

CLASS - 13 → 20/09/2023

<https://www.linkedin.com/in/manojofficialmj/>

SEARCHING & SORTING

LEVEL - 03

1. Divide two number using Binary search without using any / and % operator

PROBLEM:

Given two integers one is a dividend and the other is the divisor, we need to find the quotient when the dividend is divided by the divisor without the use of any “/” and “%” operators.

Examples 01:

Input: dividend = 10, divisor = 2
Output: 5
Explanation: $10/2 = 5$

Examples 02:

Input: dividend = 10, divisor = 3
Output: 3
Explanation: $10/3 = 3.3333\dots$ which is truncated to 3

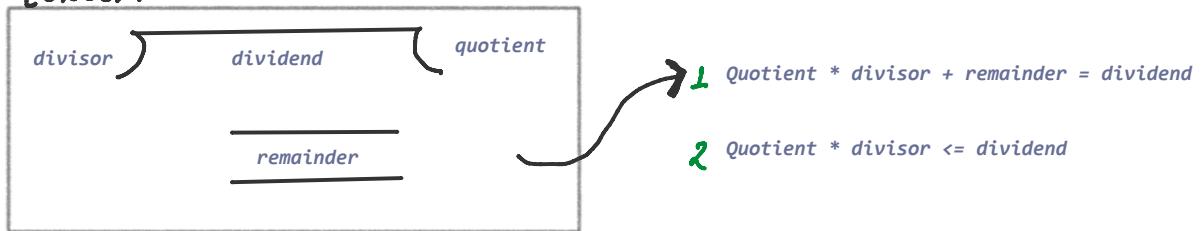
Examples 03:

Input: dividend = 10, divisor = -2
Output: -5
Explanation: $10/-2 = -5$.

Examples 04:

Input: dividend = -10, divisor = 2
Output: -5
Explanation: $-10/2 = -5$.

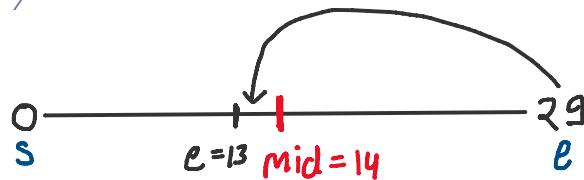
CONCEPT



DRY RUN

DRY RUN THIS EXAMPLE:

Input: dividend = 29, divisor = 7
Output: 4



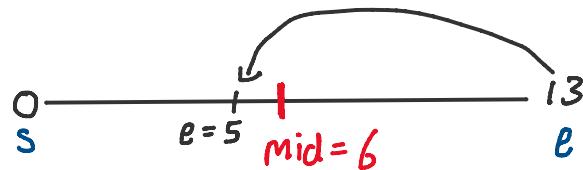
$$S = 0$$

Is 14 a possible ANS?

$$\begin{aligned}
 S &= 0 \\
 e &= 29 \\
 \text{mid} &= \frac{0+29}{2} \\
 &= 14
 \end{aligned}$$

Is 14 a possible ANS?

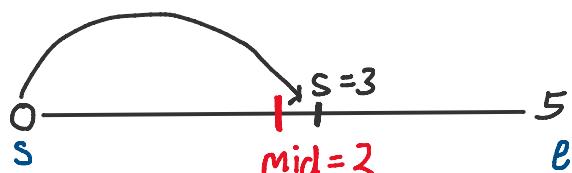
mid * division ≤ 29
 $14 * 7 \Rightarrow 98 \leq 29$ **False**
 ↳ go to left side
 $e = \text{mid} - 1$



$$\begin{aligned}
 S &= 0 \\
 e &= 13 \\
 \text{mid} &= \frac{0+13}{2} \\
 &= 6
 \end{aligned}$$

Is 6 a possible ANS?

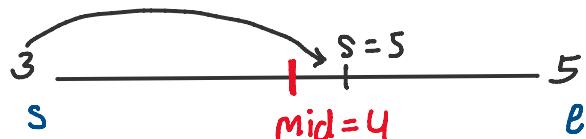
mid * division ≤ 29
 $6 * 7 \Rightarrow 42 \leq 29$ **False**
 ↳ go to left side
 $e = \text{mid} - 1$



$$\begin{aligned}
 S &= 0 \\
 e &= 5 \\
 \text{mid} &= \frac{0+5}{2} \\
 &= 2
 \end{aligned}$$

Is 2 a possible ANS?

mid * division ≤ 29
 $2 * 7 \Rightarrow 14 \leq 29$ **True**
 ↳ stone mid in ANS
 $\text{ANS} = 2$
 ↳ go to right side
 $s = \text{mid} + 1$



$$\begin{aligned}
 S &= 3 \\
 e &= 5 \\
 \text{mid} &= \underline{3+5}
 \end{aligned}$$

Is 4 a possible ANS?

mid * division ≤ 29
 $.. * .. \Rightarrow .. \leq 29$

$$\begin{aligned} e &= 5 \\ \text{mid} &= \frac{3+5}{2} \\ &= 4 \end{aligned}$$

Is 4 a possible ans?
 $4 * 7 \Rightarrow 28 \leq 29$
 True
 stone mid in Ans
 $\text{Ans} = 4$

Go to right side
 $s = \text{mid} + 1$



$$\begin{aligned} s &= 5 \\ e &= 5 \\ \text{mid} &= \frac{5+5}{2} \\ &= 5 \end{aligned}$$

Is 5 a possible Ans?
 $5 * 7 \Rightarrow 35 \leq 29$
 False
 Go to left side
 $e = \text{mid} - 1$

$(s \leq e)$ $s = 5 \leq e = 4$ No RUK jao
END

```

// Program 01: Divide two number using Binary search without using any / and % operator
#include<iostream>
using namespace std;

// Divide two number using Binary search
int getQuotient(int dividend,int divisor){
    int s = 0;
    int e = dividend;
    int mid = s+(e-s)/2;
    int ANS = -1;

    while (s<=e)
    {
        // Return quotient
        if(mid * divisor == dividend){
            return mid;
        }
        // Goto to left side to get possible quotient
        else if(mid * divisor > dividend){
            e = mid - 1;
        }
        // Store possible mid in ANS and goto to right side to get possible quotient
        else if(mid * divisor < dividend){
            ANS = mid;
            s = mid + 1;
        }

        // Updated mid
        mid = s+(e-s)/2;
    }
    return ANS;
}

// Time Complexity: O(logN), where N is dividend.
// Space Complexity: O(1)

int main(){
    int dividend = 29, divisor = -7;

    int quotient = getQuotient(abs(dividend),abs(divisor));

    // Now we need to decide ki sign konsa lagaye +VE ya -VE
    if((dividend > 0 && divisor < 0) || (dividend < 0 && divisor > 0)){
        quotient *= -1;
    }
    cout<<"Quotient: "<<quotient<<endl;
}

```

2. Binary search on nearly sorted array U.U.U.Imp

What is nearly sorted array?

Real World Example:

Lets assume ki ek Ladka hai jiska name SOHAN hai or wo gali no. 10 me rahata hai to hum bol sakte hai ki SOHAN nearly gali no. 9, 10, yaa 11 me se kisi ek gali me rahta hogा.

Nearly Sorted Array

20	10	30	50	40	70	60
0	1	2	3	4	5	6

CATCH: 3

MID-1 >= 0

-1 0 2

ans-11 ✓

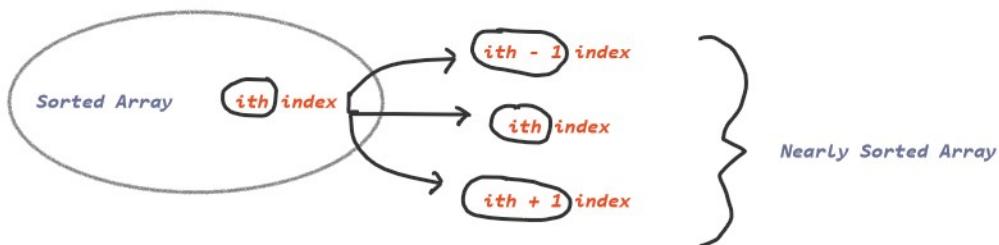
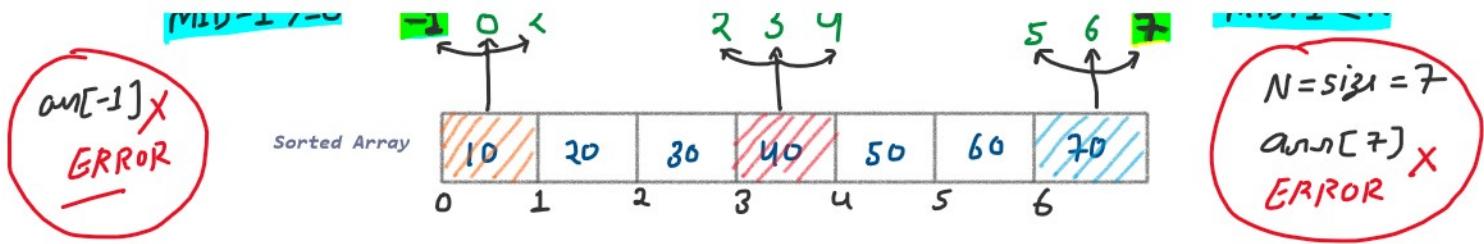
CATCH: 4

MID+1 < N

N = size = 7

2 3 4

5 6 7



Actual Example:

Sorted array me jo 40 value at index 3 par hai to nearly sorted array me 40 value ke three possible index 2, 3, yaa 4 me se kisi ek index par lie zaroor karti hogi.

Normal sorted array	Nearly sorted array
$\rightarrow \text{if } (\text{ans}[mid] == \text{target})$ { return mid; }	$\rightarrow \text{if } (\text{ans}[mid-1] == \text{target})$ CATCH-3 { return mid-1; Mid-1 >= 0 }
$\rightarrow \text{if } (\text{target} > \text{ans}[mid])$ { ↳ Right s = mid+1 }	$\rightarrow \text{if } (\text{ans}[mid] == \text{target})$ { return mid; }
ELSE { ↳ Left e = mid-1 }	$\rightarrow \text{if } (\text{ans}[mid+1] == \text{target})$ CATCH-4 { return mid+1; mid+1 < N } $\rightarrow \text{if } (\text{target} > \text{ans}[mid])$ { ↳ Right s = mid+2 CATCH-1 }

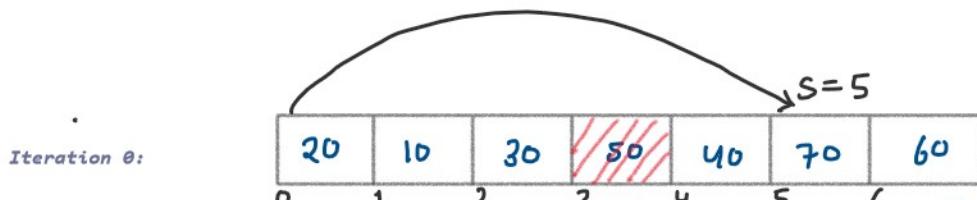
Mid+2 and Mid-1 kyun kar rahe hain?, kyunki hum 1 value ko target se 2 times compare nahi karna chahte hai

DRY RUN

Nearly Sorted Array

arr	20	10	30	50	40	70	60
	0	1	2	3	4	5	6

Target = 70
Output = 5



$$S=0$$

$$e=6$$

$$\text{mid} = \frac{0+6}{2} \\ = 3$$

$$\text{arr}[\text{mid}-1] \Rightarrow 30 == 70 \times$$

$$\text{arr}[\text{mid}] \Rightarrow 50 == 70 \times$$

$$\text{arr}[\text{mid}+1] \Rightarrow 40 == 70 \times$$

$\text{arr}[\text{mid}] < \text{target}$ ✓

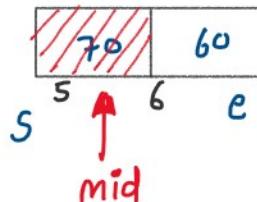
→ Go to right side

$$S = \text{mid} + 1 \times$$

Agar mid+1 karunga to iska matlb yeh hai ki me target se 40 value ko 2 times compare kar rha hu.. CATCH: 01

$$\rightarrow S = \text{mid} + 2 \checkmark$$

Iteration 1:



$$S=5$$

$$e=6$$

$$\text{mid} = \frac{5+6}{2} \\ = 5$$

$$\text{arr}[\text{mid}-1] \Rightarrow 40 == 70 \times$$

$$\text{arr}[\text{mid}] \Rightarrow 5 == 70 \checkmark$$

$$\text{arr}[\text{mid}+1] \Rightarrow 60 == 70 \times$$

→ Return MID END

```
● ● ●
// Program 02: Binary search on nearly sorted array
#include<iostream>
#include<vector>
using namespace std;
```

```

// Program 02: Binary search on nearly sorted array
#include<iostream>
#include<vector>
using namespace std;

// Find target in nearly sorted array
int searchNearlySorted(vector<int> arr,int target){
    int n=arr.size();
    int s=0;
    int e=n-1;
    int mid = s+(e-s)/2;

    while(s<=e){
        // CATCH 03 -> ARR[-1] PAR CODE FAT JAYEGA
        if(mid-1>=0 && arr[mid-1]==target){
            return mid-1;
        }
        if(arr[mid]==target){
            return mid;
        }
        // CATCH 04 -> ARR[N] PAR CODE FAT JAYEGA
        if(mid+1<n && arr[mid+1]==target){
            return mid+1;
        }
        // CATCH 01 -> GOTO RIGHT SIDE TO GET TARGET'S INDEX
        if(arr[mid]<target){
            s=mid+2;
        }
        // CATCH 02 -> GOTO LEFT SIDE TO GET TARGET'S INDEX
        else{
            e=mid-2;
        }
        mid = s+(e-s)/2;
    }
    return -1;
}

/*
Time Complexity: O(logN), where N is array's length
Space Complexity: O(1)
*/
int main(){
    vector<int> arr{20,10,30,50,40,70,60};
    int target = 100;

    int index = searchNearlySorted(arr,target);

    if(index == -1){
        cout<<"Target not found"<<endl;
    }else{
        cout<<"Target "<<target<<" found at index "<<index<<endl;
    }

    return 0;
}

/*
Example 01:
Input: arr[20,10,30,50,40,70,60] , N=7 , and target = 70
Output: Target 70 found at index 5

Example 02:
Input: arr[20,10,30,50,40,70,60] , N=7 , and target = 20
Output: Target 20 found at index 0

Example 03:
Input: arr[20,10,30,50,40,70,60] , N=7 , and target = 60
Output: Target 60 found at index 6

Example 04:
Input: arr[20,10,30,50,40,70,60] , N=7 , and target = 100
Output: Target not found
*/

```

3. Find the Number Occurring Odd Number of Times

Conditions:

1. All elements - even number of times occurs except one
2. All repeating elements - pairs repeat and pairs are not repeated
3. Ek element ek bar me 2 times nahi aa skta hai

→ **ODD Element**

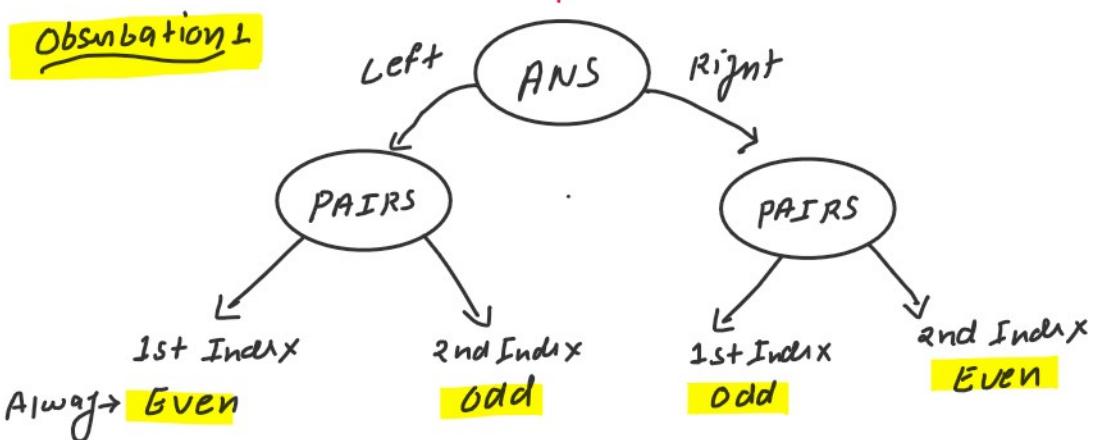
- All elements - even number of times occurs except one
- All repeating elements - pairs repeat and pairs are not repeated
- Ek element ek bar me 2 times nahi aa skta hai

OUTPUT --> Find element that occurs odd time

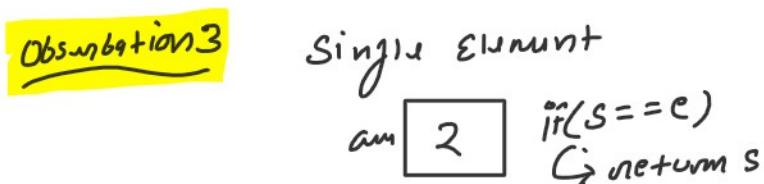
EXAMPLE 01: DRY RUN

am	10	10	2	2	5	5	2	5	5	20	20	11	11	10	10
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

ANS
↓



Observation 2
Ans ⇒ Always lies at Even Index



DRY RUN

Iteration 01

MID	10	10	2	2	5	5	2	5	5	20	20	11	11	10	10
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

$$S=0$$

if ($MID == 2$) { }

$$e=14$$

if ($am[mid] == am[mid-1]$)

$$mid = \frac{0+14}{2}$$

$$\hookrightarrow 5 == 2 \text{ FALSE}$$

$$= 7$$

E1ST

↳ go to left side

$$e = mid - 1$$

Iteration 01

								MID
10	10	2	2	5	5	2		
0	1	2	3	4	5	6		

$$\begin{aligned} s &= 0 \\ e &= 6 \\ \text{mid} &= \frac{0+6}{2} \\ &= 3 \end{aligned}$$

$\text{if } (\text{MID} \% 2 != 0) \{$
 $\quad \text{if } (\text{arr}[\text{mid}] == \text{arr}[\text{mid}-1])$
 $\quad \quad \hookrightarrow 2 == 2 \text{ True}$
 $\quad \quad \text{go to right side}$
 $\quad s = \text{mid} + 1$

Iteration 02

								MID
5	5	2						
4	5	6						

$$\begin{aligned} s &= 4 \\ e &= 6 \\ \text{mid} &= \frac{4+6}{2} \\ &= 5 \end{aligned}$$

$\text{if } (\text{MID} \% 2 != 0) \{$
 $\quad \text{if } (\text{arr}[\text{mid}] == \text{arr}[\text{mid}-1])$
 $\quad \quad \hookrightarrow 5 == 5 \text{ True}$
 $\quad \quad \text{go to right side}$
 $\quad s = \text{mid} + 1$

Iteration 03

2
6

$$\begin{aligned} s &= 6 \\ e &= 6 \end{aligned}$$

$\text{if } (s == e)$
 $\quad \hookrightarrow \text{return } s$ END

Output = 2

DRY RUN

EXAMPLE 02: DRY RUN

1	1	5	5	2	2	3	3	2	4	4
0	1	2	3	4	5	6	7	8	9	10

ANS

Iteration 0

											MID
1	1	5	5	2	2	2	3	3	2	4	4
0	1	2	3	4	5	6	7	8	9	10	

$$s = 0$$

$$e = 10$$

$$\text{mid} = \frac{0+10}{2} \\ = 5$$

if ($\text{mid} \% 2 == 0$) {

if ($\text{arr}[\text{mid}] == \text{arr}[\text{mid}-1]$)

→ 2 == 2 true
go to right side

$$s = \text{mid} + 1$$

Iteration 01

					MID
3	3	2	4	4	
6	7	8	9	10	

$$s = 6$$

$$e = 10$$

$$\text{mid} = \frac{6+10}{2} \\ = 8$$

if ($\text{mid} \% 2 == 0$) {

if ($\text{arr}[\text{mid}] == \text{arr}[\text{mid}+1]$)

→ 2 == 4 False

else

→ store mid in ANS and go to left side

$$\text{ANS} = \text{mid}$$

$$e = \text{mid}$$

END / ANS
END / peak

Iteration 02

			MID
3	3	2	
6	7	8	

$$s = 6$$

$$e = 8$$

$$\text{mid} = \frac{6+8}{2} \\ = 7$$

$$\text{ANS} = 2$$

if ($\text{mid} \% 2 != 0$) {

if ($\text{arr}[\text{mid}] == \text{arr}[\text{mid}-1]$)

3 == 3

→ go to right side

$$s = \text{mid} + 2$$

$$\left. \begin{array}{l} s=9 \\ e=8 \end{array} \right\} - (s \leq e) \quad \text{Ruk jaaao...} \quad \text{OUTPUT} = 2$$

X ↴ Returns Ans = 2 End

```

// Program 03: Find the Number Occurring Odd Number of Times
#include<iostream>
#include<vector>
using namespace std;

// Find number occurring odd number of times
int findOddOccurringElement(vector<int> arr){
    int n=arr.size();
    int s=0;
    int e=n-1;
    int mid=s+(e-s)/2;
    int ans = -1;

    while(s<=e){
        // Single element
        if(s==e){
            return s;
        }
        // when mid index is odd
        if(mid & 1){

            // Goto right side and CATCH: 01 mid-1>=0 because arr[-1] nahi ho skta hai
            if(mid-1>=0 && arr[mid]==arr[mid-1]){
                s=mid+1;
            }
            // Goto left side
            else{
                e=mid-1;
            }
        }
        // when mid index is even
        else{

            // Goto right side to get ans and CATCH: 02 mid+1<n because agar n=7 hai to arr[n] nahi ho skta hai
            if(mid+1<n && arr[mid]==arr[mid+1]){
                // CATCH: 03 we don't want to capare twice for one element
                s=mid+2;
            }
            // CATCH: 04 Store mid in ans and go to left side to get ans (ans/end || peak/ans)
            else{
                ans=mid;
                e=mid;
            }
        }
        // updated mid
        mid=s+(e-s)/2;
    }
    return ans;
}
/*
Time Complexity: O(logN), where N is array's length
Space Complexity: O(1)
*/
int main(){
    vector<int> arr{10,10,2,2,5,5,2,5,5,20,20,11,11,10,10};

    int ansIndex = findOddOccurringElement(arr);

    if(ansIndex == -1){
        cout<<ansIndex<<endl;
    }
    else{
        cout<<"Final ans is "<<arr[ansIndex]<<endl;
    }

    return 0;
}

/*
Example 01:
Input: arr[10,10,2,2,5,5,2,5,5,20,20,11,11,10,10], N=15
Output: Final ans is 2

Example 02:
Input: arr[1,1,5,5,2,2,3,3,2,4,4], N=11
Output: Final ans is 2

Example 03:
Input: arr[1,1,5,5,2,2,3,3,4,4], N=10
Output: -1

Example 04:
Input: arr[1,1,5,5,2,2,3,3,4,4,4], N=10
Output: -1
*/

```