001.	A compiler would tie the parser and the semantic action program fragments together,				
	•	ucing module.	D	huo	
	A C	one three	B D	two more than one	
002	_	polish places the operator at the ri	_		Α
002.		Postfix	grit ei B	Prefix	
	C	Infix	D	Polish	
003	_	ementation of syntax-directed translato	_		С
000.	A	input	В	output	
	C	input-output	D	parse table	
004.		directed translation allows subrou		•	Α
-		productions of a context free grammar.			
	Α .	syntax	В	semantic	
	С	syntax and semantic	D	structure	
005.	Synt	ax directed translation subroutines gen	erate	code.	Α
	Α	intermediate	В	source	
	С	object	D	error	
006.	A sy	ntax directed translation scheme is me	rely a	grammar	C
	Α	regular	В	context-sensitive	
	С	context-free	D	single	
007.	The	action is enclosed in br			В
	Α	syntax	В	semantic	
	С	both	D .	error	_
008.		•	o ado	dresses for theand one for the	Α
	resul		_		
	A	operand	В	operator	
000	C	result	D	statement	Ь
009.	_	statement is an abstract for			В
	A C	2-address Intermediatecode	B D	3-address address	
010	_	properties of an entity are called as			В
	_	values		attributes	Ь
	C	numbers	D	digits	
011	_	valuate the expression, a st		_	Α
0	A	Postfix	В	Prefix	
	C	Both	D	Polish	
012.	_	general strategy is to scan the postfix of	_		Α
	Α	left-right	В	right-left	
	С	from middle	D	from right	
013.	If the	attributes of the parent depend on the	attrib	•	D
		d as attributes.		•	
	Α	made	В	discovered	
	С	new	D	inherited	
014.		is a tree in which each leaf repres	sents	an operand and each interior node an	C
	oper	ator.			
	Α	Parse Tree	В	Semantic Tree	
	С	Syntax Tree	D	Structured Tree	
015.		is associating the attributes with the	Ξ		В
	A	rotation	В	translation	
	С	transformation	D	evolving	_
016.		address code for array reference we as		• • • • • • • • • • • • • • • • • • •	Α
	_	cripts range from 1 to some limit known	_		
	A	compile	В	run	
	С	execution	D	process	

017.		refers to the expression or expressions	s in th	ne form of three address codes.	С
	Α	value	В	place	
	С	code	D	number	
018.	Whic	h is not the way of implement the 3-ad	dress	statement.	D
		Quadruples	В	Triples	
		Indirect Triples	D	Parse Tree	
019.		record structure has 4 fields.	_		A
		Quadruples	В	Triples	
000		Indirect Triples	D	Parse Tree	
020.		nthesized numbers are used to represe			Α
	_	pointer	B D	stack	
024		queue Triples are listing pointers to tri	_	value	В
UZ 1.		Triples are listing pointers to tri selves	pies,	Tattler than ilsting the triples	D
	A	Direct	В	Indirect	
	C	Multiple	D	New	
022	_	eclared name and type incompatibilities	_		В
	A	Syntactic errors	В	Semantic errors	_
	С	Lexical Phase errors	D	Reporting errors	
023.	Minir	num distance correction is		1 3	Α
		Syntactic Phase errors	В	Semantic errors	
	С	Lexical Phase errors	D	Reporting errors	
024.	A co	mpiler needs to collect and use informa	ation a	about the names appearing in the	Α
	sour	ce program. This information is entered	l into	a data structure called a	
		·			
	Α	Symbol Table	В	Lexical analysis	
	С	Syntactic analysis	D	Records	
025.		riples uses only 3			A
		fields	В	operator	
	С	operand	D ,	instruction	_
026.		is used in the several stage			В
		Syntax Table		•	
027	C	Queue	D boo	Stack	٨
027.	IIIIOII	nation about the name is entered into	lile Sy	mbol table during and	A
		lexical and syntactic analysis.	R	lexical and code generation.	
	C	lexical and error handler.	D	lexical and code optimization.	
028.	_	entry in the symbol table is a pair of the			Α
0_0 .	A	Name and information.	В	Name and function.	-
	С	Name and Data.	D	Name and procedures.	
029.	The	accurate term for Code Optimization is			В
	Α	Intermediate Code	В	Code Improvement	
	С	Latter Optimization	D	Local Optimization.	
030.		quality of the object program is general	ly me		С
	Α	Cost	В	Time	
	С	Size or Its running time	D	Code Optimization.	
031.	Wha	t is the length of identifier for DIMPLE?			В
	Α	5	В	6	
	С	4	D	3	
032.		er discards input symbol until a		token is encountered.	В
	A	synchronizing	В	Synchronizing	
000	C	Group	D	none	_
033.		message should not be redundant in _		Companie anno	D
	А	Syntactic Phase errors	В	Semantic errors	

	C Lexical Phase errors	D	Reporting errors	
034.	Two types of data areas are			В
	Two types of data areas are A Common and stack	В	Common and equivalence.	
	C Register and stack	D	Code and equivalence.	
035.	The simplest way to implement a symbol to	able i	s as a of records,	Α
	one record per name.			
	A Linear array	В	Multidimensional array	
	C Rectangular array	D	Jagged Array.	
036.	performed within a straight	line a	and no jump.	Α
	A Local Optimization.	В	Code optimization	
	C Latter Optimization	D	Code optimization Loop optimization.	
037.	From anyone in the loop to any other, there	e is a	path of length one or more	D
	is			
	A Weakly Connected	В	B Unique Entity	
	C Multi Connected	D	Strongly Connected.	
038.	The term constant folding is used for the _			C
	A Local Optimization.	В		
	C Latter Optimization	D	Loop optimization.	
039.	The code optimization techniques consist of	of det		В
	these patterns.			
	A Errors and replacing	В	Patterns and replacing	
	C Errors and editing	D	Patterns and editing.	
040.	may be local or global.			Α
	A Code Optimization	В	Variable	
	C Sub expression	D	Patterns.	
041.	90-10 rule states that of the ti	me is	s spent in of the code.	Α
	A 90%,20%	В	80%,10%	
	C 90%,10%	D	90%,90%	
042.	The important sources of optimization are	the ic	lentification of	В
	common			
	A Regular expression	В	Sub expression	
	C expression	D	Time	
043.		symb	ol.	C
	A Nodes	В	Leaves	
	C Interior Nodes	D	Roots.	
044.	Computed results can be left in		as long as possible.	Α
	A Registers	В	Triples	
	C Indirect Triples	D		
045.	Constructing a DAG from		is a good way of determining common	С
	sub expression.			
	A 2 address statement	В	4 address statement	
	C 3 address statement	D		
046.	If some sequences of statements form arithmetic	nmet	ic progressions, we say such identifiers	В
	as			
	A Reduction	В	Induction Variables	
	C Code motion	D	Inner Loops	
047.	The replacement of an expensive operation	n by	a cheaper one is called	Α
	in strength			
	A Reduction	В	B Induction Variables	
	C Code motion	D	Inner Loops	
048.	Full form of DAG			С
	A Dynamic Acyclic Graph	В	,	
_	C Directed Acyclic Graph	D	3 - 7	_
049.	A useful data structure for automatically ar	ıalyzi	ng basic block is a	C

	Α	Dynamic Acyclic Graph	В	Data Acyclic Graph	
	С	Directed Acyclic Graph	D	Detecting Acyclic Graph	
050.	Load	Is and stores are reduced in			В
	Α	Global optimization	В	Peephole optimization	
		Latter optimization	D	Local optimization.	
051.		ch of the following comment about peep			Α
	Α	It is applied to a small part of the code	В	It can be used to optimize	
	_	and applied repeatedly	_	intermediate code	
	С	It can be applied to a portion of the	D	It is applied in the symbol table to	
		code that is not contiguous		optimize the memory requirements.	_
052.		ple jumps are reduced accordingly to _		·	С
		Local optimization.	В	Code optimization	
		Peephole optimization	D .	Latter optimization.	_
053.	_	lly the register descriptor shows that all			В
	A	Full	В	Empty	
054	C	Half-filled	D	None.	_
U 54 .		eep track of the locationis			В
		5 5	В	Address descriptor	
0EE		Allocation descriptor	D	register.	С
055.		invoke a function GETREG	**	Codo motion	C
		Code optimization	B D	Code motion Intermediate code.	
056		the code generation algorithm DAG representation of a Quadruples is			С
050.	A	Nodes	а В	 Leaves	C
	Ĉ	Tree	D	Pattern.	
057	_	ompiler optimization, operator strength r	_		В
037.		ace slow math operations with faster op			_
		acements is an illustration of operator st			
	A	Replace P + P by 2 * P or Replace 3			
	, ,	+ 4 by 7.		110pid001 02 by 1 110	
	С		D	Replace (P << 4) P by P * 15	
058.		, the bodies of the two loops are			D
		ided that they do not make any reference			
	A	Loop unrolling	В	Strength reduction	
	C	Loop concatenation	D	Loop jamming	
059.	_	graph that shows basic blocks and thei		, ,	С
	Α	DAG	В	Control graph	
	С	Flow graph	D	Hamiltonian graph	
060.	Whic	ch of the following class of statement us	sually	.	Α
		oiled?	·		
	Α	Declaration	В	Assignment statements	
	С	Input and output statements	D	Structural statements	
061.	Subs	stitution of values for names (whose val	ues a	are constants) is done in	С
	Α	Local optimization	В	Loop optimization	
	С	Constant folding	D	Strength reduction	
062.	In co	mpiler terminology reduction in strengtl	h mea	ans	D
	Α	Replacing run time computation by	В	Removing loop invariant computation	
		compile time computation		•	
	С	Removing common subexpressions	D	Replacing a costly operation by a	
				relatively cheaper one	
063.	Whic	ch of the following statements about peo	ephol	e optimization is	D
	Α	It is applied to a small part of the code	В	It can be used to optimize	
				intermediate code	
	C	To get the best out of this, it has to be	D	It can be applied to the portion of the	

		applied repeatedly		code that is not contiguous	
064.	Som	e code optimizations are carried out or	the i	ntermediate code because	Α
	Α	They enhance the portability of the	В	Program analysis is less accurate on	
		compiler to other target processors		intermediate code than on machine	
				code	
	С	The information from dataflow	D	The information from the front end	
		analysis cannot otherwise be used for		cannot otherwise be used for	
		optimization		optimization	
065.	Whic	ch one of the following is FALSE?			D
	Α	A basic block is a sequence of	В	Available expression analysis can be	
		instructions where control enters the		used for common sub expression	
		sequence at the beginning and exits		elimination.	
		at the end.			
	С	Live variable analysis can be used for	D	x = 4 5 => $x = 20$ is an example of	
		dead code elimination.		common subexpression elimination.	
066.	In th	e context of compiler design, reduction	in str	•	Α
	Α	Code optimization obtained using	В	Reduction in accuracy of the output	
		cheaper machine instructions			
	С	Reduction in the range of values of	D	Reduction in efficiency of the program	1
		input variables			
067.	Loop	unrolling is a code optimization techni			Α
	Α	That avoids tests at every iteration of	В	That improves performance by	
		the loop.		decreasing the number of instructions	,
	_		_	in a basic block.	
	С	That exchanges inner loops with	D	That reorders operations to allow	
		outer loops		multiple computations to happen in	
	_			parallel	_
068.		hole optimization is a form of	_	Landa Calada	С
	A	Loop optimization	В	Local optimization	
000	C	Constant folding	D	Data flow analysis	Ь
069.		d-code elimination in machine code opt			В
	А	Removal of all labels.	В	Removal of values that never get	
	С	Removal of function which are not	D	used. Removal of a module after its use.	
	C	involved.	D	ivernoval of a module after its use.	
070	Code	e optimization is responsibility of:			Α
070.	A	Application programmer	В	Server System	^
	C	Client System	D	Browser	
071		d-code elimination in machine code opt	_		В
0.	A	removal of all labels.	В	removal of values that never get	
	, ,			used.	
	С	removal of function which are not	D	removal of a module after its use.	
		involved.			
072.	The	identification of common sub-expression	n an	d replacement of runtime computations	A
		ompile-time computations is:		·	
	Á	Local optimization	В	Constant folding	
	С	Loop Optimization	D	Data flow analysis	
073.	One	of the purposes of using intermediate	code	in compilers is to	C
	Α	make parsing and semantic analysis	В	improve error recovery and error	
		simpler.		reporting.	
	С	increase the chances of reusing the	D	improve the register allocation.	
		machine-independent code optimizer			
	_	in other compilers.		<u> </u>	_
074.	Cons	sider the grammar rule E E1 - E2 for a	rithm	etic expressions. The code generated	В

	the f	rgeted to a CPU having a single user re irst operand to be in the register. If E1 a ession, in order to get the shortest pos	and E	2 do not have any common sub	
	A	E1 should be evaluated first	В	E2 should be evaluated first	
	C	Evaluation of E1 and E2 should	D	Order of evaluation of E1 and E2 is of	
	C	necessarily be interleaved	D		
075	Λ ar	•	vo fo	no consequence	С
075.		ammar that is both left and right recursi			C
	A C	Ambiguous Information is not sufficient to decide	В	Unambiguous	
	C		D	Context Free Grammar	
		whether it is ambiguous or			
070	^	Unambiguous.			_
076.	_	hesized attribute can easily be simulate			С
	A	LL grammar	В	Ambiguous grammar	
	C	LR grammar	D	RL grammar	_
077.	_	which of the following situations, inherit			В
	A	Evaluation of arithmetic expressions		Keeping track of variable declaration	
	С	Checking for the correct use of L-	D	Syntax Checking	
	_	values and R-values			_
078.		syntax directed translation scheme, if the			С
		tion of the values of the attributes of its			
	Α	Synthesized attribute	В	Inherited attribute	
	С	Canonical attribute	D	Register attribute.	
079.	Inco	mpatible types work with the		_	В
	Α	Syntax tree	В	Semantic analyzer	
	С	Code optimizer	D	Lexical analyzer	
080.		sider the basic block given below.a = b			A
		mum number of nodes and edges pres	ent in	the DAG representation of the above	
	basi	c block respectively are			
	Α	6 and 6	В	8 and 10	
	С	9 and 12	D	4 and 4	
081.		sider the following Syntax Directed Trai			C
	term	inals {S, A} and terminals {a,b}.SaA{pri	nt1} S	Sa{print2} SSb{print3}Using the above	
	SDT	S, the output printed by a bottom-up pa	arser,	for the input aab is:	
	Α	132	В	223	
	С	2 3 1	D	syntax error	
082.	Whic	ch of the following is not an intermediate	e cod	e form?	В
	Α	Postfix notation	В	Syntax trees	
	С	Three address codes	D	Quadruples.	
083.	Thre	e address codes can be implemented l	οу		D
	Α	Indirect triples	В	Direct triples	
	С	Quadruples	D	Indirect triples and Quadruples	
084.	Synt	ax directed translation scheme desirab	le be		C
	Α	It is based on the syntax	В	Its description is independent of any	
				implementation	
	С	It is easy to modify	D	Syntaxes	
085.	The	graph depicting the inter-dependencies	of th	e attributes of different nodes in a	В
	pars	e tree is called a			
	A	Flow graph	В	Dependency graph	
	С	Karnaughs graph	D	Steffi graph	
086.	Cho	ose the correct statements.			В
	Α	Topological sort can be used to	В	Evaluation order for a dependency	
		obtain an evaluation order of a		graph dictates the order in which the	
		dependency graph.		semantic rules are done.	
	С	Code generation depends on the	D	Simulated by poises	

order in which semantic actions are performed.

087.	A syr	ntax tree			Α
	Α	Is another name for a parser tree	В	Is a condensed form of parse tree	
	С	Should not have keywords as leaves	D	is binary search tree	
088.	Whic	th table is a permanent database that h			D
	Α	Reductions	В	Identifier table	
	С	Literal table	D	Terminal table	
089.	Acce	ess time of the symbol table will be loga	rithm		В
	Α	Linear list	В	Search tree	
	С	Hash table	D	Self-organizing list.	_
090.		th of the following symbol table implem	entati	ion is best suited if access time to be	С
		num?	_		
		Linear list	В	Search tree	
004	C	Hash table	D	self-organization list	_
091.		e address code involves	_	At the court O or I have a co	В
	A	Exactly 3 addresses	В	At the most 3 addresses	
000	С	No unary operator	D	None of the above.	_
092.		bol table can be used for	_	Overage and a development	С
	Α	Checking type compatibility	В	Suppressing duplicate error	
	_	Otava na alla satiana	_	messages	
000	C	Storage allocation	D	Type casting.	_
093.			entat	ion is based on the property of locality	ט
		ference	D	Coords troo	
	A C	Linear list	В	Search tree	
004	_	Hash table	D or ma	self-organization list	В
U34.		ch data structure in a compiler is used f cheir attributes?	OI IIIc	anaging information about variables	Ь
	_	Abstract syntax tree	В	Symbol table	
	A C	Semantic stack	D	Parse table	
0 95	_	th of the following optimization techniques	_		Α
033.	A	Removal of invariant computation	В	Elimination of induction variables	^
	C	Peephole optimization	D	Invariant computation	
096	_	cept which can be used to identify loops		invariant computation	Α
050.	A	Dominators	В В	DAG	
	C	Breadth first ordering	D	Nodes	
097.	_	leal compiler should	_	110000	В
••••	Α	Detect error	В	Detect and report error	_
	C	Detect, report and correct error	D	Convert source code.	
098.	_	phole Optimization is a form of	_		В
	Α	Loop optimization	В	Local optimization	
	С	Constant folding	D	Data flow analysis	
099.	Subs	stitution of values for names whose values	ues a	•	С
	Α	Local optimization	В	Loop optimization	
	С	Constant folding	D	Global optimization	
100.	User	s write the programs in which language	e?	·	В
	Α	Low-level Language	В	High-Level Language	
	С	Decimal-Format	D	Middle-Level Language	
101.	Does	s the compiler program translate the wh	nole s	<u> </u>	D
	Α	No	В	Depends on the Compiler	
	С	Don 't Know	D	Yes	
102.	A ba	sic block can be analyzed by a			Α
	Α	DAG	В	Graph which may involve cycles	
	С	Flow-graph	D	Flow-chart	

103.	The calle	technique of replacing run time computed	tation	by compile time computations is	Α
	A	Constant folding	В	Code hoisting	
	C	Peephole optimization	D	Invariant computation	
104.		graph that shows basic blocks and their		•	В
. •	Α	DAG	В	Flow graph	_
	C	Control graph	D	Leader Graph	
105.		uction in strength means			Α
	A	Replacing run computation by compile computation	В	Removing loop invariant computation	
	С	Removing common sub-expression	D	Replacing a costly operation by a relatively cheaper one	
106.	In a	two pass assembler, adding literals to I	literal	table and address resolution of local	D
	syml	bols are done using?			
	Α	First pass and second respectively	В	Second pass	
	С	Second pass and first respectively	D	First pass	
107.	In T	wo pass assembler the object code ger	neratio	on is done during the	Α
	A	Second pass	В	First pass	
	С	Zero pass	D	Not done by assembler	
108.	Whic	ch of the following class of statement us	sually	•	В
	com	-	,	•	
	Α	Assignment statement	В	Structural statements	
	С	Input and output statements	D	Declaration	
109.	In th	e compiler, the function of using interm	ediate	e code is:	D
	Α	to improve the register allocation	В	to increase the error reporting &	
				recovery.	
	С	to make semantic analysis easier.	D	to increase the chances of re-using the machine-independent code	
110	\//hic	ch compiler runs on one machine and g	onor	optimizer in other compilers.	В
110.		Multipass compiler			Ь
	C	Optimizing compiler	D	Onepass compiler	
111		ch optimization technique is used to rec		•	В
		Latter optimization technique		Peephole optimization technique	
		Local optimization technique	D	Code optimization technique	
112		is a program that takes a progra		·	Α
114.		uage) as an input and produces as out			
	_	uage	putu	program in another language (object	
	A	Translator	В	Assembler	
	C	Compiler	D	Interpreter	
113.	_	e source language is high-level languag	_	•	С
		uage (assembly or machine), then such			
	A	Translator	В	Assembler	
	C	Compiler	D	Interpreter	
114.	_	hich addressing mode the operand is g	_	•	В
	Α	Absolute mode	В	Immediate mode	_
	C	Indirect mode	D	Index mode	
115.	_	the machine dependent phase of the c	_		D
	A	Syntax analysis	В	Code generation	_
	C	Lexical analysis	D	Code generation	
116.	_	stem program that combines the separ		•	В
		suitable for execution is?		,	_
	A	Assembler	В	Linking loader	
				-	

	C	Cross compiler	D	Load and Go	
117.	Туре	e checking is normally done during			C
	Α	Lexical analysis	В	Syntax analysis	
	С	Syntax directed translation	D	Code generation	
118.	The	phase receives optimized in	terme	ediate codes and generates the code	D
	for e	xecution			
	Α	lexical analyzer	В	syntax analyzer	
	С	code optimizer	D	code generator	
119.	Whic	ch of the following is NOT an advantag	e of u	sing shared, dynamically linked	C
	libra	ries as compared to statically linked lib	raries	?	
	Α	Smaller sizes of executable	В	Lesser overall page fault rate in the	
				system	
	С	Faster program start-up	D	Existing programs need not be re-	
				linked to take advantage of newer	
				versions of libraries	
120.	In ar	n implementation of a compiler, portion	s of o	ne or more phases are combined into	Α
		odule called a		•	
	Α	Pass	В	Parser	
	С	Scanner	D	Set	
121.	_	actic structure can be regarded as a tre	ee wh		С
	A	Scanner	В	Parser	
		Tokens	D	Macro	
122.		phase designed to improve the in	nterm		Α
		Code optimization	В	Code Generation	
	С	Intermediate code generator		Syntax Analyzer	
123.		structure used to record the information		· ·	В
	Α	Syntactic	В	Symbol	
	C	Value	D	Tokens	
124.	_	ch of the following are labeled by opera	tor s		В
	Α	Root	В	Interior nodes	
	C	Leaves	D	Nodes	
125.	_	t is E val for string 1 * 0?	_		С
	Α	8	В	6	
	C	4	D	12	
126.	The	action of parsing the source program in	- nto pr		В
		are are the second and are the gramm			
	A	Syntax analysis	В	Lexical analysis	
	С	Interpretation analysis	D	General syntax analysis	
127.	Whic	ch one to the following false?		,	D
	Α	The code contains loop-in variant	В	There is scope of common sub-	
		computation		expression elimination in this code	
	С	There is scope strength reduction in	D	There is scope of dead code	
		this code		elimination in this code	
128.	Wha	It is the function of the syntax phase?			Α
	Α	recognize the language and to call	В	Build a literal table and an identifier	
		the appropriate action routines that		table	
		will generate the intermediate form or			
		matrix for these constructs			
	С	Build a uniform symbol table	D	Parse the source program into the	
	_	_ ind a dimenti dyffilddi tablo	_	basic elements or tokens of the	
				language	
129	Whir	ch languages necessarily need heap al	locati	<u> </u>	С
	A	Those that support recursion	В	Those that use dynamic scoping	•
	C	Those that Allow dynamic data	D	Those that use global variables	
	_	ayriairiib data	_	gionai tallanio	

structure