

MODULE 9: Queueing Networks (contd...)

LECTURE 34

Mean-Value Analysis Algorithm

Mean-Value Analysis (MVA) Algorithm

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- Mean-value analysis (MVA) is an approach for calculation of expected value performance measures (mean number and mean sojourn time in each node), without actually computing the normalizing constant $G(N)$ or the state probability distribution.
- The MVA algorithm that recursively computes the performance measures by incrementing the number of customers until the desired number is reached is built on two basic principles:
 - ▶ (Mean value theorem or Arrival theorem) The queue length observed by an arriving customer is the same as the general time queue length in a closed Jackson network with one less customer, i.e., $a_n(N) = p_n(N - 1)$.
 - ▶ Little's formula is applicable throughout the network.
- The algorithm is applicable to systems where the service times at each node are exponentially distributed.
- We also assume FCFS queueing discipline at each node.

- First principle \implies
 - Recall, for $M/M/1$, $W = \frac{1 + L}{\mu}$. That is, the average time an arriving customer must wait is the average time to serve the queue size as seen by an arriving customer plus itself.
 - For $M/M/c$ also no adjustment is needed as L is based on $a_n = p_n$.
 - The equivalent equation for our closed network, assuming that $c_i = 1 \quad \forall i$, is

$$W_i(N) = \frac{1 + L_i(N - 1)}{\mu_i},$$

where

$W_i(N)$ = mean waiting time at node i for a network with N customers

μ_i = mean service rate for the single server at node i

$L_i(N - 1)$ = mean number at node i in a network with $N - 1$ customers

- Second principle \implies

$$L_i(N) = \lambda_i(N)W_i(N).$$

- If we can find $\lambda_i(N)$, then we have a method of calculating L_i and W_i starting with an empty network and building upto one with N customers ($L_i(0) = 0, W_i(1) = \frac{1}{\mu_i}$)

To compute $\lambda_i(N)$:

- Let $D_i(N)$ represent the average delay per customer between successive visits to node i for a network with N customers. Then

$$\lambda_i(N) = \frac{N}{D_i(N)}$$

- To get $D_i(N)$: Take traffic equations, set $\nu_i = \mu_i \rho_i$. Then $\nu_i = \sum_{j=1}^k \nu_j r_{ji}$.

► Set one ν_i to 1, say, $\nu_l = 1$ and solve for others.

► Then ν_i are relative throughput through node i , i.e., $\nu_i = \frac{\lambda_i}{\lambda_l}$.

► We can write $D_l(N) = \sum_{i=1}^k \nu_i W_i(N)$.

Mean-Value Analysis (MVA) Algorithm:

Objective: To find $L_i(N)$ and $W_i(N)$ in a k -node, single-server-per-node network with routing probability matrix $R = (r_{ij})$.

- 1 Solve $\nu_i = \sum_{j=1}^k \nu_j r_{ji}$, $i = 1, 2, \dots, k$, setting one of the ν_i 's (say, ν_l) equal to 1.
- 2 Initialize $L_i(0) = 0$, $i = 1, 2, \dots, k$
- 3 For $n = 1$ to N , calculate
 - a) $W_i(n) = \frac{1 + L_i(n-1)}{\mu_i}$, $i = 1, 2, \dots, k$
 - b) $\lambda_l(n) = \frac{\mu_l}{\sum_{i=1}^k \nu_i W_i(n)}$ (assume $\nu_l = 1$)
 - c) $\lambda_i(n) = \lambda_l(n) \nu_i$, $i = 1, 2, \dots, k$ $i \neq l$
 - d) $L_i(n) = \lambda_i(n) W_i(n)$ $i = 1, 2, \dots, k$

Example (One-machine three-node network)

Consider the previously considered two-machines three-node network, but now with only one machine (to make it as a single-server-at-all-nodes network).

The traffic equations are $(\nu_1, \nu_2, \nu_3) = (\nu_1, \nu_2, \nu_3) \begin{pmatrix} 0 & 3/4 & 1/4 \\ 2/3 & 0 & 1/3 \\ 1 & 0 & 0 \end{pmatrix}$.

Choosing $\nu_2 = 1$ (i.e., $l = 2$ here), we get $\nu_1 = 4/3$ and $\nu_3 = 2/3$.

Now, $i = 1, 2, 3$ and $n = N = 1$.

For step-3(a), $W_1(1) = (1 + L_1(0))/\lambda = 1/\lambda = 1/2$, $W_2(1) = 1/\mu_2 = 1$, $W_3(1) = 1/\mu_3 = 1/3$.

For step-3(b), $\lambda_2(1) = \frac{1}{\sum_{i=1}^3 \nu_i W_i(1)} = \frac{9}{17}$.

For step-3(c), $\lambda_1(1) = \nu_1 \lambda_2(1) = 12/17$, $\lambda_3(1) = \nu_3 \lambda_2(1) = 6/17$.

For step-3(d), $L_1(1) = \lambda_1(1)W_1(1) = 6/17$, $L_2(1) = 9/17$, $L_3(1) = 2/17$.

Since here $N = 1$, we are done.

Verification: Recall that the solution of traffic equations was $\rho_1 = 2/3$, $\rho_2 = 1$, $\rho_3 = 2/9$ and $p_{n_1, n_2, n_3} = (1/G(1))(2/3)^{n_1} (2/9)^{n_3}$. This implies that $G(1) = 17/9$ and $p_{100} = 6/17$, $p_{010} = 9/17$, $p_{001} = 2/17$ which in turn gives the above values of L_i 's and W_i 's.

Mean-Value Analysis (MVA) Algorithm (Multi-Server Case):

Denote the marginal probability of j at node i in an n -customer system by $p_i(j, n)$.

And, set $\alpha_i(j) = \begin{cases} j, & j \leq c_i \\ c_i, & j \geq c_i \end{cases}$.

Objective: To find $L_i(N)$ and $W_i(N)$ in a k -node, c_i -servers-at-node- i network with routing probability matrix $R = (r_{ij})$.

1 Solve $\nu_i = \sum_{j=1}^k \nu_j r_{ji}$, $i = 1, 2, \dots, k$, setting one of the ν_i 's (say, ν_l) equal to 1.

2 Initialize $L_i(0) = 0, p_i(0, 0) = 1, p_i(j, 0) = 0, j \neq 0, i = 1, 2, \dots, k$

3 For $n = 1$ to N , calculate

$$\text{a) } W_i(n) = \frac{1}{c_i \mu_i} \left(1 + L_i(n-1) + \sum_{j=0}^{c_i-2} (c_i - 1 - j) p_i(j, n-1) \right), \quad i = 1, 2, \dots, k$$

$$\text{b) } \lambda_l(n) = \frac{n}{\sum_{i=1}^k \nu_i W_i(n)} \quad (\text{assume } \nu_l = 1)$$

$$\text{c) } \lambda_i(n) = \lambda_l(n) \nu_i, \quad i = 1, 2, \dots, k \quad i \neq l$$

$$\text{d) } L_i(n) = \lambda_i(n) W_i(n), \quad i = 1, 2, \dots, k$$

$$\text{e) } p_i(j, n) = \frac{\lambda_i(n)}{\alpha_i(j) \mu_i} p_i(j-1, n-1), \quad j = 1, 2, \dots, n; \quad i = 1, 2, \dots, k$$

- *Exercise: Take the two-machines three-node closed network example considered earlier and apply the MVA algorithm (multi-server case). You can also carry out a verification of the obtained results.*
- Comparisons (Naive, Convolution and MVA):
 - ▶ All three methods are easy for small problems
 - ▶ Buzen's and MVA are computationally superior (over naive) with respect to efficiency (storage and speed) and stability, for larger problems.
 - ▶ For a single-server-at-all-nodes networks, MVA is superior if we want only the mean waiting times and mean system sizes at each node.
 - ▶ If we have multiserver nodes or if we want marginal probability distributions at the nodes, then it is not clear if MVA is really better than Buzen procedure.