



\*) If  $a$  and  $b$  are both  $0$ , return "yes".

\*) If one is  $0$  and the other is  $\text{succ } n$ , return "no".

\*) If  $a = \text{succ } m$  and  $b = \text{succ } n$ , then return the answer to "does  $m = n$ ?"

Our job now is to *prove* that this algorithm always gives the correct answer. The proof that  $0 = 0$  is `rfl`. The proof that  $0 \neq \text{succ } n$  is `zero_ne_succ n`, and the

$m\ n : \mathbb{N}$

Assumptions:

$h : m \neq n$

Goal:

$\text{succ } m \neq \text{succ } n$

contrapose!  $h$

✖ Retry

Active Goal

Objects:

$m\ n : \mathbb{N}$

Assumptions:

$h : \text{succ } m = \text{succ } n$

Goal:

$m = n$

apply `succ_inj` at  $h$

✖ Retry

Active Goal

Objects:

$m\ n : \mathbb{N}$

Assumptions:

$h : m = n$

Goal:

$m = n$

exact  $h$

✖ Retry

level completed! 🎉

**succ\_inj** ✖

$(a\ b : \mathbb{N}) (h : \text{MyNat.succ } a = \text{MyNat.succ } b) : a = b$

## Statement

If  $a$  and  $b$  are numbers, then `succ_inj a b` is the proof that  $(\text{succ}(a) = \text{succ}(b)) \implies a = b$ .

## More technical details

There are other ways to think about `succ_inj`.

You can think about `succ_inj` itself as a function which takes two numbers  $a$  and  $b$  as input, and outputs a proof of  $(\text{succ}(a) = \text{succ}(b)) \implies a = b$ .