Compiler Design RA1911003010143 Abhisher kumar 4ssignments * Machine Independent optimization A deep optimization -> To apply loop optzynation, we must fourt detect loops. Jor detecting loops we canhol flow analysis (cfn) using program flow graph (PF4) To find If 4 we need to find basic blocks. Optimization HLL IGG GFA (1FG) 5,13 C.0 Note: for loop is not visible at IGG. At Ily even if, Switch Stalement have gots! hence, -> Gele indirect a loup BBLJ

I finding the basic blocks find the leaders in the program then a basic block will Start from one leader to the next leader. 1 Identifyry leader in basic block: 1) of Statement is a leader 2) Statement that in target of Conductional ex. y(...) gato (Stalement) gato (stalent () 5 luder 3) Statement that follows immedially a Conditional or unconditional statement is leader Convert to Baddress Code -> ex: Code: 0 1 = 1 4 B, 6 i=2 fact (x) for (i=2;ix=n;i++) B 14 (17x)9060(9) 7BL O t=f*(i-) S = ti; | 133. return of, (9) golo (3)

D goto (return) & By

optimizations on leap (a) frequency reduction morning the code from high frequency to low frequency region is Called Code motion. ex! -A = Sin(u)/con(u) +i'; = Sins (h)/cosin); (oundy once calcular while (1'(5000) unralling: Taking lesser iterations loop while (110) うり(世) -0 N[1+] = 0;) combining p hence reducing

for (1=0 ; i210 xi+) for (J=0; I(10, It+) n Eight Fog for (i = 0; icto; i++) for (J=0; -J<10; J++) 1 2[1,1] = 0; x[i,]] =0; obtimized above sols, but actually only three laddress Code are used at this Stage folding: replacing an expression that can be computed at compile time by rts value Redundancy elimination: - (DAG - directed acycle graph) graph)

operation by Cheaper one.

(x. B = 4 × 2 3 b = ACC)

A = BtC

D = 2+ 0+3+C

A = A+O & Demenate Trina

n = n + 1 | Shewt.