

## Sorting - 2

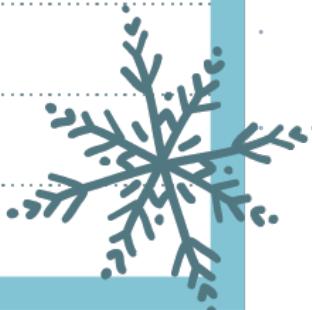


### Agenda :

- Sort 0 & 1
- Pivot Partition
- Quick Sort
- Randomised Quick Sort
- Comparator
- Sorting based on Count of Factors
- Largest Number



Good Morning  
(Starting at 7:05)



## Sort 0 & 1

< Question > : Given an integer array of zeros and ones in random order, segregate 0's on left and 1's on right.

A[ ] = 

0	1	0	0	1	1	0	1	0
---	---	---	---	---	---	---	---	---

o/p: 

0	0	0	0	0	1	1	1	1
---	---	---	---	---	---	---	---	---

### Idea -1

Use Inbuilt Sorting Algo

T.C:  $O(N \log N)$

### Idea -2

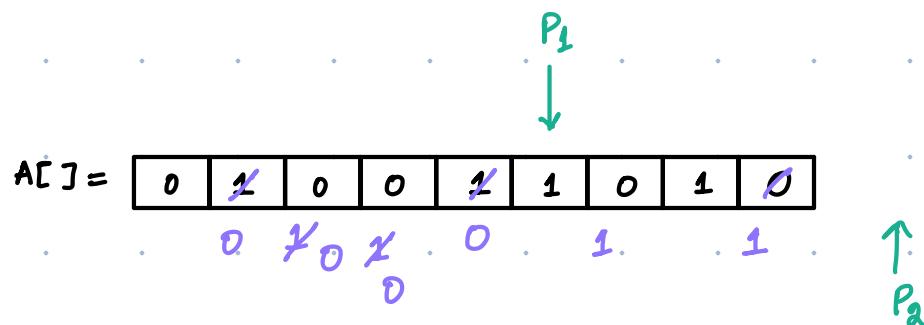
Count Sort

- Iterate over array & count freq. of 0's & 1's.  $\rightarrow O(N)$
- Iterate again & set 0's & 1's.  $\rightarrow O(N)$

T.C:  $O(N)$

**Idea-3**

Two Pointers

 $P_1$  : where 0. should be placed $P_2$  : to iterate over the array..

&lt;/&gt; Code

```
void sort01( int[ ] arr ) {  
    P1 = 0, P2 = 0;  
    while ( P2 < n ) {  
        if ( arr[ P2 ] == 0 ) {  
            Swap( arr[ P1 ], arr[ P2 ] );  
            P1++;  
        }  
        P2++;  
    }  
}
```

T.C: O(N)

S.C: O(1)



## Pivot Partition

**< Question >** : Given an integer array, consider last elements as pivot, rearrange the elements such that for all  $i$  :

- If  $A[i] < p$  then it should be present on left side of array.
- If  $A[i] \geq p$  then it should be present on right side of array.

$1 \leq N \leq 10^5$

$A[] = [54 | 26 | 93 | 17 | 77 | 31 | 44 | 55 | 20]$

$|$   
 $[17 | 54 | 26 | 93 | 77 | 44 | 55 | 31 | 20]$

$|$   
 $[17 | 26 | 44 | 31 | 20 | 93 | 77 | 54 | 55]$

$|$

$A[] = [54 | 26 | 93 | 17 | 77 | 20 | 31 | 55 | 40]$

$[26 | 20 | 31 | 17 | 54 | 93 | 77 | 55 | 40]$

$[20 | 17 | 26 | 31 | 93 | 40 | 54 | 77 | 55]$



Idea -1

Sort the array using inbuilt sorting.

T.C:  $O(N \log N)$



Idea -2 Two Pointers

$P_1$  : where ele < pivot should be placed.

$P_2$  : to iterate over the array.



< / > Code

```
void pivotPartition (int [ ] arr) {
```

Pivot = arr[n-1];

$$P_1 = 0 \text{ , } P_2 = 0 \text{ ; }$$

while. ( $P_2 < n$ ) {

if (arr[p2] < pivot) {

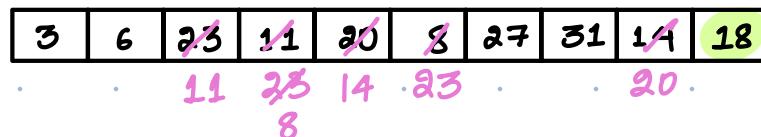
Swap (arr[P<sub>2</sub>], arr[P<sub>1</sub>]);

$P_1 + + \circ$

P<sub>2</sub> + + :

۴

P<sub>1</sub>



Pivot = 18

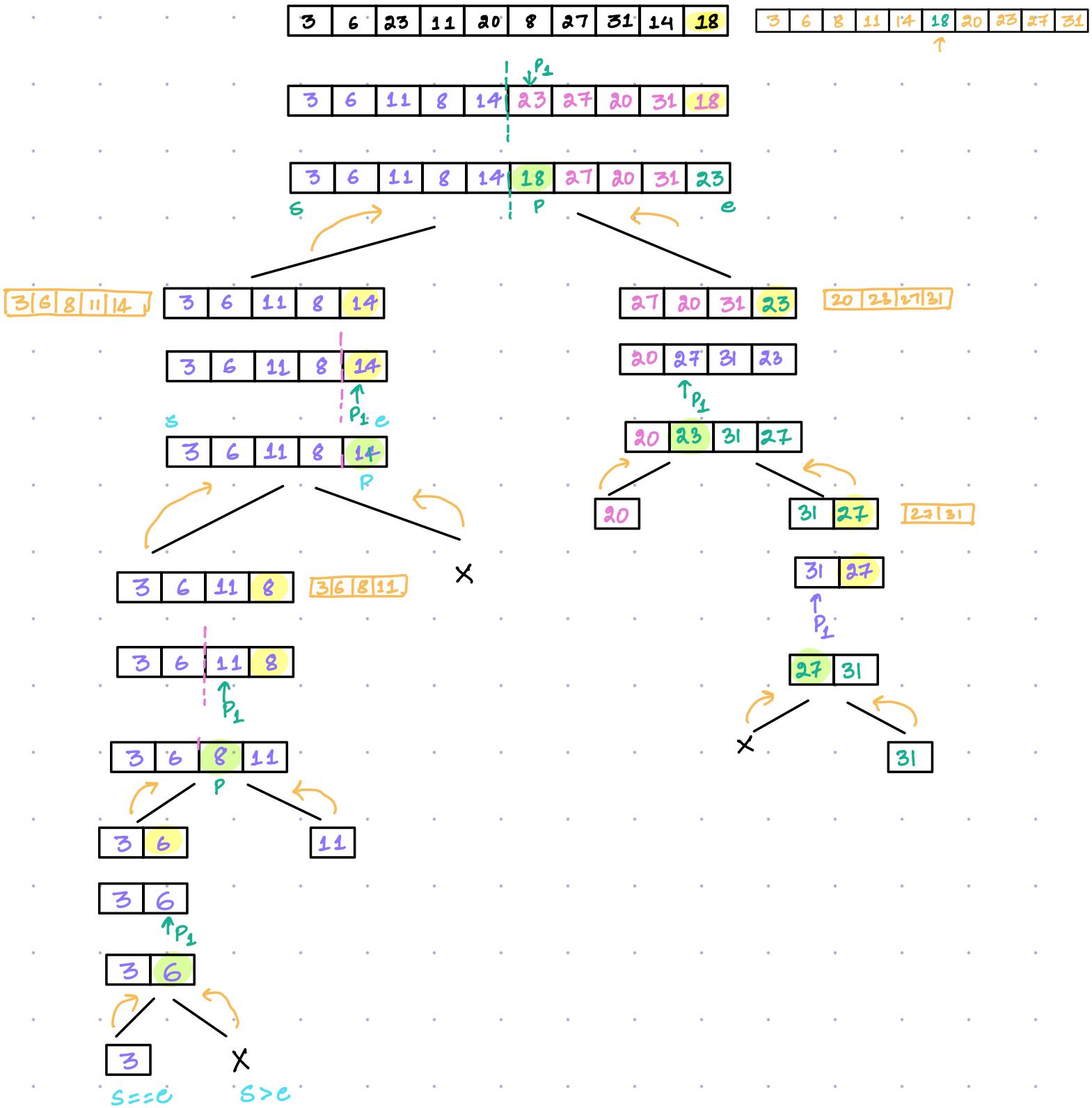
A sequence of numbers from 3 to 18 is shown in a row of boxes. A green dashed vertical line with a green arrow labeled  $P_1$  points to the box containing the number 14.

↑  
P<sub>a</sub>

3	6	11	8	14	18	27	31	20	23
---	---	----	---	----	----	----	----	----	----



# Quick Sort



Approach:

Smaller element on left side of array  
↑ greater or equal elements on right  
side of array

- Choose a pivot & partition the array.
- Bring pivot element to its correct position in the sorted array (swap pivot element &  $\text{arr}[P_i]$ ).
- Repeat the process on the 2 subarrays.

$[S, P-1]$ .

$[P+1, e]$ .



&lt;/&gt; Code

```
void Quicksort (int [ ] A, int s, int e) {  
    if (s > e) return; // Base Case.  
    int pivotIndex = Partition (A, s, e);  
    Quicksort (A, s, pivotIndex - 1);  
    Quicksort (A, pivotIndex + 1, e);  
}
```

```
int Partition (int [ ] arr, int s, int e) {
```

$P_1 = s \Rightarrow P_2 = s$ , pivot = arr[e];

while ( $P_2 \leq e$ ) {

if ( $arr[P_2] < pivot$ ) {

swap (arr[P<sub>1</sub>], arr[P<sub>2</sub>]);

P<sub>1</sub>++;

P<sub>2</sub>++;

}

// Bring pivot element to its correct position

swap (arr[P<sub>1</sub>], arr[e]);

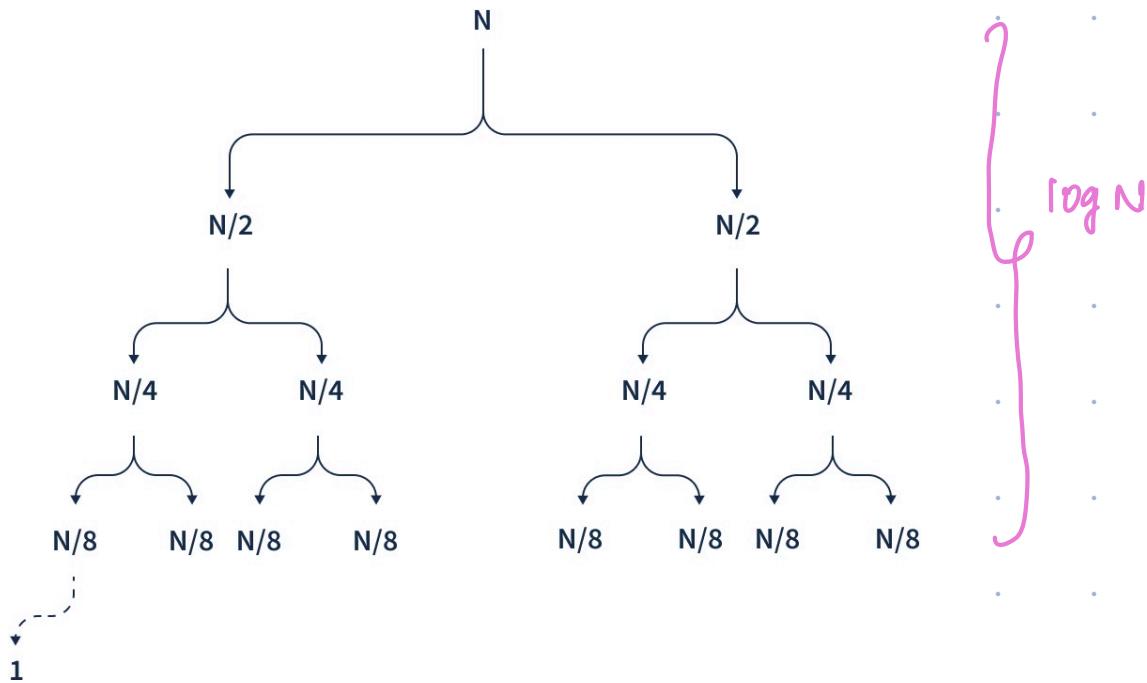
return P<sub>1</sub>;

y

O(N)

## Time Complexity Analysis of Quick Sort

**Best - Case :** When pivot divides the array into 2 equal size subarrays (approx.)



T.C : Total No of calls \* Time per call

$$= \log N * N$$

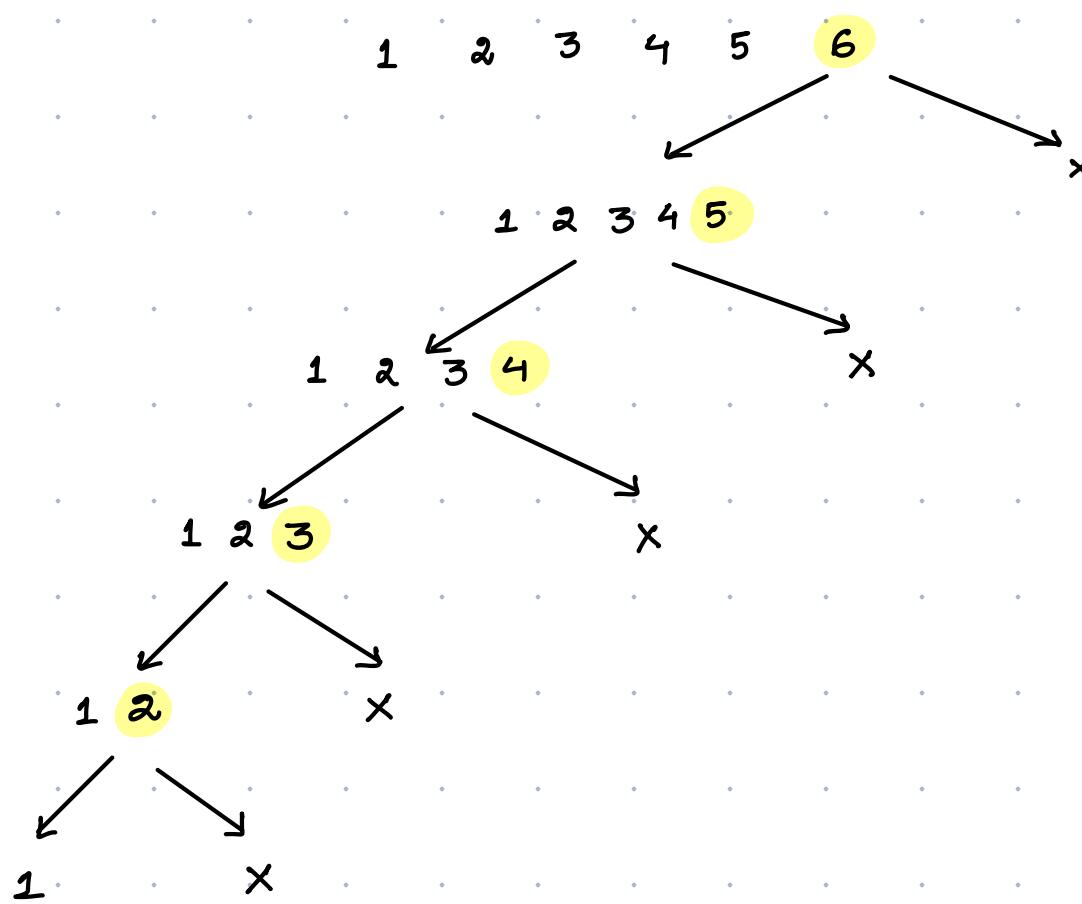
**T.C :  $O(N \log N)$**

S.C : Max Stack Space

**S.C :  $O(\log N)$**



## Worst - Case



T.C : Total no of func calls \* Time per call

$$= N \cdot N$$

T.C :  $O(N^2)$

S.C : Max Stack Space

S.C :  $O(N)$



# Randomised Quick Sort

- Technique where we randomly select the pivot not necessarily 1<sup>st</sup> element or last element.
- Helps to get away with worst time complexity, as probability of always the random element being MIN/MAX is very low.

Given N elements,

- Probability that a random element is minimum :-  $\frac{1}{N}$

- Next time with remaining N-1 elements,

Probability that a random element is minimum :-  $\frac{1}{N-1}$

- Next time with remaining N-2 elements,

Probability that a random element is minimum :-  $\frac{1}{N-2}$

- Next time with remaining N-3 elements,

Probability that a random element is minimum :-  $\frac{1}{N-3}$

$$\frac{1}{N} * \frac{1}{N-1} * \frac{1}{N-2} \cdots \cdots \cdots 1 = \frac{1}{N!}$$

(very small)

Randomised quick sort helps us to achieve avg. T.C of  $O(N \log N)$  most of the times.



**Comparator** - Function that compares 2 values & returns a result indicating whether the values are equal less than or greater than each other.

- Used in Sorting Algorithm to compare elements in a datastructure & arrange them in specified order.

### Java, Python, JS, Ruby

- If we want first argument to come before second : **return -ve**
- If we want second argument to come before first : **return +ve**
- If both are same : **return 0**

### C++

- If we want first argument to come before second : **return true**  
else : **return false**



## Sorting based on Factors

< Question > : Given arr[N], sort the data in ascending order of count of factors.

If count of factors are equal, then sort the elements on the basis of their values.

arr - [ 9 , 3 , 10 , 6 , 4 ]

0 1 2 3 4

$1 \leq N \leq 10^5$

$1 \leq A[i] \leq 10^4$

9 → 1, 3, 9 . . . . . 3

3 → 1, 3 . . . . . 2

Op : [ 3 , 4 , 9 , 6 , 10 ]

10 → 1, 2, 5, 10 . . . . . 4

6 → 1, 2, 3, 6 . . . . . 4

4 → 1, 2, 4 . . . . . 3

QUIZ

10	4	5	13	1
----	---	---	----	---

10 → 1, 2, 5, 10 . . . . . 4

4 → 1, 2, 4 . . . . . 3

Op : [ 1 , 5 , 13 , 4 , 10 ]

5 → 1, 5 . . . . . 2

13 → 1, 13 . . . . . 2

1 → 1 . . . . . 1



&lt;/&gt; Code

C++

```
bool compare(int val1, int val2) {  
    int f1 = count_OF_factors(val1);  
    int f2 = count_OF_factors(val2);  
  
    if (f1 == f2) {  
        if (val1 < val2) // val1 should come before val2  
            return true;  
        else  
            return false;  
    }  
  
    else if (f1 < f2) // val1 should come before val2  
        return true;  
    return false; // val2 should come before val1  
}
```

```
vector<int> FactorBasedSorting (vector<int> A) {  
    sort(A.begin(), A.end(), compare);  
    return A;  
}
```



## Java

```
public ArrayList<Integer> solve(ArrayList<Integer> A) {  
    Collections.sort(A, new Comparator<Integer>(){  
        @Override  
        public int compare(Integer v1, Integer v2){  
            if(factors(v1) == factors(v2)){  
                if(v1 < v2) return -1;  
                else if(v2 < v1) return 1;  
                return 0;  
            }  
            else if(factors(v1) < factors(v2)){  
                return -1;  
            }  
            return 1;  
        }  
    });  
    return A;  
}
```

## Python

```
def compare(v1, v2):  
    if(factors(v1) == factors(v2)):  
        if(v1 < v2):  
            return -1;  
        if(v2 < v1):  
            return 1;  
        else  
            return 0;  
    elif (factors(v1) < factors(v2)):  
        return -1;  
    else  
        return 1;  
  
class Solution:  
    def solve(self, A):  
        A = sorted(A, key = functools.cmp_to_key(compare))  
        return A
```



## Largest Number

Given a list of non-negative integers, arrange them such that they form the largest number and return it. Since the result may be very large, so you need to return a string instead of an integer.

[10, 2]

Op: "210"

[3, 30, 34, 5, 9]

Op: "9534330"

Sort in des order X



[10, 5, 2, 8, 200]

Op: "85220010"



### Approach

X

3

Y

30

String X+Y  
"330"

>

Y+X  
"303"

9

"95"

5

"59"

```
bool compare(int a, int b) {
```

```
    if (to_string(a) + to_string(b) > to_string(b) + to_string(a))  
        return true;  
    else  
        return false;
```

}

```
string largestNumber(int arr[]) {
```

```
    sort(arr.begin(), arr.end(), compare);
```

```
    if (arr[0] == '0') return "0";
```

```
    string ans = " ";
```

```
    for (int i = 0; i < arr.length(); i++) {
```

```
        ans += to_string(arr[i]);
```

y

```
    return ans;
```

y

0, 9, 2, 3, 4 → 9, 4, 3, 2, 0

0, 0, 0, 0, 0 → 0, 0, 0, 0, 0

[3, 30, 34, 5, 9] → [9 5 34 3 30]

ans: "9534330"

3 34 30

33 334 330

















