

# Parallel Sorting Algorithms

- 1 Sorting in C and C++
  - using `qsort` in C
  - using STL `sort` in C++
- 2 Bucket Sort for Distributed Memory
  - bucket sort in parallel
  - communication versus computation
- 3 Quicksort for Shared Memory
  - partitioning numbers
  - quicksort with OpenMP
  - parallel sort with Intel TBB

MCS 572 Lecture 12  
Introduction to Supercomputing  
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## using `qsort`

C provides an implementation of quicksort. The prototype is

```
void qsort ( void *base, size_t count, size_t size,  
            int (*compar)(const void *element1,  
                          const void *element2) );
```

`qsort` sorts an array whose first element is pointed to by `base` and contains `count` elements, of the given `size`.

The function `compar` returns

- `-1` if `element1 < element2`,
- `0` if `element1 = element2`,
- `+1` if `element1 > element2`.

We will apply `qsort` to sort a random sequence of doubles.

# generating and writing numbers

```
void random_numbers ( int n, double a[n] )
{
    int i;
    for(i=0; i<n; i++)
        a[i] = ((double) rand())/RAND_MAX;
}
```

```
void write_numbers ( int n, double a[n] )
{
    int i;
    for(i=0; i<n; i++) printf("%.15e\n", a[i]);
}
```

## using qsort

```
int compare ( const void *e1, const void *e2 )
{
    double *i1 = (double*)e1;
    double *i2 = (double*)e2;
    return ((*i1 < *i2) ? -1 : (*i1 > *i2) ? +1 : 0);
}
```

in the function `main()`:

```
double *a;
a = (double*)calloc(n, sizeof(double));
random_numbers(n, a);

qsort((void*)a, (size_t)n, sizeof(double), compare);
```

## code to time `qsort`

We use the command line to enter the dimension and to toggle off the output.

To measure the CPU time for sorting:

```
clock_t tstart, tstop;
tstart = clock();
qsort((void*)a, (size_t)n, sizeof(double), compare);
tstop = clock();
printf("time elapsed : %.4lf seconds\n",
      (tstop - tstart)/((double) CLOCKS_PER_SEC));
```

## timing `qsort` on 3.47GHz Intel Xeon

```
$ time /tmp/time_qsort 1000000 0
time elapsed : 0.2100 seconds
real        0m0.231s
user        0m0.225s
sys         0m0.006s
$ time /tmp/time_qsort 10000000 0
time elapsed : 2.5700 seconds
real        0m2.683s
user        0m2.650s
sys         0m0.033s
$ time /tmp/time_qsort 100000000 0
time elapsed : 29.5600 seconds
real        0m30.641s
user        0m30.409s
sys         0m0.226s
```

Observe:  $O(n \log_2(n))$  is almost linear in  $n$ .

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## using the STL container `vector`

```
#include <vector>
using namespace std;

vector<double> random_vector ( int n );
// returns a vector of n random doubles

vector<double> random_vector ( int n )
{
    vector<double> v;
    for(int i=0; i<n; i++)
    {
        double r = (double) rand();
        r = r/RAND_MAX;
        v.push_back(r);
    }
    return v;
}
```

## writing STL vectors

```
#include <iostream>
#include <iomanip>
#include <vector>
using namespace std;

void write_vector ( vector<double> v );
// writes the vector v

void write_vector ( vector<double> v )
{
    for(int i=0; i<v.size(); i++)
        cout << scientific
            << setprecision(15)
            << v[i] << endl;
}
```

## using the STL sort

```
#include <vector>
#include <algorithm>
using namespace std;

struct less_than // defines "<"
{
    bool operator()(const double& a,
                    const double& b)
    {
        return (a < b);
    }
};
```

in the main program:

```
sort(v.begin(), v.end(), less_than());
```

## timing STL `sort` on 3.47GHz Intel Xeon

```
$ time /tmp/time_stl_sort 1000000 0
time elapsed : 0.36 seconds
real        0m0.376s
user        0m0.371s
sys         0m0.004s
$ time /tmp/time_stl_sort 10000000 0
time elapsed : 4.09 seconds
real        0m4.309s
user        0m4.275s
sys         0m0.033s
$ time /tmp/time_stl_sort 100000000 0
time elapsed : 46.5 seconds
real        0m48.610s
user        0m48.336s
sys         0m0.267s
```

Different distributions may cause timings to fluctuate.

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# bucket sort

Given are  $n$  numbers, suppose all are in  $[0, 1]$ .

The algorithm using  $p$  buckets proceeds in two steps:

- Partition numbers  $x$  into  $p$  buckets:  
 $x \in [i/p, (i+1)/p[ \Rightarrow x \in (i+1)\text{th bucket.}$
- Sort all  $p$  buckets.

The cost to partition the numbers into  $p$  buckets is  $O(n \log_2(p))$ .

Note: radix sort uses most significant bits to partition.

In the best case: every bucket contains  $n/p$  numbers.

The cost of Quicksort is  $O(n/p \log_2(n/p))$  per bucket.

Sorting  $p$  buckets takes  $O(n \log_2(n/p))$ .

Total cost is  $O(n(\log_2(p) + \log_2(n/p)))$ .

# parallel bucket sort

On  $p$  processors, all nodes sort:

- 1 Root node distributes numbers: processor  $i$  gets  $i$ th bucket.
- 2 Processor  $i$  sorts  $i$ th bucket.
- 3 Root node collects sorted buckets from processors.

Is it worth it? Recall: serial cost is  $n(\log_2(p) + \log_2(n/p))$ .

Cost of parallel algorithm:

- $n \log_2(p)$  to place numbers into buckets,
- $n/p \log_2(n/p)$  to sort buckets.

$$\begin{aligned} \text{speedup} &= \frac{n(\log_2(p) + \log_2(n/p))}{n(\log_2(p) + \log_2(n/p)/p)} \\ &= \frac{1 + L}{1 + L/p} = \frac{1 + L}{(p + L)/p} = \frac{p}{p + L}(1 + L), \quad L = \frac{\log_2(n/p)}{\log_2(p)}. \end{aligned}$$

## comparing to quicksort

$$\begin{aligned}\text{speedup} &= \frac{n \log_2(n)}{n(\log_2(p) + n/p \log_2(n/p))} \\ &= \frac{\log_2(n)}{\log_2(p) + 1/p(\log_2(n) - \log_2(p))} \\ &= \frac{\log_2(n)}{1/p(\log_2(n) + (1 - 1/p) \log_2(p))}\end{aligned}$$

Example:  $n = 2^{20}$ ,  $\log_2(n) = 20$ ,  $p = 2^2$ ,  $\log_2(p) = 2$ ,

$$\begin{aligned}\text{speedup} &= \frac{20}{1/4(20) + (1 - 1/4)2} \\ &= \frac{20}{5 + 3/2} = \frac{40}{13} \approx 3.08.\end{aligned}$$



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# communication and computation

The scatter of  $n$  data elements costs  $t_{\text{start up}} + nt_{\text{data}}$ , where  $t_{\text{data}}$  is the cost of sending 1 data element.

For distributing and collecting of all buckets, the total communication time is  $2p \left( t_{\text{start up}} + \frac{n}{p} t_{\text{data}} \right)$ .

The computation/communication ratio is

$$\frac{(n \log_2(p) + n/p \log_2(n/p)) t_{\text{compare}}}{2p \left( t_{\text{start up}} + \frac{n}{p} t_{\text{data}} \right)}$$

where  $t_{\text{compare}}$  is the cost for one comparison.

# the computation/communication ratio

The computation/communication ratio is

$$\frac{(n \log_2(p) + n/p \log_2(n/p)) t_{\text{compare}}}{2p \left( t_{\text{start up}} + \frac{n}{p} t_{\text{data}} \right)}$$

where  $t_{\text{compare}}$  is the cost for one comparison.

We view this ratio for  $n \gg p$ , for fixed  $p$ , so:

$$\frac{n}{p} \log_2 \left( \frac{n}{p} \right) = \frac{n}{p} (\log_2(n) - \log_2(p)) \approx \frac{n}{p} \log_2(n).$$

The ratio then becomes  $\frac{n}{p} \log_2(n) t_{\text{compare}} \gg 2n t_{\text{data}}$ .

Thus  $\log_2(n)$  must be sufficiently high...

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## a recursive algorithm

```
void quicksort ( double *v, int start, int end ) {  
    if(start < end) {  
        int pivot;  
        partition(v, start, end, &pivot);  
        quicksort(v, start, pivot-1);  
        quicksort(v, pivot+1, end);  
    }  
}
```

where partition has the prototype:

```
void partition  
( double *v, int lower, int upper, int *pivot );  
/* precondition: upper - lower > 0  
 * takes v[lower] as pivot and interchanges elements:  
 * v[i] <= v[pivot] for all i < pivot, and  
 * v[i] > v[pivot] for all i > pivot,  
 * where lower <= pivot <= upper. */
```

## a partition function

```
void partition
( double *v, int lower, int upper, int *pivot )
{
    double x = v[lower];
    int up = lower+1;    /* index will go up */
    int down = upper;    /* index will go down */
    while(up < down)
    {
        while((up < down) && (v[up] <= x)) up++;
        while((up < down) && (v[down] > x)) down--;
        if(up == down) break;
        double tmp = v[up];
        v[up] = v[down]; v[down] = tmp;
    }
    if(v[up] > x) up--;
    v[lower] = v[up]; v[up] = x;
    *pivot = up;
}
```

## partition and qsort in main()

```
int lower = 0;
int upper = n-1;
int pivot = 0;
if(n > 1) partition(v, lower, upper, &pivot);

if(pivot != 0)
    qsort((void*)v, (size_t)pivot,
          sizeof(double), compare);

if(pivot != n)
    qsort((void*)&v[pivot+1], (size_t)(n-pivot-1),
          sizeof(double), compare);
```

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## a parallel region in `main()`

```
omp_set_num_threads(2);  
#pragma omp parallel  
{  
    if(pivot != 0)  
        qsort((void*)v, (size_t)pivot,  
              sizeof(double), compare);  
    if(pivot != n)  
        qsort((void*)&v[pivot+1], (size_t)(n-pivot-1),  
              sizeof(double), compare);  
}
```

## on dual core Mac OS X at 2.26 GHz

```
$ time /tmp/time_qsort 10000000 0  
time elapsed : 4.0575 seconds
```

```
real      0m4.299s  
user      0m4.229s  
sys       0m0.068s
```

```
$ time /tmp/part_qsort_omp 10000000 0  
pivot = 4721964
```

```
-> sorting the first half : 4721964 numbers  
-> sorting the second half : 5278035 numbers
```

```
real      0m3.794s  
user      0m7.117s  
sys       0m0.066s
```

Speed up:  $4.299/3.794 = 1.133$ , or 13.3% faster with one extra core.

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## using `parallel_sort` of the Intel TBB

At the top of the program, add the lines

```
#include "tbb/parallel_sort.h"
```

```
using namespace tbb;
```

To sort a number of random doubles:

```
int n;  
double *v;  
v = (double*)calloc(n, sizeof(double));  
random_numbers(n, v);  
parallel_sort(v, v+n);
```

## an interactive test run

```
$ /tmp/tbb_sort 4 1
4 random numbers :
3.696845319912231e-01
7.545582678888730e-01
6.707372915329120e-01
3.402865237278335e-01
the sorted numbers :
3.402865237278335e-01
3.696845319912231e-01
6.707372915329120e-01
7.545582678888730e-01
$
```

## timing parallel runs

```
$ time /tmp/tbb_sort 10000000 0
```

```
real      0m0.479s
```

```
user      0m4.605s
```

```
sys       0m0.168s
```

```
$ time /tmp/tbb_sort 100000000 0
```

```
real      0m4.734s
```

```
user      0m51.063s
```

```
sys       0m0.386s
```

```
$ time /tmp/tbb_sort 1000000000 0
```

```
real      0m47.400s
```

```
user      9m32.713s
```

```
sys       0m2.073s
```

```
$
```

## recommended reading

- Edgar Solomonik and Laxmikant V. Kale: **Highly Scalable Parallel Sorting**. In the proceedings of the IEEE International Parallel and Distributed Processing Symposium (IPDPS), 2010.
- Mirko Rahn, Peter Sanders, and Johannes Singler: **Scalable Distributed-Memory External Sorting**. In the proceedings of the 26th IEEE International Conference on Data Engineering (ICDE), pages 685-688, IEEE, 2010.
- Davide Pasetto and Albert Akhriev: **A Comparative Study of Parallel Sort Algorithms**. In SPLASH'11, the proceedings of the ACM international conference companion on object oriented programming systems languages and applications, pages 203-204, ACM 2011.

# Summary + Exercises

In the book of Wilkinson and Allen, bucket sort is described in §4.2.1 and chapter 10 is entirely devoted to sorting algorithms.

## Exercises:

- 1 Consider the fan out scatter and fan in gather operations and investigate how these operations will reduce the communication cost and improve the computation/communication ratio in bucket sort of  $n$  numbers on  $p$  processors.
- 2 Instead of OpenMP, use Pthreads to run Quicksort on two cores.
- 3 Instead of OpenMP, use the Intel Threading Building Blocks to run Quicksort on two cores.