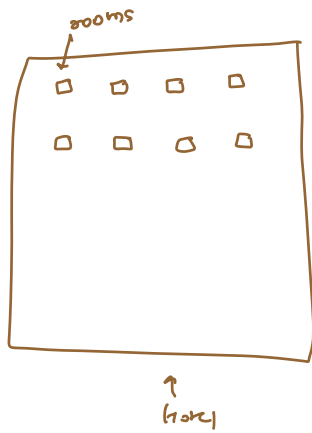


Hashing - 1



100 rooms (Register)



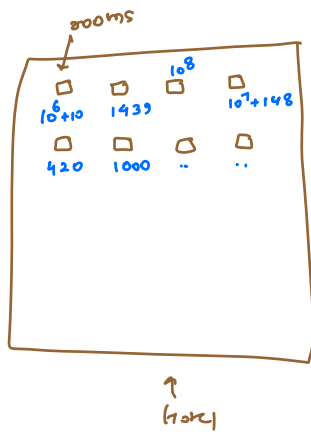
1000 rooms

Boolean arr [1001]



70th room

arr[70]



1000 special numbers.



Each

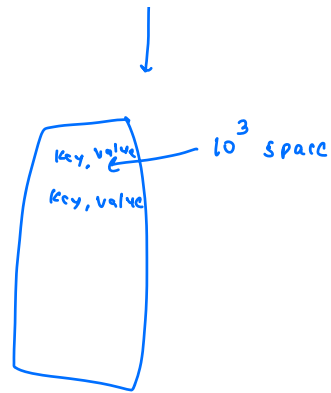
number

$[1 - 10^9]$

arr [10⁹] → 10⁹ space

10² is actually used

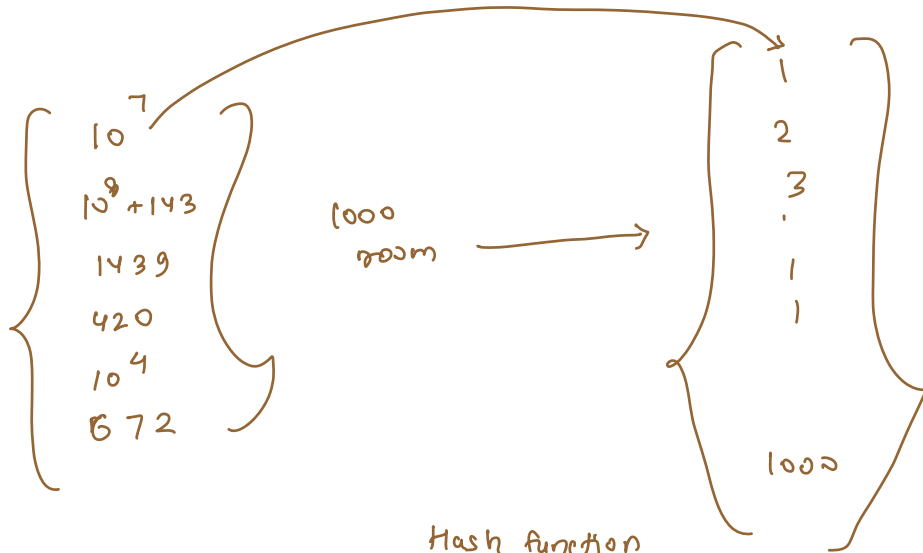
Hashing



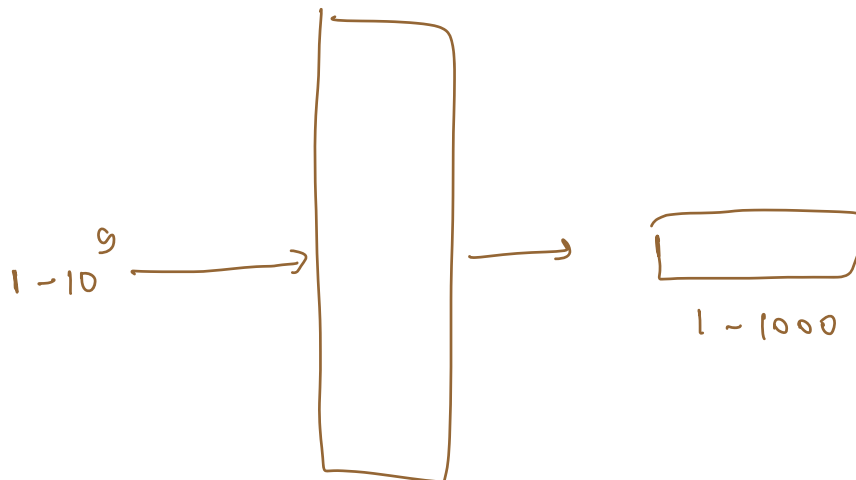
Tc: $O(1)$

Insertion
deletion
searching

HashMap
Hashed

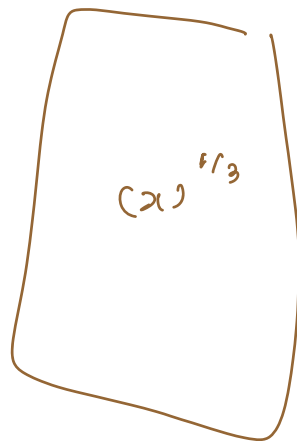


Hash function



$$y = f(x)$$

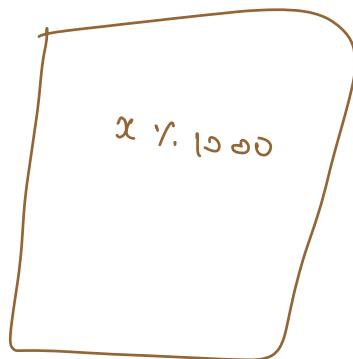
$1 - 10^9$



$1 - 10^3$



$1 - 10^9$



$0 - 999$



1

$1 \% 1000$



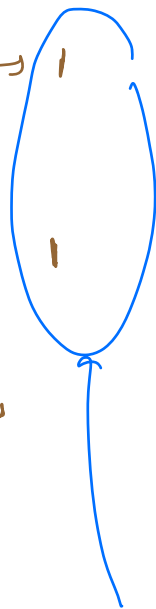
1

1001

$1001 \% 1000$



1



Collisions

Chaining

$A = [2, 5, 10, 7, 3, 4, 7, 12, 17, 12, 1]$

↓
x.5
0-4

0 → 5, 10

1 → 1

2 → 2, 7, 7, 12, 17, 12

3 → 3

4 → 4

worst case operation
↓
Tc: $O(N)$

1. Hashfunction very good. ST chances of collision is v. rare.

10000 values

map [1 ~ 1000] in range



$[1 - 10^5]$

101 pigeons

100 holes

Open addressing

T.C: $O(1)$ ✓ Each function
in Hash function

Avg / T.C. practical

Theoretically

(worst & best)

HashMap

$O(1)$

HashSet

$O(1)$

$O(N)$

$O(N)$

ordered set
ordered map

self-balancing
trees

Avg

worst

C++

Java

$O(\log N)$

$O(\log N)$

map

TreeMap

$O(\log N)$

$O(\log N)$

set

TreeSet

Problems

$$i \neq j$$

index

Q. Given an array $[N]$. Find pair (i, j)

s.t. $A[i] == A[j]$ & $|j - i|$ is minimum

return $|j - i|$

Ex

0	1	2	3	4	5	6	7	8
1	2	3	6	1	2	3	2	1

ans

4 $[0, 4]$

2 $[5, 7]$ \rightarrow

ans = 2

BF

check each pair

if values are equal

update ans

Tc: $O(N^2)$

Sc: $O(1)$

0 1 2 3 4 5 6 7 8
 1 2 3 6 1 2 3 2 1

1 → ~~8~~
 2 → ~~7~~
 3 → ~~6~~
 6 → 3

1 - - - 1 - - - 1 - - - 1 - - - 1

ans = ~~4~~
2

↓ ↓
 HashMap < int, int > hm

ans = INT_max

for (i = 0; i < n; i++)

{
 if (arr[i] in hm)
 {
 ans = min(ans, i - hm[arr[i]])
 }
 update/add arr[i] → i
 }

return ans

TC: O(N)
 SC: O(N)

Problems

Q. Given an array $[N]$. Find pair (i, j)
 s.t. $A[i] == A[j]$ & $|j-i|$ is maximum,
 $i \neq j$
 index

Ex

0	1	2	3	4	5	6	7	8
1	2	3	6	1	2	3	2	1

ans = 8 - 0 = 8

1 _ _ _ 1 _ _ 1 _ _ _ 1 _ 1

↓ ↓
HashMap < int, int > hm

ans = 0

for (i = 0 ; i < n ; i++)

{
 if (arr[i] in hm)
 {
 ans = max(ans, i - hm[arr[i]])
 }
 else
 {
 add arr[i] → i
 }
}

return ans

TC: $O(N)$

SC: $O(N)$

Q. Given array, Find longest subarray with

sum = 0.

return len of subarray.

	0	1	2	3	4	5	6	7	8	9	10
arr :	2	1	2	-3	4	3	1	-8	6	-2	1
PF[] :	2	3	5	2	6	9	10	2	8	6	7

$$\text{sum}[L-R] = \text{PF}[R] - \text{PF}[L-1] = 0$$

↓

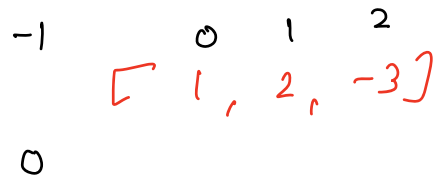
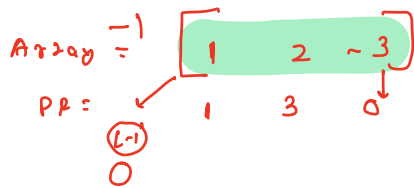
$$\text{PF}[R] = \text{PF}[L-1]$$

	0	1	2	3	4	5	6	7	8	9	10
arr :	2	1	2	-3	4	3	1	-8	6	-2	1
PF[] :	2	3	5	2	6	9	10	2	8	6	7

	0	1	2	3	4	5	6	7	8	9	10
arr :	2	1	2	-3	4	3	1	-8	6	-2	1
PF[] :	2	3	5	2	6	9	10	2	8	6	7

1

previous question



HashMap < int, int > hm

hm.add { 0, -1 }

ans = 0

for (i = 0 ; i < n ; i++)

{

if (^{PF} ~~arr~~[i] in hm)

{

ans = max(ans, i - hm[^{PF} ~~arr~~[i]])

else

{

add ^{PF} ~~arr~~[i] → i

}

}

return ans

TC: O(N)

SC: O(N)

Break
8:31 8:35

6-att

Q. Given an array $[N]$. Find length of longest subsequence which can be reordered into consecutive elements

arr = 101, 4, 3, 6, 10, 20, 11, 5, 100

ans = 4

100, 100 \rightarrow 100, 101 (2)

10, 11 \rightarrow 10, 11 (2)

20 \rightarrow 20 (1)

4, 3, 6, 5 \rightarrow 3, 4, 5, 6 (4)

Idea 1: Sort the array

3, 4, 5, 6, 10, 11, 20, 100, 101
 └───┘ └──┘ └─┘ └───┘
 4 2 1 2

ans = 4

Tc: $O(N \log N + N)$: $O(N \log N)$

Sc: $O(1)$

Idea 2: add hashmap

array: 101, 4, 3, 6, 10, 20, 11, 5, 100
 ↓

101, 4,
 3, 6, 10,
 20, 11, 5,
 100

101 → x

1

4 → 5, 6 → x

2

3 → 4, 5, 6, → x

4

→ ans = 4

6 → x

1

10 → 11, x

2

20 → x

1

11 → x

1

5 → 6, x

2

100 → 101, x

2

Tc: $O(N^2)$

Sc: $O(N)$

↓
hs

arr[] = 101, 4, 3, 6, 10, 20, 11, 5, 100 ↓

101, 4,
3, 6, 10,
20, 11, 5,
100

101 → X

(1)

4 → X

(1)

3 → 4, 5, 6,

(4) → ans = 4

6 → X

(1)

10 → 11

(2)

20 → X

(1)

11 → X

5 → X

(1)

100 → 101, X

(2)

~~Tc: $O(N^2)$~~

~~Sc: $O(N)$~~

→

Tc: $O(N^2)$

Sc: $O(N)$

arr = 4, 4, 4, 4, 4, 4, 4, 4, 5, 6, 7, 8, 9, 10, 11, 12

4	
5	
6	12
7	
11	10
8	
9	

4 → 5, 6, 7, 8, 9, 10, 11, 12

4 → 5, 6, 7, 8, 9, 10, 11, 12

4 → _____

4 → _____

4 → _____

4 → _____

4 → _____

Array → array

(without duplicate)

Put all element in hs

hs

ans = 0

for(iterate of hs.keySet())

x

if (x-1 is not in hs)

cnt = 1

curr = x

while (curr+1 in hs)

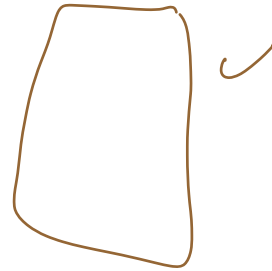
{
cnt++
curr++

ans = max(ans, cnt)

else

nothing

return ans



TC: $O(N + N)$; $O(N)$

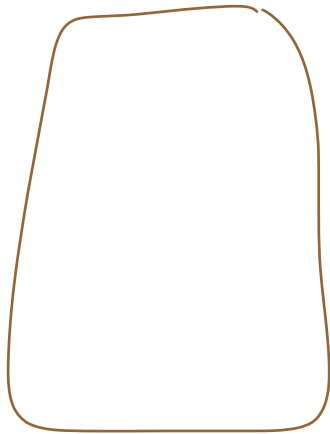
SC: $O(N)$

X

X

4 → 5 → 6 → 7 → 8 → x

3 → 4 → 5 → 6 → 7 → 8 → x



4 interview

45 mins + 15 min

DSA

Googleyness