

→ we will be solving problems. (puzzle + math + coding)

→ Relevant to interviews

→ Do not make notes in class.

→ Revise every class day weekdays.

→ Every class → HW.

→ LeetCode

200
1500
20

20
↓
understand
while solving
problem

200
actually

1500
think

Most Recommended

99%

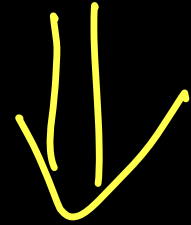
1% → fast & fun

C++ and Java

they are fast

Show their code

Online →
 20% → Don't care
 30% → can't do
 50% → leg



Python and JS

app driver

Do OSA with JS

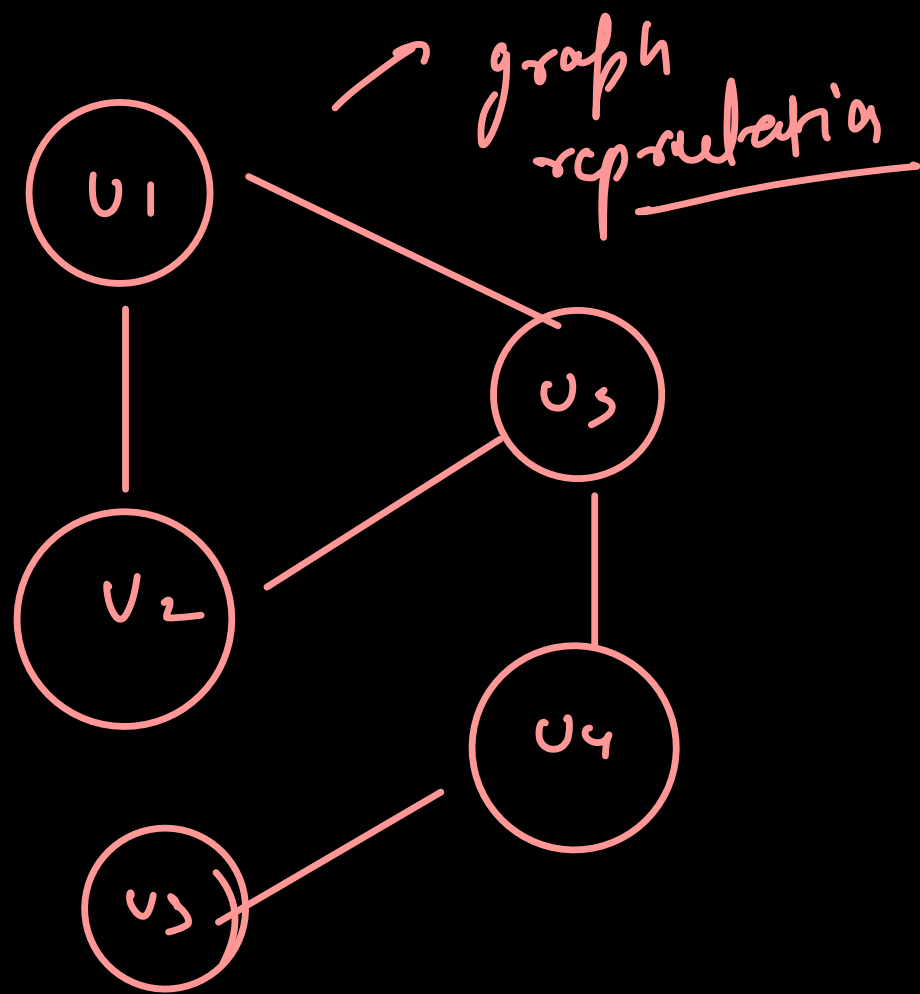
→ 7-8 hr for tutorial or Java

(By next week) Made this watch that

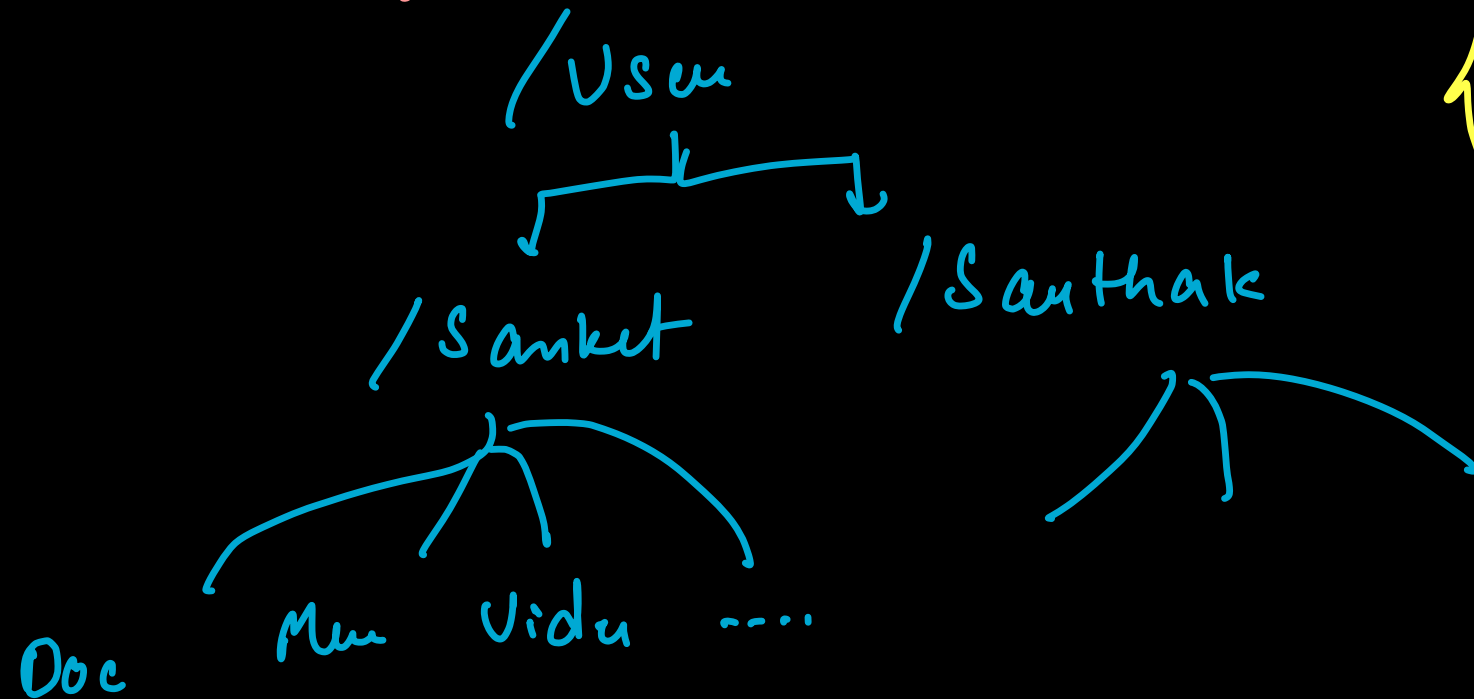
(DSA)

Data Structures & Algorithms

↳ Data Structures →



folder



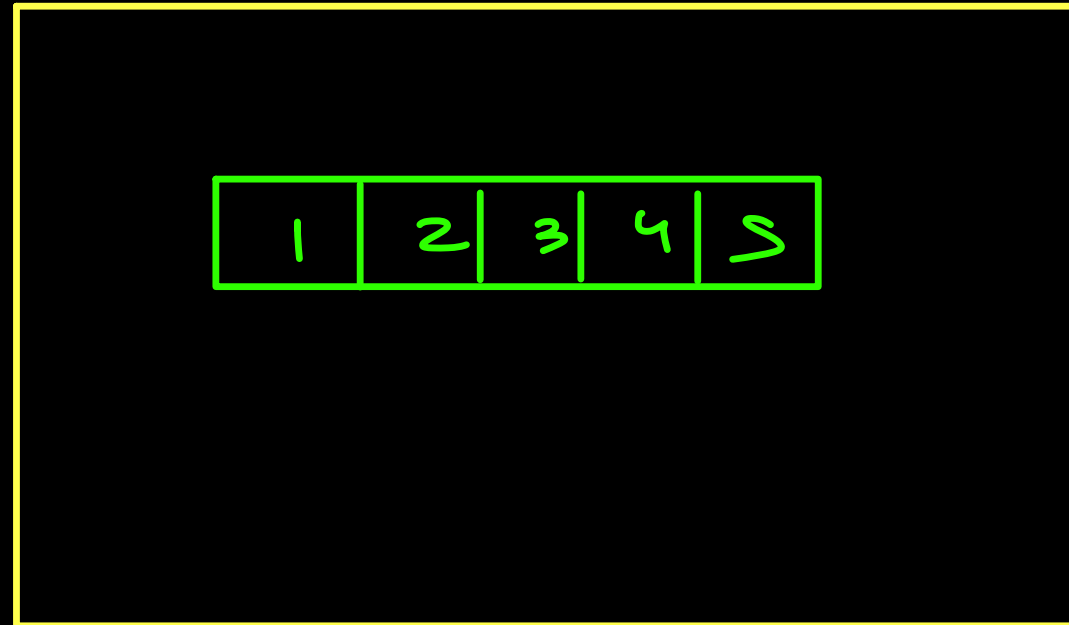
hierarchy

Data structures are storage models which helps us to store data in efficient manner for diff use cases.

↳ arrays

[1, 2, 3, 4, 5]

→ linear time



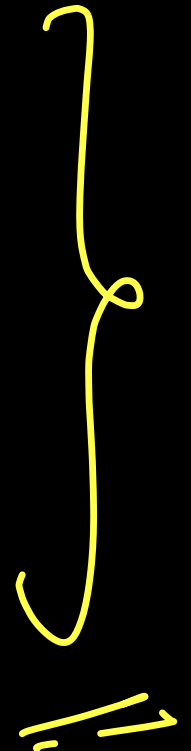
arrays
linked list

stacks

queue

recursion

etc



Algorithms → Defined steps that helps us to
achieve a task is called algorithm.

Problem-1

→ 2 sum

nums → $\begin{matrix} 0 & 1 & 2 & 3 \\ [2, & 7, & 11, & 15] \end{matrix}$ $\xrightarrow{\text{idx}}$

from this array find a pair of elements
that sums up to

target.

(Do not use same
element twice)

return array with
indices of pair.

target → 9

ans → [0, 1]

$2 \leq \text{length of array} \leq 10^4$

$-10^9 \leq \text{element of array} \leq 10^9$

$-10^9 \leq \text{target} \leq 10^9$

TLE → time limit exceeded

1 sec to run

In one sec approx

10^8 instructions

we can exceed.

Pair → target
 $[a_0, a_1, a_2, a_3, \dots, a_{n-1}]$

worst solution / brute force

why not prepare all possible pairs & check them

Sum

$\begin{matrix} & i & & j \\ & \downarrow & & \downarrow \\ 0 & 1 & 2 & 3 \end{matrix}$ → index
 $[2, 7, 11, 15]$

$j \rightarrow$ 2nd elem of pair

$i \rightarrow$ first value of pair

(2, 7)

(7, 11)

(11, 15)

(2

(2, 11)

(7, 15)

(2, 15)

we can fix the first element of the pair &
then combine it with the remaining elements. —
 $n \leftarrow \text{length of array}$

```
for (i = 0; i ≤ n - 2; i++) {  
    for (j = i + 1; j < n; j++) {  
        print(arr[i], arr[j])  
    }  
}
```

print all
pairs of elements

n-2 times

1st

outer

for (i=0; i ≤ n-2 ; i++) {

for (j=i+1; j < n ; j++) {

if (num[i] + num[j] == target) {

return [i, j];

}

}

in every internal loop
iteration we execute
3-4 instructions

}

Outer

i=0

→

internal loop

→

n-1 times

i=1

→

internal loop

→

n-2 times

i=2

→

" "

→

n-3 times

⋮

↓
Total iterations

$$(n-1) + (n-2) + (n-3) + \dots + 2 + 1$$

Sum of first $(n-1)$ natural no \rightarrow $\frac{(n)(n-1)}{2}$

\swarrow
approx

Total instructions is approx $\frac{2 \times (n)(n-1)}{2}$

$$\rightarrow \underline{2(n)(n-1)}$$

$$\rightarrow \underline{2 \times 10^4 (10^4 - 1)} \approx \underline{\underline{10^8}}$$

Note \rightarrow Never randomly solve array & strings problems.

DSA → notes

↓

techniques

↓

Examples

→ Trigonometry

Exercises

10¹²

5-7

> 100-200

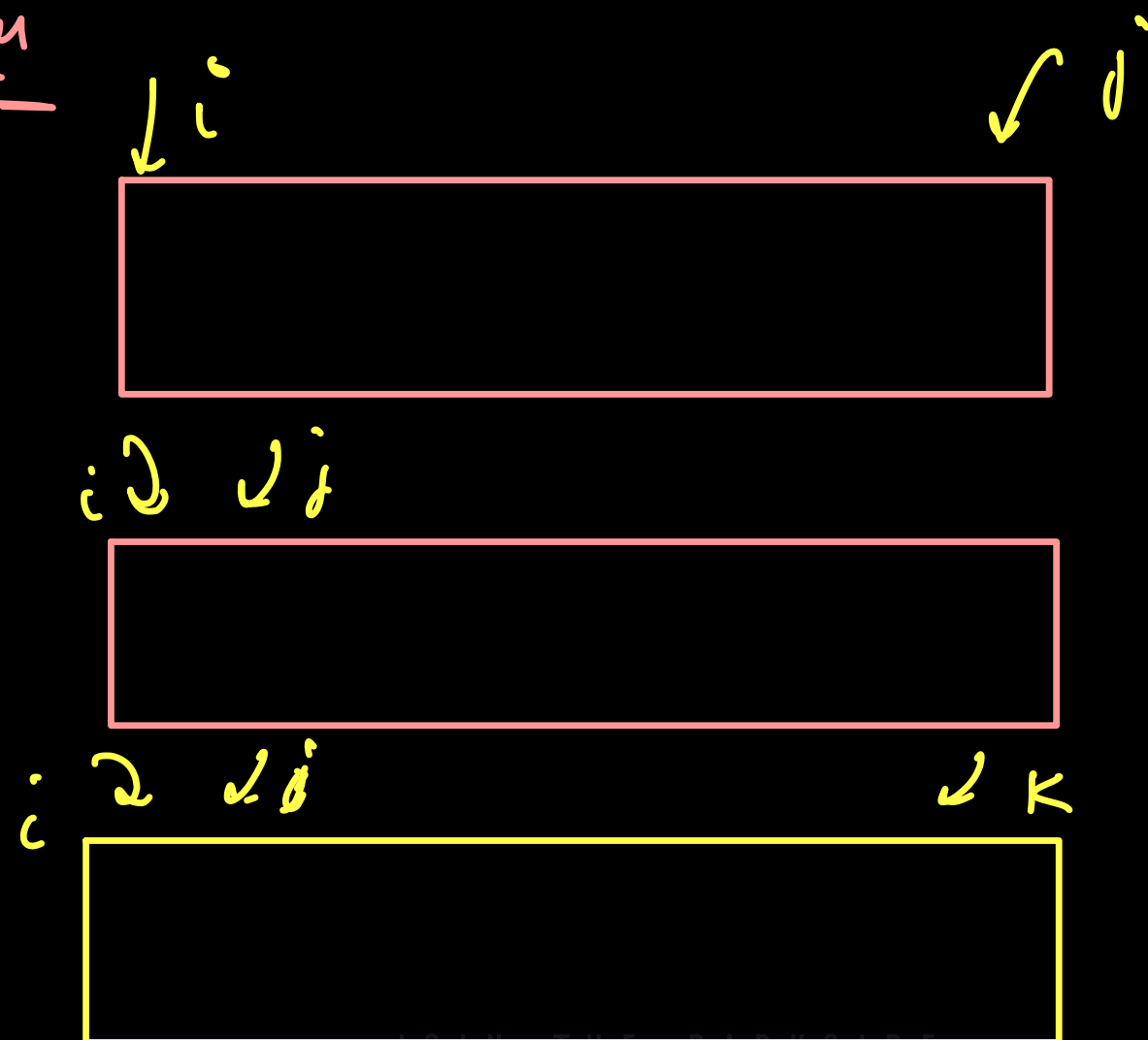
new

|||||

Two pointers (technique)

- prepare → 2-3 variable → pointers
- you place these variable on different parts of the input.
- now then

while moving dou
ng to discard
some part of
data set



2 → Problem - 344 (Reverse a string)

$\rightarrow S = ["h", "e", "l", "l", "o"] \quad ["o", "l", "l", "e"]$
 \downarrow
array
 $\rightarrow ["o", "l", "l", "e", "h"]$

Backflow → because we want data of array in reverse order, we can read the array in reverse order. And, one by one store the data in a new array.

```
n = s.length;  
ans = [];
```

```
for (i = n - 1; i >= 0; i--) {  
    ans.push(s[i]);
```

```
}
```

```
return ans;
```

Can we improve ??

→ we are taking an extra space in memory because
of the ^{ans.} array. We can avoid it.

$i=0$
 $j=n-1$

$[h, e, l, l, o]$

$[0, i] \rightarrow \text{left region}$
 $[j, n-1] \rightarrow \text{right region}$

already
sorted

$[i+1, j-1] \rightarrow \text{this region needs to be sorted}$

→ [t , e , k , ^ , a , s] _{0 1 2 3 4 5} n=6

[^]_{temp}

i = 0

j = n-1

while (i <= j) {

→ temp = s[i]

s[i] = s[j]

s[j] = temp

i++;

j--;

}

} Swapping

$$\left. \begin{array}{l} a = 10 \\ b = 20 \end{array} \right\}$$

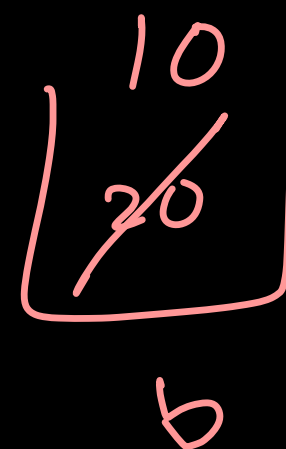
bitwise

arithmetic

$$a = a + b$$

$$b = a - b$$

$$a = a - b$$



Problem - 169

can we optimise ??

target \Rightarrow 15

2^i i
[1, 3, 5, 7, 11, 16, 22]

$i = 0$
 $j = n - 1$

in each iteration you are
eliminating one element

7

total
iter $\rightarrow \underline{\underline{n}}$

$$\underline{\underline{4 \times n}}$$

$$\underline{\underline{4 \times n}}$$

$$\underline{\underline{n \leq 3 \times 10^4}}$$

$$4 \times 3 \times 10^4$$

$$\rightarrow 12 \times 10^4 \approx \underline{\underline{10^5}}$$

while ($i < j$) {

if ($\text{nums}[i] + \text{nums}[j] == \text{target}$)
return $[i+1, j+1]$

else if ($\text{nums}[i] + \text{nums}[j] > \text{target}$)
j--;

else
i++;

}