

DSA(SKM)

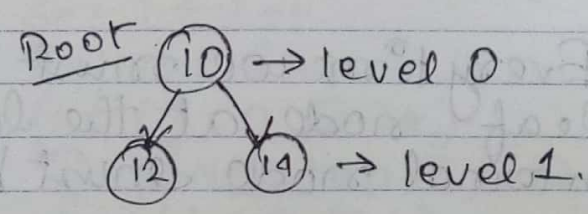
Tree

* non-linear DS.

linear	non linear
① sequential	① non-sequential.
② Only one path is followed.	② Multiple path is followed or possible.

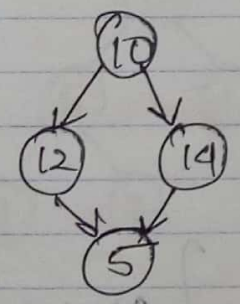
$$T(1) + T(2) + \dots + T(n) = \text{Forest}$$

Tree has to be hierarchical.

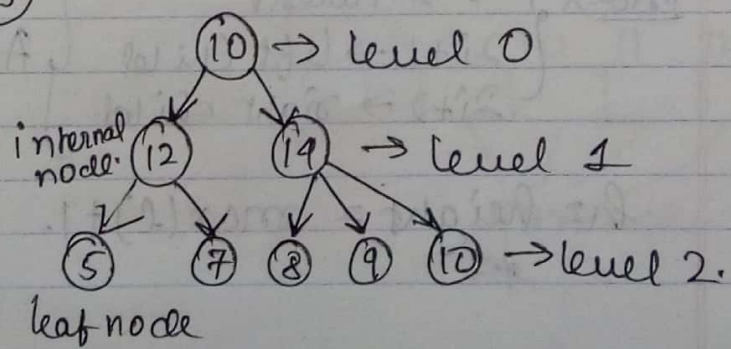
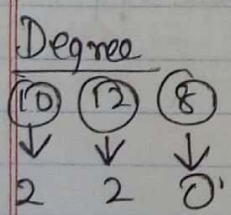


Tree \rightarrow no loop allowed.

Tree is a collection of nodes, connecting with themselves, maintaining a hierarchical form, & no loop



Not a tree, but a graph.



Internal node \rightarrow degree ≥ 1 .
leaf node \rightarrow degree $= 0$.

Siblings \rightarrow whose parents are same.

height of (node) $12 \rightarrow$ Root of that node $\rightarrow 1$.

Depth of $12 \rightarrow 3 \rightarrow$ (node to leaf).
Height of node \rightarrow Root to that node.
Height of tree \rightarrow (max level) $+ 1$.

Binary tree

Total no. of nodes $= 2^n - 1$ (Applicable only for complete binary tree)

Full \rightarrow Every time we must have the leaf node at the last level, & internal node must have degree 2.

Complete \rightarrow Every time we must have the leaf node at the last level, & the leaf node should be filled from left to right.

B T using Array

index $\left\{ \begin{array}{l} i \rightarrow \text{Parent} \\ 2i+1 \rightarrow \text{left child} \\ 2i+2 \rightarrow \text{right child} \end{array} \right\}$ Applicable iff $i \geq 0$.

$h = \text{height} = \max(l) + 1$.

If no child want to insert, it should be -1.

per level we are inserting 2^i elements.

Algorithm:- Create Binary tree.

Step 1: Start.

Step 2: Read n as height of Binary Tree.

Step 3: Print "Enter Root data", set $i \leftarrow 0$.

Step 4: Read $a[i]$ as Root data.

Step 5: if $(a[i] \neq -1)$.

Step 6: Print "Enter the left child of", $a[i]$.

Step 7: Read $a[2i+1]$ as left child.

Step 8: Print "Enter the Right child of", $a[i]$.

Step 9: Read $a[2i+2]$ as Right child.

Step 10: Set $i \leftarrow i+1$.

Step 11: Repeat Step 5 to 11.

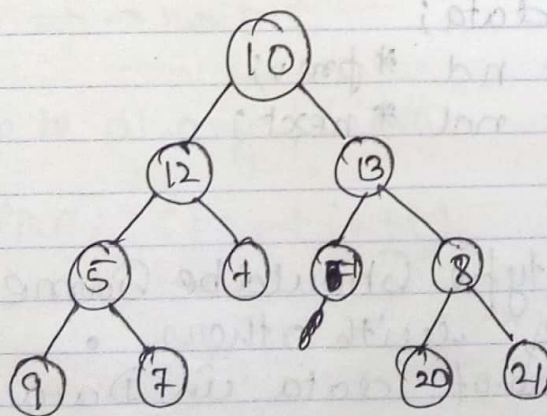
for $(j=0 \text{ to } 2^k)$.

Step 12: Repeat Step 5 to 12.

for $(k=0 \text{ to } h)$.

Step 13: Return a .

Step 14: Stop

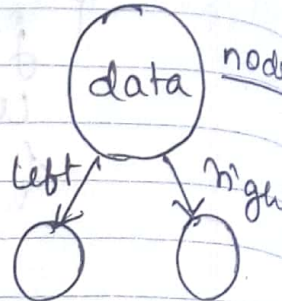


i	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
a	10	12	13	5	-1	-1	8	9	7					20	21

garbage.

Binary Tree Construction using Linked List

```
typedef struct Bnode
{
    int data;
    struct Bnode *left;
    struct Bnode *right;
}
```



```
};
bnode * CreateBinaryTree();
void main()
```

```
{
    bnode * root;
    root = NULL;
    root = CreateBinaryTree();
}
```

```
bnode * CreateBinaryTree()
```

```
{
    int n; bnode * nw;
    printf("Enter node data");
    scanf("%d", &n);
    if (n == -1)
        return NULL;
    nw = (bnode *) malloc(sizeof(bnode));
    nw->data = n;
    printf("Enter the left child of %d", n);
    nw->left = CreateBinaryTree();
    printf("Enter the right child of %d", n);
    nw->right = CreateBinaryTree();
    return nw;
}
```


Topic :-

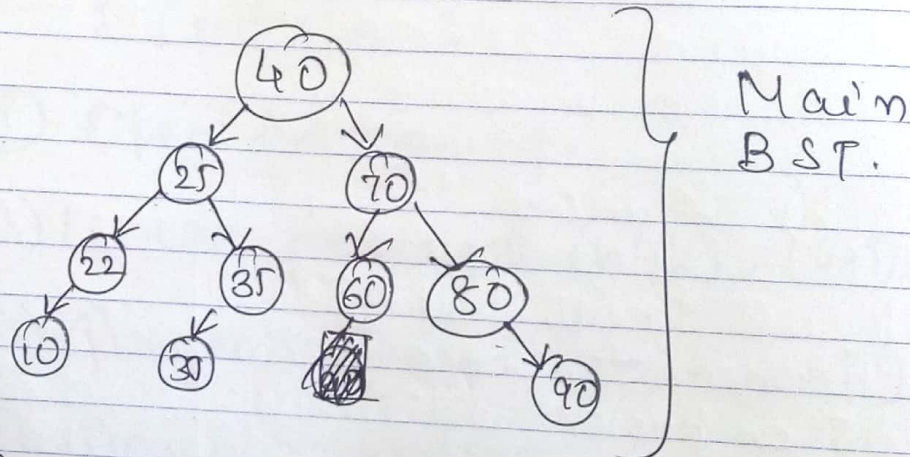
- ① Create a BST.
- ② Delete a node from BST.

Binary Search Tree :- (Order of $\log n$).

(For Insertion) Rule :-

- ① value of left subtree is less than the value of Parent node.
- ② Value of Right subtree is greater than the value of Parent node.

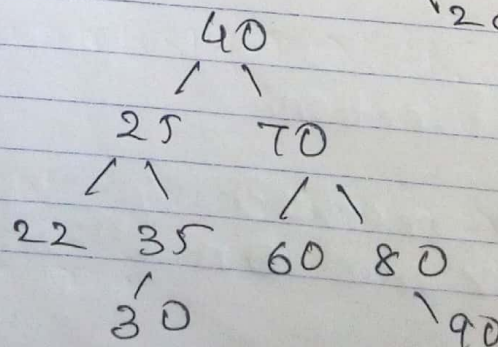
BST using 40, 25, 70, 22, 35, 60,
80, 90, 10, 30.

For Deletion :-

Rule :-

- ① leaf node \rightarrow Remove that.
- ② Non-leaf node \leftarrow 1 child
2 child.

Delete 10

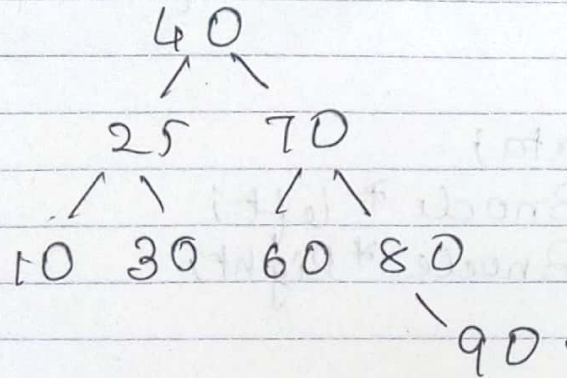


If delete 90, then simply remove 90, as like 10.

Rule → Non leaf

- 1 child → Attach the child to the parent of the to be deleted node.
- 2 child → Find the Inorder successor to be deleted node. Replace with that.

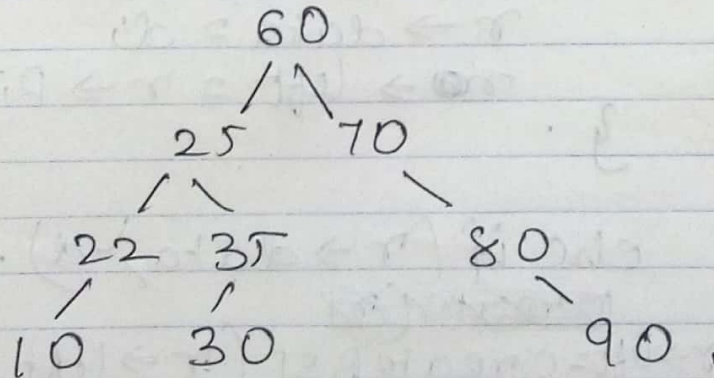
Delete 22



Inorder →

10	22	25	30	35	40	60	70	80	90
----	----	----	----	----	----	----	----	----	----

Delete → 40.

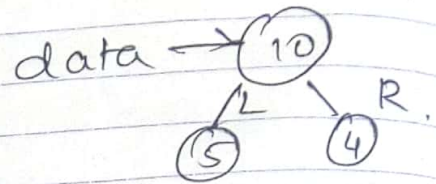


When removing parent node, then the node that will be parent is the smallest number in the right path.

Construct a BST :-

- using recursive
- using non-recursive.

Recursive:-



```
typedef struct Bnode {
```

```
    int data;  
    struct Bnode *left;  
    struct Bnode *right;
```

```
} bnode;
```

~~createBST~~
bnode* createBST(bnode *r, int d).

```
{  
    if (r == NULL).
```

```
{  
    r = (bnode*) malloc(sizeof(bnode));
```

```
    r->data = d;
```

```
    r->left = r->right = NULL;
```

```
}
```

```
else if (r->data > d).
```

~~createBST~~

```
r->left = createBST(r->left, d);
```

```
else
```

```
r->right = createBST(r->right, d);
```

```
return r;
```

```
}
```

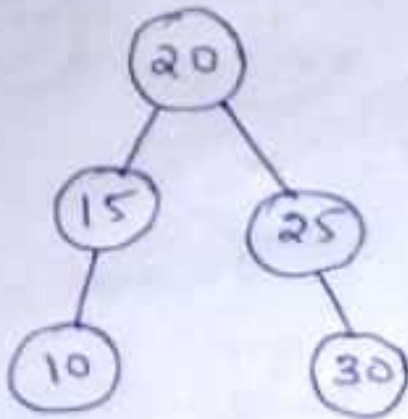


```
void main()
{
    bnode *Root;
    Root = NULL;
    int n;
    do { scanf ("%d", &n);
        Root = CreateBST(Root, n);
        printf ("Do u want to insert?");
        ch = getch();
    } while (ch != 'Y' || ch != 'N')
```

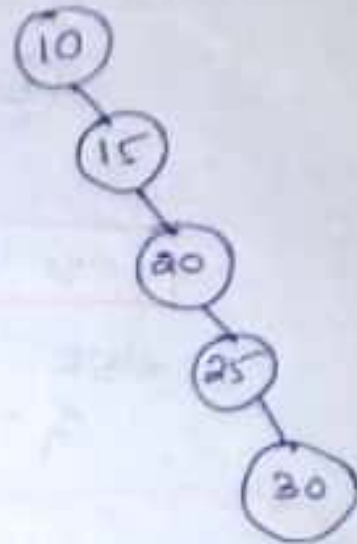

Possible Question:

→ construct a height balance tree.

Height Balanced Tree



Vs



- 1) searching efficiency. [in some case]
- 2) less time required.

3 comparisons
[for shorter
height tree]

5 comparisons
[for longer height
tree.]

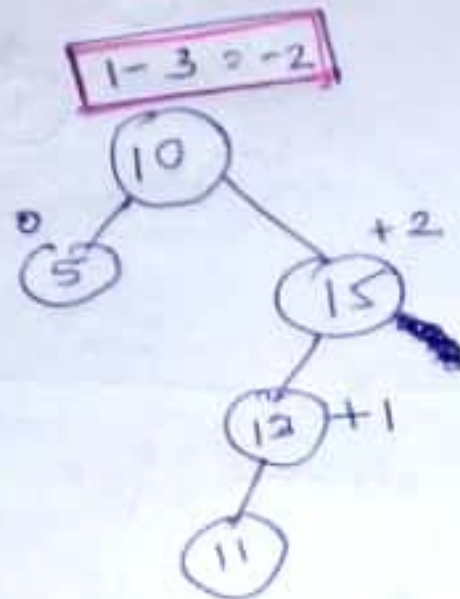
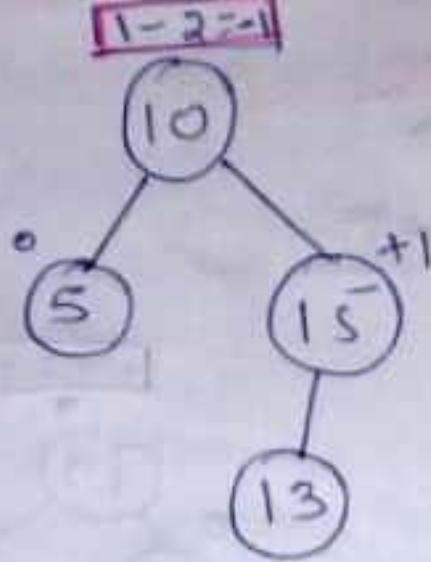
Using: Balance Factor.

Balance Factor of a node.

Height of left sub-tree

Height of right sub-tree.

permissible values of Balance Factor = $(+1, 0, -1) = \text{left} - \text{right}$



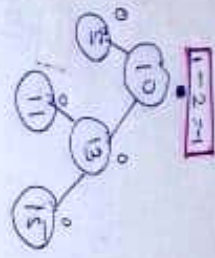
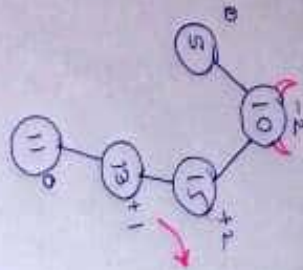
Rotation.

if not permissible balanced factor = then **affected**.

Position of Rotation is the nearest affected node with respect to the last inserted node. due to which the tree has become unbalanced Tree.

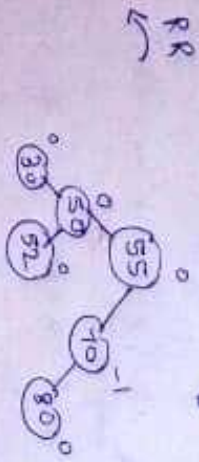
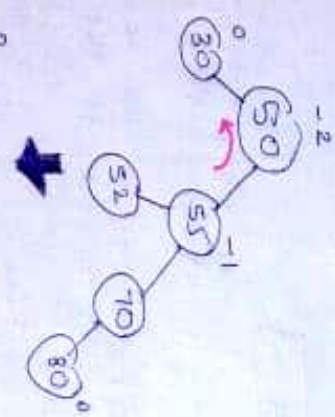
Direction of the inserted node
with respect to →
 the affected node.

Rotation:
 LL →
 from left.

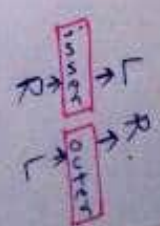
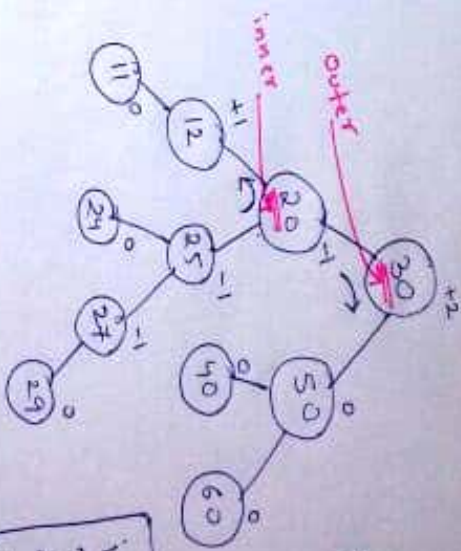
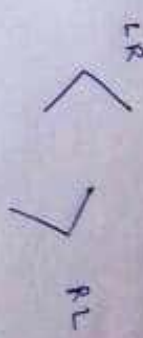


|| Select largest chain for calc ||

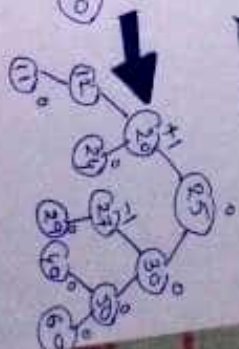
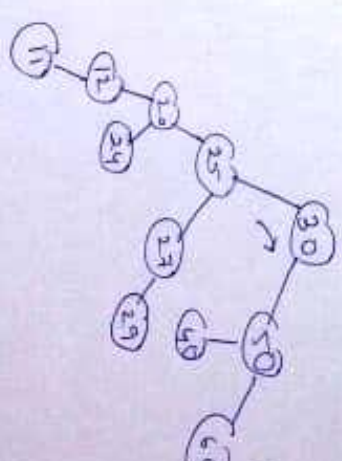
insert 85



→ Double Rotation:



insert mode in the direction change

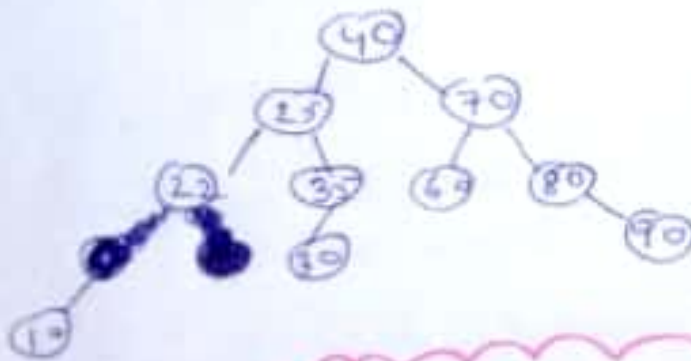


- ①. Create a Binary Search Tree.
- ②. Delete a node from Binary Search Tree.

Rule:-

1. Value of left subtree is less than the value of parent node.
2. Value of Right subtree is greater than value of parent node.

40, 25, 70, 22, 35, 60, 80, 90, 10, 30

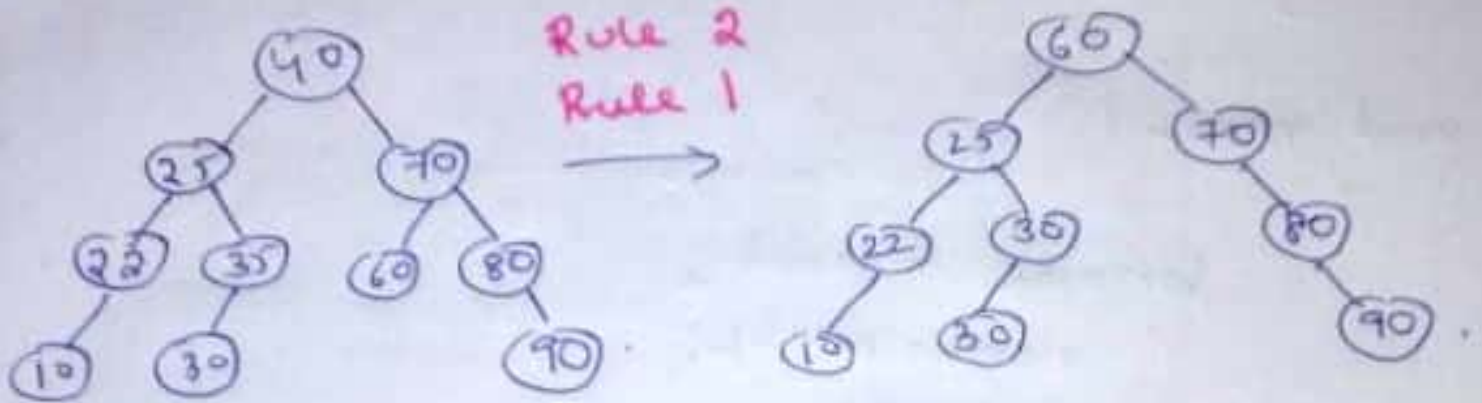


Two types of node

- 1) leaf node → remove that.
- 2) non leaf node → 1 child.
→ 2 children.

1 child \rightarrow Attach the child to the parent of the to be deleted node.

2 child \rightarrow Find the in order successor of to be deleted node and replace with that. (L A/Right)



Construct BST: \rightarrow Recursive
 \rightarrow Non - Recursive

```
typedef struct Bnode {
    int data;
    struct Bnode * left;
    struct Bnode * right;
} bnode;
```

Recursive method:

```
bnode * create BST (bnode * r, int d)
{
    if (r == NULL)
    {
        r = (bnode *) malloc (sizeof (bnode));
        r->data = d;
        r->left = r->right = NULL;
    }
}
```

```

        else
            if (r->data > d)
                r->left = CreateBST(r->left, d);
            else
                r->right = CreateBST(r->right, d);
        return r;
    }

```

```

void main ()

```

```

{

```

```

    bnode * root;

```

```

    root = NULL;

```

```

    int x;

```

```

    do { scanf ("%d", &x);

```

```

        root = CreateBST (root, x);

```

```

        printf (" Do you want to
                insert");

```

```

        ch = getch();

```

```

    } while (ch != 'n' || ch != 'N');

```

```

}

```