

7:35AM

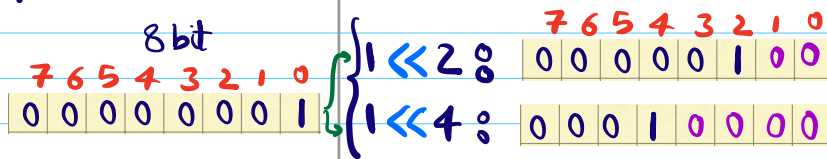
Topics - << applications in bit manipulation  
 - check bit  
 - unset bit  
 - count of set bit

- negative binary

## Bit Manipulation 2

Usage of <<

shifting the "1"



<< ①  
 ^ ②  
 & ③

- checking and/or setting specific bits

①  
 $N / (1 \ll i)$

$N = 45 \rightarrow$

	5	4	3	2	1	0
$1 \ll 2$	0	0	0	1	0	0
OR	1	0	1	1	0	1
	$\rightarrow 45$					

	5	4	3	2	1	0
$1 \ll 4$	0	1	0	0	0	0
OR	1	1	1	1	0	1
	$\rightarrow 61$					

$N / (1 \ll i) \begin{cases} N \text{ if } i\text{th bit is set in } N \\ N + (1 \ll i) \text{ if } i\text{th bit is unset} \end{cases}$

②  
 $N \wedge (1 \ll i)$

$N = 45$

	5	4	3	2	1	0
$1 \ll 2$	0	0	0	1	0	0
XOR	1	0	1	0	0	1
	$\neq 45$					

	5	4	3	2	1	0
$1 \ll 4$	0	1	0	0	0	0
XOR	1	1	1	0	1	1
	$\neq 45$					

$N \wedge (1 \ll i) \rightarrow$  Flips  $i$ th bit

③  
 $N \& (1 \ll i)$

$N = 45$

	5	4	3	2	1	0
$1 \ll 3$	0	0	1	0	0	0
AND	0	0	1	0	0	0
	$\neq 0$					

	5	4	3	2	1	0
$1 \ll 1$	0	0	0	0	1	0
AND	0	0	0	0	0	0
	$= 0$					

$N \& (1 \ll i) \begin{cases} 1 \ll i (\neq 0) \text{ if } i\text{th bit is set in } N \\ 0 \text{ if } i\text{th bit is unset in } N \end{cases}$

P1 Unset  $i$ th bit of a number if it is set  
otherwise no change

ex  $N = 45$   $\overset{5}{1} \overset{4}{0} \overset{3}{1} \overset{2}{1} \overset{1}{0} \overset{0}{1}$   
 $i = 2$   $101001$   
 $i = 4$   $101101$

1.1  
check  
 $i$ th Bit

```

if (check Bit(N, i)) N = N ^ (1 << i)
else N = N → not needed
return N

```

```

bool check Bit(N, i) {
    return N & (1 << i) != 0
}
or
return N / (1 << i) == N
return N ^ (1 << i) < N → why does this work?

```

if  $i$ th bit is set result will be smaller  
 if  $i$ th bit is unset result will be larger

true

(MSB) most Significant bit  
 8bit  
 Least Significant bit (LSB)  
 isOdd  
 bool func1(N) {  
 return N & 1  
 }  
 get Least Significant  
 return  $n \& 1 == 1$

P2 Count the number of set bits in N  $\rightarrow$  is positive

N = 101101 ans = 4



int  $\rightarrow$  32 bit

double 64 bit  
long 64 bit

long long 128 bit  
long double 128 bit

TC:  
 $O(\text{BITS})$

```
int countBit1(int n){
    ans = 0
    for(int i=0; i<32; i++){
        if(checkBit(n,i)) ans += 1
    }
    return ans;
}
```

Quiz  
TC:  $O(\log_2 n)$   
SC:  $O(1)$

```
int countBit2(int n){
    ans = 0
    while(n > 0){
        if(n & 1 == 1) ans += 1
        n = n >> 1
    }
    return ans;
}
```

$N \gg 1$

$\log_2 n$

	7	6	5	4	3	2	1	0
	0	0	0	0	1	0	1	0
ans	0	0	0	0	0	1	0	1
1	0	0	0	0	0	0	1	0
2	0	0	0	0	0	0	0	1
	0	0	0	0	0	0	0	0
	0	0	0	0	0	0	0	0
	0	0	0	0	0	0	0	0
	0	0	0	0	0	0	0	0
	0	0	0	0	0	0	0	0

## Negative binary numbers

intro:

Most significant bit

$$\begin{array}{r} \textcircled{7} \phantom{0000} \\ + \phantom{0000} \\ \hline 8 \phantom{0000} \end{array}$$


(a)  $01111111 = 2^7 - 1 = 128 - 1 = 127$

$\frac{10000000}{76543210} \rightarrow 2^7$        $n+1 = 2^i$   
 $n = 2^i - 1$

$2^0 + 2^1 + 2^2 + 2^3 + \dots + 2^6$

$\frac{a(r^i - 1)}{r - 1}$

31 ... 32 10



$2^{31} > 2^{30} + 2^{29} + \dots + 2^2 + 2^1 + 2^0$

$$(-45)_{10} = (?)_2$$

$$45 + 3 = 0 \quad \checkmark$$

B will be -45 if

- 1- Convert absolute number to binary ✓
- 2- invert all bit 1's complement ✓
- 3- add 1 to the inverted number

## 2's Complement

→ why like this?  
why this works?

8 bit  
1's Comp.

	7	6	5	4	3	2	1	0	
	0	0	1	0	1	1	0	1	(+5)
	1	1	0	1	0	0	1	0	

+ -

MSB | 1 → negative  
0 → positive

B ← 2's Camp.  
(45)

✓ 76 54 3 2 1 0

$$= 2^7 + 2^6 + 2^4 + 2^1 + 2^0 = 211$$

$$= -1 \times 2^7 + 2^6 + 2^4 + 2^1 + 2^0 = -128 +$$

$$(64 + 16 + 2 + 1)$$

$$83$$

$$-45$$

## Quiz

$$(-3)_{10} = (?)_2$$

$$(+3)_{10} \rightarrow (\quad)_2$$

$$+3 = (00000011)_2$$

111 1100

A horizontal number line is shown. It has a green tick mark on the left labeled '0' and a black tick mark on the right labeled '1'.

1 1 1 1 0 1  
7 6 5 4 3 2 1 0

$$-1x^7 + 2x^6 + 2x^5 + 2x^4 + 2x^3 + 2x^2 + 1$$

$$-128 + \underbrace{(64 + 32 + 16 + 8 + 4 + 1)}_{125} = -3$$

Quiz  $(45)_{10} \rightarrow (45)_2 + (-45)_{10} \rightarrow (-45)_2$

①

0	0	1	0	1	1	0	1
1	1	0	1	0	0	1	1
0	0	0	0	0	0	0	0

$\rightarrow 0$   $N+B=0$

$(45-12)_2 = (45)_2 + (-12)_2 = (33)_2$

$$\begin{array}{r} 0000 \ 1100 \\ 1111 \ 0011 \\ \hline 1111 \ 0100 \end{array}$$

0	0	1	0	1	1	0	1
1	1	1	1	0	1	0	0
0	0	1	0	0	0	0	1

5 4 3 2 1 0  
 $2^5 + 2^0 = 33$

Range of int

MSB 31 is reserved for sign

Max  $0 \overbrace{11111 \dots 1}^{31} = 2^{31} - 1 \approx +2 \times 10^9$

Min  $1 \overbrace{000 \dots 0}^{31} = 2^{31} \times -1 = -2^{31} \approx -2 \times 10^9$

Range of long

Max  $0 \overbrace{1111 \dots 1}^{63} = 2^{63} - 1 \approx +9 \times 10^{18}$

Min  $1 \overbrace{0000 \dots 0}^{63} = -2^{63} \approx -9 \times 10^{18}$

4

P3 Calculate sum of all elements in array  $a[]$

Wrong →

```
int n = a.length;
int sum = 0;
for (i = 0; i < n; i++) {
    sum += a[i];
}
return sum;
```

→ wrong

Constraints

$$1 \leq N \leq 10^5$$

$$1 \leq a[i] \leq 10^6$$

$$a = \{10^6, 10^6, \dots, 10^6\}$$

$\underbrace{\hspace{10em}}_{10^5}$

$$\text{sum} = 10^{11}$$

P4 For given two integers  $a$  and  $b$ , return  $a * b$ .

```
int f(int a, int b) {
    int ans = 0;
    ans = a * b;
    return ans;
}
```

$$a \leq 2 * 10^9$$

$$b \leq 2 * 10^9$$

$$\begin{aligned} &= 2 \times 10^9 \times 2 \times 10^9 \\ &= 4 \times 10^{18} \end{aligned}$$

```
long f(int a, int b) {
    long ans = 0;
    ans = a * b;
    return ans;
}
```

C, C++, Java

←  
101000

```
long f (int a, int b) {  
    long ans = 0  
    ans = (long)(a * b) X  
    ret ans  
}
```

---

```
long f (int a, int b) {  
    long ans = 0  
    ans = ((long)(a)) * b ✓  
    ret ans  
}
```