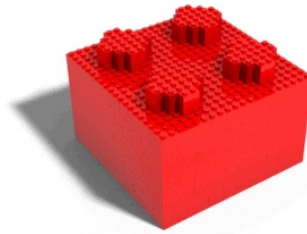


Recursion?

function calling
itself / subproblem



- master theorem

- Topics:
- Quizes
 - kth symbol easy
 - kth symbol hard
 - Towers of hanoi

wikipedia

Optional, check online resources if interested!

Master Theorem:

$$T(N) = T\left(\frac{N}{2}\right) + 1$$

if $\rightarrow T(n) = a T(n/b) + O(n^d)$

$a \geq 1$ $b > 1$ $d \geq 0$

a	b	d
1	2	0

$$b^d = 2^0 = 1 = a$$

solutions:

- ① $b^d > a \rightarrow TC = O(n^d)$
- ② $b^d = a \rightarrow TC = O(n^d \log_b(n))$
- ③ $b^d < a \rightarrow TC = O(n^{\log_b(a)})$

$$TC = O(n^{\log_2(n)})$$

Quiz



Quiz

output 3 2 1

```
void solve(int n= 3){
    if(n==0) ret
    print(n)
    solve(n-1)
}
```

```
void solve(int n= 3){
    if(n==0) ret
    print(n)
    solve(n-1)
}
```

```
void solve(int n= 2){
    if(n==0) ret
    print(n)
    solve(n-1)
}
```

```
void solve(int n= 1){
    if(n==0) ret
    print(n)
    solve(n-1)
}
```

```
void solve(int n= 0){
    if(n==0) ret
    print(n)
    solve(n-1)
}
```

Quiz

output -3 -4 -5 ...

```
void solve(int n= -3){
    if(n==0) ret
    print(n)
    solve(n-1)
}
```

```
void solve(int n= -3){
    if(n==0) ret
    print(n)
    solve(n-1)
}
```

```
void solve(int n= -4){
    if(n==0) ret
    print(n)
    solve(n-1)
}
```

```
void solve(int n= -5){
    if(n==0) ret
    print(n)
    solve(n-1)
}
```

```
void solve(int n= -6){
    if(n==0) ret
    print(n)
    solve(n-1)
}
```

stack overflow

Quiz

```
void solve(int n= 3){
    if(n==0) ret
    solve(n-1)
    print(n)
}
```

output 1 2 3

```
void solve(int n= 3){
    if(n==0) ret
    solve(n-1)
    print(n)
}
```

```
void solve(int n= 2){
    if(n==0) ret
    solve(n-1)
    print(n)
}
```

```
void solve(int n= 1){
    if(n==0) ret
    solve(n-1)
    print(n)
}
```

```
void solve(int n= 0){
    if(n==0) ret
    solve(n-1)
    print(n)
}
```

P1 kth symbol (easy!)

Problem Description

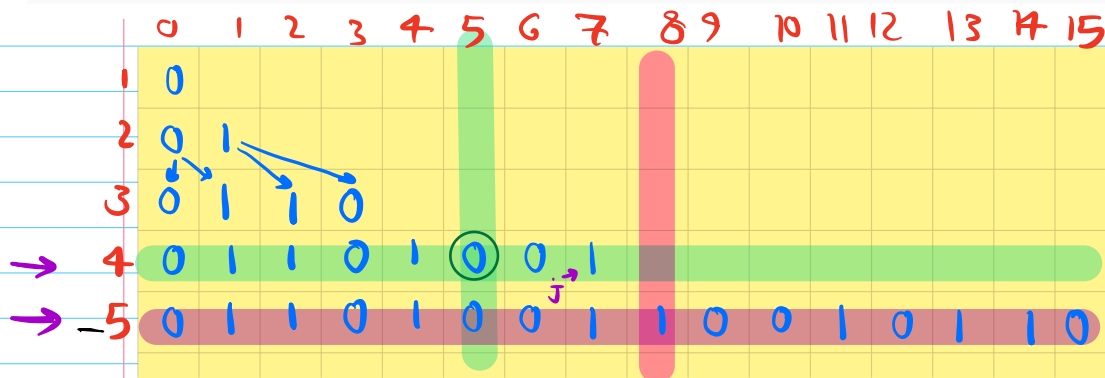
On the first row, we write a 0. Now in every subsequent row, we look at the previous row and replace each occurrence of 0 with 01, and each occurrence of 1 with 10.

Given row number A and index B, return the Bth indexed symbol in row A. (The values of B are 0-indexed.).

Problem Constraints

$1 \leq A \leq 20$

$0 \leq B < 2^{A-1}$



ex1 $\frac{A}{4} \frac{B}{5} \rightarrow 0 \leftarrow \text{ans}$

ex2 $\frac{A}{5} \frac{B}{8} \rightarrow 1 \leftarrow \text{ans}$

TC
 $O(2^A)$

$2^{A-1} + 2^{A-2} + \dots + 2^0$
 $O(2^A)$

row# col# 0-based

```

int kthNum(int A, int B){
    curRow = ArrayList<int>; curRow.append(0);
    for(i=2; i<=A; i++){
        nextRow = ArrayList<int>();
        for(j=0; j<curRow.len; j++){
            if(curRow[j] == 0){
                nextRow.append(0);
                nextRow.append(1);
            }
            else{
                nextRow.append(1);
                nextRow.append(0);
            }
        }
        curRow = nextRow;
    }
    return curRow[B];
}

```

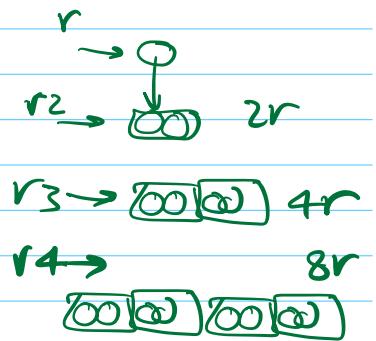
TC

```

ArrayList<int> AthRow(int A){
    if(A==1) return new ArrayList<int>(0);
    List<int> prevRow = AthRow(A-1);
    List<int> r = new ArrayList<int>();
    for(i=0; i<prevRow.len; i++){
        if(prevRow[i] == 0){
            r.append(0);
            r.append(1);
        }
        else{
            r.append(1);
            r.append(0);
        }
    }
    return r;
}

```

SCB
 $O(A + 2^{A-1} + 2^{A-2} + \dots + 2^0)$
 stack
 $O(2^A)$



$1 + 2 + 4 + 8 + 16 + \dots + 2^{A-1}$

P 2^kth symbol (hard)

Problem Description

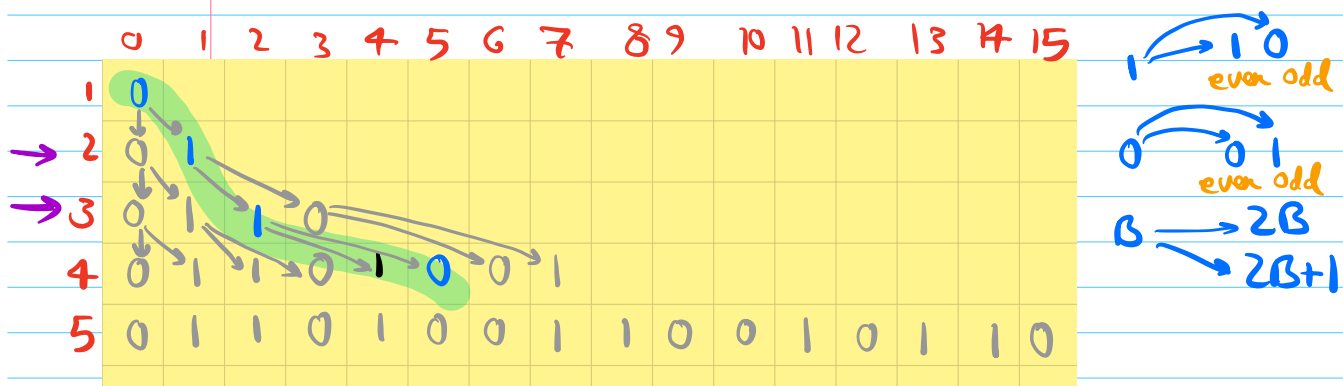
On the first row, we write a 0. Now in every subsequent row, we look at the previous row and replace each occurrence of 0 with 01, and each occurrence of 1 with 10.

Given row number A and index B, return the Bth indexed symbol in row A. (The values of B are 0-indexed.).

Problem Constraints

$$1 \leq A \leq 10^5$$

$$0 \leq B \leq \min(2^{A-1} - 1, 10^{18})$$



TC: $O(A)$

only the green path

SC: $O(A)$

```
int kthNum2(int A, long B){
    if(A==1) return 0;
    P = kthNum2(row-1, B/2);
    if(B%2==0){
        return P;
    }
    return 1-P;
}
```

main → return kthNum2(A, B)

A B
4 5
4-1=3 $\lfloor \frac{5}{2} \rfloor = 2$
row → row-1
col → $\lfloor \frac{col}{2} \rfloor$
parent info

if parent is (1)
if (B%2==0) → 1
else → 0
if parent is 0
if (B%2==0) → 0
else → 1

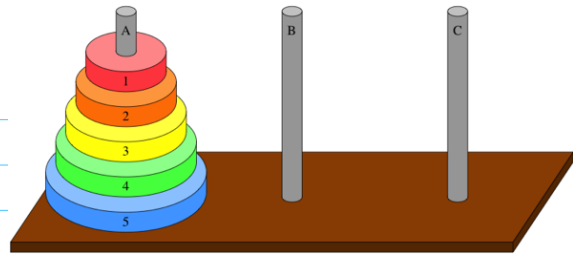
return (parent + B%2)/2

row col
A B
4 5
3 2
2 1
1 0

Towers of Hanoi

P1 three towers A, B, C

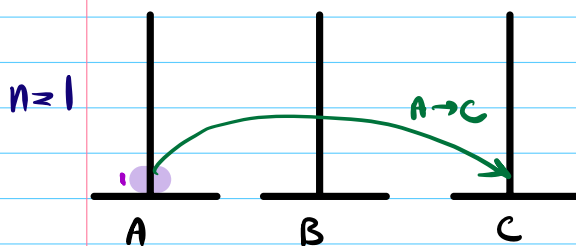
n discs on tower A;
different sizes, smaller on top of larger



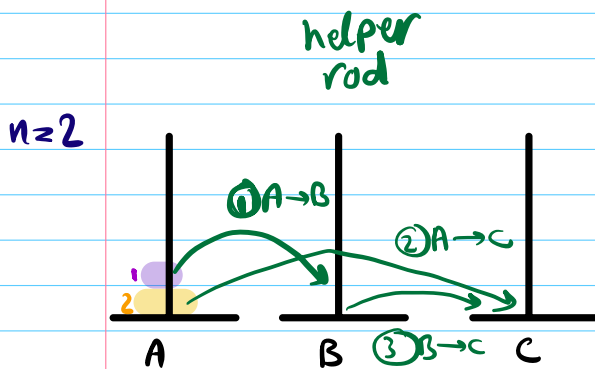
Goal: Move all discs from tower A to C.

Constraint: 1- we can move one disc at a time
2- Large disc cannot be placed on small disc

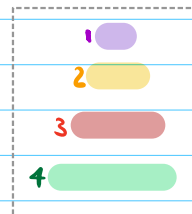
- Print the moves s.t. all the discs move from A to C in minimum steps.



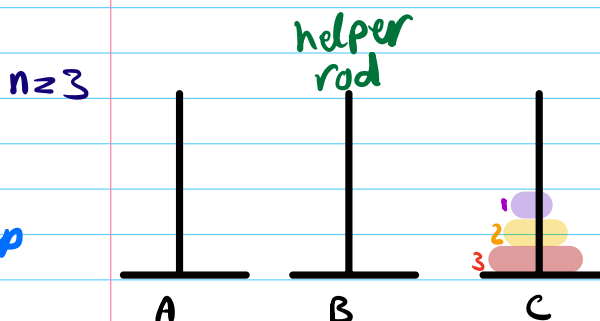
output:
A → C



output:



disc bag!

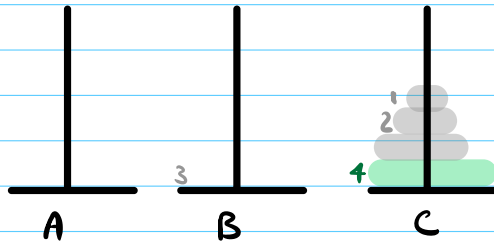


output:

7 step



$n=4$



char \rightarrow rod labels

int \leftarrow #disc
 Source \rightarrow dest \rightarrow helper

```
void tower(n, A, C, B) {
  if (n == 1) { print(1, A  $\rightarrow$  C); ret }
  tower(n-1, A, B, C)
  print(n, A  $\rightarrow$  C)
  tower(n-1, B, C, A)
}
```

$n=3$

$n-1$



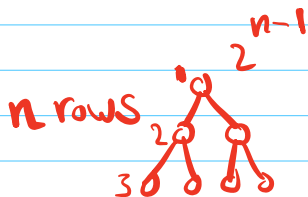
$$T(n) = T(n-1) + 1 + T(n-1)$$

$$= 2T(n-1) + 1$$

$$= 4T(n-2) + 1 + 1$$

$$= 8T(n-3) + 1 + 1 + 1$$

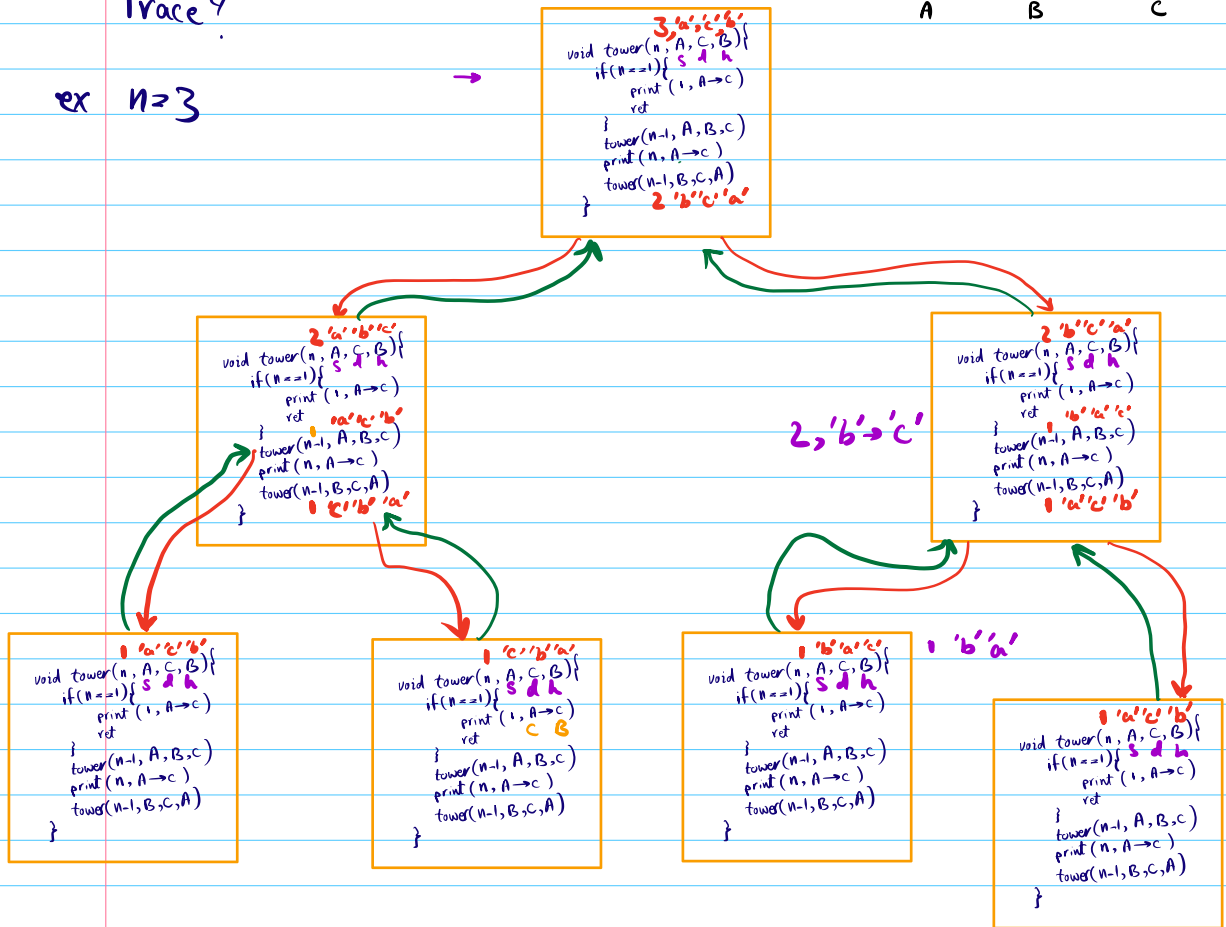
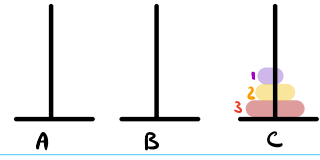
\vdots



$$TC: O(2^n)$$

Trace?

ex $n=3$



TC

optional

SC \rightarrow HW