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Ans. 1 >

Pseudocode for linear search

```
for (i = 0 to n)
{
```

```
    if (arr[i] == value) // element found
```

```
}
```

Ans. 2 >

```
void insertion(int arr[], int n)
```

```
{
```

```
    if (n <= 1)
```

```
        return;
```

```
    insertion(arr, n-1);
```

```
    int nth = arr[n-1];
```

```
    int j = n-2;
```

```
    while (j >= 0 && arr[j] > nth)
```

```
{
```

```
        arr[j+1] = arr[j];
```

```
        j--;
```

```
}
```

```
    arr[j+1] = nth;
```

```
}
```

for ( $i=1$  to  $n$ )

{

key  $\leftarrow A[i]$

$j \leftarrow i-1$ ;

while ( $j \geq 0$  and  $A[j] > \text{key}$ )

{

$A[j+1] \leftarrow A[j]$

$j \leftarrow j-1$

}

$A[j+1] \leftarrow \text{key}$

}

Insertion sort is online sorting because it doesn't know the whole input, more input can be inserted with the insertion sorting is running.

Ans. 3 > Complexity  $\rightarrow$

Selection sorting  $\rightarrow$

Best  $\rightarrow O(n^2)$

Worst  $\rightarrow O(n^2)$

Average  $\rightarrow O(n^2)$

Bubble sort  $\rightarrow$

Best  $\rightarrow O(n)$

Worst  $\rightarrow O(n^2)$

Average  $\rightarrow O(n^2)$



	Best	Worst	Average
Insertion sort $\rightarrow$	$O(n)$	$O(n^2)$	$O(n^2)$
Heap sort $\rightarrow$	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$
Quick sort $\rightarrow$	$O(n \log n)$	$O(n^2)$	$O(n \log n)$
Merge sort $\rightarrow$	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$

Ans. 7)

Inplace Sorting

Stable Sorting

Online Sorting

Bubble

Merge sort

Insertion

Selection

Bubble sort

Insertion

Insertion

Quick sort

Count sort

Heap sort

Ans. 5)

```
int binary (int arr[], int l, int r, int x)
```

```
{
```

```
    if (r >= l)
```

```
    {
```

```
        int mid = l + (r - l) / 2;
```

```
        if (arr[mid] == x)
```

```
            return mid;
```

```
        else if (arr[mid] > x)
```

```
            return binary (arr, l, mid - 1, x);
```

```
        else
```

```
            return binary (arr, mid + 1, r, x);
```

```
    }  
    return -1;
```

```
    }  
    int binary (int arr [], int l, int r, int x)  
    {  
        while (l <= r)  
        {  
            int m = l + (r - l) / 2;  
            if (arr[m] == x)  
                return m;  
            else if (arr[m] > x)  
                r = m - 1;  
            else  
                l = m + 1;  
        }  
        return -1;  
    }
```

Time complexity of

linear search  $\Rightarrow O(n)$

Binary search  $\Rightarrow O(\log n)$

Ans. 6 >

$$T(n) = T(n/2) + 1$$

where  $T(n)$  is the time required for binary search in an array of size 'n'



Ans. 7

```
int find (A[], n, k)
{
    sort (A, n)
    for (i = 0 to n-1)
    {
        n = binary search (A, 0, n-1, k - A[i])
        if (n)
            return 1
    }
    return -1
}
```

$$\begin{aligned}\text{Time complexity} &= O(n \log n) + n \cdot O(\log n) \\ &= O(n \log n)\end{aligned}$$

Ans. 8

- Quick sort is the fastest ~~sort~~ general purpose sort
- In most practical ~~situ~~ situation, quick sort is the method of choice. If stability is important and space is available, merge sort might be best

Ans. 9

A pair  $(a[i], a[j])$  is said to be inversion if  $a[i] > a[j]$

In  $arr[] = \{7, 21, 31, 6, 10, 1, 20, 6, 4, 15\}$

total no. of inversion are 31, using merge sort

Ans. 10)

The worst case time complexity of quick sort is  $O(n^2)$ .

This case occurs when the picked pivot is always an extreme (smallest or largest) element

This happens when input array is sorted or reverse sorted.

- The best case of quick sort is when we wisely select pivot as a mean element.

Ans. 11)

Recurrence relation of

$$\text{Merge sort} \rightarrow T(n) = 2T(n/2) + n$$

$$\text{Quick sort} \rightarrow T(n) = 2T(n/2) + n$$

- Merge sort is more efficient and works faster than quick sort in case of larger array size or dataset.
- Worst case complexity for quick sort is  $O(n^2)$  whereas  $O(n \log n)$  for merge sort.

Ans. 12) Stable selection sort  $\rightarrow$

```
void stableSelection (int arr[], int n)
{
    for (int i=0; i<n-1; i++)
    {
        int min = i;
```



```

for (int j = i+1; j < n; j++)
{
    if (arr[min] > arr[j])
        min = j;
}
int key = arr[min];
while (min > i)
{
    arr[min] = arr[min-1];
    min--;
}
arr[i] = key;
}
}

```

Ans. 13) Modify bubble sort ->

```

void bubble (int a[], int n)
{
    for (int i = 0; i < n; i++)
    {
        int swaps = 0;
        for (int j = 0; j < n-1-i; j++)
        {
            if (a[j] > a[j+1])
            {
                int t = a[j];
                a[j] = a[j+1];
                a[j+1] = t;
                swaps++;
            }
        }
    }
}

```

if (swaps == 0)

break;

y

y