

1 Acceptance Rejection Method

Introduced by Von-Neumann, this method is among the most widely applicable mechanism for generating random samples. This method generates samples from a target distribution, say F , by first generating candidates from a more convenient distribution, say G and then rejecting some samples generated from G and accepting the rest. The rejection mechanism is designed so that the accepted samples are indeed distributed according to the target distribution.

Suppose we wish to generate samples from a PDF f . Let g be a PDF and we know the technique to generate samples from it. Also, assume that the following condition holds. For some real constant $c \geq 1$,

$$f(x) \leq cg(x) \quad \text{for all } x \in \mathbb{R}.$$

In the acceptance rejection method, we generate a sample X from g and accept the sample with probability $f(X)/cg(X)$. This can be implemented by sampling U from $U(0, 1)$. If X is rejected, a new candidate is sampled from g and the acceptance test applied again. The process repeats until the acceptance test is passed and the accepted value is returned as a sample from f . Thus, we have Algorithm 1. The Theorem 1 provides the justification of the algorithm.

Algorithm 1 Acceptance Rejection Method

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1: repeat
2:   generate  $X$  from distribution  $g$ .
3:   generate  $U$  from  $U(0, 1)$ .
4: until  $U \leq f(X)/cg(X)$ 
5: return  $X$ 
  
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Theorem 1. *Let f and g be two PDFs such that*

$$f(x) \leq cg(x) \quad \text{for all } x \in \mathbb{R} \text{ and for some } c \geq 1.$$

Then X generated by Algorithm 1 has PDF f .

Proof: Fix $x \in \mathbb{R}$. Let Y be the random variable having PDF g . Now, the CDF of X is

$$\begin{aligned}
 P(X \leq x) &= P\left(X \leq x, U \leq \frac{f(Y)}{cg(Y)}\right) + P\left(X \leq x, U > \frac{f(Y)}{cg(Y)}\right) \\
 &= P(X \leq x, cg(Y)U \leq f(Y)) + P(X \leq x | cg(Y)U > f(Y)) P(cg(Y)U > f(Y)) \\
 &= \int_{-\infty}^{\infty} P(X \leq x, cg(Y)U \leq f(Y) | Y = t) g(t) dt + P(X \leq x) P(cg(Y)U > f(Y)).
 \end{aligned}$$

Now, the probability of acceptance is

$$\begin{aligned}
P\left(U \leq \frac{f(Y)}{cg(Y)}\right) &= \int_{-\infty}^{\infty} P(cg(Y)U \leq f(Y)|Y = t) g(t)dt \\
&= \int_{-\infty}^{\infty} P\left(U \leq \frac{f(t)}{cg(t)}|Y = t\right) g(t)dt \\
&= \int_{-\infty}^{\infty} P\left(U \leq \frac{f(t)}{cg(t)}\right) g(t)dt \\
&= \int_{-\infty}^{\infty} \frac{f(t)}{cg(t)} \times g(t)dt \\
&= \frac{1}{c} \int_{-\infty}^{\infty} f(t)dt \\
&= \frac{1}{c}.
\end{aligned}$$

Thus, we have

$$\frac{1}{c}P(X \leq x) = \int_{-\infty}^x P\left(U \leq \frac{f(t)}{cg(t)}\right) g(t)dt = \frac{1}{c} \int_{-\infty}^x f(t)dt \implies P(X \leq x) = \int_{-\infty}^x f(t)dt.$$

Hence, using the definition of PDF, it is clear that the PDF of X is f . □

Example 1 (Gamma Distribution). Consider the Gamma distribution with shape parameter $\alpha > 0$ and scale parameter $\beta > 0$. The PDF of the distribution is

$$f(x) = \frac{\beta^\alpha}{\Gamma(\alpha)} x^{\alpha-1} e^{-\beta x} \quad \text{for } x > 0.$$

Note that $X \sim \text{Gamma}(\alpha, 1)$, then $Y = \frac{X}{\beta} \sim \text{Gamma}(\alpha, \beta)$. Therefore, if we can generate from $\text{Gamma}(\alpha, 1)$, we can easily generate from $\text{Gamma}(\alpha, \beta)$, and hence, in this example, we will assume that $\beta = 1$. Thus, our target PDF is

$$f(x) = \frac{1}{\Gamma(\alpha)} x^{\alpha-1} e^{-x} \quad \text{for } x > 0.$$

Now, we will consider three cases based on the values of α .

Case I: $0 < \alpha < 1$

Note that $f(x)$ approaches to ∞ as $x \rightarrow 0$. Also, $f(x) \rightarrow 0$ as $x \rightarrow \infty$. Here, we try to bound f using a integrable function into two parts. For $0 < x < 1$,

$$f(x) \leq \frac{1}{\Gamma(\alpha)} x^{\alpha-1}.$$

For $x \geq 1$,

$$f(x) \leq \frac{1}{\Gamma(\alpha)} e^{-x}.$$

Thus, we take

$$g(x) = \begin{cases} \frac{x^{\alpha-1}}{A} & \text{if } 0 < x < 1 \\ \frac{e^{-x}}{A} & \text{if } x \geq 1, \end{cases}$$

where $A = \frac{1}{\alpha} + \frac{1}{e}$ and $c = \frac{A}{\Gamma(\alpha)}$. Then,

$$f(x) \leq cg(x),$$

where g is a PDF as given above. The corresponding CDF is

$$G(x) = \begin{cases} \frac{x^\alpha}{\alpha A} & \text{if } 0 < x < 1 \\ 1 - \frac{e^{-x}}{A} & \text{if } x \geq 1. \end{cases}$$

Now, using inversion method, we can very easily generate random numbers from g . Note that G^{-1} is given by

$$G^{-1}(u) = \begin{cases} -\ln A - \ln(1 - u) & \text{if } 0 < u < \frac{1}{\alpha A} \\ (\alpha A u)^{1/\alpha} & \text{if } \frac{1}{\alpha A} \leq u < 1. \end{cases}$$

Thus, we have Algorithm 2 to generate random number from $Gamma(\alpha, 1)$ distribution.

Algorithm 2 Generation from $Gamma(\alpha, 1)$ for $0 < \alpha < 1$

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1: repeat
2:   generate  $U_1$  from  $U(0, 1)$ 
3:   Set  $X = G^{-1}(U_1)$ 
4:   generate  $U_2$  from  $U(0, 1)$ 
5: until  $cg(X)U_2 \leq f(X)$ 
6: return  $X$ 

```

Case II: $\alpha \geq 1$ and α is an integer

Notice that $Y_1 + Y_2 + \dots + Y_n \sim Gamma(n, 1)$ if $Y_1, Y_2, \dots, Y_n \stackrel{i.i.d.}{\sim} Exp(1)$. Thus, we will set $n = \alpha$ and have the Algorithm 3.

Algorithm 3 Generation from $Gamma(\alpha, 1)$ when α is a positive integer

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1:  $n = \alpha$ 
2:  $Y = 0$ 
3: while  $n \neq 0$  do
4:   generate  $U$  from  $U(0, 1)$ 
5:   Set  $X = -\ln(U)$ 
6:    $Y = Y + X$ 
7:    $n = n - 1$ 
8: end while
9: return  $Y$ 

```

Case II: $\alpha > 1$ and α is not an integer

Note that we can write $\alpha = \lfloor \alpha \rfloor + \{\alpha\}$, where $\{x\}$ denotes the fractional part of the positive real

number x . Also, $X + Y \sim \text{Gamma}(\alpha_1 + \alpha_2, \beta)$ if $X \sim \text{Gamma}(\alpha_1, \beta)$, $Y \sim \text{Gamma}(\alpha_2, \beta)$, and X and Y are independent. Thus, we have the Algorithm 4.

Algorithm 4 Generation from $\text{Gamma}(\alpha, 1)$ when $\alpha > 1$ and α is not an integer

- 1: generate Y from $\text{Gamma}(\{\alpha\}, 1)$ using Algorithm 2
 - 2: generate X from $\text{Gamma}(\lfloor \alpha \rfloor, 1)$ using Algorithm 3
 - 3: $Z = X + Y$
 - 4: **return** Z
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Example 2 (Beta Distribution). The PDF of Beta distribution with parameters $\alpha_1, \alpha_2 > 0$ is given by

$$f(x) = \frac{1}{B(\alpha_1, \alpha_2)} x^{\alpha_1-1} (1-x)^{\alpha_2-1} \quad 0 < x < 1,$$

where

$$B(\alpha_1, \alpha_2) = \int_0^1 x^{\alpha_1-1} (1-x)^{\alpha_2-1} dx = \frac{\Gamma(\alpha_1)\Gamma(\alpha_2)}{\Gamma(\alpha_1 + \alpha_2)}.$$

Varying the parameters α_1 and α_2 results in a variety of shapes, making this a versatile family of distribution. For example, the case, $\alpha_1 = \alpha_2 = 1/2$ is the arcsine distribution. If $\alpha_1 = \alpha_2 = 1$, we have uniform distribution on $(0, 1)$. If $\alpha_1, \alpha_2 \geq 1$ and at least one of the parameter exceeds 1, the beta density is unimodal and achieves its maximum at $\frac{\alpha_1 - 1}{\alpha_1 + \alpha_2 - 2}$. Let c be the value of the density f at this point, *i.e.*,

$$c = f\left(\frac{\alpha_1 - 1}{\alpha_1 + \alpha_2 - 2}\right).$$

Then $f(x) \leq c$ for all $x \in \mathbb{R}$. For the purpose of acceptance rejection method, we may choose g to be the PDF of uniform distribution over $(0, 1)$, *i.e.*,

$$g(x) = 1 \quad \text{for } 0 < x < 1.$$

Therefore, we can use the Algorithm 5 to generate random numbers from a Beta distribution with parameters $\alpha_1 \geq 1$ and $\alpha_2 \geq 1$, where $\alpha_1 + \alpha_2 > 2$.

Algorithm 5 Generation from $\text{Beta}(\alpha_1, \alpha_2)$ with $\alpha_1 \geq 1$, $\alpha_2 \geq 1$ and $\alpha_1 + \alpha_2 > 2$.

- 1: **repeat**
 - 2: generate U_1 and U_2 from $U(0, 1)$
 - 3: **until** $cU_2 \leq f(U_1)$
 - 4: **return** U_1
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Note that this choice of g will not work if $\alpha_1 < 1$ or $\alpha_2 < 1$. In this case the PDF of Beta distribution is unbounded either at $x = 0$ (for $\alpha_1 < 1$) or at $x = 1$ (for $\alpha_2 < 1$). ||

Example 3 (Beta Distribution). Note that if $X \sim \text{Gamma}(\alpha_1, \beta)$, $Y \sim \text{Gamma}(\alpha_2, \beta)$ and X and Y are independent, then $\frac{X}{X+Y} \sim \text{Beta}(\alpha_1, \alpha_2)$ (Why?). Thus, Algorithm 6 can be used to generate from $\text{Beta}(\alpha_1, \alpha_2)$ distribution.

Algorithm 6 Generation from $\text{Beta}(\alpha_1, \alpha_2)$ distribution

1: generate X from $\text{Gamma}(\alpha_1, \beta)$

2: generate Y from $\text{Gamma}(\alpha_2, \beta)$

3: $Z = \frac{X}{X+Y}$

4: **return** Z

Unlike the previous example, there is no extra restriction on the parameters α_1 and α_2 here. Also, we can take any value for $\beta > 0$ to implement this algorithm. ||