001.	A Sy	inthesized attribute is an attribute whos	e van	ue at a parse tree node depends on	C
	Α	Attributes at the siblings only	В	Attributes at parent node only	
	С	Attributes at children nodes only	D	Attributes at children nodes and parents nodes.	
002.	In a	bottom up evaluation of a syntax direct	ion d	•	С
	A	Always be evaluated	В	Be evaluated only if the definition is L attributed	
	С	Evaluation only done if the definition has synthesized attributes	D		
003.	Con	sider the given below SDT. P1: S -> MI	۷ {S.۱	val= M.val + N.val} P2: M -> PQ {M.val	C
	= P.	val * Q.val and P.val =Q.val} Select the		•	
	Α	Both P1 and P2 are S-attributed.	В	P1 is S-attributed and P2 is L-attributed.	
	С	P1 is L-attributed but P2 is not L-attributed.	D	Both P1 and P2 are L- attributed	
004.	Synt	thesized attribute can be easily simulate	_		D
	Α	LL grammar	В	Ambiguous grammar	
	C	CLR grammar	D	LR grammar	_
005.	_	errors that can be pointed out by the co			D
	A	Internal errors	В	Logical errors	
000	C	Semantic errors	D	Syntax errors	
006.		rited attribute is a natural choice in	В	Correct use of L and B values	Α
	A	Tracking declaration of a variable		Correct use of L and R values	
007	C	Tracking the declaration of L values	D	Correct use of R values	С
υυ <i>τ</i> .	A	nantic errors can be detected at Compile time only	В	Run-time only	C
	C	Both Compile and Run time	D	Loading time	
സമ		ch of the following groups is/are token t			С
000.	A	Syntax Analyzer	B	Semantic Analyser	C
	C	Lexical Analyzer	D	Intermediate Code Generation	
009		tax directed translation can be based or		memediate dode deneration	С
000.	A	Syntax Tree	''	Parse Tree	J
	C	Syntax tree as well as Parse Tree.	D	Abstract Syntax Tree	
010.		ations for associating semantic rules with			D
	Α	Syntax Directed Definition	В	Translation Scheme	
	С	Postfix Notation	D	Syntax Directed Definition and	
				Translation Scheme	
011.	The	interdependencies among inherited an	d syn	thesized attributes at nodes in parse	В
		can be depicted by		•	
	Α	DAG	В	Dependency graph	
	С	Interdependency graph	D	Wait for graph	
012.	Inhe	rited attribute can easily be simulated b	y an	OW	Α
	Α	LL grammar	В	Ambiguous Grammar	
	С	LR grammar	D	Unambiguous Grammar.	
013.	Synt	tax Directed Translation (SDD) that invo	olves	only synthesized attributes is	Α
	calle	ed			
	Α	S-Attributed	В	L-Attributed	
	С	K-Attributed	D	R- Attributed	
014.	Whi	ch one of the following is not method fo	r eva	luating semantic rules? NSWER	C
	Α	Parse tree methods	В	Oblivious method	
	С	Syntax -based methods	D	Rule-based methods	
015.	In sy	nthesize attribute, we can evaluate in_		order.	Α

	A	Bottom-up	В	i op-aown	
	С	Preorder	D	postorder.	
016.	Synt	ax Directed Translation (SDD) without	side e	effects is called	C
	Α	Context free grammar	В	Operator grammar	
	С	Attribute grammar	D	Context sensitive grammar	
017.	Whe	en is the type checking usually done?		•	Α
	Α	During syntax directed translation	В	During lexical analysis	
	C	During code optimization	D	During syntax analysis	
018	_	ch of the following component is import			D
010.	_		В	Symbol Table	D
	A C	Lex			
040	_	Yacc	D	Type Checking	_
019.		e compiler, the function of using interm			С
	Α	toincrease the error reporting &	В	tomake semantic analysis easier	
		recovery			
	С	toincrease the chances of re-using	D	toimprove the register allocation	
		the machine-independent code			
		optimizerin other compilers			
020.	Whic	ch one of the following is not type of int	erme	diate code representations?	С
	Α	Syntax tree	В	Postfix	
	C	Preorder	D	Three address code	
021	_	herited attribute, we can evaluate in	_		В
UZ 1.		WER		Older.Ovv	ט
	_		D	Top down	
	A	Bottom-up	В	Top-down	
000	C	Post order	D	Preorder	
022.		ch of the following is essential for conve	erting	an infix expression to the postfix from	Α
		iently?			
		An operator stack	В	An operand stack	
	С	An operand stack and an operator	D	A parse tree	
		stack			
023.	Fron	n the following production with semanti	c rule	E.val is E E1 + T {	D
		I = E1.val + T.val}		·	
	Α	L-Attribute	В	Inherited Attribute	
	С	Syntax Attribute	D	Synthesized Attribute	
024		It is true about Syntax Directed Definition			С
02 11	A	Syntax Directed Definitions +	B	Syntax Directed Definitions + CFG =	
	/ \	Semantic rules = CFG	D	Semantic rules	
	С		D		
	C	CFG + Semantic rules = Syntax	D	CSG+ Syntax Directed	
005		Directed Definitions		Definitions=Semantic Rules	^
025.		_ constructs the desired target progra	m tror	n the intermediate representation of	С
	_	source program.	_		
	A	Analysis Part	В	Lexical Part	
	С	Synthesis Part	D	Syntactic Part	
026.	In a	single pass assembler, most of the for	ward i	references can be avoided by putting	В
	the r	estriction.			
	Α	on the number of strings/life reacts	В	code segment to be defined after	
		· ·		data segment	
	С	on unconditional jumps	D	on conditional jumps	
027.	_	k-patching is useful for handling	_		D
U	A	uncondinal jumps	В	conditional jump	
	Ĉ	backward references	D	forward references	
വാഠ	_				D
U Z 0.		ch of the following component is import		•	ט
	A	Yacc	В	Lex	
000	С	Symbol Table	D	Type Checking	_
029.		I ype Checking is defined a	as the	type checking being done at run time.	C

	Α	Static	В	Control	
	С	Dynamic	D	Syntactic	
030.	Tran	slating the expression given below into	quad	druple representation, how many	D
	oper	ations are required? (i*j) + (e+f) * (a*b+	c).		
	Α	5	В	2	
	С	4	D	6	
031.	Qua	druple is a record structure with four fie	lds		Α
	Α	op, arg1, arg2 and result	В	op1, op2, arg2 and result	
	С	arg1, arg2, result and op	D	arg1,arg2,arg3 and result	
032.	Grap	oh used to represent semantic network	is		В
	-	Undirected Graph	В	Directed Graph	
	С	Directed Acyclic Graph (DAG)	D	Directed Complete Graph	
033.	Cons	sider line number 3 of the following C -	progr	·	C
	n, i+	_			
	A	No compilation error	В	only lexical errors	
	С	only syntactic errors	D	both lexical and syntactic errors	
034.	_	s an example of	_		Α
•••	Α	postfix notation	В	abstract syntax tree	
	C	three address code	D	parse tree	
035.	_	ch of the following representation is a ty		•	D
000.	A	Quadruples	В	Triples	
	C	Indirect triples	D	All of the above	
036	_	ch statement is an abstract form of inter	_		Α
000.	A	3- address	В	2-address	^
	C	address	D	Intermediate code	
037	_	tify the function which generates three-	_		В
031.	A	gen_code()	В	emit()	_
	C	new_label()	D	lookup()	
บรร	_	three address code is a combination of		• •	В
030.	A	True	В	False	ט
	C	True or False	D	True and False	
U30	_	tify the incorrect statement about three	_		D
033.	A	It is not used by the optimizing	В	The instruction was hard to translate	ט
	\wedge	compilers.	0	into assembly language.	
	С	•	D	All the mentioned	
	C		U	All the mentioned	
040	\//ha	has at least three operands.	ont?		D
U 4 U.		t is the function of the storage assignm	В	Assign storage to all temperary	ט
	Α	Assign storage to all variables	Ь	Assign storage to all temporary	
		referenced in the source program		locations that are necessary for intermediate results	
	<u></u>	Assign storage to literale, and to	D		
	С	Assign storage to literals, and to	D	All of the mentioned	
		ensure that the storage is allocated			
		and appropriate locations are			
044	14/1 -	initialized			_
041.		t is garbage?	_	Allered a Leteran Leteran	В
	Α	Unallocated storage	В	Allocated storage whose access	
	_	All control	_	paths are destroyed?	
0.40	С	Allocated storage	D	Uninitialized storage	_
042.		_	e tor	Boolean expressions and flow-of-	С
	_	rol statements in one pass.	_	_	
	A	Procedure call	В	Type expression	
	C	Back-patching	D	Type equivalence	_
043.	Thre	e address code statements are typically	y imp	lemented in the compiler	D
	ลร				

	Α	hash tables	В	symbol tables	
	С	linked lists	D	records	
044.		have pointers to triple.			D
	Α	Triples	В	Quadruples	
	C _	TAC	D	Indirect triples	
045.		riples, a record withfields	s repr	esent each of the Three Address	С
		e(TAC) statements.	_	_	
	A	4	В	2	
0.40	С	3	D	1	
046.				s the procedure name, parentheses	Α
		actual parameter names or values, and			
	A	Procedure Call	В	Calling Sequence	
047	C	Callee Sequence	D locate	Calling Function	۸
047.		memory allocation, storage is al			A
	_	blished and destroyed during the execu	лион с В	Static	
	A C	Dynamic Automatic	D D	Heap.	
0/18	_	ch field is not present in activation reco	_	Πeap.	В
040.	A	Saved machine status	B	Register allocation	ט
	C	Optional control link	D	Temporaries	
049.		•		ed or deallocated at arbitrary points	D
040.		ng its execution.	inocai	ica of acamocated at arbitrary points	
	A	Static	В	Dynamic	
	C	Automatic	D	Program Controlled	
050.	_	languages that need heap allocation in			D
	Α	Those that use Global Variable	В	Those that use Dynamic Scoping	
	С	Those that support Recursion	D	Those that allow Dynamic Data	
				Structures	
051.	Iden	tify the Technique used to replace Run	-Time	Computations with Compile-Time	С
	Com	putations.			
	Α	Peephole Optimization	В	Invariant Computations	
	С	Constant Folding	D	Code Hoisting	
052.	Dyna	amic memory allocation is implemented	l usin		D
	Α	Array	В	Stacks	
	С	Heap	D	Stacks and Heap	
053.	The	memory for variable is allocated before	the e	execution of a program is called	Α
		_allocation.	_	_	
	A	Static	В	Dynamic	
054	С	Automatic	D	Program Controlled	
054.			-	ntrol enters and leaves activations.	Α
	A	Activation	В	Parse	
OEE	C	Syntax	D	Semantic	ь
U 35.	_	ctivation tree each node represent	В	Activation of a procedure	В
	A C	Activation of main program Activation of a function call	D D	Activation of a procedure	
056	_			Activation of a record	D
050.	A	ch of the following fields are of activatio Return value	В	Local data	ט
	C	Temporaries	D	All of the above SHOW ANSWE	
057		ch is not part of runtime memory subdiv			С
υσ <i>ι</i> .	A	Stack	В	r Heap	J
	C	Access link	D	Static data	
058	_	size field of activation record can be de	_		В
550.	A	Run Time	В	Compile Time	٠
	C	Compile Time and Run Time	D	Compile Time or Run Tim	
	-	r		r	

059.	Whi	ch of the following symbol table impleme	entati	on makes efficient use of memory?	С
	Α	List	В	Search Tree	
	С	Hash Table	D	Self-Organizing List	
060.	In w	hich storage allocation strategy size is r			Α
	Α	Static Allocation	В	Dynamic Allocation	
	С	Stack Allocation	D	Heap Allocation.	_
		ctivation record, optional control link poi			Α
				Activation record of callee	
000	C	Activation record of procedure			_
062.		ch of the following is used in various sta	_		С
	A C	Records SymbolTable	B D	Program Table	
063		field of actual parameter in activation re	_		В
003.		edure	coru	is used by willon	Ь
	A		В	Called procedure	
	C	0 1	D	Calling sequence	
064.		node for a is the parent of node for b if			В
	Α	If lifetime of a occurs before lifetime of			_
		b		b	
	С	If lifetime of b occurs before lifetime of	D		
		a		a	
065.	If the	e occurrence of name in procedure is in	the s	scope of declaration within the	Α
	proc	edure then it is said to be			
	Α	Local	В	Temporary	
	С	Global	D	Class	
066.		by reference also called as	·		С
		,	В	Call-by-location	
~~=	C	Call-by-address and location		Call-by-value	_
067.		hich allocation, names are bound to sto	rage	as program is	D
		piled Static	В	Hoon	
	A C	Dynamic	D	Heap Stack	
	_	graph that shows basic blocks and thei	_		С
000.	1110	graph that shows basic blocks and the	1 3000	cessor relationship is called	J
	A	DAG	В	Control Graph	
	C	Flow Chart	D	Hamilton graph	
069.	An o	ptimizing compiler		3 4	Α
	Α	Optimized the code	В	Is optimized to occupy less space	
	С	Is optimized to take less time for	D	Optimized to occupy less space and	
		execution		less time for execution	
070.	Som	e code optimizations are carried out on	the i	ntermediate code because	В
	Α	They enhance the portability of the	В	Program analysis is name accurate	
		complier to other target processors		on intermediate code than on	
				machine code	
	С	The information from data flow	D	The information from the front end	
		analysis cannot otherwise be used for		cannot otherwise be used for	
		optimization		optimization	_
071.		never a procedure is executed, its activ	ation/	record is stored on the stack, also	В
		vn as?	Ь	Control ato alc	
	A	Access Stack	В	Control stack	
072	C	Formal Stack	D	Return Stack	С
u/Z.	Δ	location of memory (address) where an		k-value	C

	С	I-value	D	t-value	
073.	A ba	sic block can be analyzed by		_	В
	Α	Graph with Cycles	В	Directed Acyclic Graph(DAG)	
	С	Flow Graph	D	Directed Graph	
074.	The	value of which variable is updated insic	de the	loop by a loop-invariant value?	D
	Α	Invariable	В	Strength	
	С	Loop	D	Induction	
075.	Subs	stitution of values for names (whose va	lues a	are constants) is done	C
	in	·			
	Α	Local optimization	В	Loop optimization	
	С	Constant folding	D	Strength reduction	
076.	Whic	ch of the following comment about peep	o-hole	e optimization is true?	Α
	Α	It is applied to small part of the code	В	It can be used to optimize	
		and applied repeatedly		intermediate code	
	С	It can be applied to a portion of the	D	It is applied in symbol table to	
		code that is not contiguous		optimize the memory requirements	
077.	In co	impiler terminology reduction in strengt	h me		D
	Α	Replacing run time computation by	В	Removing loop invariant computation	
		compile time computation			
	С	Removing common sub expressions	D	Replacing a costly operation by a	
				relatively cheaper one	
078.	DAG	representation of a basic block allows			Α
	Α	Automatic detection of local common	В	Detection of induction variables	
	_	sub expressions	_		
	C	Automatic detection of loop variant	D .	Detection of dead code elimination	_
079.	_	acement of an expensive operation by			D
	A	Common Sub Expression Elimination		Loop-Invariant Computation	
000	С	Code Motion	D	Reduction in Strength	_
080.		phole optimization	_	Land Order and	С
	A	Loop Optimization	В	Local Optimization	
004	C	Constant folding	D	Data Flow analysis	_
UO1.	_	ch graph describes the basic block and DAG	_		В
	A C	Control graph	B D	Flow graph Hamilton graph	
ດຂວ		ch graph describes the basic block and		• .	С
002.	A	Control graph	В	DAG	C
	C	Flow graph	D	Hamilton graph	
083	_	e can be optimized at		Transition graph	Α
000.	A	Source from user	В	Target code	•
	С	Intermediate code	D	Assembler	
084.	Optin	mization can be categorized broadly int	:0	types.	Α
	Α΄	2	В	3	
	С	4	D	5	
085.	Code	e optimization is responsibility of			В
	Α	Application Programmer	В	System Programmer	
	С	Operating System	D	Database Administrator	
086.	Dead	d-code elimination in machine code opt	imiza	tion refers to	В
	Α	Removal of all labels.	В	Removal of values that never get	
				used.	
	С	Removal of function which are not	D	Removal of a module after its use.	
		involved.			
087.	The	technique of live variable analysis is us	ed fo		С
	Α	Type checking	В	Code generation	
	С	Code optimization	D	Parsing	

088.	Which of the following symbol table impler	nenta	tion is based on the property of locality	В
	of reference?	_	0 1/4	
	A Linear list	В	Self-organizing list	
	C Search tree	.D	Hash table	_
089.	Which optimization technique is used to re			В
	A Latteroptimization technique	В	• •	
	C Localoptimization technique	D	Codeoptimization technique	
090.	Input to code generator is			В
	A Source code	В	Intermediate code	
	C Target code	D	Assembly code	_
091.	Which method merges the multiple loops i			C
	A Constant Folding	В	Loop rolling	
	C Loop fusion or jamming			
092.	The optimization which avoids test at ever	-		В
	A Loop jamming		Loop unrolling	
	C Constant folding	D	Loop Invariant	
093.	The optimization technique which is typica			D
	A Removal of invariant computation			
	C Constant folding		All of these	
094.	Local and loop optimization in turn provide			Α
	A Data flow analysis	В	3	
	C Pee hole optimization	_	3	
095.	Who is responsible for the creation of the	symbo	ol table?	В
	A Assembler	В	Compiler	
	C Interpreter	D	Loader	
096.	x * 2 can be replaced by x << 1 is an exam	nple o	f?	C
	A Algebraic expression simplification	В		
	C Strength reduction	D		
097.	The following code is an example of? void	add_	ten(int x) { return x + 10; printf(""value	В
	of x is %d"", x); }			
	A Redundant instruction elimination	В	Unreachable code	
	C Flow of control optimization	D	None of the above	
098.	In Algebraic expression simplification, a =	a + 1	can simply be replaced by?	В
	A A	В	INC a	
	C DEC a	D	MUL a	
099.	Which of the following class of statement u	usuall	y produces no executable code when	Α
	compiled?			
	A Declaration	В	Assignment statements	
	C Input and output statements	D	Structural statement	
100.	is the final phase of compile	er		В
	A Semantic analysis	В	Code generation	
	C Target code generation	D	Syntax analysis	
101.	Consider the following intermediate progra	am in t	three address code p = a - b q = p * c p	В
	= u * v q = p + q Which one of the following	g corr	esponds to a static single assignment	
	form of the above code?	_		
	A p1=a-b	В	n3 - a - h	
	b1=8-0		p3 = a - b	
	q1 = p1 * c		q4 = p3 * c	
	p1 = u * c		p4 = u * c	
	q1 = p1 + q1		q5 = p4 + q4	
	С	D		

		p3 = u * c		p2 = u * c	
		q2 = p4 + q3		q2 = p + q	
102.		e code optimizations are carried out or	the i	ntermediate code	A
	A	They enhance the portability of the compiler to other target processors	В	Program analysis is more accurate on intermediate code than on machine code	
	С	The information from dataflow analysis cannot otherwise be used for optimization		The information from the front end cannot otherwise be used for optimization.	
103.	Whic	ch of the following is related to synthesi	s pha	se?	В
	Α	Syntax analysis	В	Code generation	
	С	Lexical analysis	D	Semantic analysis	
104.	The	input to the code generator is a		<u>_</u> .	NC
					AN SV ER GI' EN
	A C	Sequence of tree at lexical level Sequence of assembly language instruction	B D	Sequence of tree at semantic level Sequence of machine idioms	
105.	Wha	t is Machine Code			С
	Α	Serial number of the CPU	В	Instructions and data in human readable form	
	С	Instructions and data in binary	D	Instructions and data in assembly code mnemonics	
106.		many descriptors are used for track bo			Α
	addr	esses (location of values) while genera			
	A	2	В	3	
407	С	4	D .	5	_
107.		are used to keep track of men	nory id	ocations where the values of identifiers	В
	_	stored.	D	Address descriptor	
	A	Register descriptor	В	Address descriptor	
100	Code	Memory descriptor	D on to	Process descriptor	D
100.		e generator uses function	טוו נט	uctermine the status of available	D
	A	sters and the location of name values. setReg	В	cinReg	
	C	pfReg	D	•	
100		acing the expression 4*2.14 by 8.56 is	_	getReg	Α
103.	A	Constant Folding	B	Induction Variable	^
	C	Strength reduction	D	Code reduction	
110	_	gister descriptor	ט	Code reduction	В
	A	Keeps track of the usage of a register	В	Keeps track of what is currently in each register	
	С	Keeps track of what is currently in each address	D	Keeps track of the address of each register	
111	The	cost of following instruction sequence i	sMO\	· ·	С
	A	3	B	2	_
	C	6	D	4	
112.	_	ddress descriptor is maintained for			Α
		· · · · · · · · · · · · · · · · · · ·			

p1 = a - b q1 = p2 * c

113.	A C How	Each name in a block Each register variable many descriptors are used for track bo	B D oth the	Each symbol in the program Each memory variable e registers (for availability) and	A
	_	resses (location of values) while genera		_	
	A	2	В	3	
111	C A fro	4	D and a	5	С
114.		igment of code that resides in the loop in the loop in the loop.	and d	omputes the same value at each	C
	A	Induction analysis	В	Strength reduction	
	Ĉ	Loop-invariant code	D	Dead code elimination.	
115		compiler can make use of memory hier	_		Α
	A	TRUE	В	FALSE	^
	C	Can be true or false	D	Can notsa	
116.	The	cost of the instruction MUL #1,R1 is			Α
	Α	2	В	3	
	С	1	D	6	
117.	Forn	nula for Instruction cost			Α
	Α	1+cost for source and destination	В	1-cost for source and destination	
		address modes		address modes	
	С	1+cost for source or destination	D	1-cost for source or destination	
		address modes		address modes	
118.		cost of following instruction sequence is			C
	Assu	uming R0, R1 and R2 contain the addre	esses	of a, b, and c	
	Α	4	В	3	
	С	2	D	1	
119.		ess descriptorkeeps track of the location			Α
	Α	Current value of that variable can be	В	Current register of that variable can	
		found.	_	be found.	
	С	Current address of that variable can	D	Current memory of that variable can	
		be found.		be found.	_
120.		simple code generation algorithm the G	EIR	EG() function is used	В
	to	Datum the leasting	_	Determine the leasting	
	A	Return the location	В	Determine the location	
121	C	Find the previous location	D	Find the next location	^
121.	A	tify the odd statement in the list a = a+b	В	hlil – a	Α
	C	a = a+b c = &b	D	b[i] = a a = *c	
122	_	ch algorithm invokes a function GETRE	_	a = 0	D
122.	A	Code motion algorithm	В	Code optimization algorithm	
	C	Intermediate code	D	Code generation algorithm	
123.	_			of basic blocks, helps to see the flow	Α
		alues flowing among the basic blocks, a		•	
	Α	DAG	В	CAG	
	С	SAG	D	PAG	
124.	Мар	ping names in the source program to a	ddres	ses of data object in run time memory	В
	-	e comparatively by the front end and the			
	Α	Target programs	В	Memory management	
	С	Instruction selection	D	Register allocation	
125.	Addr	ress modes involving registers have co	st		D
	Α	One	В	Two	
	С	Three	D	Zero	
126.	Whic	ch one is the form of target program			В
	Α	High level language	В	Absolute machine language	
	С	Low Level language	D	Middle level language	

127.	Initia	ally all register descriptors value is			В
	Α	Zero	В	Empty	
	С	Non Empty	. D	One	_
128.		criptors are necessary for variable load		· · · · · · · · · · · · · · · · · · ·	D
		essary for variable load and store decis	sion. L	Descriptors are necessary for	
	vana A	ables Load or Store decision	В	Load and Move decision	
	Ĉ	Store and Move decision	D	Load and Store decision	
129.	_	ister allocation depends on:	D	Load and otore decision	Α
0.	A	Cost of loads/stores needed.	В	Cost of loads/moves needed.	•
	С	Cost of instructions/stores needed.	D	Cost of moves/stores needed.	
130.	Glob	pal register allocation can be solved			С
	Α	NP-Complete problems	В	Cache Memory Problems	
	С	Graph colourings problems	D	NP-Hard Problems	
131.				ich one of the following corresponds to	В
		atic single assignment from the above	•		
	Α	x1= a - b y1= x2* c x3= u * v y2= x4+	В	x1= a - b y2= x1* c x2= u * v y3= x2+	
	С	y3	D	y2 x1- 2	
132.	C	p = a - b q = p * c p = u * v q = p + q	D	x1= a - b y1= p * c x2= u * v y2= p + q	D
132.			for (i =	0, i <n; i++)<="" th=""><th>ט</th></n;>	ט
			{		
			for (j=0; j <n; j++)<="" th=""><th></th></n;>	
			1 ;5	(i%2)	
			{	(1/02)	
				x += (4*j + 5*i);	
				y+= (7+4*j);	
			1 3		
			}		
	_	sider the following C code segment.	% _		
	Α	The code contains loop invariant	В	There is scope of common sub-	
		computation			
	C	·	D	expression elimination in this code	
	С	There is scope of strength reduction	D	There is scope of dead code	
133		There is scope of strength reduction in this code		There is scope of dead code elimination in this code	B
133.	The	There is scope of strength reduction in this code pass number for each of the following	activit	There is scope of dead code elimination in this code ties 1. Object code generation 2.	В
133.	The	There is scope of strength reduction in this code pass number for each of the following als added to literal table 3. Listing prince	activit	There is scope of dead code elimination in this code ties 1. Object code generation 2. Address resolution of local symbols	В
133.	The Liter	There is scope of strength reduction in this code pass number for each of the following	activit ted 4.	There is scope of dead code elimination in this code ties 1. Object code generation 2.	В
	The Liter A C	There is scope of strength reduction in this code pass number for each of the following als added to literal table 3. Listing print 1, 2, 1, 2	activit ted 4. B D	There is scope of dead code elimination in this code ties 1. Object code generation 2. Address resolution of local symbols 2, 1, 2, 1 1, 2, 2, 2	В
	The Liter A C Addi A	There is scope of strength reduction in this code pass number for each of the following rals added to literal table 3. Listing print 1, 2, 1, 2 2, 1, 1, 2 ress modes involving memory location One	activit ted 4. B D or lite B	There is scope of dead code elimination in this code ties 1. Object code generation 2. Address resolution of local symbols 2, 1, 2, 1 1, 2, 2, 2 eral have cost Two	
134.	The Liter A C Addi A C	There is scope of strength reduction in this code pass number for each of the following als added to literal table 3. Listing print 1, 2, 1, 2 2, 1, 1, 2 ress modes involving memory location One Three	activit ted 4. B D or lite B D	There is scope of dead code elimination in this code ties 1. Object code generation 2. Address resolution of local symbols 2, 1, 2, 1 1, 2, 2, 2 eral have cost Two Zero	A
134.	The Liter A C Addi A C Whice	There is scope of strength reduction in this code pass number for each of the following rals added to literal table 3. Listing print 1, 2, 1, 2, 2, 1, 1, 2 ress modes involving memory location. One Three ch of the following class of statement up to the statement of the following class of statement up to the following class of st	activit ted 4. B D or lite B D	There is scope of dead code elimination in this code ties 1. Object code generation 2. Address resolution of local symbols 2, 1, 2, 1 1, 2, 2, 2 eral have cost Two Zero	
134.	The Liter A C Addi A C Whice	There is scope of strength reduction in this code pass number for each of the following als added to literal table 3. Listing print 1, 2, 1, 2, 2, 1, 1, 2 ress modes involving memory location. One Three ch of the following class of statement upiled?	activit ted 4. B D or lite B D sually	There is scope of dead code elimination in this code ties 1. Object code generation 2. Address resolution of local symbols 2, 1, 2, 1 1, 2, 2, 2 eral have cost Two Zero produces no executable code when	A
134.	The Liter A C Addi A C Whice com A	There is scope of strength reduction in this code pass number for each of the following rals added to literal table 3. Listing print 1, 2, 1, 2 2, 1, 1, 2 ress modes involving memory location One Three ch of the following class of statement upiled? Declaration	activit ted 4. B D or lite B D sually	There is scope of dead code elimination in this code ties 1. Object code generation 2. Address resolution of local symbols 2, 1, 2, 1 1, 2, 2, 2 eral have cost Two Zero produces no executable code when Assignment statements	A
134. 135.	The Liter A C Addi A C Whice com A C	There is scope of strength reduction in this code pass number for each of the following rals added to literal table 3. Listing print 1, 2, 1, 2 2, 1, 1, 2 ress modes involving memory location. One Three ch of the following class of statement upiled? Declaration Input and output statements	activit ted 4. B D or lite B D sually B	There is scope of dead code elimination in this code ties 1. Object code generation 2. Address resolution of local symbols 2, 1, 2, 1 1, 2, 2, 2 eral have cost Two Zero produces no executable code when Assignment statements Structural statement	A
134. 135.	The Liter A C Addi A C Whio com A C In a	There is scope of strength reduction in this code pass number for each of the following als added to literal table 3. Listing print 1, 2, 1, 2 2, 1, 1, 2 ress modes involving memory location One Three ch of the following class of statement upiled? Declaration Input and output statements resident- OS computer, which of the following class of	activit ted 4. B D or lite B D sually B	There is scope of dead code elimination in this code ties 1. Object code generation 2. Address resolution of local symbols 2, 1, 2, 1 1, 2, 2, 2 eral have cost Two Zero produces no executable code when Assignment statements Structural statement	A
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134. 135.	The Liter A C Addi A C Whio com A C In a	There is scope of strength reduction in this code pass number for each of the following als added to literal table 3. Listing print 1, 2, 1, 2 2, 1, 1, 2 ress modes involving memory location One Three ch of the following class of statement upiled? Declaration Input and output statements resident- OS computer, which of the following class of	activit ted 4. B D or lite B D sually B	There is scope of dead code elimination in this code ties 1. Object code generation 2. Address resolution of local symbols 2, 1, 2, 1 1, 2, 2, 2 eral have cost Two Zero produces no executable code when Assignment statements Structural statement ag system software must reside in the	A
134. 135. 136.	The Liter A C Addi C Whice C In a mair A C	There is scope of strength reduction in this code pass number for each of the following rals added to literal table 3. Listing print 1, 2, 1, 2 2, 1, 1, 2 ress modes involving memory location One Three ch of the following class of statement upiled? Declaration Input and output statements resident- OS computer, which of the formemory under all situations? Assembler Loader	activited 4. B D or lite B D sually B D olllowin	There is scope of dead code elimination in this code ties 1. Object code generation 2. Address resolution of local symbols 2, 1, 2, 1 1, 2, 2, 2 eral have cost Two Zero produces no executable code when Assignment statements Structural statement ag system software must reside in the Linker Compiler	A
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134. 135. 136.	The Liter A C Addi A C Whice Com A C In a mair A C The and	There is scope of strength reduction in this code pass number for each of the following rals added to literal table 3. Listing print 1, 2, 1, 2 2, 1, 1, 2 ress modes involving memory location. One Three ch of the following class of statement upiled? Declaration Input and output statements resident- OS computer, which of the form memory under all situations? Assembler Loader expression(a*b)* cop where '	activited 4. B D or lite B D sually B D ollowin B D on &# an be</th><th>There is scope of dead code elimination in this code ties 1. Object code generation 2. Address resolution of local symbols 2, 1, 2, 1 1, 2, 2, 2 tral have cost Two Zero produces no executable code when Assignment statements Structural statement ag system software must reside in the Linker Compiler #39 is one of '+','*'</th><th>A A</th></tr><tr><th>134. 135. 136.</th><th>The Liter A C Addi A C Whice Com A C In a mair A C The and</th><th>There is scope of strength reduction in this code pass number for each of the following rals added to literal table 3. Listing print 1, 2, 1, 2 2, 1, 1, 2 ress modes involving memory location. One Three ch of the following class of statement upiled? Declaration Input and output statements resident- OS computer, which of the form memory under all situations? Assembler Loader expression(a*b)* cop where '↑ ' (exponentiation) compared to the following class of statement upiled?</th><th>activited 4. B D or lite B D sually B D ollowin B D on &# an be</th><th>There is scope of dead code elimination in this code ties 1. Object code generation 2. Address resolution of local symbols 2, 1, 2, 1 1, 2, 2, 2 tral have cost Two Zero produces no executable code when Assignment statements Structural statement ag system software must reside in the Linker Compiler #39 is one of '+','*'</th><th>A A</th></tr></tbody></table>		

storing 138. reduces the evaluation frequency of expression. В Α Variable propagation В Code motion Dead code elimination C D Induction variable **139.** Generate the code for Indexed Assignment statement a : = b[i] ? Α MOV b(Ri), R MOV b, a(Ri) Α В C MOV b, a[Ri] MOV b[Ri], R D **140.** Register allocation is only within a basic block. It follows ___ approach. D Bottom up Left to right Α В С Right to left D Top-down **141.** Consider the following code segment. x = u - t; y = x * v; x = y + w; y = t - z; y = x * vΑ y;The minimum number of total variables required to convert the above code segment to static single assignment form isNote: This question was asked as Numerical Answer Type. 10 8 Α В C 9 D 6 142. Relative to the program translated by a compiler, the same program when interpreted В runs? Faster В Slower Α At the same speed D May be faster or slower **143.** Advantage of Local register allocation is D Does not consider non-uniform В few values reside in registers distribution of uses C Does consider non-uniform D Heavily used values reside in distribution of uses registers **144.** To find out the live range(s) of each variable and the area(s) where the variable is C used/defined allocation is needed. Α Local Memory В C Global D Register