

	First()	Follow()
$S \rightarrow cc$	$\{c, d\}$	$\{\$ \}$ <del><math>\{ \}</math></del>
$c \rightarrow c d$	$\{c, d\}$	$\{c, d, \$ \}$

LL(1) Predictive Parsing Table :-

	\$	c	d
S		$S \rightarrow cc$	$S \rightarrow cc$
c		$c \rightarrow c d$	$c \rightarrow d$

There are no multiple entries in the predictive parsing table. So we can say that it is a valid LL(1) parse type grammar.

Also, an E-free LL(1) grammar can be SLR(1) grammar, in turn can also be LALR(1) & LR(1) as well.