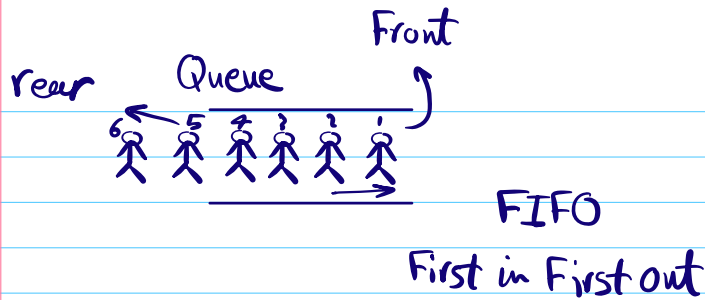


LIFO  
Stack  


Queues  
Intro.



Topics

- 1 - Qs
- 2 - Qs with LL
- 3 - Q with SS
- 4 - Perfect num
- 5 - largest item in window of size B

Queue  
Operations

enqueue Insert a new object from rear

dequeue Removes & returns the object from front end

isEmpty whether any item is in Q

Peak  $\leftarrow$  front get the front element

rear get the rear element

Similar  
to Stock  
Peak



# P1 Implement queue using Linked List:

enqueue Q.eng(7)

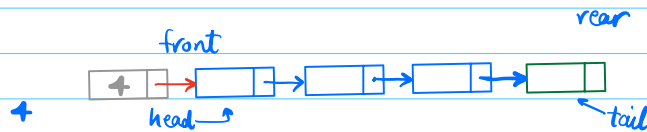
6 3 \* +

Q.eng(3)

rear front

Q.eng(6)

dequeue Q.dq()



new  
node

O(1)

```
void enqueue(int x) {
    nn = new node(x);
    if (head == null) { // isEmpty()
        head = nn;
        tail = nn;
    } else {
        1 tail.next = nn;
        2 tail = nn;
    }
}
```

```
int dequeue() {
    if (head == null) { // isEmpty()
        error, ex, lag;
        ret;
    } else { // Q is not empty
        x = head.data;
        head = head.next;
        if (head == null) { // isEmpty()
            tail = null;
        }
        ret x;
    }
}
```

Q.isEmpty

```
bool isEmpty() {
    ret (head == null) // O(1)
}
```

```
int front() {
    if (isEmpty()) ret ex, error, ...;
    else ret head.data;
} // O(1)
```

equivalently

```
int rear() {
    if (tail == null) { // isEmpty()
        ret ex, error, lag etc;
    } else {
        ret tail.data;
    }
} // O(1)
```

## P2 Implement Queue using 2 stacks.

front  
rear



Q.eng(4)

Q.eng(3)

Q.dq()

Q.dq()

Q.eng(6)

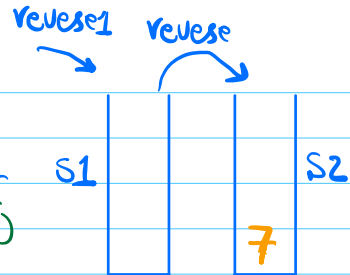
Q.dq()

rear  
outputs

front

6 3 4  
③ ② ①

Q



~~1, 2, 3, 4, 5, 6, 7~~

```
void enqueue(int x){
```

```
    st1.push(x)
```

```
}
```

```
bool isEmpty(){
```

```
    return (st1.isEmpty() && st2.isEmpty())
```

```
}
```

```
int dequeue(){
```

```
    if (isEmpty()) return ex, error, lag etc.
```

```
    → if (st2.isEmpty()) { // re Fill only if empty
```

```
        while (!st1.isEmpty()) {
            st2.push(st1.pop())
        }
```

```
    }
```

```
    return st2.pop() o(1)
```

```
}
```

avg  $O(1)$

eng(1)

eng(2)

eng(3)

deq() → ①

eng(4)

eng(5)

deq() → ②

deq() → ③

deq() → ④

eng(6)

eng(7)

deq() → ⑤

deq() → ⑥

$O(1)$

$O(1)$

$\frac{1}{n} \{ O(1) \}$

$O(n)$

∇ item average asymptotic analysis

2push+ 2pop

for each item

break?

Perfect number: any number that is composed of 1 and/or 2

P3 Find  $n^{\text{th}}$  number with digits only 1 or 2 or both.

1  
2

1, 2, 11, 12, 21, 22, 111, 112, 121, 122, ...

1st 2nd 3rd 4 5 6 7 8 9 10

1, 2, 3, ..., 9, 10,

11, 12, 13, 14, ..., 19, 20

21, 22, 23, ... 30

31 ... 40

100 101 ... 111 112 113 ...

9  
10

6th  $\rightarrow$  ans = 22

8th  $\rightarrow$  ans = 112

$j^{\text{th}}$

idea 1:

list  $i^{\text{th}}$  Number  
for  $i=1, 2, \dots$ , upper-bound

check if digits are only 1 and/or 2

idea 2:

use Q

if ( $n^{\text{th}} \leq 2$ ) ret nth

init. Q

q.enqueue(1); q.enqueue(2)

for ( $i=3$ ;  $i \leq n^{\text{th}}$ ;  $i+=2$ ) {

x = q.dequeue();

y = 10 \* x + 1

z = 10 \* x + 2

if ( $i == n^{\text{th}}$ ) ret y;

if ( $i+1 == n^{\text{th}}$ ) ret z;

q.enqueue(y); q.enqueue(z)

}

f  
x, x, 11, 12 21 22

smaller nums  $\rightarrow$  to large  
nums

$\log_{10} x$

x.toString() + "1"

+ "2"

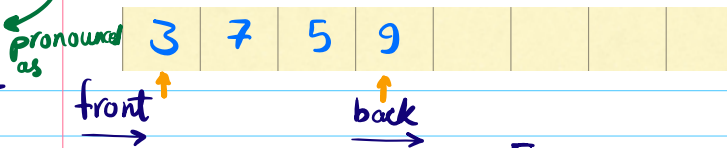
if x, y, z overflow

TC: HW

$O(n)$

$\downarrow$   
nth

dequeue  
(deck)  
deque  
intro  
cheque



- ① add front → ③ remove front
- ② add back → ④ remove back

① & ④ → Normal Q

① & ③ → Stack

P4 Given an int array  $a$ , find max in all sliding window of size  $B$ .  
output is an array for windows of size  $B$ .

ex

0	1	2	3	4	5	6	7
1	3	-1	-3	5	3	6	7

$B=3$

$$8 - B + 1 = 6$$

$O(B)$

$B$

$(n - B + 1)B$

output: { 3, 3, 5, 5, 6, 7 }

ex

0	1	2	3	4	5	6	7	8
1	2	3	4	2	7	1	3	6

$B=6$

$$9 - B + 1 = 4$$

$B$

output: { 7, 7, 7, 7 }

idea 1

TC:  $O(n \times B)$

idea 2

size PQ keep it  $B$ ,  $O(n \log B)$  optional. later in heap  
max heap

idea 3

0	1	2	3	4	5	6	7
1	3	-1	-3	5	3	6	7

step

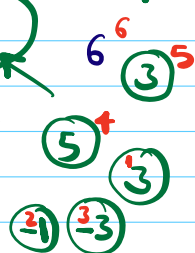
Start index

$A[i]$

\* front of deque is always max of window  $B$

$B$

deque



HW

$O(N+B)$

1, 1, 5, -1, -2, -3, 4, ...

// q is my deque

n = A.length

int start = 0 // start index of window B

```
for (i = 0; i < B; i++) {  
    while (!q.isEmpty() && A[q.back()] < A[i])
```

```
        q.removeBack();
```

```
    q.addBack(i);
```

```
}  
ans.add(A[q.front()]);
```

```
for (i = B; i < n; i++) {
```

```
    // A[i]
```

```
    while (!q.isEmpty() && A[q.back()] < A[i]) {
```

```
        q.removeBack();
```

```
    }
```

```
    q.addBack(i)
```

```
    // remove left most from deque only if it was  
    // affecting max
```

```
    if (start > q.front()) q.removeFront();
```

```
    ans.add(A[q.front()]);
```

```
    start++
```

```
return ans;
```

while  
might  
be  
easier  
to understand

