Strings:

- Intro, ASCII codes
- Flip
- sort char array
- Longest Palindromic Substring

Strings def. : array of chars.

characters & ASCII values Google ASCII table

-32 +32

a-z 26	0	'A' → 65	'a' → 97	'0' → 48
	1	'B' → 66	'b' → 98	'1' → 49
	2	'C' → 67	⋮	⋮
ch-65	⋮			'9' → 57
ch-'A'	25	'Z' → 90	'z' → 122	'10' → x 49 48

-32 +32

↳ index

↳ not single char

C/C++

8 bit 1 byte

~~UTF-8~~

~~UTF-16~~ ...

Java	Python
println((int)'A') 65	print(ord('A')) 65
println((char)66) B	print(chr(66)) B
println('B'>'A') true	print('B'>'A') True
println('C'<'A') false	print('C'<'A') false
println('C'-'A') 2	print(ord('C')-ord('A')) 2

32

'a'-'A'

ex string s="abcd"

0 1 2 3

index print(s[0]) → a s.charAt(0)

print(s[2]) → c

P1 Given a char array of only alphabet chars

toggle every char Capital ⇄ Small

ex AnaConDa
aNAcONdA

no built in functions → 'a' ... 'z' ↔ 'A' ... 'Z'

Quiz*
TC: O(n)
SC: O(1)

```
char[] toggle(char S[]){
    int n = S.len
    for(i=0; i<n; i++){
        if('a'<=S[i] && S[i]<='z'){
            S[i] = (char)(S[i] - 32)      -> ('a'-'A')
        }else{ // capital
            S[i] = (char)(S[i] + 32)
                                          ('a'-'A')
        }
    }
    ret S
}
```

32
0010 0000

Part 2: optional
use XOR & Bitmanipulation

'A' → 65 0100 0001 (7 6 5 4 3 2 1 0)
'B' → 66 0100 0010 (7 6 5 4 3 2 1 0)
⋮
'Z' → 90 0101 1010

'a' → 97 0110 0001
'b' → 98 0110 0010
⋮
'z' → 122 0111 1010

1 ^ 1 = 0
0 ^ 1 = 1

S[i] = S[i] ^ 32 ← solution one line
(1<<5)

P2 Given a char array s[], that contains only lower case alphabet letters, sort s[] chars in alphabetic order

no of chars in arr $1 <= n <= 10^5$

$'a' <= ch[i] <= 'z'$

ex ch[]: { d, a, b, a, c, d, z, d }

a, a, b, c, d, d, d, z

idea1: Arrays.sort(ch) ✓

TC: n log(n)

SC: O(1)

exist[ch - 'a']

ch[]: { d, a, b, a, c, d, z, d }

⚠ idea2: {a, a, b, c, d, d, d, z}

freq count

freq[index]++

	char	count
0	'a'	2
1	'b'	1
2	'c'	1
3	'd'	3
4	'e'	0
⋮		
25	'z'	1

{a, a, b, c, d, d, d, z}

i

k

Quiz

TC: O(n)

SC: O(1)

```
void sortString(char S[]){
    int n = S.Len
    int freq[26] = {0}
    for(i=0; i<n; i++){
n       index = S[i] - 'a'
        freq[index]+=1
    }
    k=0
    for(i=0; i<26; i++)
        char ch = i + 'a'
        for(j=0; j<freq[i]; j++){
n           S[k] = ch; k++;
        }
    }
    ret S;
}
```

O(n+n)

Substring ⟶ Sub array

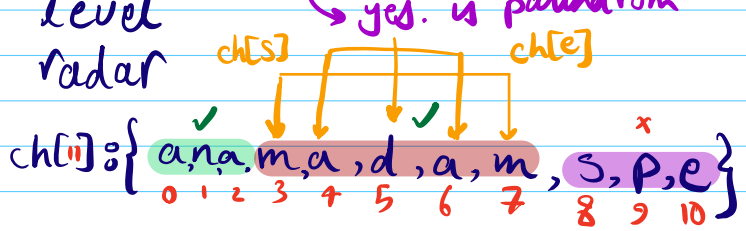
1. Continouse part of string
2. Full string is substring
3. Single char can be substring

how to identify a substring?

s, e → start index, end index

P3 Check if given subString is palindrom or not? → true/false

ex mom dad madam xyyx civic level radar "x" → yes. is palindrom



```
bool isPalindrom(char ch[], int s, int e){
    while(s<e){
        if( ch[s] != ch[e]) ret false;
        else{ // ch[s] == ch[e]
            s++
            e--
        }
    }
    ret true
}
```

TC: O(n)

SC: O(1)

$\frac{e-s}{2} \rightarrow \frac{n}{2}$

output int

P4 Given string, calculate length of longest Palindrom substring

Quiz* a b a c a b
0 1 2 3 4 5
ans = 5

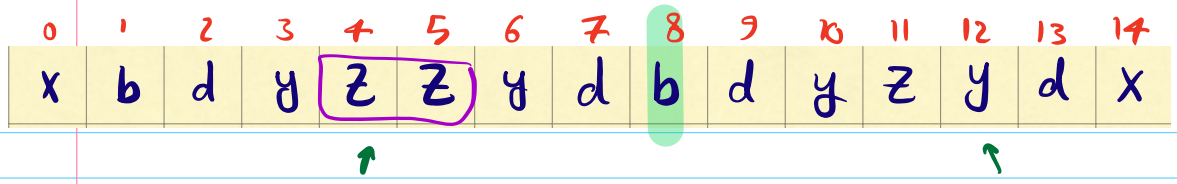
Quiz* a b c d e
0 1 2 3 4
ans = 1

Quiz*
TC: $O(n^3)$
SC: $O(1)$

```
int longestPalindrom1(char ch[]){
    int n=ch.Len    ans=0
    int ans = 0;
    for(i=0; i<n; i++){
        for(j=i; j<n; j++){
            if( isPalindrom(ch,i,j) ){   // O(n) [a,b]
                ans = Max(ans, j-i+1)
                                  a-b+1
            }
        }
    }
    ret ans
}
```

idea

ex



Case 1: pallindrom has odd length with the center

given $i \rightarrow s$, e

stretch as much as stay pallindrom

madam

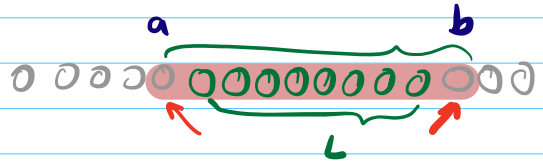
Case 2: pallindrom has even length with the center

pair is given $(i, i+1)$

stretch as much as stay pallindrom

```
int expand(char S[], int P1, int P2){
    while(P1>=0 && P2<S.Len && S[P1]==S[P2])
        P1--; P2++;
    }
    ret P2-P1-1
}
```

O(n)



L2 - 2 = L

P2 - P1 + 1 - 2 = P2 - P1 - 1

b - a + 1

Quiz*

TC: O(n²)

SC: O(1)

```
int longestPalindrome2(char ch[]){
    n = ch.len
    ans = 1
    for(i=0; i<n; i++){   // odd case ①
        P1 = i; P2 = i
        ans = max(ans, expand(ch, P1, P2))
    }
    for(i=0; i<n-1; i++){   // even case ②
        P1 = i; P2 = i+1;
        ans = max(ans, expand(ch, P1, P2))
    }
    ret ans
}
```

n

O(n²)

O(n²)