Merge Sort : Divide and Conquer

input: an array of unsorted elements

output: an array of sorted elements

small problem: If an array contains only a single element in an array then it is itself a sorted array and that we consider as the small problem.

Function call -> Preorder

Function execute -> Postorder

Merge Procedure:

Worst case number of comparisons in Merge Procedure :

Number of comparisons: 7

$$m + n - 1$$

m = number of elements in sorted subarray

1 n = number of elements in sorted subarray

$$4 + 4 - 1 = 7$$

Best case number of comparison in Merge Procedure

```
1 2 3 4 5 6 7 8
```

```
(10,5) = 5 1st comparison
(10,6) = 6 2nd comparison
```

General formula for best case scenario of merge procedure

```
: min(m,n)
```

m = number of elements in sorted subarray 1

n = number of elements in sorted subarray 2

Overall Time Complexity of Merge Procedure:

Number of moves in best and in worst case scenario = m + n

Time complexity = Number of moves + Number of

comparisons =
$$O(m + n)$$

Implementation:

```
def merge(arr, I, m, r):
    // Find sizes of two subarrays to be merged
    n1 = m - I + 1
    n2 = r - m
    // Create temp arrays
    array1 = []
    array2=[]
    // Copy data to temp arrays
```

```
for i in range(n1):
        array.append(arr[l+i])
forin range(n2):
               arrary2.append(arr[m+1+j])
// Initial indexes of first and second subarrays
i = 0, j = 0
// Initial index of merged subarray array
k = 1
while(i<n1 && j<n2):
       if array1[i]<=array[j]:</pre>
               arr[k]=array1[i]:
               j++
       else:
               arr[k]=array2[j]
               j++
       k++
// Copy remaining elements of array1[] if any
while (i < n1):
        arr[k] = array1[i]
        j++
       k++
// Copy remaining elements of array2[] if any
while (j < n2):
        arr[k] = array2[j]
        j++
       k++
```

Note: MergeSort is an outplace sorting algorithm because here we are using a new array to store the elements after doing comparisons.

Mergesort Algorithm:

```
MergeSort(arr,i,j): #arr array , i -starting index, j - ending index
       // small problem
       if(i == j):
              return arr[i]
       // big problem
       else:
              mid = (i + j)/2
                                                   // Divide
                                                                          0(1)
                                                   // Conquer
               MergeSort(arr,i,mid);
                                                   // Left side tree
                                                                          T(n/2)
               MergeSort(arr,mid+1,j);
                                                   // Right side tree
                                                                          T(n/2)
               MergeProcedure(arr,i,mid,mid+1,j); // combine
                                                                          O(n)
```

Overall Time Complexity:

Best, average and worst case scenario

$$O(1) + 2T(n/2) + O(n) = 2T(n/2) + O(n)$$

= $O(n log n)$