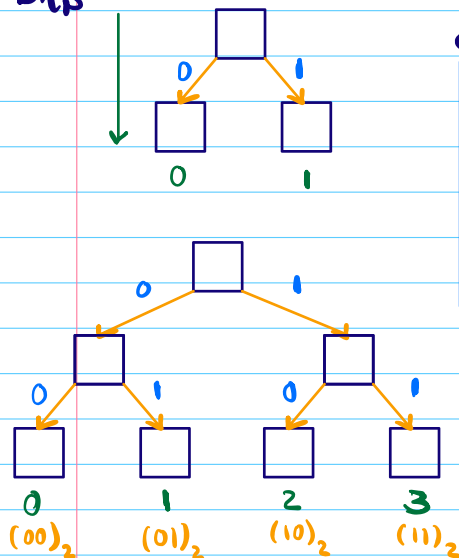


## Tries of Bits

0-1



0  $\rightarrow$  (0)<sub>2</sub>  
1  $\rightarrow$  (1)<sub>2</sub>

3  $\rightarrow$  (11)<sub>2</sub>

```
class Node {
    int data;
    Node bit[];
    Node() {
        bit = new Node[2];
    }
}
```

## Topics - Tries of Bits

- max XOR pair
- max XOR subarray
- BT  $\rightarrow$  list

1-b

$$\begin{aligned} \text{max} \\ (1^b \wedge x) = 1 \\ \text{max} \\ (0^b \wedge x) = 1 \end{aligned}$$

max number  $\rightarrow n$

$$n = 2^H - 1$$

3210  
1010

$$((x \gg 2) \& 1)$$

$$((x \gg i) \& 1) \rightarrow \text{give me bit } i \text{ from } x$$

bit manipulation

$$2n = n \ll 1$$

$$n \gg 1 = n/2$$

P1 Given an array of integers A, find the max value of  $A[i] \wedge A[j]$ , for all possible  $i, j$

ex  $A = \{9, 8, 10, 7\}$

1001 1000 1010 0111

$$9 \wedge 8 = 0001$$

$$7 \wedge 8 = 1111 \rightarrow 15$$

$$\begin{aligned} x \wedge x &= 0 \\ y \wedge x &= x \end{aligned}$$

$$A[i] \wedge A[j]$$

bruteforce

$i, j$

$$TC: O(n^2)$$

$$SC: O(1)$$

ans = 0

for  $i = 0 \rightarrow n-1$

for  $j = i+1 \rightarrow n-1$

$$\text{ans} = \max(\text{ans}, a[i] \wedge a[j])$$

$$\begin{cases} x \wedge x = 0 & i \neq j \\ x \wedge y = y \wedge x & j > i \end{cases}$$

most significant bit

Note

MSB is important!

① 1000 > 0111 ②

intuition & sol.

in 3 step

- Solve manually
- Solve with trie
- Code

{9, 7, 5, 11, 8}

		3	2	1	0
ex 9		1	0	0	1
7		0	1	1	1
5		0	1	0	1
11		1	0	1	1
8		1	0	0	0
		1	1	1	1

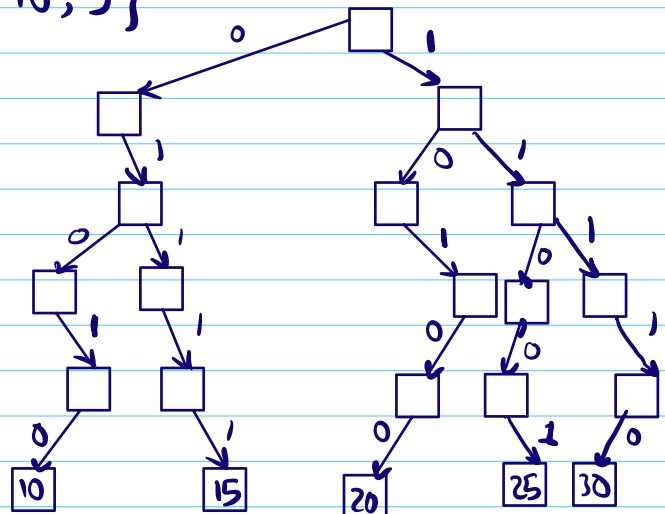
{7, 5}

{9, 11, 8}

ex  $A = \{20, 30, 15, 25, 10, 5\}$

max  
ans

		4	3	2	1	0
0	20	1	0	1	0	0
10	30	1	1	1	1	0
27	15	0	1	1	1	1
27 <sup>21</sup>	25	1	1	0	0	1
30	10	0	1	0	1	0
30 <sup>28</sup>	5	0	0	1	0	1



Code

$O(b)$

$\sim O(1)$

```
void insert(root, X) {
    cur = root
    for (i = 4; i >= 0; i--) {
        b = ((X >> i) & 1)
        if (cur.bit[b] == null) {
            cur.bit[b] = new Node(1)
        }
        cur = cur.bit[b]
    }
    cur.data = X
}
```

ret max XOR in trie  
so far

```
int search(root, X) {
    for (i = 4; i >= 0; i--) {
        b = ((X >> i) & 1)
        t = 1 - b
        if (cur.bit[t] != null)
            cur = cur.bit[t]
        else
            cur = cur.bit[b]
    }
    return X ^ cur.data
}
```

```
main() {
```

```
    ans = a[0] ^ a[1]
```

```
    insert(root, a[0])
```

```
    insert(root, a[1])
```

```
    for (i = 2; i < n; i++) {
```

```
        ans = Max(ans, search(root, a[i]))
```

```
        insert(root, a[i])
```

```
    }
```

```
    return ans
}
```

TC :  $O(n)$

SC :  $O(b \times n) \sim O(n)$

$O(n)$

bit man

P2 Given an integer array, find max subarray XOR of the given array.

ex  $A = \{1, 6, 2\}$

1 0

6 0

2 0

1, 6  $\rightarrow 1 \wedge 6 = 7$

6, 2  $\rightarrow 2 \wedge 6 = 4$

1, 6, 2  $\rightarrow 1 \wedge 6 \wedge 2 = 5$

ans

SC:  $O(1)$

TC:

$O(n^3)$

for  $i = 0 \rightarrow n$

for  $j = i+1 \rightarrow n$

XOR( $a[i], \dots, a[j]$ )  $O(n)$

prefix sum

$$P[i] = a[i] + P[i-1] \quad ①$$

$$\forall i, j, \text{sum}(a[i], \dots, a[j]) \quad ②$$
$$= \begin{cases} P[j] - P[i-1] \\ \text{if } i \geq 0 \rightarrow \text{sum} = P[j] \end{cases}$$

def. prefix XOR

$$a[i] \wedge a[i-1] \wedge \dots \wedge a[0]$$

$$P[i] = a[i] \wedge P[i-1]$$

$$\text{XOR}(a[i] \wedge \dots \wedge a[j]) =$$

$$a[i] \wedge a[i+1] \wedge \dots \wedge a[j] \wedge a[0] \wedge a[1] \wedge a[2] \wedge \dots \wedge a[i-1]$$

$$a[0] \wedge a[1] \wedge a[2] \wedge \dots \wedge a[i-1]$$

$$① \quad X \wedge X = 0$$

$$② \quad 0 \wedge Y = \underbrace{X \wedge X} \wedge Y = Y$$

$$= P[j] \wedge P[i-1]$$

$i \geq 0?$   $\Delta$

$$\text{XOR}(a[i] \dots a[j]) = P[j] \wedge P[i-1]$$

edge case  $i=0$

maximize

$\geq$

maximize

1. Calculate PXOR

TC

SC

2. find  $P[i \wedge \dots \wedge j]$

PS XOR

same nest loop

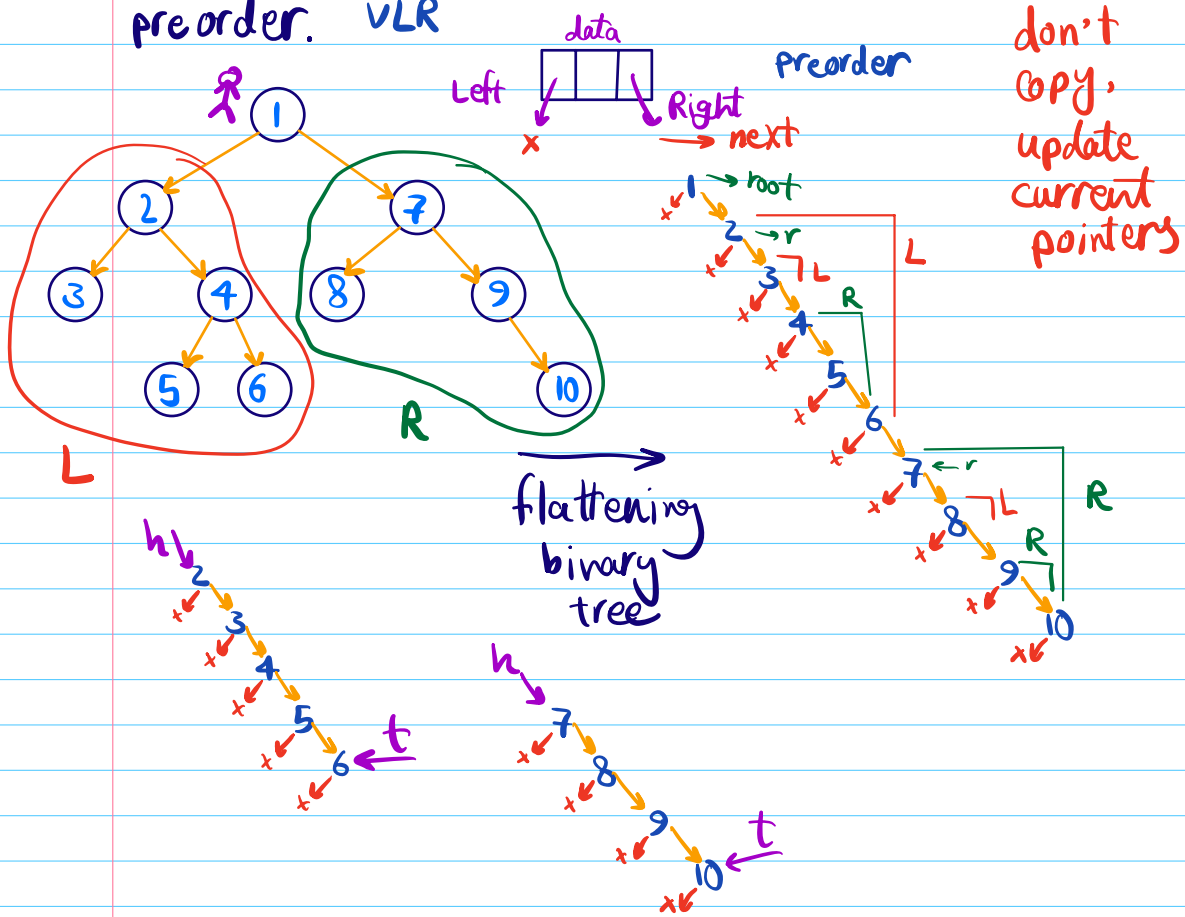
TC:  $O(n + n^2) \sim O(n^2)$  SC:  $O(n)$

code assignment

95% similar

PS max sub array

P3 Convert a given binary tree to linked list in preorder. VLR



```

(head, tail) flatten(root) {
    if (root == null) ret (head, null, tail, null) < Pair >
    // root

```

```

    Pair ← L = flatten(root.left)
    Pair ← R = flatten(root.right)
    root.left = null

```

```

    if (L.head == null && R.head == null) {
        ① ret (root, root)
    }

```

```

    else if (L.head == null) {
        ② root.right = R.head
        ret (root, R.tail)
    }

```

```

    else if (R.head == null) {
        3 root.right = L.head
        ret (root, L.tail)
    }
    else {

```

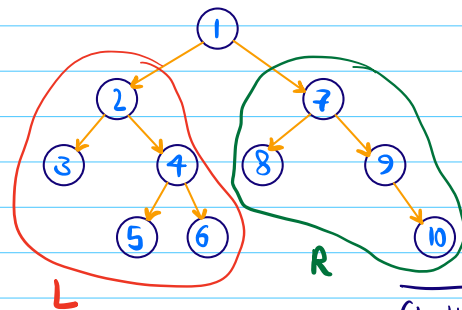
```

        ④ root.right = L.head
        L.tail.right = R.head
        ret (root, R.tail)
    }
}

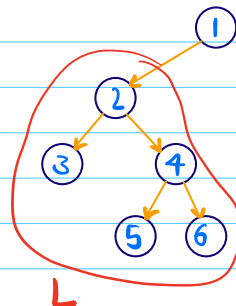
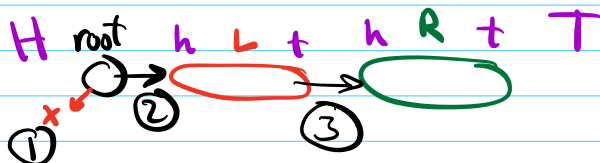
```

4 different cases

	L.head	R.head	
isnull?	T	T	①
	T	F	②
	F	T	③
	F	F	④



flattening  
binary  
tree



TC  $O(n)$   
SC  $O(\text{height})$