Google: 74% (desktop), 92% (mobile)
 - ➤ Bing: 11.4% (desktop), 1% (mobile)
 - ➤ Yahoo: 2.2% (desktop), 1.3% (mobile)
- ➤ 2023 (all platforms):
 - ➤ Google (92%), Bing (3%), Yahoo (1,2%)

SERVING A SEARCH QUERY

(Google)

PROBLEMS

- ➤ Need to use energy-efficient, low-cost CPUs
 - ➤ low peak performance

- ➤ Need to guarantee fast response times
 - ➤ to keep users happy

- ➤ Need to parallelize queries
 - ➤ since there are tens of billions instructions per query to execute

HOW PARALLEL?

➤ Rough estimate:

➤ **34,000** queries per second and say, > 8 seconds per query (sequential execution)

➤ Hence,

➤ one would need to involve > **34,000** computers (with 8 cores each)

➤ Notes:

- ➤ a very rough estimate since, e.g., caching of results might reduce the number of computers, etc
- ➤ might be more limited by memory than CPU speed
- ➤ is this the average or the peak queries per second?

www.google.com/search?
 q=systems+engineering

2. Browser resolves www.google.com

www.google.com is an alias for www.l.google.com.
www.l.google.com has address 64.233.183.104
www.l.google.com has address 64.233.183.147
www.l.google.com has address 64.233.183.99
www.l.google.com has address 64.233.183.103

URI - Uniform Resource Identifier

www.google.com/search? q=systems+engineering

Browser resolves www.google.com

www.google.com is an alias for www.l.google.com. www.l.google.com has address 64.233.183.104 www.l.google.com has address 64.233.183.147 www.l.google.com has address 64.233.183.99 www.l.google.com has address 64.233.183.103 (located in the US)

URI - Uniform Resource Identifier

www.google.com/search?
 q=systems+engineering

Query parameters are key-value pairs (e.g., k=v) added to the end of a URI, typically after a question mark

2. Browser resolves www.google.com

```
www.google.com is an alias for www.l.google.com.
www.l.google.com has address 64.233.183.104
www.l.google.com has address 64.233.183.147
www.l.google.com has address 64.233.183.99
www.l.google.com has address 64.233.183.103
```

www.google.com/search?
 q=syst/ms+engineering

Need to translate www.google.com to an IP address - using DNS

2. Browser resolves www.google.com

www.google.com is an alias for www.l.google.com.
www.l.google.com has address 64.233.183.104
www.l.google.com has address 64.233.183.147
www.l.google.com has address 64.233.183.99
www.l.google.com has address 64.233.183.103

www.google.com/search?
 q=systems+engineering

www.google.com is a CNAME that maps to www.l.google.com

2. Browser resolves www.google.com

www.google.com is an alias for www.l.google.com.
www.l.google.com has address 64.233.183.104
www.l.google.com has address 64.233.183.147
www.l.google.com has address 64.233.183.99
www.l.google.com has address 64.233.183.103

www.google.com/search?
 q=systems+engineering

www.l.google.com maps to multiple IP addresses

Browser resolves www.google.com

www.google.com is an alias for www.l.google.com.

www.l.google.com has address 64.233.183.104

www.l.google.com has address 64.233.183.147

www.l.google.com has address 64.233.183.99

www.l.google.com has address 64.233.183.103

ROUND-ROBIN DNS

➤ DNS rotate results:

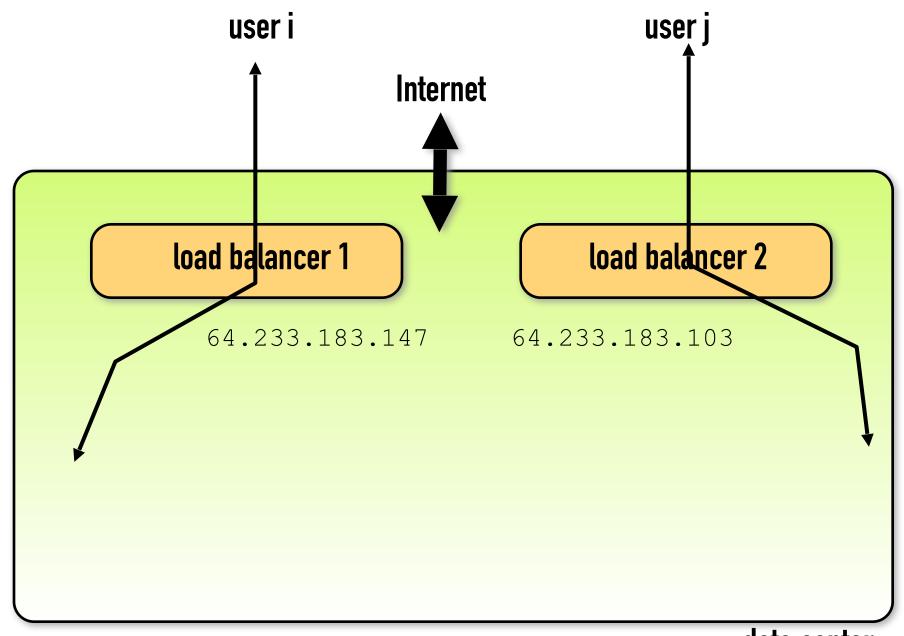
```
www.l.google.com has address 64.233.183.147

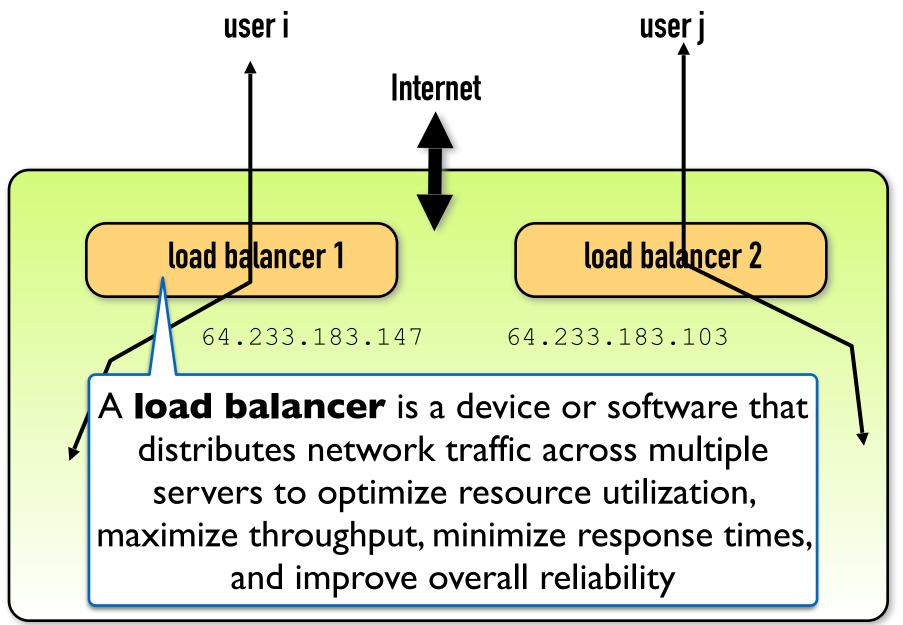
www.l.google.com has address 64.233.183.99

www.l.google.com has address 64.233.183.103

www.l.google.com has address 64.233.183.104
```

IP addresses rotate in round-robin fashion





IN 2009, 2010:

Browser resolves www.google.com

```
www.google.com is an alias for www.l.google.com
```

```
www.l.google.com has address 74.125.39.103
www.l.google.com has address 74.125.39.105
www.l.google.com has address 74.125.39.106
www.l.google.com has address 74.125.39.147
www.l.google.com has address 74.125.39.104
www.l.google.com has address 74.125.39.99
```

(located in Berlin)

5 MINUTES LATER:

Browser resolves www.google.com

```
www.google.com is an alias for www.l.google.com
www.l.google.com has address 209.85.129.103

www.l.google.com has address 209.85.129.147

www.l.google.com has address 209.85.129.147

www.l.google.com has address 209.85.129.99
```

(located in Mountain View, CA)

DNS

- ➤ Google's DNS servers used to map <u>www.google.com</u> on IP addresses based on
 - geographical location of user
 - -> minimize round-trip times
 - ➤ load of the individual Google clusters
 - -> coarse grain load balancing

2013/4

```
www.google.com has address 173.194.69.106
www.google.com has address 173.194.69.99
www.google.com has address 173.194.69.147
www.google.com has address 173.194.69.104
www.google.com has address 173.194.69.105
www.google.com has address 173.194.69.103
www.google.com has IPv6 address
2a00:1450:4008:c01::63
```

(located in Mountain View, CA?)

2013/4

```
www.google.com has address 173.194.69.106
www.google.com has address 173.194.69.99
www.google.com has address 173.194.69.147
www.google.com has address 173.194.69.104
www.google.com has address 173.194.69.105
www.google.com has address 173.194.69.103
www.google.com has IPv6 address
2a00:1450:4008:c01::63
```

DNS also returns an **IPv6** address

(located in Mountain View, CA?)

2019–2024: SINGLE ADDRESS

www.google.com has address 216.58.207.68

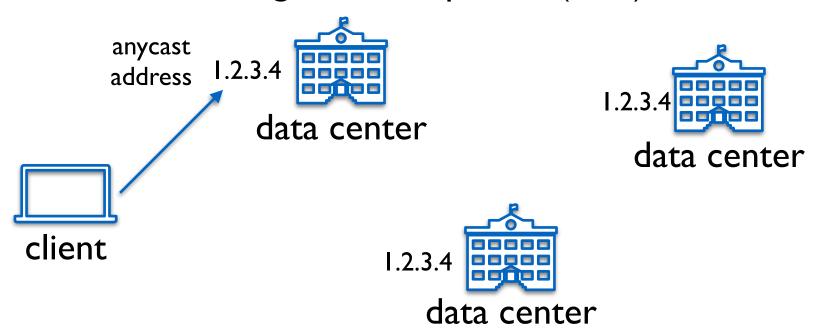
www.google.com has IPv6 address

2a00:1450:400e:800::2004

IPv6 address the same but IPv4 has changed in 2024

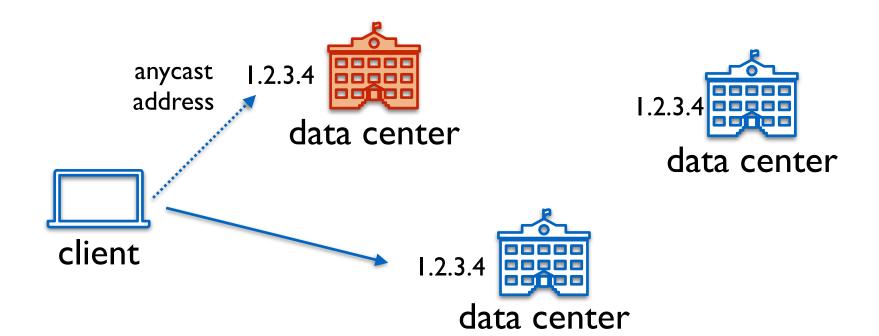
Anycast

- Ensure that load is routed to closest data center (a ,, I to I routing scheme")
 - minimizing Round Trip Time (RTT)



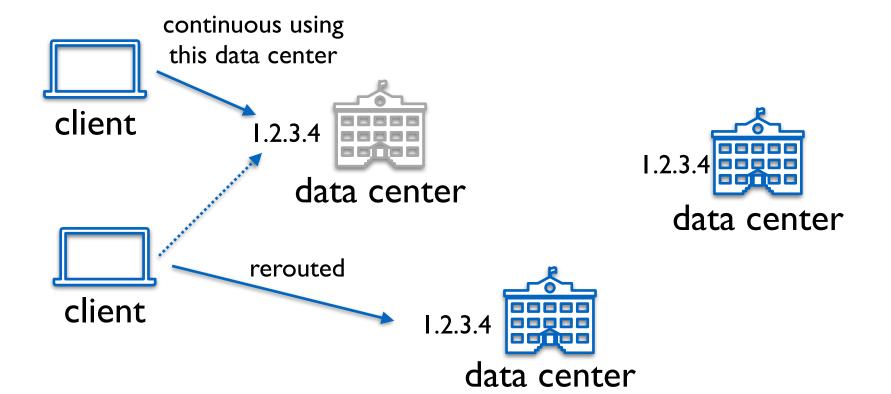
Rerouting

 In case closest data center is overloaded, traffic is rerouted to another datacenter



Load Balancing

 One might only reroute parts of the traffic to keep data centers from being overloaded



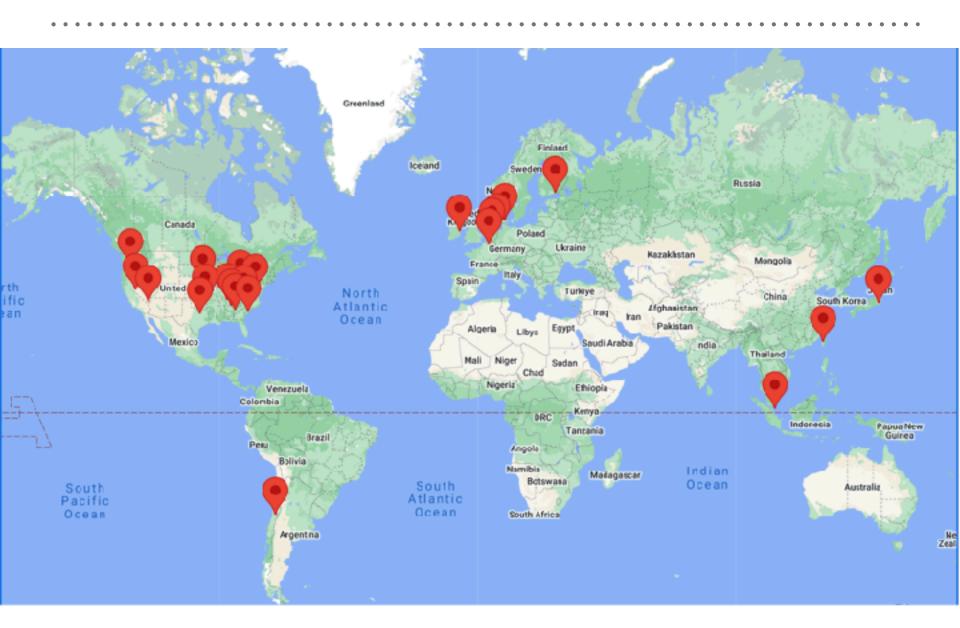
Technical Challenges

 Data to compute answer must be accessible in all data centers

Google search:

- need to replicate the index for searching
- personalized search one needs to keep personalized data accessible in all data centers?
- Mail / Documents / ...:
 - need to keep data available in all data centers?

GOOGLE DATA CENTERS - GEOGRAPHICALLY DISTRIBUTED



GOOGLE'S DATACENTERS

- ➤ Google has geographically distributed
 - data centers consisting of thousands of servers each
 - ➤ 2009 guess: about 1million servers total now ??



(Google data center, Oregon US, (C) New York Times)

GOOGLE'S DATACENTERS

Look at electricity consumption?

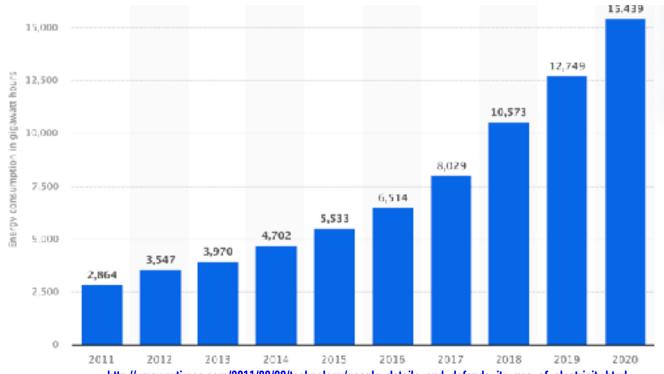
- ➤ Google has geographically distributed
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(Google data center, Oregon US, (C) New York Times)

POWER CONSUMPTION!

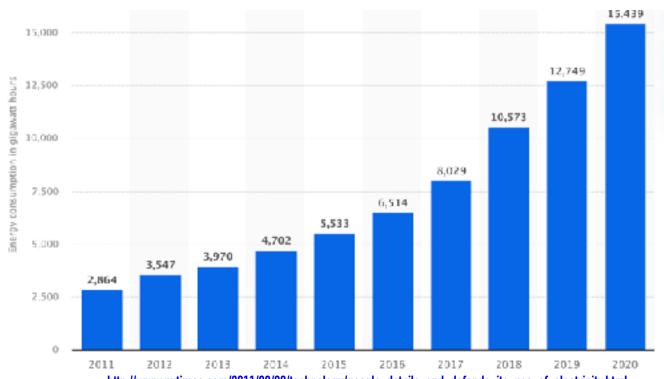
- > Server computer:
 - ➤ about 200-400Watts
- ➤ Electricity consumption of Google in 2020



POWER CONSUMPTION!

2023: Google and Microsoft each consumed more than 24 TWh

- ➤ Server computer:
 - about 200-400Watts
- ➤ Electricity consumption of Google in 2020

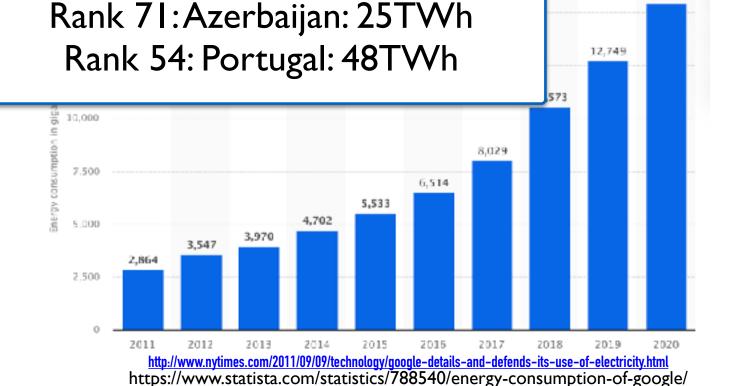


POWER CONSUMPTION!

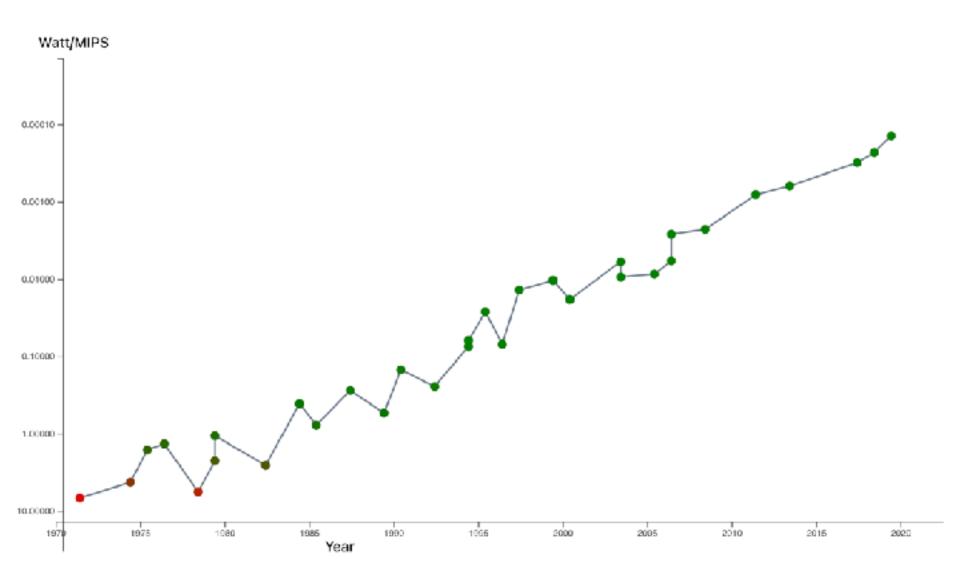
2023: Google and Microsoft each consumed more than 24 TWh

15.439

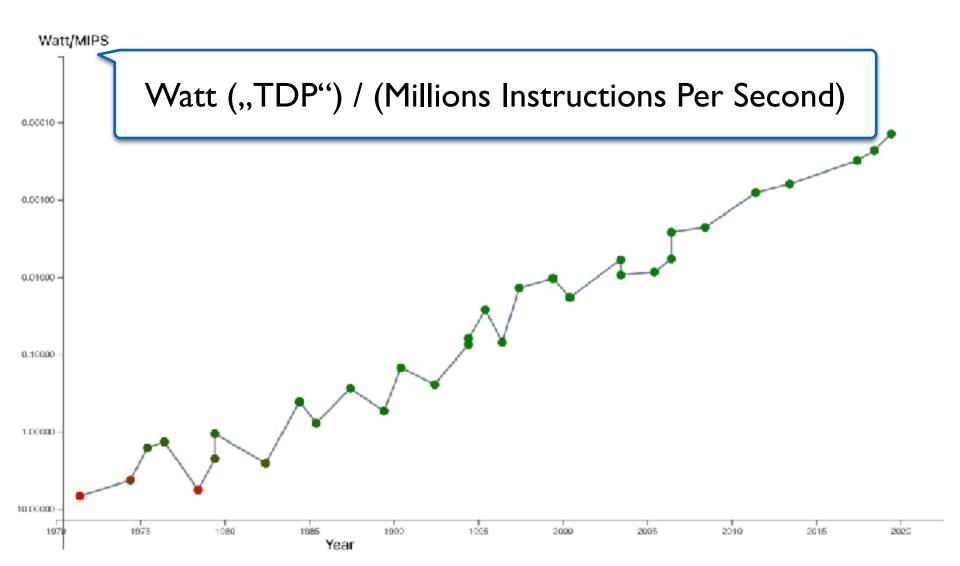
- Server computer:
 - ➤ about 200-400Watts
- ➤ Electricity consumption of Google in 207



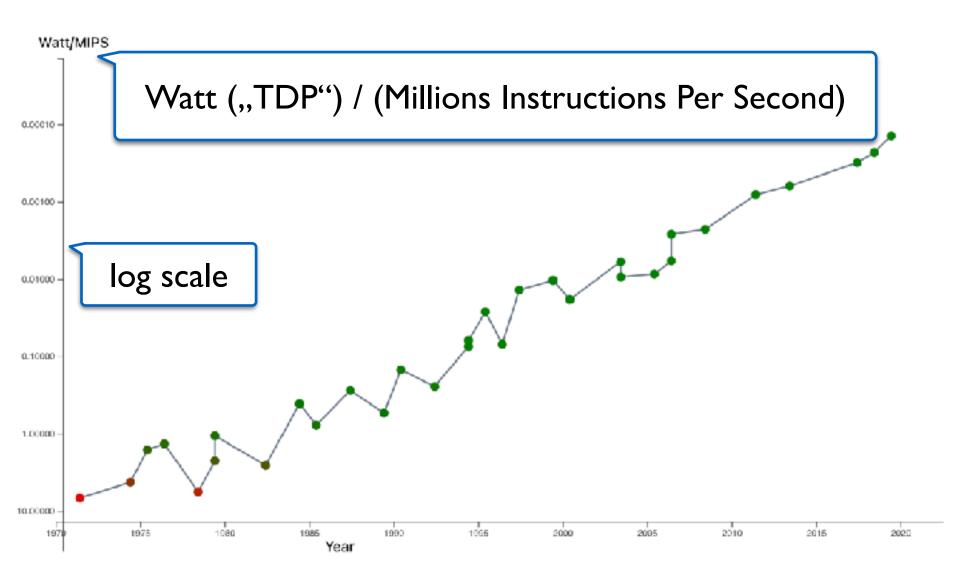
Energy Efficiency Increases



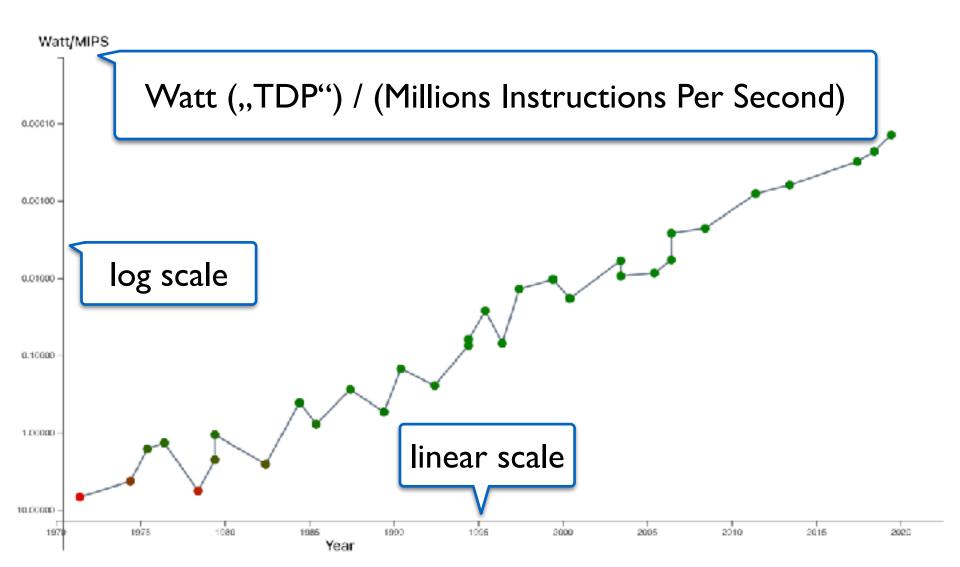
Energy Efficiency Increases



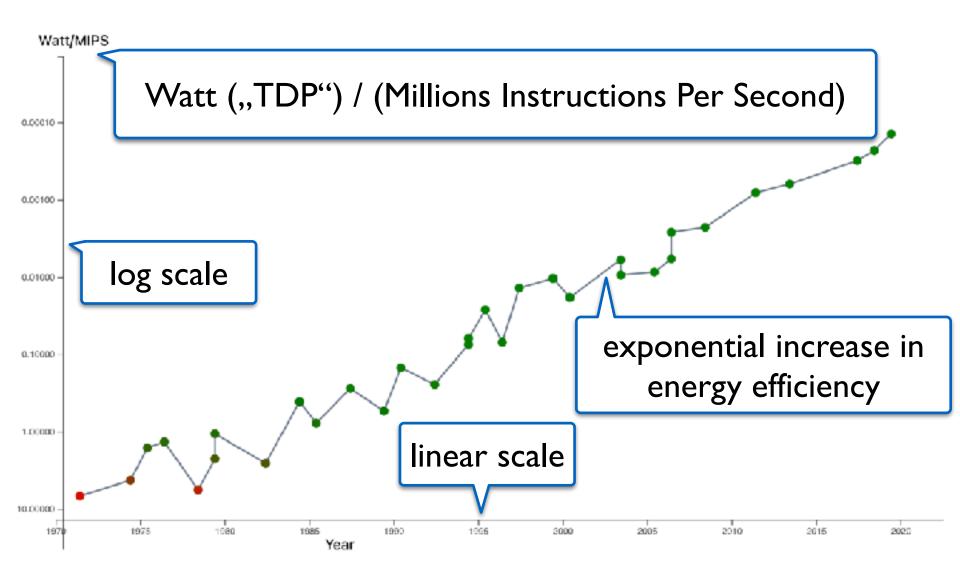
Energy Efficiency Increases



Energy Efficiency Increases



Energy Efficiency Increases



Note: Exponential Increase in Computing Power

- over last decades -

Note: Exponential Increase in Computing Power

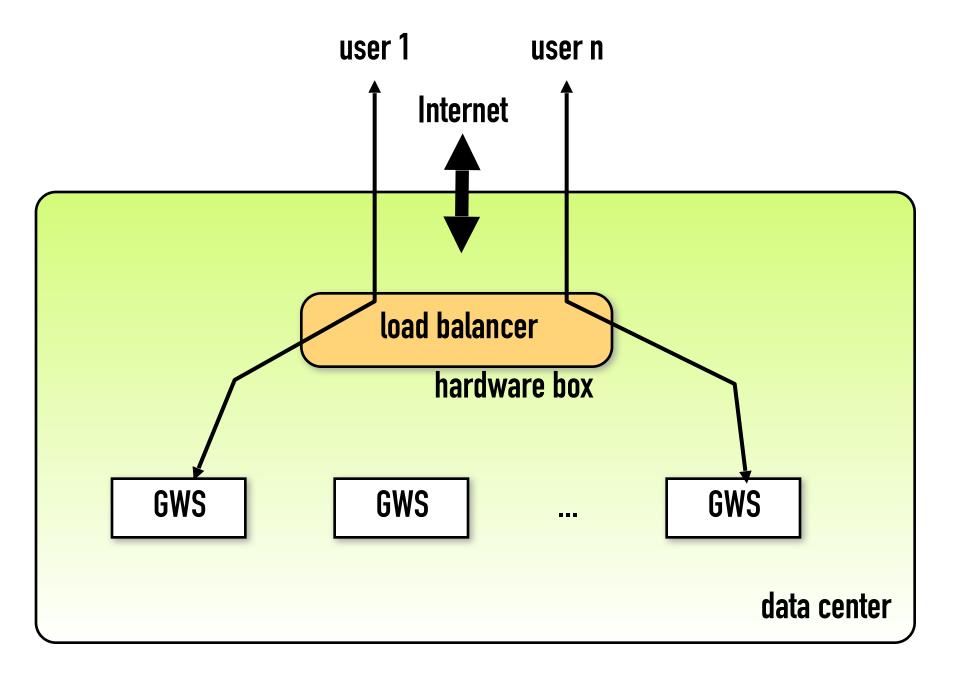
- over last decades -

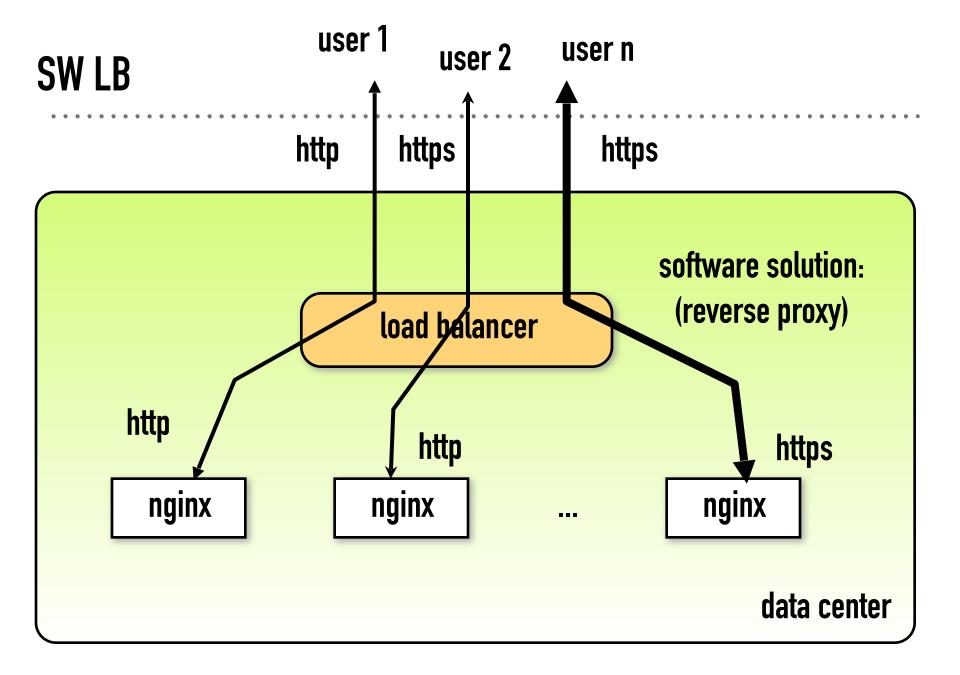
Al workloads use lots of computing power...

QUERY (3)

Browser (Chrome, Firefox, ...) sends http request to IP address

- ➤ Hardware load balancer distributes requests amongst Google Web Servers (GWS)
- ➤ GWS coordinates execution of a request





SEQUENTIAL QUERY PROCESSING

- ➤ GWS queries index server
- ➤ Index server contains:
 - inverted index: (word , list(URL,score))
- ➤ Index server computes hit list for
 - ➤ each query word, and
 - > computes intersection of individual hit lists
- ➤ Computes score of the documents
 - ➤ in the intersection

INDEX SERVER: HIT LISTS

systems

URLx, 1000

URLy, 800

...

URLz, 600

engineering

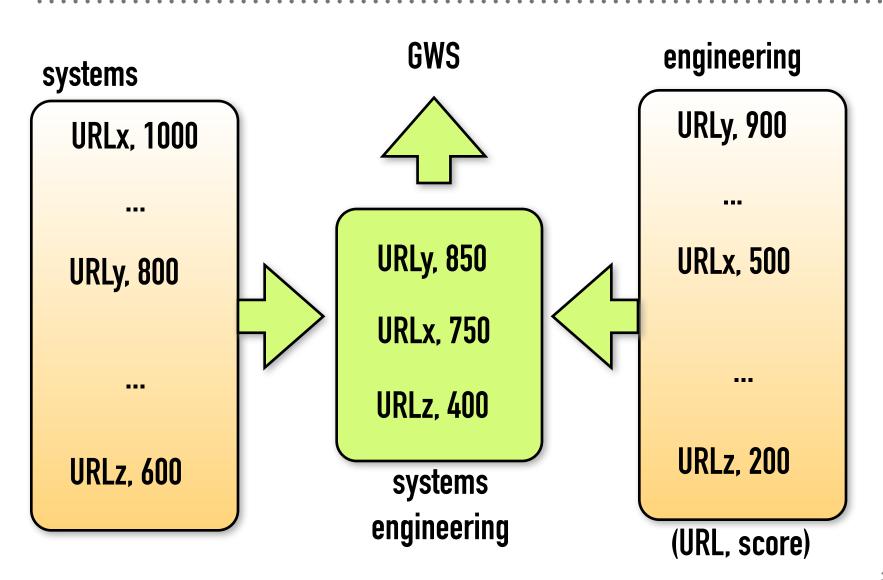
URLy, 900

URLx, 500

URLz, 200

(URL, score)

COMBINE HIT LISTS



PROBLEMS?

- ➤ Billions of web pages!
 - ➤ index does not fit into single machine

- ➤ Very high # of requests per second
 - single index server cannot serve all requests
 - > even if index would fit on single machine

PARTITION!

- ➤ Index is randomly partitioned into
 - ➤ Index shards
 - each shard contains index for a disjoint subset of URLs
 - ➤ e.g., use hash function to partition URLs
- > Pool of servers serves each shard
 - requests are broadcast to all shards
 - ➤ load balancer assigns request to one or more servers in a shard

URLa, URLb, URLc, URLd, URLe, URLf, URLg, URLg,

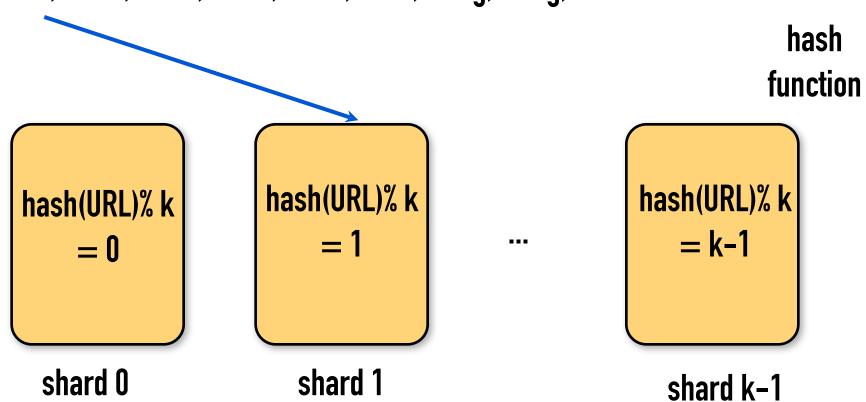
hash function

hash(URL)% k
= 0
shard 0

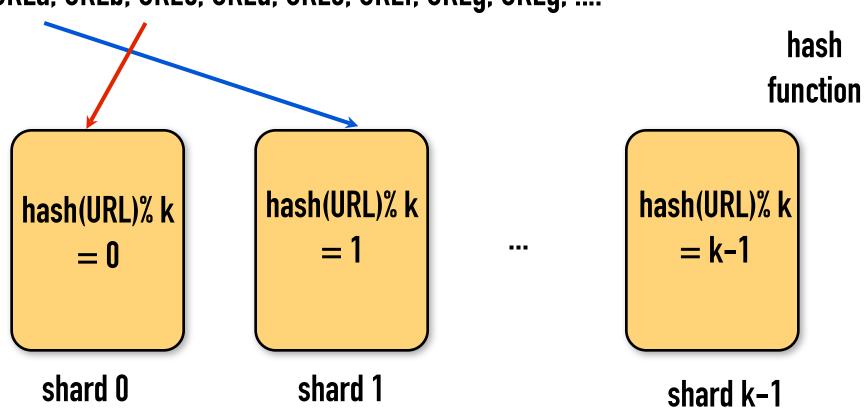
hash(URL)% k = 1 shard 1 hash(URL)% k = k-1

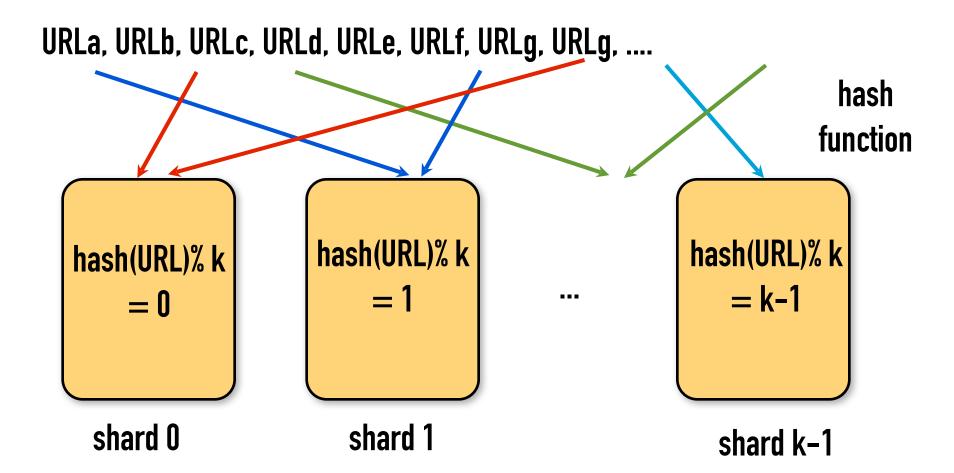
shard k-1

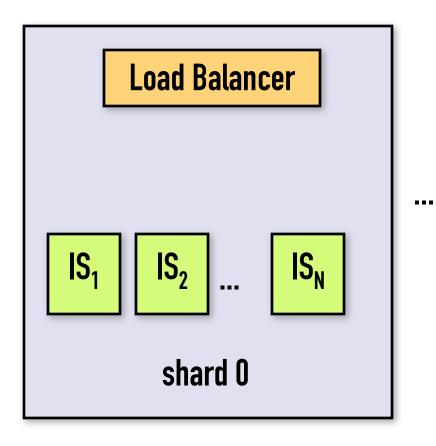
URLa, URLb, URLc, URLd, URLe, URLf, URLg, URLg,

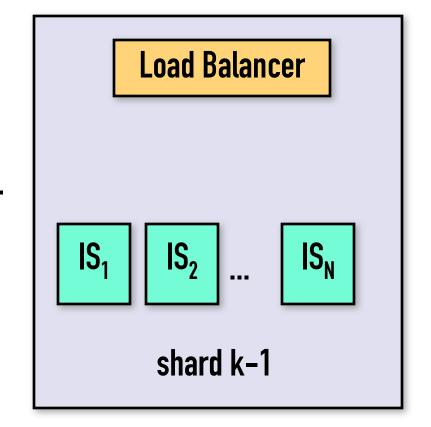


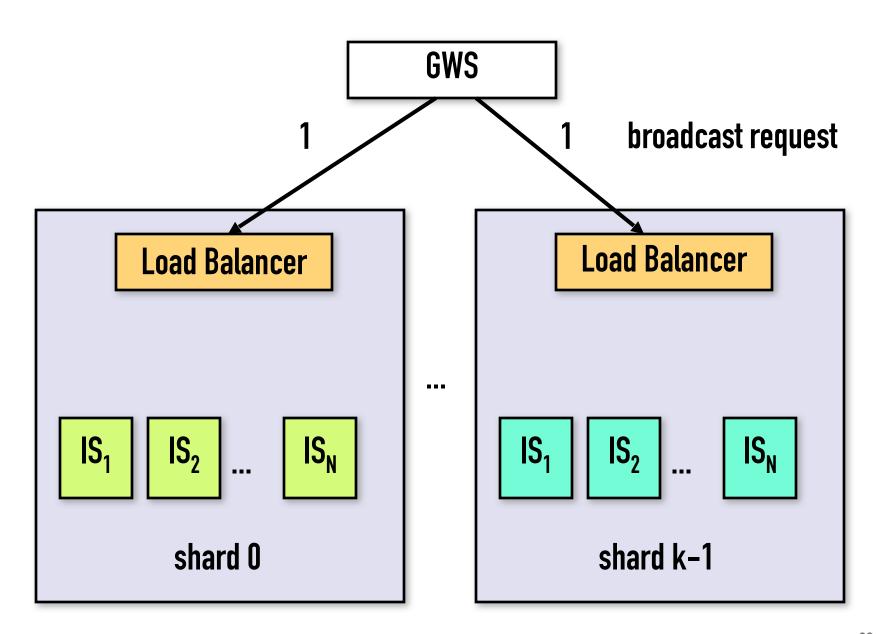
URLa, URLb, URLc, URLd, URLe, URLf, URLg, URLg,

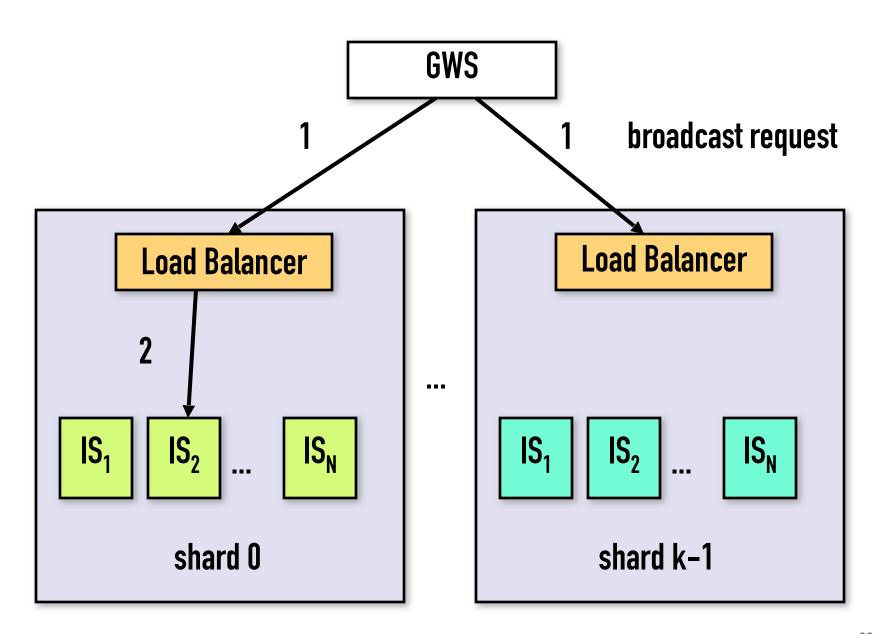


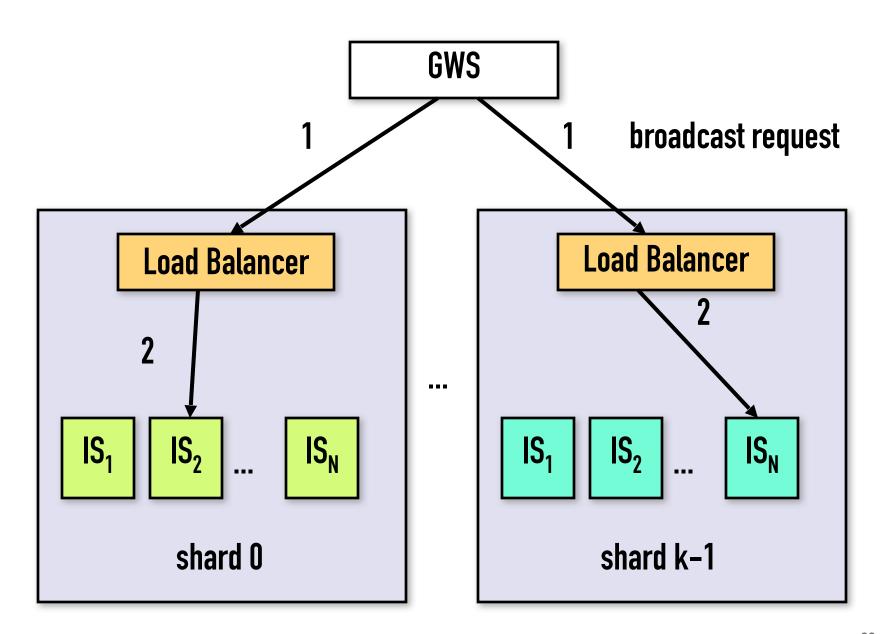


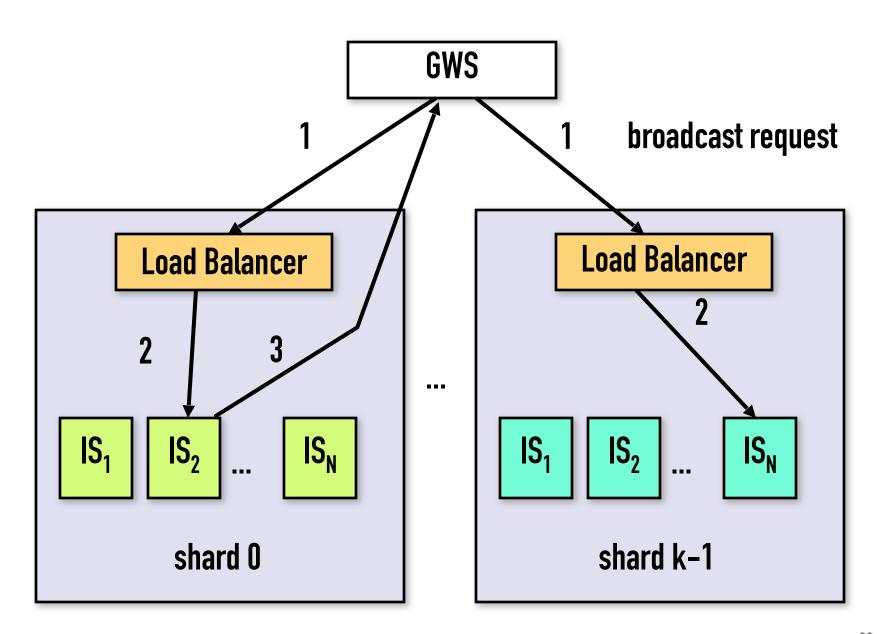


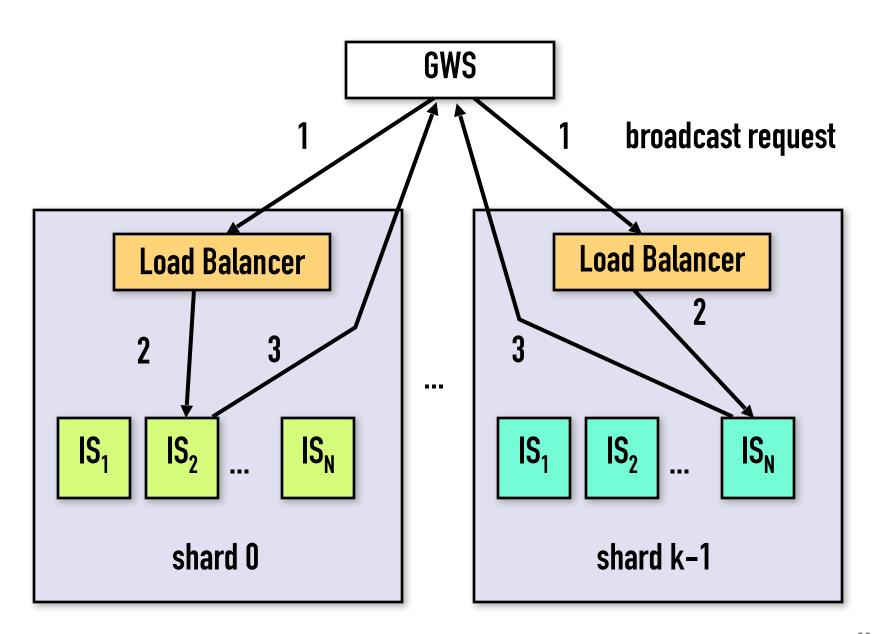




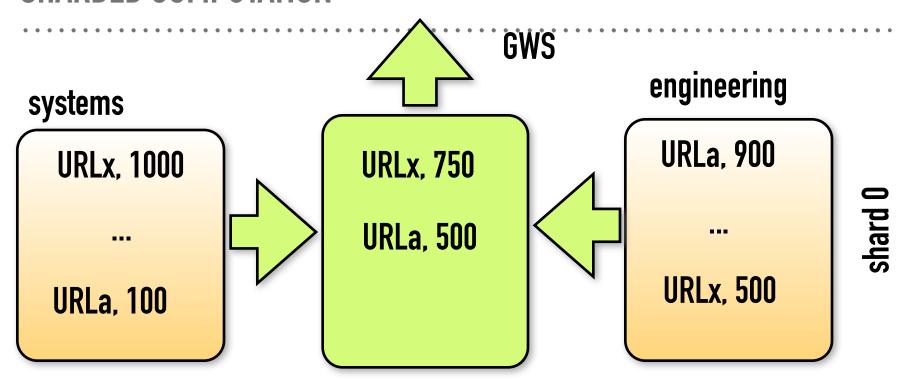




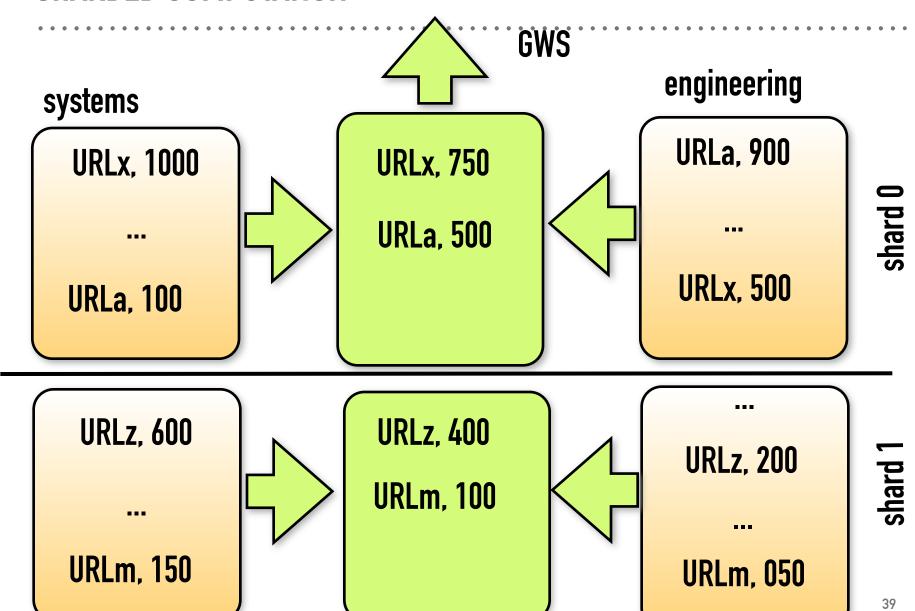




SHARDED COMPUTATION



SHARDED COMPUTATION



SCALABILITY

- Increasing number of requests / second
 - ➤ add more GWS & more replicas per shard
- "Imperfect" hash function:
 - ➤ use different number of replicas per shard
- ➤ Increasing size of index
 - ➤ add more shards
- ➤ Reaching limit of a data center
 - ➤ add new data center using DNS / network load balancer

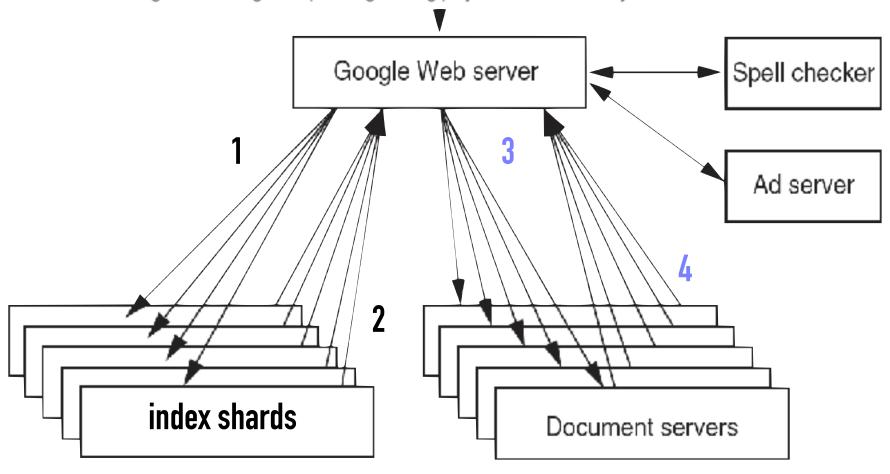
2ND PHASE: RESULT PROCESSING

- ➤ Phase 1:
 - ordered list of document ids from index shards
 - merge sort of the lists

- ➤ Phase 2: Result processing
 - retrieve all documents in list
 - compute title, url, text snippet

(title) Systems engineering - Wikipedia, the free encyclopedia en.wikipedia.org/wiki/Systems_engineering -

(snippet) Systems engineering is an interdisciplinary field of engineering that focuses on how to design and manage complex engineering projects over their life cycles.



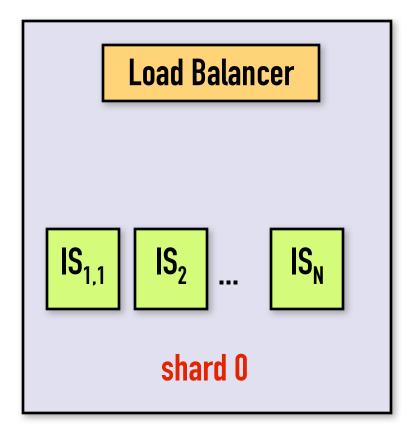
PARALLELIZE DOCUMENT RETRIEVAL

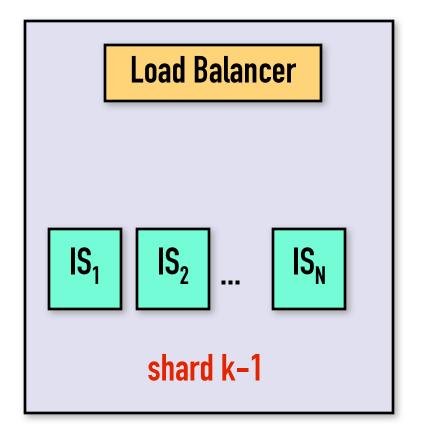
- ➤ Parallelize by:
 - randomly assign documents to shards
 - each shard is served by a pool of servers
 - request sent to a document server in a shard via a load balancer

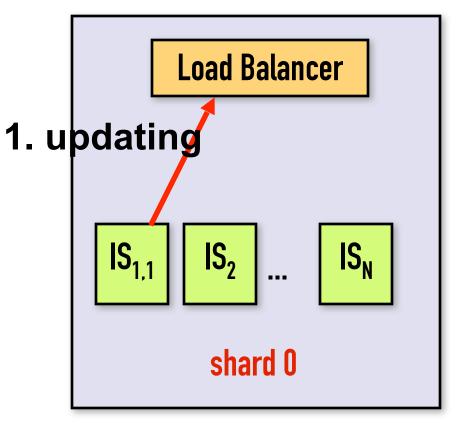
REPLICATION FOR FAULT-TOLERANCE & CAPACITY

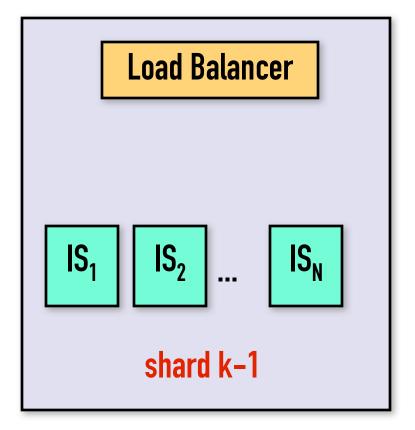
- ➤ Synchronization index shards:
 - > can be avoided because data is read only
 - but one needs to be able to update the index!
- ➤ Practical requirements:
 - ➤ no downtime by update of index
 - rebalancing of mapping of URLs to shards
 - needs to be supported

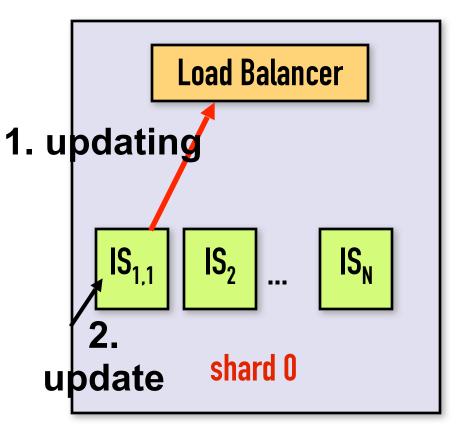
- ➤ To update index server
 - divert all requests to other servers in the pool

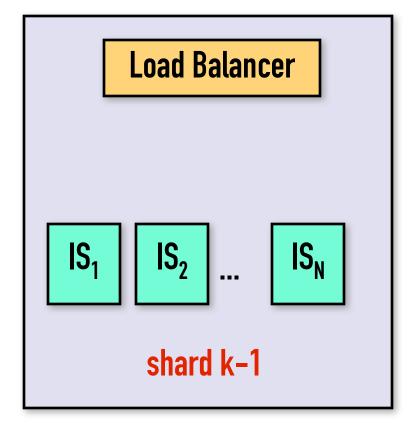


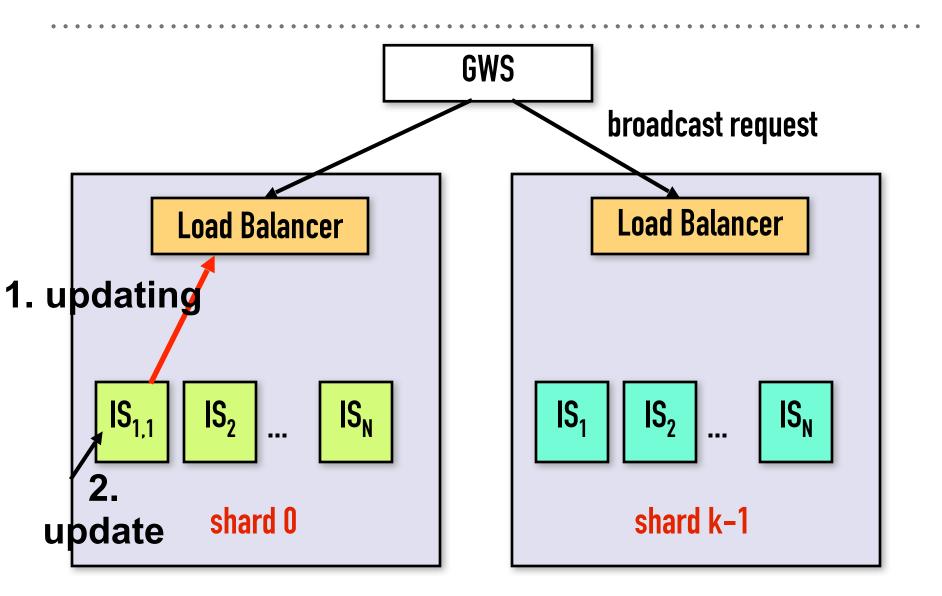


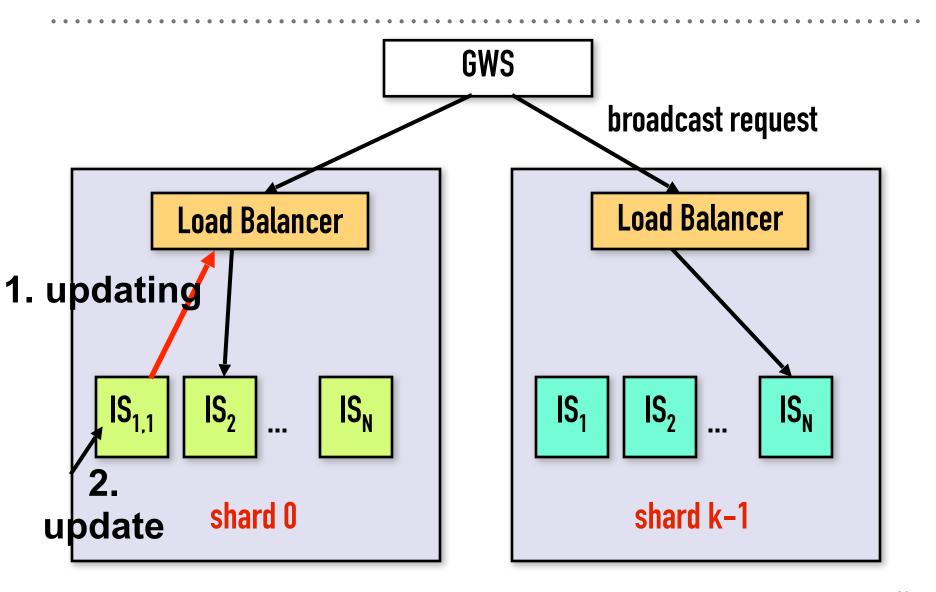












HOMEWORK:

- ➤ What if mapping of URLs to index servers changes?
 - ➤ e.g., could have double entries, or
 - > omitted entries

- ➤ How to deal with this?
 - without additional overhead in index servers?

REAL-TIME PROBLEM

- ➤ What if index is continuously updated?
 - > e.g., news sites, tweets must be indexed in near real-time
 - ➤ how can we update index efficiently?

SCALABILITY LESSONS

Scalability:

- ➤ use the inherent parallelism in the application
 - e.g., retrieve doc list in parallel & inexpensive merge
- ➤ use multiple clusters to divide the load
- ➤ Increase the number of shards
 - > to scale with an increase of the index
 - ➤ to accommodate slower CPUs
- Increase the number of servers per pool
 - to increase throughput of system

REFERENCES

- [1] L. A. Barroso, J. Dean, and U. Hölzle. Web search for a planet: The google cluster architecture. IEEE Micro, 23(2):22–28, 2003.
- [2] Luiz André Barroso and Urs Hölzle, "The Datacenter as a Computer", 2013 by Morgan & Claypool. (available online)