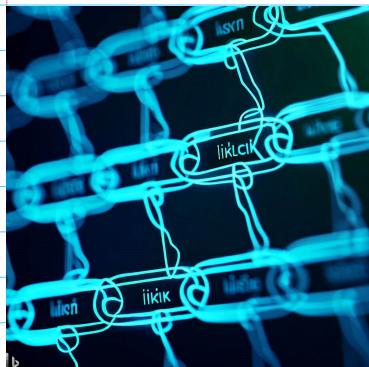


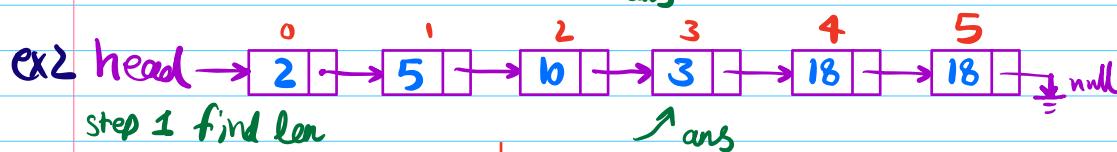
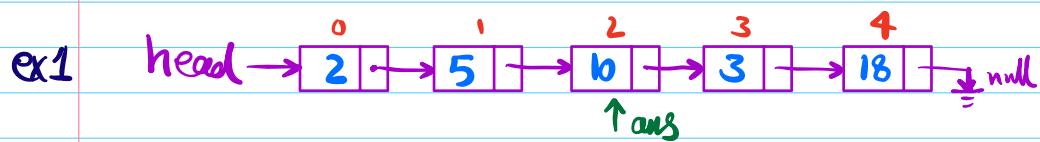
Generated by Bing AI image creator
<https://www.bing.com/images/create>
Query: "linked list code algorithm visualization"



Topics

- ✓ 1 - mid element in LL
- ✓ 2 - merge two sorted LL
- ✗ 3 - merge sort LL
- ✗ 4 - cycle?
- ✗ 5 - beginning of cycle

P1 Find middle element in a Linked List. ret node pointer



idea1

len = 0

cur = head

while (cur != null){

 len++;

 cur = cur.next;

}

$O(n)$



2 min

Step 2 go $\frac{len}{2}$ from head

cur = head

for ($i = 1$; $i \leq \frac{len}{2}$; $i++$)

 cur = cur.next

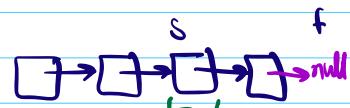
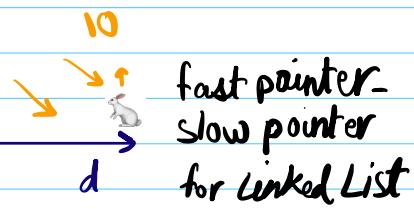
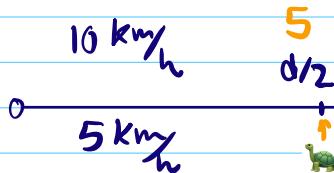
}

ret cur

$O(\frac{n}{2}) \sim O(n)$

TC: $O(n)$ SC: $O(1)$

idea2



slow fast

Code

```

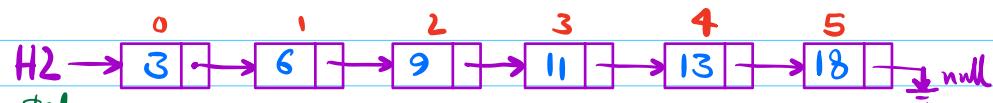
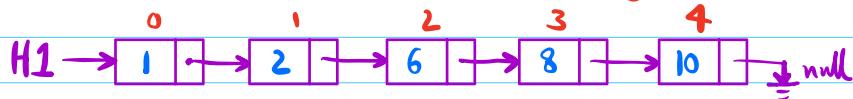
 $\Delta$  {
    s = head      f = head
    while (f != null && f.next != null) {
        s = s.next
        f = f.next.next
    }
    ret s
}
    
```

? Also what about slow

TC: $O(\frac{n}{2}) \approx O(n)$ SC: $O(1)$

P2 Merge two sorted linked List into a single sorted list.

Constraint: Do not create new list, only update pointers



3 min

Sorted

Head 1 → 2 → 3 → 6 → 6 → 8 → 9 → 10 → 11 → 13 → 18

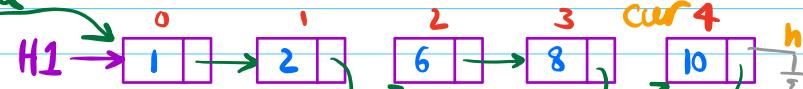
Quiz 🍏

TC: $O(m+n)$

SC: $O(1)$

vs Arr

SC: $O(m+n)$



sorted

Head



$h_1 = H_1$

$h_2 = H_2$

Δ if ($h_1 == null$) ret h_2
if ($h_2 == null$) ret h_1

set
head

{ if ($h_1.data < h_2.data$) {
 head = h_1
 $h_1 = h_1.next$
} else {
 head = h_2
 $h_2 = h_2.next$
}}

main logic

cur = head
while ($H_2 != null$ & $H_1 != null$) {
 if ($H_1.data \leq H_2.data$) {
 cur.next = H_1
 $H_1 = H_1.next$
 } else {
 cur.next = H_2
 $H_2 = H_2.next$
 }
 cur = cur.next
}

if ($H_1 != null$) cur.next = h_1
if ($H_2 != null$) cur.next = H_2

which
list
has more
items

var head is for merged list

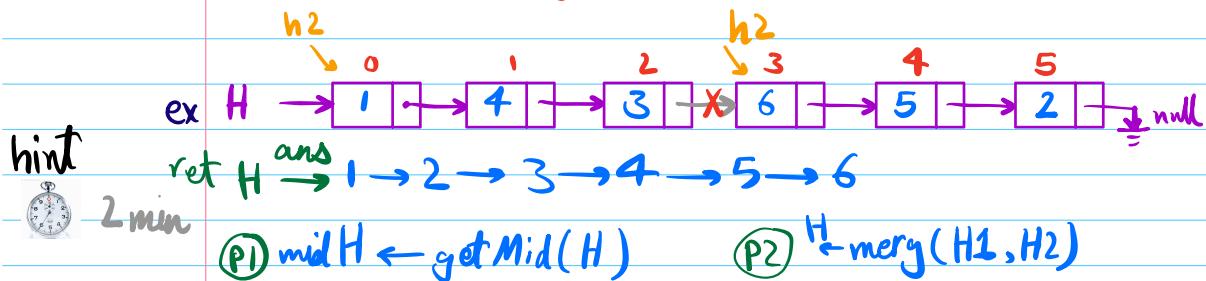
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Query: "merge sort array linked list algorithm"



P3 Sort a given linked list using merge sort.

Constraint: Do not create new list, only update pointers



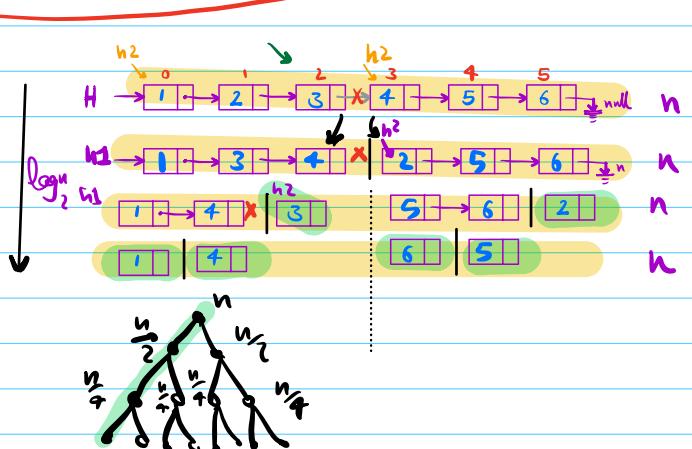
Quiz
TC: $O(n \log n)$
SC: $O(\log n)$

VS
Arr
TC: $O(n \log n)$
SC: $O(n)$

node $\text{mergesort}(\text{head})\{$
if ($\text{head} == \text{null}$ || $\text{head.next} == \text{null}$) ret head
 $\underline{h2 = head}$ $\underline{\text{len} = 0}$ $\underline{\text{len} = 1}$
 $\underline{\text{node mid} = \text{getMid}(h)}$ $O(n)$
 $\underline{h1 = h; h2 = mid.next}$ Yes needed
A $\underline{\text{mid.next} = \text{null}}$ → split (do we need this?)
cut it

$h1 = \text{mergesort}(h1);$
 $h2 = \text{mergesort}(h2);$
ret $\text{merg}(h1, h2)$

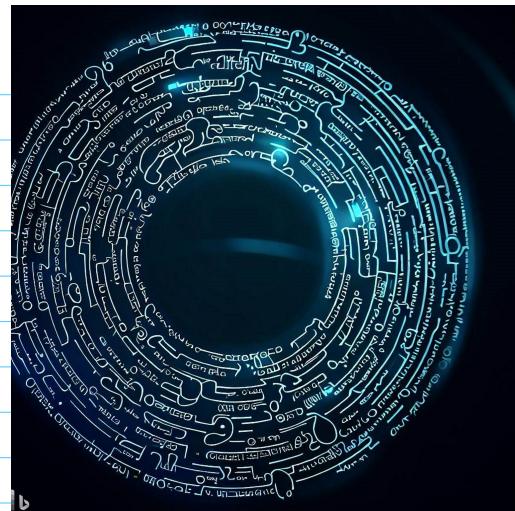
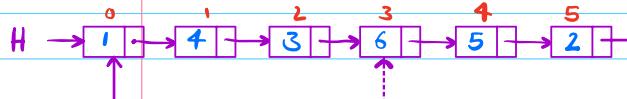
$O(1)$
 $O(1)$
 $O(1)$
 $O(1)$
Stack



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<https://www.bing.com/images/create>

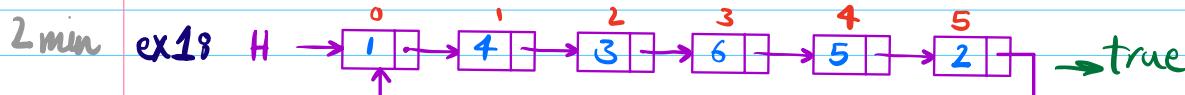
Query: "circular linked list code algorithm visualization"

circle linked list

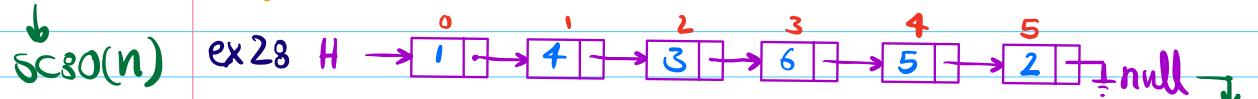


true/false

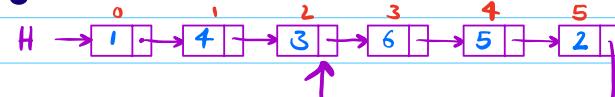
P4 check if a given linked list has a cycle.



hashset<Node>



ex3



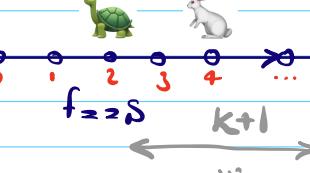
false

(I)

statement
 if cycle ⇒
 $f == s$

f and s
 will meet
 on a node/why?

math Proof
 beyond scope



$3 \rightarrow 4 \rightarrow 5 \rightarrow 6$

$\rightarrow 7 \rightarrow 8 \rightarrow 9 \rightarrow 10 \rightarrow 11$

n size of cycle

$k > 1$

max after

n

→

Code

```

 $s = \text{head}; f = \text{head}$ 
 $\Delta \text{ while } (f \neq \text{null} \& \& f.\text{next} \neq \text{null})\{$ 
     $s = s.\text{next}$ 
     $f = f.\text{next}.\text{next}$ 
     $\text{if } (s == f) \text{ ret true}$ 
}
 $\text{ret false}$ 

```

pointer compare

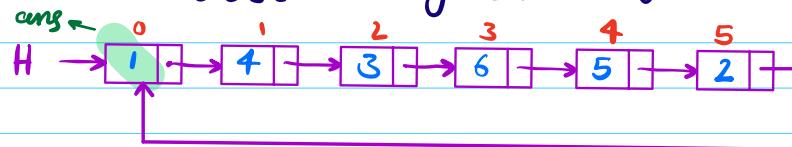
TC₈

O(n)

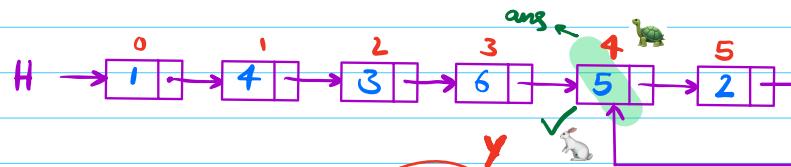
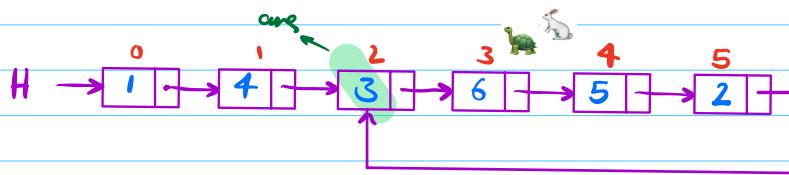
SC80(1)

* domain knowledge; you are not expected to infer all in interview

④ P5 Given a linked list with cycle, find the start of cycle

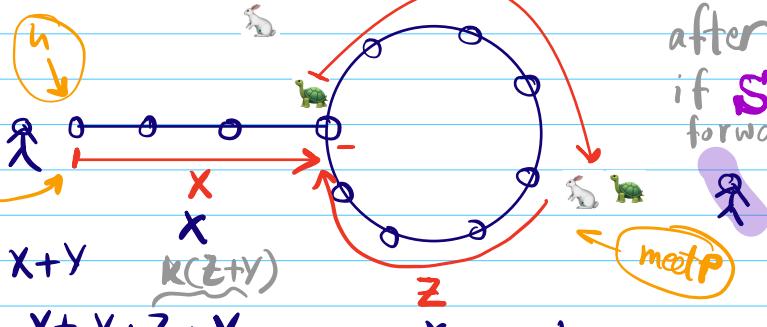


SC80(1)



2 min

(begin)



after T unit of time
if S went 3 steps (dot)
forward, how many steps
did f go?

$$2S = 1f$$

meet {

- $X+Y$
- $X+Y+Z+Y$
- $X+Y+Z+Y$

length of cycle

$$(X+Y)2 = X+Y+Z+Y \text{ at meeting point}$$

$$\cancel{X+X+Y+Y} = \cancel{X+X+Y+Y}$$

$$X = Z$$

$$X = Z + \frac{(k-1)(Z+V)}{C_1}$$

code

```
S = head; f = head
while (f != null && f.next != null) {
    S = S.next
    f = f.next.next
    if (S == f) break ✓ they met!
```

} meetP = S | x point forward to meet
begin = h the start of cycle

```
while (meetP != begin) {
    meetP = meetP.next
    begin = begin.next
}
```

ret meetP → definitely start of cycle!

$O(n)$ gTC

$O(1)$ gSC