



## Interview Problems 2

Topics: 0, 1, add ✓  
 - 0, 1, swap ✓  
 - triplets  
 - Josephuse

Google  
 TC

Q1 Given an array of 0 & 1s, we can replace exactly one of the 0s with a 1, return the count of max consecutive 1's in the array:

ex  $a[] = \{1, 1, 0, 1, 1, 0, 1, 1\}$   $\text{ans} = 5$   $\{1, 1, 1\}$

<sub>0 1 2 3 4 5 6 7</sub>

Quiz

$a[] = \{1, 1, 0, 1, 1, 0, 1, 1, 1\}$   $\text{ans} = 6$

<sub>0 1 2 3 4 5 6 7 8</sub>

Quiz

$a[] = \{0, 1, 1, 1, 0, 1, 1, 0, 1, 1, 0\}$   $\{1, 0, 0, 1\}$

<sub>0 1 2 3 4 5 6 7 8 9 10</sub>

$\underbrace{0+1+3}_4$      $\underbrace{3+1+2}_6$      $\underbrace{2+1+2}_5$      $\underbrace{2+1+0}_3$

← lcount 0 rcount →

$\text{ans} = \max(\text{ans}, \text{lcount} + 1 + \text{rcount})$

ex  $a[] = \{1, 1, 1, 0, 1, 1, 0, 1, 1, 1, 1, 0, 0, 1, 1, 0, 1, 1\}$

<sub>0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17</sub>

$\underbrace{3+1+2}_4$      $\underbrace{2+1+4}_7$      $\underbrace{4+1+0}_6$      $\underbrace{0+1+2}_3$      $\underbrace{2+1+2}_5$

$\cancel{4}$     7    6    3    5

i  
j

Quiz

TC:  $O(n)$

SC:  $O(1)$

```
int CountMaxOnes(int a[]) {
```

{0,0,0}

ans = 0

① Count all 1s. #1s == n ret n  $O(n)$

② for (i=0; i<n; i++) {

if (a[i] == 0) {

lcount ← look left Count 1s till 0 or start of arr

rcount ← look right Count 1s till 0 or end of arr

ans = max(ans, lcount + 1 + rcount)

}

}

③ ret ans

}



a[]: { 1, 1, 1, 0, 1, 1, 0, 1, 1, 1, 1, 0, 0, 1, 1, 0, 1, 1 }

$O(4n) = O(n)$

$O(n + n + n + n) = O(4n)$

P2 Given an array of 0 & 1s, we can **swap** exactly one of the 0s with a 1. Return the count of max consecutive 1's in the array:

ex a: { 1, 0, 1, 1, 0, 1 }      0, 0, 1, 1, 1, 1  
4 ones

```
010110101
01010010
101001
```

Quiz { 1, 1, 0, 1, 1, 1 }      0, 1, 1, 1, 1, 1  
ans = 5

ex a[]: { 1, 1, 0, 1, 1, 0, 1, 1, 1 }

{ 0, 0, 1, 1, 1, 0, 1, 1, 1 }      { I ↓count + 1 + rcount  
 II ↓count + rcount }

```
int CountMaxOnes(int a[]) {
    ans = 0
```

① Count all 1s. #1s == n ret n // pre calculate # of 1s

② for (i = 0; i < n; i++) {

if (a[i] == 0) {

lcount = look left Count 1s till 0 or start of arr

rcount = look right Count 1s till 0 or end of arr

neighborOnes = lcount + rcount

if (neighborOnes == #1s) chainLen = lcount + rcount

else chainLen = lcount + rcount + 1

ans = max(ans, chainLen)

}

③ ret ans

}

TC

SC



P3 Given an array  $a[n]$  Count the number of triplets.

triplets def.  $\begin{cases} i < j < k \\ a[i] < a[j] < a[k] \end{cases}$

ex  $a = \{1, 2, 3, 4, 5\}$   $ans = 10$

$(1, 2, 3), (1, 2, 4), (1, 2, 5), (1, 3, 4), (1, 3, 5), (1, 4, 5)$   
 $(2, 3, 4), (2, 3, 5), (2, 4, 5)$   
 $(3, 4, 5)$

Quiz  $a[] = \{2, 6, 9, 4, 10\}$

Count = 5  
ans

i	j	k	a[i]	a[j]	a[k]
0	1	2	2	6	9
0	1	4	2	6	10
0	2	4	2	9	10
0	3	4	2	4	10
1	2	4	6	9	10

```

int CountTriplet(int a[]) {
    n = a.len; Count = 0;
    for (i = 0; i < n; i++) {
        for (j = i+1; j < n; j++) {
            for (k = j+1; k < n; k++) {
                if (a[i] < a[j] && a[j] < a[k]) {
                    Count++;
                }
            } // k
        } // j
    } // i
    return Count;
}

```

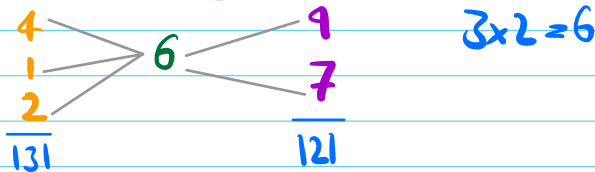
TC:  $O(n^3)$

SC:  $O(1)$



ex

$a[] = \{4, 1, 2, 6, 9, 7, 2\}$



hint  
Contribution  
technique

① 1 to  $n-1$

① iterate over all elements, consider  $a[i]$  as center

① for loop left of  $c$  count if  $a[j] < c \rightarrow l$

① for loop right of  $c$  count if  $a[j] > c \rightarrow r$

ans +=  $l \times r$

② ret ans

TC:  $O(n^2)$

SC:  $O(1)$

why  
time

Complexity

is  $O(n^2)$

```
for(i=1; i<n-1; i++){
    for(j=i-1; j>=0; j--) if...
    for(j=i+1; j<n; j++) if...
}
```

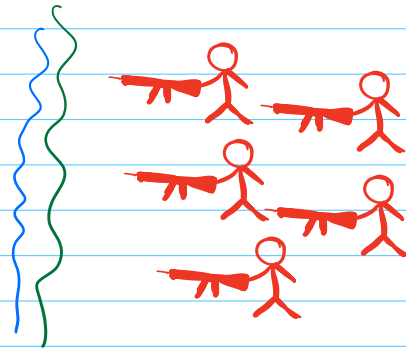
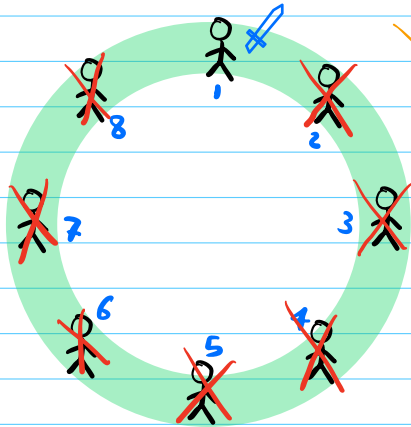
$n$   $n \times n$   $x+y \sim n$

Puzzle problem, not common in interviews!

P4 Josephus problem:

\*find which location in the circle will stay alive

input:  
N



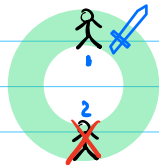
N=1



ans=1

Quiz

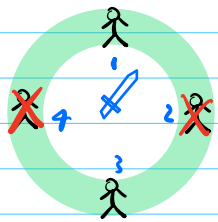
N=2



ans=1

Quiz\*

N=4

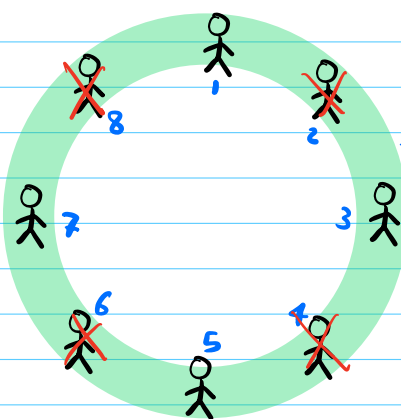


1 round

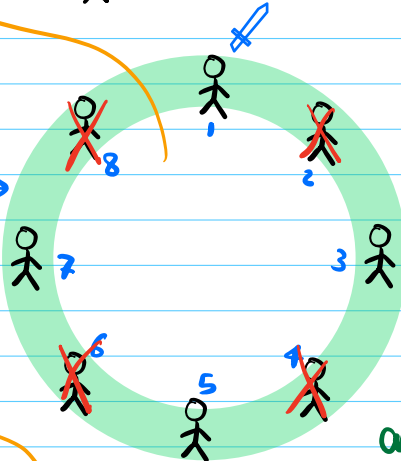


ans=1

N=8



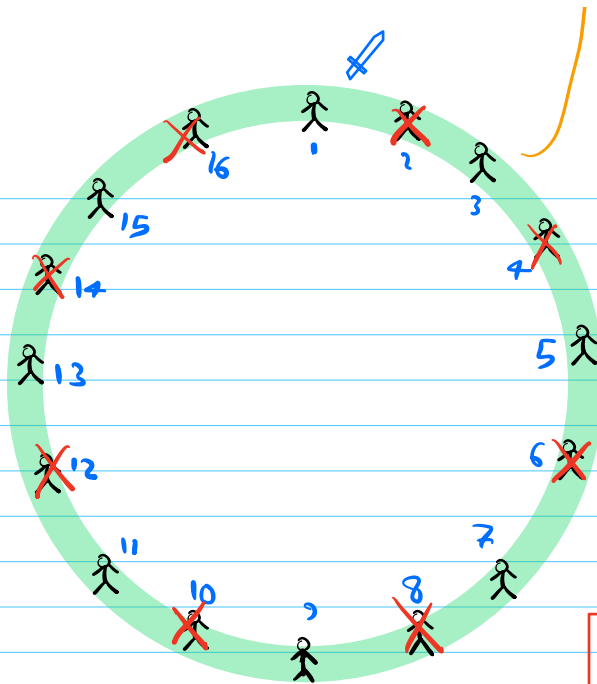
1 round



ans=1

Quiz

$$N=16$$



$$ans = 1$$

$$N = 2^n \quad 2^{n-1}$$

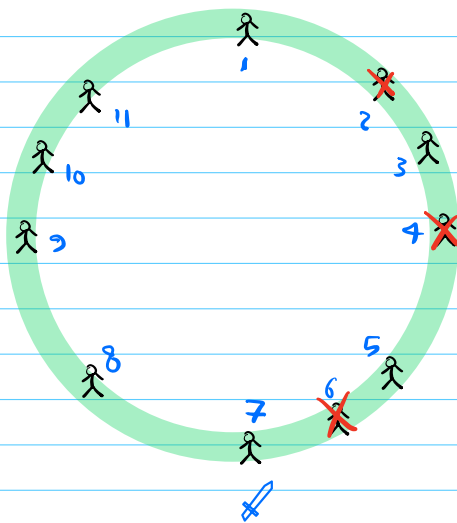
$$N \xrightarrow{\text{1st round}} \frac{N}{2}$$

proof by induction

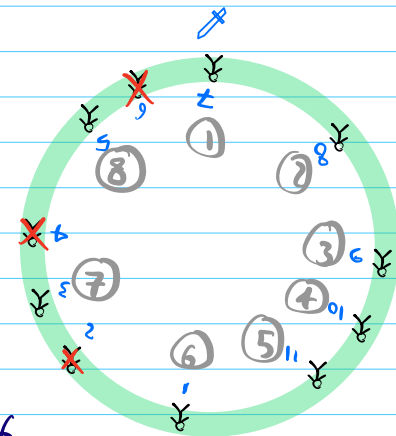
for all  $N = 2^n$   
Josephus problem ans.  
in 1

Quiz

$$N=11$$



remaining  
 $8 = 2^3$



$$8 < 11 < 16 \\ 2^3 < 11 < 2^4 \quad \lfloor \log_2 11 \rfloor \rightarrow 3$$

step 1 - find the closest power of 2  $N = 11 - 8 = 3$  **k**

$$a(1) \quad N - 2^{\lfloor \log_2 N \rfloor} = k$$

$$11 - 2^{\lfloor \log_2 11 \rfloor} = 3$$

step 2 -  $2k + 1 \rightarrow ans$

$$2 \times 3 + 1 = 7$$

$$100 - 64 = 36$$

$$36 \times 2 + 1 = 73$$