- Intro
- Fractional knapsak
- Free Cars
- Candy
- Max Jobs

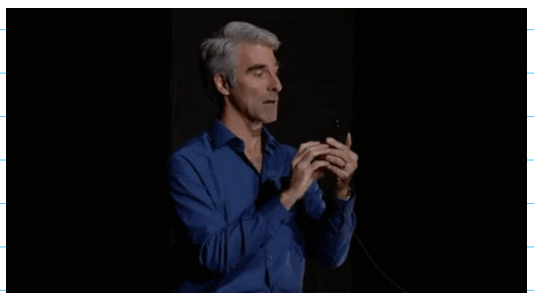
Greedy Algorithm (short sightet in a positive way!)

Choose the best at current

iphone price:

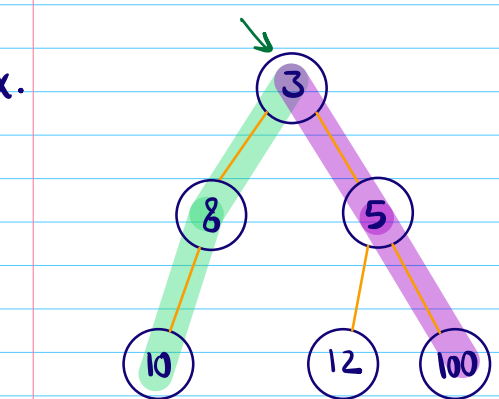
Amazon	$950
Apple	$999.99
eBay	$989.99

greedy



iphone announcement event

ex.



sum=21

ans=0

sum=108

Find a path that sum(nodes) is max

greedy aproach doesn't work all the time.

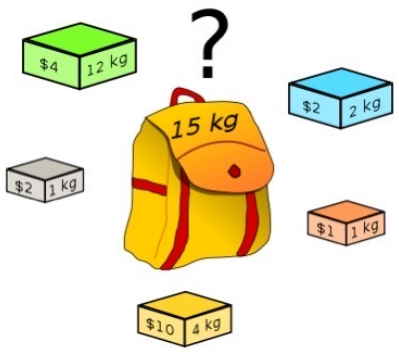
Sometime it does.

max $

P1 Fractional knapsacks

protein in grams

Food	protein Content in the current amount	Protein/kg
Tomato	200 for 20kg	10
Apple	180 for 15kg	12
Onion	250 for 50kg	5
Chicken	150 for 10kg	15
Potato	200 for 25kg	8
Mango	132 for 12kg	11
Seafood	100 for 5kg	20



16kg 1kg 20 → 5kg 6kg

maximiz protein in take if bag capacity is 7 kg.

```
double MaxProtein(int wt[], int protein[], int k){
    List<(int, int, double)> a = new ...
    for(i=0; i<wt.Len; i++){                          n
        p = (wt[i], protein[i], protein[i]*1.0/wt[i])   (density)
        a.Add(p)
    }
    Array.Sort(a, (t1, t2) => t1.density.cmp(t2.density))
    double ans=0
    for(i=0; i<a.Len; i++){                           n
        (w, p, d) = a[i]
        if(w <= k){ ans += p; k=k-w }
        else { ans += d*k; break; }
    }
    ret ans
}
```

C# tc: O(nlogn) SC: O(n)

ans ← 100

10kg, 150, 15

10kg 150

2kg ?

2kg

Seafood → 5kg

0_1 Knapsack problem → what is? Future session → greedy gives wrong answer

Free cars!

P2 There is a limited time sale on toys. Sale is 100% off!! (Car / Free)

A[i] → sale end time for ith toy

B[i] → Beauty of ith toy.

time starts from t=0, it takes 1 unit of time to buy one toy & toy can only be bought if t < A[i]. (t: current time)

buy toys s.t. sum of beauty of toys is maximized.

deal expiration time

ex A = { 3, 1, 3, 2, 3 }

beauty

B = { 6, 5, 3, 1, 9 }
(1 2 3 4 5)

t	toy#	b
0	5	9
1	1	6
2	3	3

① 3 — 18

t	toy#	b
0	2	5
1	5	9
2	1	6
3		

② — 20

Quiz

A = { 1 2 }

B = { 3 1500 }
(1 2)

t	toy#	b
0	2	1500
1	X	

①

t	toy#	b
0	1	3
1	2	1500

② 1503

ex

	1	2	3	4	5	6	7	8
A = {	1	3	3	~~3~~	5	5	5	8 }
B = {	5	2	7	~~x~~	4	3	8	1 }

time stamp

b

t	A end time	B
0	1	5
1	~~3~~ 5	~~2~~ 8
2	3	7
3	5	4
4	5	3
5	5	8

Pair & Custom cmp

Code

A&B are sorted based on time ascending O(nlogn)

```
totalBeauty=0 ; t=0 ;
for(i=0;i<n;i++){
{ if(t>=A[i]){
    if((B[i] > root of minheap)){
      b = minHeap.RemoveMin()
      totalBeauty -= b
      totalBeauty += B[i]
    } minHeap.insert(B[i])
  }else{ t<A[i]
      minHeap.insert(B[i])
      t++
      totalBeauty += B[i]
  }
}
```

check heap is not empty

log(n) pop

regret/redeem

min heap: 5, 8, 7, 4, 3

n

TC: O(nlogn)

SC: O(n)

	1	2	3	4	5	6	7	8
A={	1	3	3	3	5	5	5	8 }
B={	5	2	7	1	4	3	8	1 }

↑ i

t=0 1 2 3 4 5

B=0 5 7 14 18 24

27

(grades)

P3 There are N students with marks, teacher has to give them candies s.t.

a) Every student should have at least one candy.

b) Students with more marks than neighbors have more candies than them. → immediate neighbors

Find min candies to distribute.

ex

7 ← A = {1, 5, 2, 1} (indices 0 1 2 3)

candies: 1 3 2 1

Quiz A = {4, 4, 4, 4, 4} (indices 0 1 2 3 4)

candies: 1 1 1 1 1

ans = 5

ex

7 ← A = {8, 10, 6, 2} (indices 0 1 2 3)

candies: 1 3 2 1

i-1 i Pass 1 ←

Pass 2 →

Quiz A = {1, 6, 3, 1, 10, 12, 20, 5, 2}

candies: 1 2 1 2 1 2 3 4 2 1

ans = 19

c[] ← 3 →

① for all i, c[i] = 1

② for all i, if (A[i-1] < A[i]) ⇒ if (c[i-1] >= c[i]) c[i] = c[i-1] + 1

(A[i] is the grade)

TC: O(n)

SC: O(n)

for i = 1 → n-1

Pass 1 →

③ for all i, if (A[i] > A[i+1]) ⇒ if (c[i] <= c[i+1]) c[i] = c[i+1] + 1

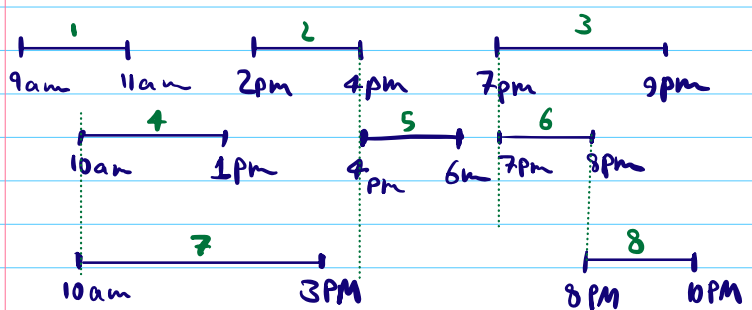
for i = n-2 → 0

Pass 2 ←

Facebook

P4 Given n jobs, with their start & end time, find max jobs that can be completed if only one job can be done at a time.

ex



Quiz

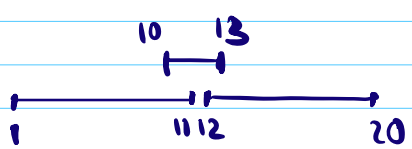
$S = \{1, 5, 8, 7, 13, 12\}$ (1 2 3 4 5 6)

$e = \{2, 10, 10, 11, 12, 20\}$

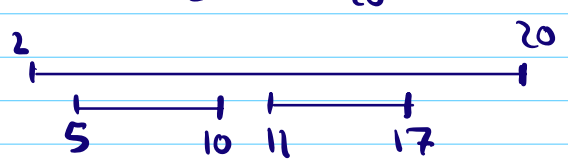
ans = 3

duration[i] = e[i] - s[i]

idea 1 duration ✗



idea 2 start time? ✗



idea 3 end time → sort ascending
start early + less duration

Code

arr<(s,e)> a = new ...

for (i = 0 → n-1) a.add(s[i], e[i])

a.sort(cmp)

// sort based on end time

```
ans = 1
lastEnd = a[0].e
for i → 1 to (N-1){
    if(a[i].s >= lastEnd){
        ans++;
        lastEnd = a[i].e
    }
    else NoOp
}
ret ans;
```

	0	1	2	3	4	5
s =	1	5	8	7	13	12
e =	2	10	10	11	19	20

lastEnd

i

ans = 3