7:05AM

Interview tips:
failing an interview session is part of process

Prefix sum
↓
Dynamic Programming
last three weeks of Advanced DSA

Topics:
1. Sum Query
2. Prefix Sum example
break?
3. Prefix Sum
4. Eq. index
5. Count of even numbers

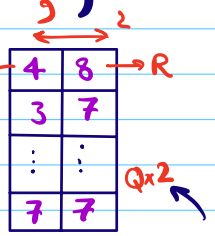
why multiple queries? why not one? ← [L, R]

P1 Given an array of N elements and Q queries, for each query calculate sum of all elements in range [L, R]

Constraints: $L \le R$ $\quad 1 \le N, Q \le 10^5$

ex a: {-3, 6, 2, 4, 5, 2, 8, -9, 3, 1}
indices: 0 1 2 3 4 5 6 7 8 9

index	L	R	ans
	4	8	9
	3	7	10
	1	3	12
	0	4	14
	7	7	-9



1000 sec
1. O(Q x n)
2. Q n
10^5 10^5
3. 100 ops
4. 1GHz
5. 1 sec

$10^5 \times 10^5 = 10^{10}$
10^2

10^9 ← 1GHz

$\frac{1 \times 10^{12}}{10^9} < 10^3$

10^{12} operation

$10^2 \rightarrow 2$
$10^9 \rightarrow 1$ sec
10^3

Quiz
TC: O(Q x N)
SC: O(1)

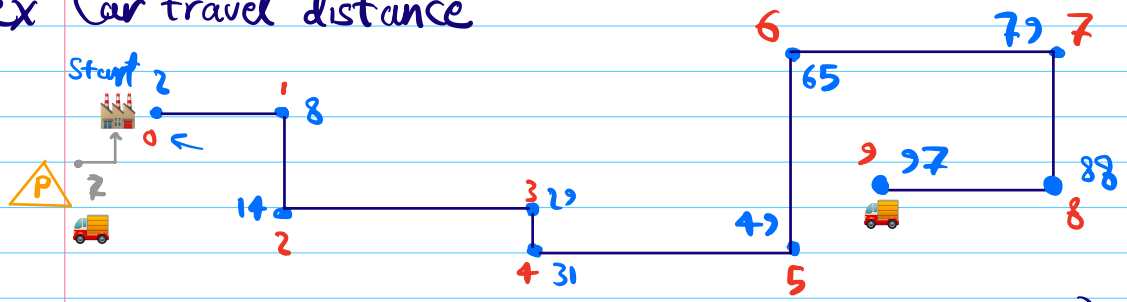
$10^5 \times 10^5 = 10^{10}$
10 - 100 → 10^2

Code:

```
void querySum(int a[], int q[][]){
    Q = q.Len // # of queries
    for(i=0; i<Q; i++){
        L = q[i][0]
        R = q[i][1]
        sum = 0
        for(j=L; j<=R; j++){
            sum += a[j]
        }
        print(sum)
    }
}
```

detour

ex Car travel distance



trip computer (odometer) = {2, 8, 14, 29, 31, 49, 65, 79, 88, 97} ← a
0 1 2 3 4 5 6 7 8 9

query : ① Last trip distanse : 97-88=9km O(1)

Quiz

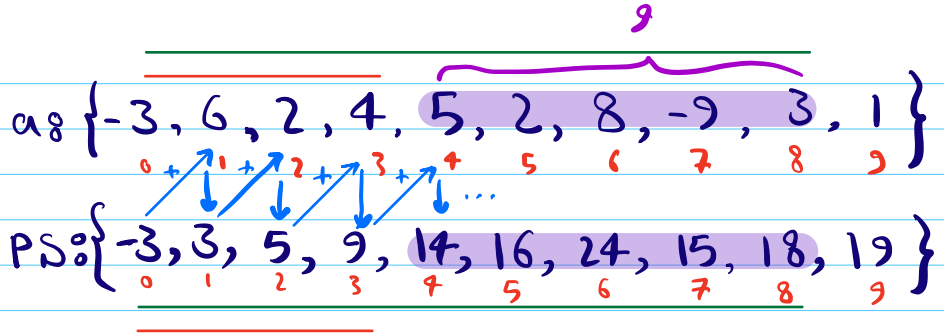
dest.
② 6th delivery, trip distance: 65-49=16km ✓ O(1)
5-to-6

dest.
③ 5th to 9th delivery, total distance: 97-49=48km ✓ O(1)

dest.
④ 1st to 5th delivery, total distance: 49-8=41km ✓ O(1)

- total distance from dest ith to dest jth ?
a[j]-a[i]

optimized ideas



— ⊖ — = ▭

Queries L	R	ans		
4	8	9	9 ✓	PS[8] - PS[3] = 9
3	7	10	10 ✓	PS[7] - PS[3-1] = 15-5 = 10
1	3	12	12 ✓	PS[3] - PS[1-1] = 9-(-3) = 12
0	4	14	14 ✓	~~PS[4] - PS[0-1] =~~ P[4] = 14
7	7	-9	-9 ✓	PS[7] - PS[7-1] = 15-24 = -9

PS[-1] → exception ✗

⊛ sum(L,R) { PS[R] - PS[L-1] if L>0 ; PS[R] if L==0 }

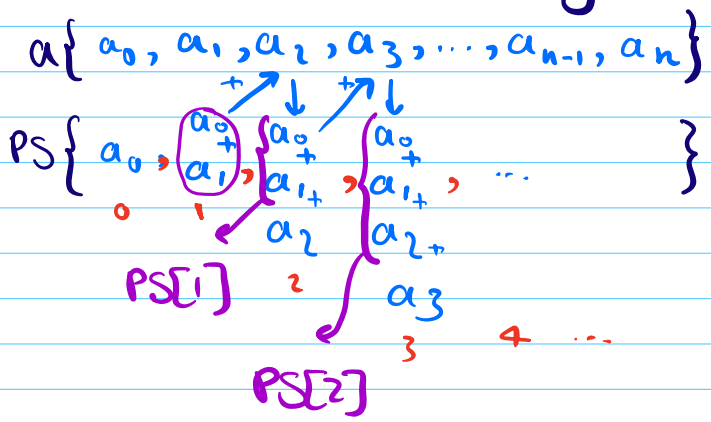
O(1)

TC: O(1)

one more time :)

How to construct prefix sum array?

Quiz ✓



$PS[i] = PS[i-1] + a[i]$

```
PS[0]=a[0]
for(i=1;i<n;i++){
    PS[i]=PS[i-1]+a[i]
}
```

TC: O(N)

$Q \times O(1) \approx O(Q)$

total $O(N+Q)$ $10^5+10^5 = 2\times10^5$

$O(Q \times N)$ 10^{10}

Optimized code for P1

SC: O(n)

Quiz

TC: O(n) + Q × O(1)

O(n+Q)

Can we optimize SC?

if we can change the original array you may use input a[] as PS[]

```
void querySum2(int a[], int q[][]){
    n = a.Len
    PS = new int[n]
    PS[0] = a[0]
    for(i=1; i<n; i++){
        PS[i] = PS[i-1] + a[i]
    }
    Q = q.Len
    for(i=0; i<Q; i++){
        L = q[i][0]
        R = q[i][1]
        if(L==0) print(PS[R])
        else print(PS[R] - PS[L-1])
    }
}
```

Calculate PS: O(n)

be aware of overflow

Q

formula (*): O(1)

P2 Equilibrium Index:

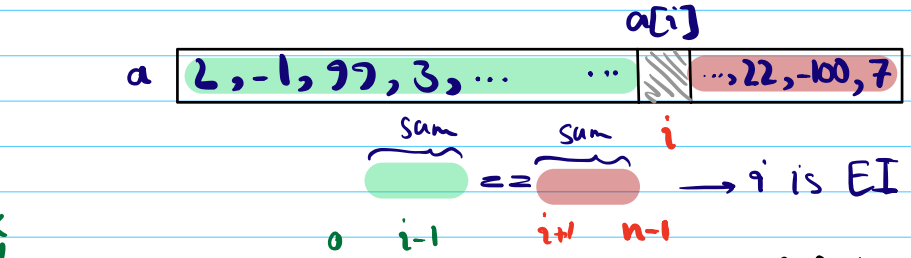
ret int

Given an array of N elements, Count the number of equilibrium indexes.

↳ what is?

equilibrium index(EI):

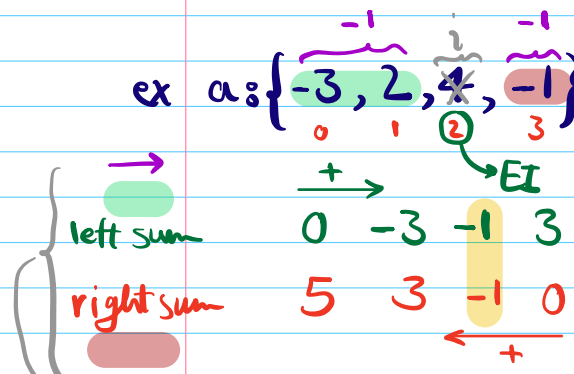
⚠ Sum of all elements on left of ith index == Sum of all elements on right of ith index



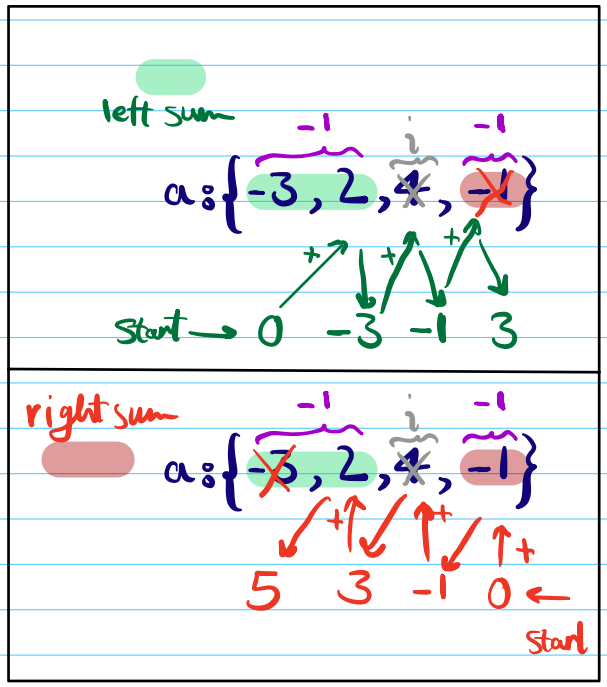
edge case
↳ Abhishek mentioned
→ when [ ] is o len
or [ ] is o len

how to get left/right sum

ex a: {-3, 2, 4, -1} ans = 1
0 1 2 3
EI



(*) we do not use left sum & right in optimized solution; just to understand EI



(*) Left sum & right sum are not exactly equal to prefix sum & postfix sum

heads up: EI is an index, not the element itself.

Quiz

a: {-7, 1, 5, 2, -4, 3, 0}
   0 1 2 3 4 5 6

ans = 2

Quiz

0

{ prefix sum
{ postfix sum

Code:

```
int CountEI(int a[]){
    n = a.Len        ans = 0
    PS = new int[n]
    PS[0] = a[0]
    for(i=1; i<n; i++){
        PS[i] = PS[i-1] + a[i]
    }                                   } Prefix sum

    for(i=0; i<n; i++){                 [i+1, n-1]
      if(PS[i-1] == P[n-1] - P[i]){       P[n-1] - P[(i+1)-1]
        ans += 1
      }
    }
    ret ans
}
```

from [0, i-1]

Quiz

SC: O(n)

TC: O(n+n)
 = O(n)

bug

HW

P3 Given an array of N elements and Q queries, for each query [L, R], find count of even numbers in a given range.

keep the original array intact.

ex

	0	1	2	3	4	5	6	7	8	9
as	2	4	3	7	9	8	6	5	4	9
eo	1	1	0	0	0	1	1	0	1	0

	L	R	ans
x	4	8	3
x	3	9	3
x	0	4	2

Code

n+m+q

TC: O(n+q)

SC: O(2n) = O(n)

```
void CountEven(int a[], int q[][]){
    e = new int[a.Len]        n = a.Len
    for(i=0; i<n; i++){
        if(a[i]%2==0) e[i]=1
        else e[i]=0
    }
    PS = new int[n]
    PS[0] = a[0]
    for(i=1; i<n; i++){
        PS[i] = PS[i-1] + e[i]
    }
    Q = q.Len
    for(i=0; i<Q; i++){
        L = q[i][0]
        R = q[i][1]
        if(L==0) print(PS[R])
        else print(PS[R]-PS[L-1])
    }
}
```

Calculate PS: O(n)

be aware of overflow

Q: formula (*) O(1)