CS343 - Operating Systems

Module-9B

Distributed Systems



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Overview

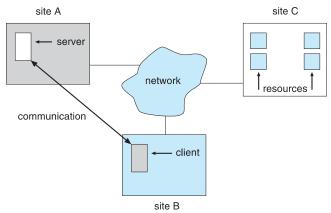
- Network Operating Systems vs Distributed Systems
- Network Structure
- Communication Structure and Protocols
- Design Issues
- Distributed File System

Objectives

- To provide a high-level overview of distributed systems and the networks that interconnect them
- ❖ To discuss the general structure of distributed operating systems
- ❖ To explain general communication structure and communication protocols
- To describe issues concerning the design of distributed systems

Distributed System

- Distributed system is collection of loosely coupled processors interconnected by a communications network
- Processors variously called nodes, computers, machines, hosts
 - Site is location of the processor
 - Generally a server has a resource a client node at a different site wants to use



Distributed System

- Resource sharing
 - Sharing and printing files at remote sites
 - Processing information in a distributed database
 - Using remote specialized hardware devices
- Computation speedup load sharing or job migration
- Reliability detect and recover from site failure, function transfer, reintegrate failed site
- Communication message passing
- ❖ Better user interfaces and easier maintenance by moving from large system to multiple smaller systems performing distributed computing

Types of Distributed Operating Systems

Network Operating Systems

Distributed Operating Systems

Network-Operating Systems

- Users are aware of multiplicity of machines
- Access to resources of various machines is done explicitly by:
 - Remote logging into the appropriate remote machine (telnet, ssh)
 - Remote Desktop (Microsoft Windows)
 - Transferring data from remote machines to local machines, via the File Transfer Protocol (FTP) mechanism
- Users must change paradigms establish a session, give network-based commands
 - More difficult for users

Distributed-Operating Systems

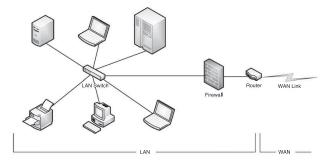
- Users not aware of multiplicity of machines
 - Access to remote resources similar to access to local resources
- Data Migration transfer data by transferring entire file, or transferring only those portions of the file necessary for the immediate task
- Computation Migration transfer the computation, rather than the data, across the system
 - Via remote procedure calls (RPCs)
 - Via messaging system

Distributed-Operating Systems

- ❖ Process Migration run an entire process, or parts of it, at different sites
- ❖ Load balancing distribute processes across network to even the workload
- Computation speedup subprocesses can run concurrently on different sites
- ❖ Hardware preference process execution may require specialized processor
- ❖ Software preference required software may be run at only a particular site
- ❖ Data access run process remotely, rather than transfer all data locally

Network Structure

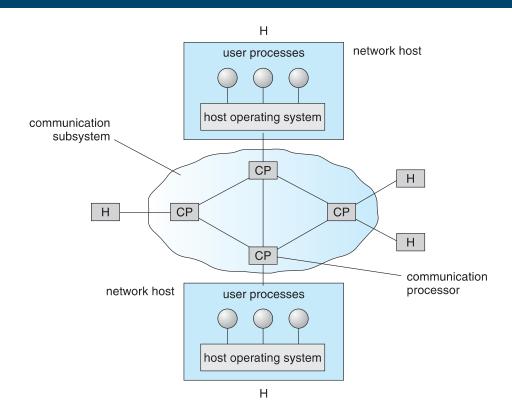
- Local-Area Network (LAN) designed to cover small geographical area
 - Multiple topologies like star or ring
 - Speeds from 1Mbps per second to 40 Gbps for fastest Ethernet over twisted pair copper or optical fibre
 - Consists of multiple computers (mainframes through mobile devices), peripherals (printers, storage arrays), routers (specialized network communication processors) providing access to other networks



Network Structure

- ❖ Wide-Area Network (WAN) links geographically separated sites
 - Point-to-point connections over long-haul lines
 - Implemented via connection processors known as routers
 - Internet WAN enables hosts world wide to communicate
 - Hosts differ in all dimensions but WAN allows communications
 - WANs and LANs interconnect, similar to cell phone network:
 - Cell phones use radio waves to cell towers
 - ❖Towers connect to other towers and hubs

Communication Processors in a Wide-Area Network



Communication Structure – Design Issues

- **❖** Naming and name resolution.
 - How do two processes locate each other to communicate?
- Routing strategies
 - How are messages sent through the network?
- Connection strategies
 - How do two processes send a sequence of messages?
- Contention
 - The network is a shared resource, so how do we resolve conflicting demands for its use?

Naming and Name Resolution

- Name systems in the network
- Address messages with the process-id
- Identify processes on remote systems by <host-name, identifier> pair
- Domain name system (DNS) specifies the naming structure of the hosts, as well as name to address resolution (Internet)

- Fixed routing A path from A to B is specified in advance; path changes only if a hardware failure disables it
 - Since the shortest path is usually chosen, communication costs are minimized
 - Fixed routing cannot adapt to load changes
 - Ensures that messages will be delivered in the order in which they were sent
- Virtual routing- A path from A to B is fixed for the duration of one session. Different sessions involving messages from A to B may have different paths
 - Partial remedy to adapting to load changes
 - Ensures that messages will be delivered in the order in which they were sent

- Dynamic routing The path used to send a message from site A to site B is chosen only when a message is sent
 - Usually a site sends a message to another site on the link least used at that particular time
 - Adapts to load changes by avoiding routing messages on heavily used path
 - Messages may arrive out of order managed by a sequence number

- Router is communications processor responsible for routing messages
- Must have at least 2 network connections
- Maybe special purpose or just function running on host
- Checks its tables to determine where destination host is, where to send messages
 - Static routing table only changed manually
 - Dynamic routing table changed via routing protocol

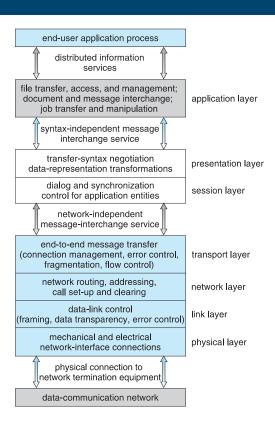
- More recently, routing managed by intelligent software more intelligently than routing protocols
 - OpenFlow is device-independent, allowing developers to introduce network efficiencies by decoupling data-routing decisions from underlying network devices
- Messages vary in length simplified design breaks them into packets (or frames, or datagrams)
- Connectionless message is just one packet
 - Otherwise need a connection to get a multi-packet message from source to destination

Connection Strategies

- Circuit switching A permanent physical link is established for the duration of the communication (i.e., telephone system)
- Message switching A temporary link is established for the duration of one message transfer (i.e., post-office mailing system)
- Packet switching Messages of variable length are divided into fixed-length packets which are sent to the destination
 - Each packet may take a different path through the network
 - The packets must be reassembled into messages as they arrive
- Circuit switching requires setup time, but incurs less overhead for shipping each message, and may waste network bandwidth
- Message and packet switching require less setup time, but incur more overhead per message

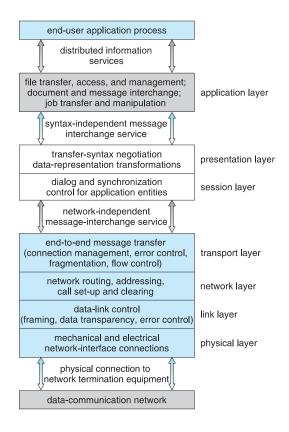
Communication Protocol

- Layer 1: Physical layer handles the mechanical and electrical details of the physical transmission of a bit stream
- ❖ Layer 2: Data-link layer handles the frames, or fixed-length parts of packets, including any error detection and recovery that occurred in the physical layer



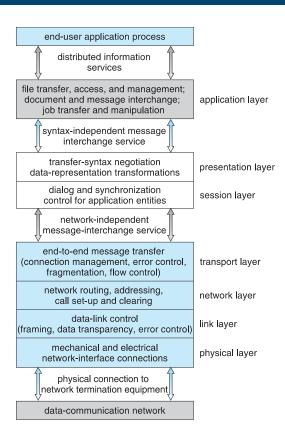
Communication Protocol

- ❖ Layer 3: Network layer provides connections and routes packets in the communication network, including handling the address of outgoing packets, decoding the address of incoming packets, and maintaining routing information for proper response to changing load levels
- ❖ Layer 4: Transport layer responsible for lowlevel network access and for message transfer between clients, including partitioning messages into packets, maintaining packet order, controlling flow

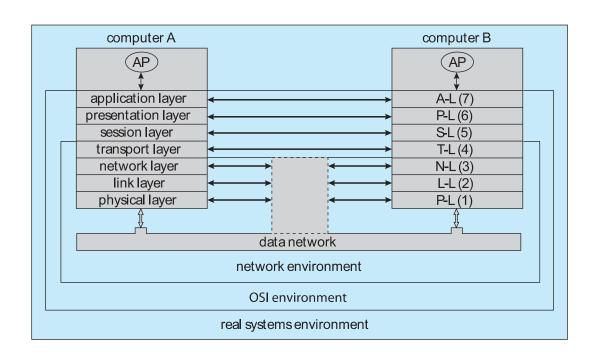


Communication Protocol

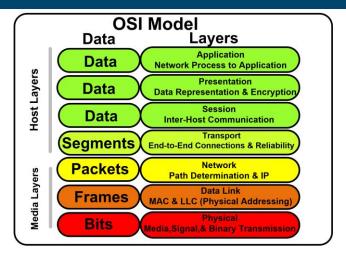
- Layer 5: Session layer implements sessions, or process-to-process communications protocols
- ❖ Layer 6: Presentation layer resolves the differences in formats among the various sites in the network, including character conversions
- ❖ Layer 7: Application layer interacts directly with the users, deals with file transfer, remote-login protocols and electronic mail, as well as schemas for distributed databases

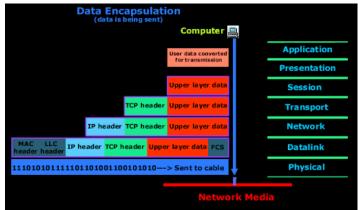


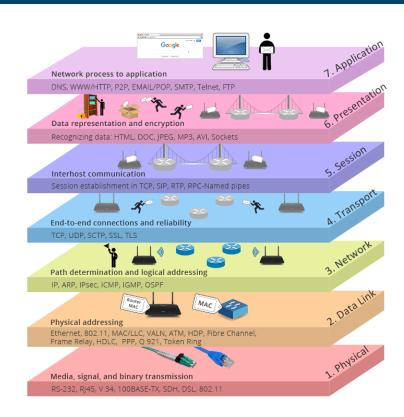
Communication Via ISO Network Model



Communication Via ISO Network Model







Robustness- Failure Detection

- Detecting hardware failure is difficult
- To detect a link failure, a heartbeat protocol can be used
- ❖ Assume Site A and Site B have established a link
 - At fixed intervals, each site will exchange an I-am-up message indicating that they are up and running
- ❖ If Site A does not receive a message within the fixed interval, it assumes either (a) the other site is not up or (b) the message was lost
- Site A can now send an Are-you-up? message to Site B
- If Site A does not receive a reply, it can repeat the message or try an alternate route to Site B

Robustness- Failure Detection

- If Site A does not ultimately receive a reply from Site B, it concludes some type of failure has occurred
- ❖ Types of failures: Site B is down The direct link between A and B is down The alternate link from A to B is down The message has been lost
- ❖ However, Site A cannot determine exactly why the failure has occurred

Robustness-Reconfiguration

- When Site A determines a failure has occurred, it must reconfigure the system:
 - If the link from A to B has failed, this must be broadcast to every site in the system
 - If a site has failed, every other site must also be notified indicating that the services offered by the failed site are no longer available
- When the link or the site becomes available again, this information must again be broadcast to all other sites

Design Issues

- Transparency the distributed system should appear as a conventional, centralized system to the user
- ❖ Fault tolerance the distributed system should continue to function in the face of failure
- Scalability as demands increase, the system should easily accept the addition of new resources to accommodate the increased demand

Distributed File System

- Distributed file system (DFS) a distributed implementation of the classical time-sharing model of a file system, where multiple users share files and storage resources
- ❖ A DFS manages set of dispersed storage devices
- Overall storage space managed by a DFS is composed of different, remotely located, smaller storage spaces
- There is usually a correspondence between constituent storage spaces and sets of files
- Challenges include:
 - Naming and Transparency
 - ❖ Remote File Access

DFS Structure

- Service software entity running on one or more machines and providing a particular type of function to a priori unknown clients
- ❖ Server service software running on a single machine
- Client process that can invoke a service using a set of operations that forms its client interface
- ❖ A client interface for a file service is formed by a set of primitive file operations (create, delete, read, write)
- Client interface of a DFS should be transparent, i.e., not distinguish between local and remote files

Naming and Transparency

- Naming mapping between logical and physical objects
- Multilevel mapping abstraction of a file that hides the details of how and where on the disk the file is actually stored
- A transparent DFS hides the location where in the network the file is stored
- Location transparency file name does not reveal the file's physical storage location
- Location independence file name does not need to be changed when the file's physical storage location changes
- For a file being replicated in several sites, the mapping returns a set of the locations of this file's replicas; both the existence of multiple copies and their location are hidden

Remote File Access

- Remote-service mechanism reduce network traffic by retaining recently accessed disk blocks in a cache, so that repeated accesses to the same information can be handled locally
 - If needed data not already cached, a copy of data is brought from the server to the user
 - Accesses are performed on the cached copy
 - ❖ Files identified with one master copy residing at the server machine, but copies of (parts of) the file are scattered in different caches
 - Cache-consistency problem keeping the cached copies consistent with the master file

Cache Update Policy

- Write-through write data through to disk as soon as they are placed on any cache
- Delayed-write (write-back) modifications written to the cache and then written through to the server later
 - Write accesses complete quickly; some data may be overwritten before they are written back, and so need never be written at all
 - ❖ Poor reliability; unwritten data will be lost when a user machine crashes
 - Variation scan cache at regular intervals and flush blocks that have been modified since the last scan
 - ❖ write-on-close, writes data back to the server when the file is closed
 - Best for files that are open for long periods and frequently modified

Consistency

- Is locally cached copy of the data consistent with the master copy?
- Client-initiated approach
 - Client initiates a validity check
 - Server checks whether the local data are consistent with the master copy
- Server-initiated approach
 - Server records, for each client, the (parts of) files it caches
 - When server detects a potential inconsistency, it must react



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