Finding the keys for a Relation R, from FDs

* Any combination that uniquely identify all attributes of a relation is a key or sapukey. (SK)

*Minimal Superkey (SK), whose no proper subset can't be used for unique identification of all attributes is a cardidate key (CK).

Process!

- 1. From given FDs, findout which attribute is not identified by any other attributes ar we can say the attribute which doesn't appear on the R.H.S of any FD. If found, they are essentials for CK.
- 9. If any such attribute(X) exist, then it must be included in Ck. say such attributes are X. (X can be single as multiple) I check whether X+ can include all the attributes of R. If Yes then it means that we can identify all attributes of R with X, and its a candidate key (Ck). Any other cambination containly X as subset can be a Sk but can't be Ck.
 - JJ X alone can't identify all attributes. Next step is to try all possible combinations containing X as subset, starting from the minimum length. If any such combination found include it in set of CK, and try other combinations.
- 3. If no such attributes found from step 1, then try all combinations of attributes that can identify all attributes of relation R. starting from length ane.

NOTE: If any attribute or set of attributes is selected as CK, then no other CK can include existing CKs as subset.

(2)

FDs= {A+BC, CD+E, E+C, D+AEH, ABH+BD, DH+BC}

* As all the attributes are identifiable, so try combinations -

At = ABC
Bt = B

$$c^+ = c$$

 $D^+ = DAEHBC = R$
 $E^+ = EC$
 $H^+ = H$ $\neq R$
 $\neq R$
 $\neq R$

* Try other combinations with rest of attributes -(without including D as its a CK)

$$(AB)^{+} = ABC$$
 $(AC)^{+} = ABC$
 $(AE)^{+} = AEBC$
 $\neq R$
 $(AE)^{+} = AEBC$
 $\Rightarrow R$
 $(AE)^{+} = ABCDE = R$
 $\Rightarrow R$
 \Rightarrow

* Try another Combinations not Containing D, AH as subsets-

$$(ABC)^{\dagger} = ABC$$

$$(ACE)^{\dagger} = ACEB$$

$$(HBC)^{\dagger} = HBC$$

$$(HCE)^{\dagger} = HCE$$

$$(BCE)^{\dagger} = BCE$$

$$(ABE)^{\dagger} = ABEC$$

* Try another combinations !
(ABCE) += ABCE | +R

(HBCE) += HBCE | +R

O ROB, CTOEF)

A and E are not identified by any other attributes. So they must be included in every CK.

(AE) = AECFDB =R

SO (AB) alove can identify all other attributes. So its the only CK for R. Any other key combination must include AE as subset and they can be SK but count be CK.

Q R(A,B,C,D) E,F,G,H) [GATE Question]

FDs={CH+G, A+BC, B+CFH, E+A, F+EG}

Dt=D //so D alove con't identify R

Try combinations - (DA), (DB), DC, (DE) (DE) DG, DH

(DA)+= (DE)+= & (DB)+=(DF)+= R

 $(DG)^{\dagger}, (DG)^{\dagger}, (DH)^{\dagger} \neq R$

So set of Ck includes !- [DA, DE]

As any other combination can't be formed without including DA or DE.

3) R (W, X, Y, Z) [GATE Question]

FDs = {Z -> W, Y -> XZ, WX -> Y}

As all attributes are identifiable. So try Combinations !-

 $X^{+} = X \neq R$

w+= w ≠R

any other combination of length 3 or more can't be journed without including 4, WX or XZ.

| Candidate keys = {Y, WX, XZ}

Ty+ = YXZW = R/W Try other combinations not including y :- options of length 21- $Z^{+}=ZW \neq R \quad (WX, WZ, XZ)$

(WX) = WXYZ=RV

 $(wz)^{\dagger} = wz \neq R$

(XZ)+= XZWY =RL

Finding keys from given FDs (4) R(A,B,C,B,E)

FDs = {AB -> CD, D -> A, BC -> DE }

As B is not identified by any attributes- $B^{+}=B \neq R$, B alone coult be a CK

for R

Try combinations: BA, BC, BD, BE $(AB)^{\pm}(BC)^{\dagger} = (BD)^{\dagger} = ABCDE = R$ But $(BE)^{\dagger} = BE \neq R$

So, CK = { AB, BC, BD3

3 R(A,B,C,D,E,F,G,H)

FDs={A>C, A>DE, B>F, F>GH}
Taking A and B as they are unidentified!
(AB)+= ABCDEFGH=R

So (AB) is the only CK.

@ R (A,B) S, (D) E, E, S, H)

FD= [AB - C, BD - EF, AD - G, A - H]

Taking A, B and Di-

(ABD) = ABDCEFGH=R

So (ABD) is the only CK

(7) R(A, K, Q, Ď, Ě) [Gate Question] CK= {BC, CD}

FDs = {BC -> ADE, D -> B}

Taking C, c+= C ≠ R

So try combinations 1- (AC, BC) (D) CE)

(BC) = (CD) = BCADE = R

 $(AC)^{+}, (CE)^{+} \neq R$

Try combination including C, but not BC or CD:only one Combination is possible -> (ACE)

 $(ACE)^{+} = ACE \neq R$

Equivalence of Sets of Functional dependencies: definition: - For two sets of FDS 1- E and F. There is a FD x -> y in E. as E= 2x -> y] @ calculate X+ with respect to F 3 check whether this X+ includes the attributes in Y 9 If this is the case for every FD in E, then F Covers E. or ESF (5) E and F set are equivalent if E CF and F SE Example 1:- R(ACDEH) Check whether IF= [A-C, AC-D, E-AD, E-H] F=E or not E= {A-CD, E-AH} Step 1:- Check whether all the dependenties in B Can be inferred from F or not! 11 FD FLON E * A+CD At = ACD // with respect to F As it CD or all the attributes of R.H.S of A -> CD, So this FD is coved in F *> E > AH ET = EADHC // with respect to F as AH S Et, So E AH is coved in F * So we can say E EF on E is covered in F. Step2!- Check whether FCE or not? *> E>H *> A>C E+=EAHCD //pom E At= OCDA 1 // with respect to E HS Et V CEA+ V Step 3 - AS ECF

* + AC + D (AC) = ACD //from E D = (AC) / *> E > AD E+= EAHCD //from E

ADSET

and F SE So E=F

Example 2! Equivalence of Set of Functional Dependence 6

F= ¿A+C, C+DI, E+AB]

E= [A+C, C+DI, EC+AB, E+c]

A>C, C>DI are common for both, so need not to check for them;

Step 1:- check for E = F or not?

* EC -> AB

(EC) = EABCDI //inferred from F

IN AB = (EC) +, so E(-) AB is covered in F

E>C E+= EABCDI // inferred from F As C ∈ E+, then E>C coverbed in F * This means that E is cover in For E⊆F

Step 2:- check for FCE or not?

* E -> AB Et = ECABDI //inferred from E

* Now F CE

Step3: As $E \subseteq F$ and $F \subseteq E$ then F = E

Example 3:- Given below are two sets of FDs for [monathe] C a Relation R(A,B,C,D,E). Are they equivalent? $F = \{A \rightarrow B, AB \rightarrow C, D \rightarrow AC, D \rightarrow E\}$ $E = \{A \rightarrow BC, D \rightarrow AE\}$

Step 1: - check for E = F or not?

* A > BC

At = ABC // inferred from F :: A > BC Good in F

* D > AE

Dt = DACE // inferred from F]: D > AE Govered in F

* E = F

Stepd: Check for FCE or net:7

* A-7B
// This dependency is directly covered in E=[A-BC]

* AB > C (AB) = ABC //inferred from E CE (AB) +, i.e. AB > C covered in E

* D > AC Dt = DAEBC // infured from E AC C Dt, ie. D > AC covered in E

* D→E

// This dependency is directly covered in E={D→AE}

* All FDs of F are Covered in E, F SE

* All FDs of F are Covered in E, F SE

Mind I was Mr. in the party

 $Stp31- E \subseteq F$ and $F \subseteq E$ then F = E

Scanned by CamScanner

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Minimal Set/Minimum Cover/Canonical Cover of FDs: - 8
                R(A,B,D)
       BAA
       D>A
      C+BA
  Step 1.1- one element on R.H.S
 Step 2 -> Check for extra attributes cm LH-S
        may be replaced with H-D or B-D
             Cr BA
             So remove AB >D and Add A >D
              (AB)+= ABD
       with A-D; At = AD
             (AB)+ + A+, So A>D can't replace AB>D
       WITH BOD: STEBDA
               (AB)+ = B+, SO B-D can replace AB-D
        Soi- often Step 21-
                B> D
  Step31- check for redundant FDs1-
                 D > A -sessential
                 B-D-essential
 , (a) is B > A redundat or net!
           without B + A, Bt = BDA, so B > A is redundant
    (b) is D>A redundant !-
             Dt = DA
         without D > A, Dt = D, so essential
   (C) 15 B-D redundant
            Bt = BD
       WHOWB -D, BT=B, so usertial
 Final - D-A and B-D
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Finding Minimal Cover/Canonical Cover

R(A, B, C)

FD: {A>BC, B>C, A>B, AB>C}

Step1 1-A→B A→C B→C A→B

Step2:- Check for extra attributes on L.H.S.

For AB > C

taking (AB) = ABC

taking A > C | SO | AB > C = A > C

L> A+= ABC

Taking B > C] so Ly B+= BC- AB > C + B > C

So we can replace AB -> C

Final Set often Stepal-

A->C
B-> C
A-> B

duplicate

duplicate

A-> C redundant
B-> C

Step 3!- Check for redundant FDs

(a) is $A \rightarrow B$ redundantswith $A \rightarrow B$, $A + = ABCT \neq$ without $A \rightarrow B$, A + = AC $A \rightarrow B$ is essential

(b) is A-C redundant!

with A-C, A+= ABCJ

without A-C, A+= ABCJ

So A-C is redundant

(c) is B→C redundant:with B→C, B+= BC] ≠ without B→C, B+= BJ ≠ So B→C is essential

Resultant Set!

A->B B->C 3

Finding Minimal Set of "FDs:
R(A,B,C,D,E)

FDs= {A>BC, CD>E, B>D, E>A]

Stept: FDs = {A>B, A>C, CD>E, B>D, E>A}

Step 2 - Check extra attributes on L. H.S !-

with $CD \rightarrow E$, $(CD)^{\dagger} = CDEAB \leftarrow CD \rightarrow E$ with $C \rightarrow E$, $C^{\dagger} = CEABO \leftarrow CD \rightarrow E$ with $D \rightarrow E$, $D^{\dagger} = DEABC \leftarrow C \rightarrow E$ or $D \rightarrow E$

FDs= [A-B, A+C, C+E, B+D, E+A]

Step3: - Check for redundant FDs1-

(1) is A -> B redundant with A -> B, A = ABCDE7 without A -> B, A = ACE_7 So A >B is essential

(b) is A > C redundant
with A > C, A = ABCDE] #
without A > C, A = ABD] #
So; A > C is essential

(c) the C→E redundant with C→E, C+=CE] ≠ without C→E, C+=CJ ≠ So, C→E is essential

(d) is B>D redundant
with B>D, B+=BD]
without B>D, B+=BJ

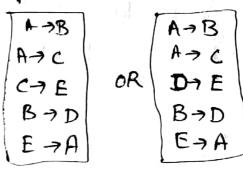
So B>D is ensential

(e) is En A redundant
with En A, Et=EABCD] 7
without En A, Et= E

30 En A is essential

A>B V A>C V C>E V B>D V E>A

so final FDs Set!-



(4) Anding Minimal Covering R(A,B,C,D,E)

FD= [A+B, AB+C, D+AC, D+E]

Step11- [A-B, AB-C, D-A, D-C, D-E]

Step 2: - Check extra symbols at L.H.S. with $AB \rightarrow C$, $(AB)^{+}=(ABC)$ with $A \rightarrow C$, $A^{+}=ACB$, so, $AB \rightarrow C=A \rightarrow C$

with B=c, B=B, so AB=c = B=c

Replace, AB > C with A > C

FD= {A-B, A-C, D-A, D-C, D-E}

Step3 :- Check for redundant FDs1-

- (a) with $A \rightarrow B$, $A^{\dagger} = ABC$ \neq without $A \rightarrow B$, $A^{\dagger} = AC$ \neq $A \rightarrow B = essential$
- (b) with A > C, A = ABC] +
 without A > C, A = AB] +
 A> C= essential

Final Set 1—
$$A \rightarrow B$$

 $A \rightarrow C$
 $D \rightarrow A$
 $D \rightarrow E$

- (d) with D > C, Dt= DABCE] without D > C, Dt= DAEBC _____

 So D > C & redundant
- (c) with D=E, D+= DAEBC] = without D=E, D+= DABC] =

 So D=E is essential.

 $DR \xrightarrow{A \to BC}$

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(5) R(w, x, y,z) (Finding Minimal Cover)
  FD= { X > w, wZ -> XY, Y > wxZ}
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decompose R.H.S:-X->W WZ > X WZ>Y Y -> W XFY 4-72

Step 2! - Check for extra elements on L.H.S: WZ7 XY (wz)+=(wzxy) * Removing w from LHS, B>XY Z+= (ZXYW) as west=zt, so z > xy can replace WZ > XY

* Removing Z from R.H.S, W > XY

After Step 3!-

 $X \rightarrow W$ X+W 1 **Zラメ** WAXI マライ ORT Y >W W->YI y -> X W EY Y > X intaking Y -> Z Y > Z this set

X->W~ Step3:-. MAX -Y-7W YAXV ソラスト

Step 3 - Check for redundant FDs: (a) is X > w redundant: with Xt= XWYZ without X > W, X+=(X) 20 x > w is essential

(b) is w → x redundant! with w+= (wxzy) without W > X, w = (wyxz) So W → X is redundant.

(c) is way redundant: w+=(wyxz) without w=>y, w+= (w),-So essential.

So waxy can also replace wzaxy (d) is Yaw redundantiwith Y=(YWXZ) without Yow, Yt= (YXZW) so you is redundant.

> (e) is yax redundantiwith yt= (YXZW) without Y = X, Y = (42) so essential

1) is Y - Z redundant! with Yt=YXZW without Y\$2, Y=(YXW) So Y-> Z is essential

Final! - set of FD!-X->W

Way YAXZ (6) Fincting Minimal Cover 1- R(A,B,C,D,E,I)

FD= {A+C, AB+C, C+DI, CD+I, EC+AB, EI+C}

Step11- FDs={A+C, AB+C, C+D, C+I, CD+I, EC+A, EC+B, EI+C}

Step2!- Check for each attributes on L.H.S:
(a) AB+C, (AB) = ABCDIT/ (a) EC+AB, (EC)=ECABDIT/=

with B+C B+S T

- (a) AB > C, (AB) = ABCDI with B>C, B = B X # with A > C, At=ACDI X # So, AB>C is required.
- (b) $CD \rightarrow I$, $(CD)^{\dagger} = CDI$ with $C \rightarrow I$, $C^{\dagger} = CDI$ with $D \rightarrow I$, $D^{\dagger} = DI$ So $CD \rightarrow I$ replaced with $C \rightarrow I$
- (c) EC→ AB, (EC) = ECABDIT = with E→AB, E= EABCDIT = #
 with C→AB, C= CABDI

 So replace EC→AB with E→AB

(e) EI -C, (EI) = EI (DAB) = with E > C, E = EGDIAB = with I + C, I = ICD

So replace EI - G with E - G

See check for dedundant attributs on right. H.S:

FDS={AB>CV

C-) I' C-) I' C-) I' redukcy E-) Ar

E-BI

(a)(AB) = ABCDI = SO AB > C B without (AB) = ABCDI = redundant

- (b) At = ACDI] of so A > C is entirely without At = A, J + so A > C is entirely
- (c) C= (+= CDI) + so (-) D= extention
- (a) without (+) I, C+= CD, So essential
- (e) E->A, Etz EABCDI, SO= 7 + So escusive without E->A, Etz (EBCDI) + So escusive

(1) where >B, Et=EALDI & soesuntil

(9) without E>C, Et= EABCDI, so redundent

Fhal: A > C C > DI E > AB