

ON THE HARDY-RAMANUJAN THEOREM

UNDERSTANDING GLOBAL & LOCAL BEHAVIOUR OF PRIME OMEGA FUNCTION

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1. **Objects:** Arithmetic functions (sequences) $f : \mathbb{N} \rightarrow \mathbb{R}$. Ex: $f(n) = n^2, \log(n), d(n), \omega(n)$.
2. **Big Oh-Notation:** For real-valued function f , a non-negative function g defined on $[a, \infty)$, $a \geq 0$ and $x_0 \geq a$ we say

$$f(x) = O(g(x)) \text{ or } f(x) \ll g(x) \text{ for all } x \geq x_0,$$

if $\exists C > 0$ such that $\forall x \geq x_0$,

$$|f(x)| \leq Cg(x).$$

If x_0 is not specified, then the notation $f(x) \ll g(x)$ assumes the existence of some appropriately large positive real for which the above statement is true.

3. **Little Oh-Notation:** For real-valued function f , and a positive function g defined on $[a, \infty)$, $a \geq 0$ we say

$$f(x) = o(g(x)) \text{ as } x \rightarrow \infty,$$

$$\text{if } \lim_{x \rightarrow \infty} \frac{f(x)}{g(x)} = 0.$$

4. **Asymptotic \sim :** For positive real-valued function f and g defined on $[a, \infty)$, $a \geq 0$ we say

$$f(x) \sim g(x) \text{ as } x \rightarrow \infty$$

$$\text{if } f(x) = g(x) + o(g(x)).$$

THE FUNCTION $\omega(n)$

Define $\omega(n)$ to be *the number of distinct prime factors* of n :

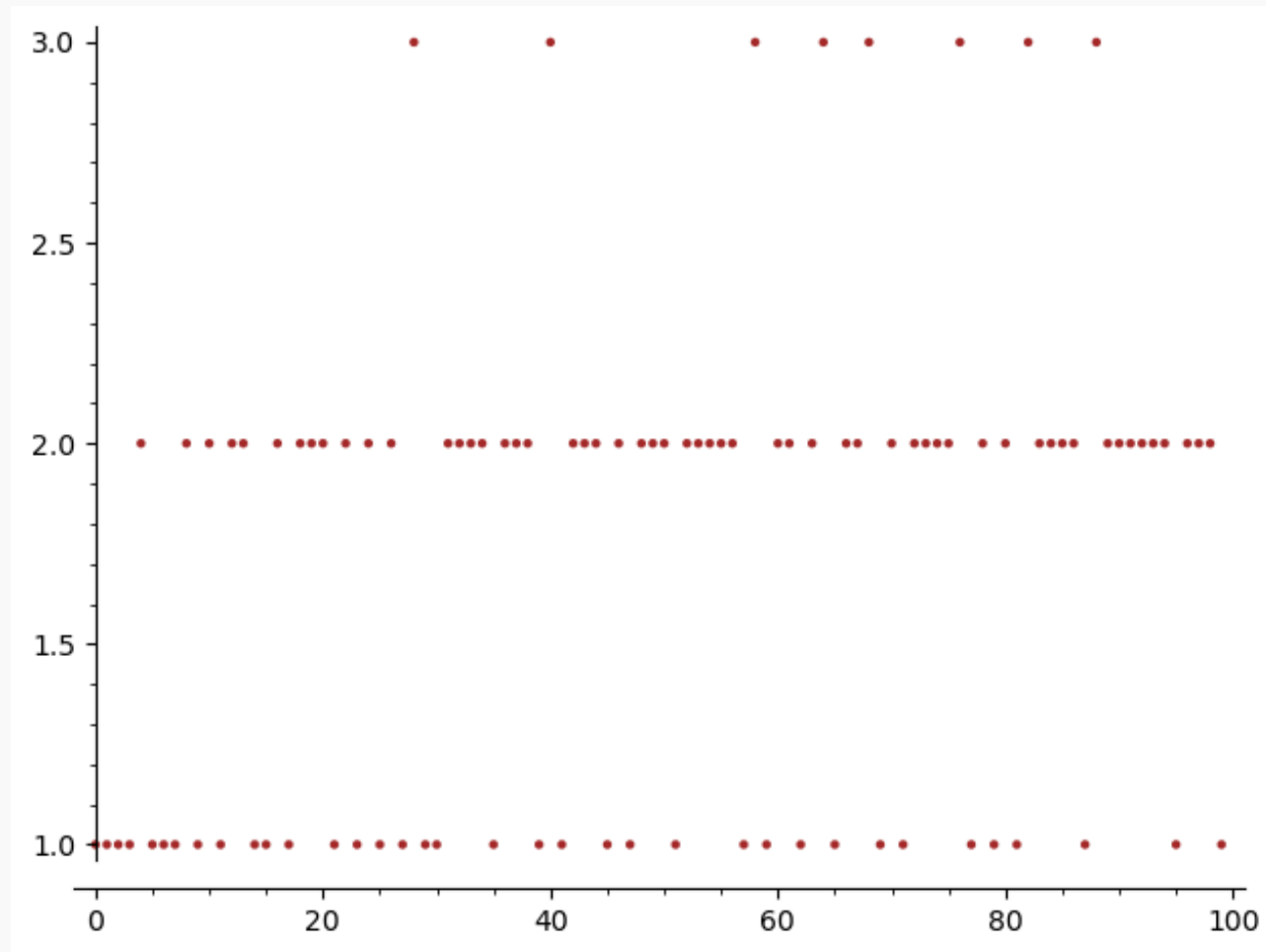
$$\omega(n) := \sum_{p|n} 1,$$

for $n \in \mathbb{N}$.

Example:

- $\omega(1) = 0$
- $\omega(2) = 1$
- $\omega(p^k) = 1$ for any prime p and $k \in \mathbb{N}$
- $\omega(200560490130) = \omega(2 \cdot 3 \cdot 5 \cdot \dots \cdot 31) = 11$
- $\omega(200560490131) = 1$

VISUAL: PLOT OF $\omega(n)$



x : natural number n ; y : $\omega(n)$

Understand the following for $\omega(n)$:

- Global behaviour in terms of *average order*.
- Local behaviour via the notion of *normal order*.

AVERAGES



WHAT IS AVERAGE?

A real-valued function g (generally monotonic) is called *average order* of the arithmetic function f , if

$$\frac{f(1) + f(2) + \dots + f(N)}{N} = \frac{1}{N} \sum_{n \leq N} f(n) \sim g(N) \text{ as } N \rightarrow \infty.$$

Theorem 0.1: Average of $\omega(n)$ is $\log \log(n)$, *i.e.*,

$$E(N) = \frac{1}{N} \sum_{n \leq N} \omega(n) = \log \log N + O(1).$$

AVERAGE OF $\omega(n)$

Theorem 0.2: Average of $\omega(n)$ is $\log \log(n)$, *i.e.*,

$$E(N) = \frac{1}{N} \sum_{n \leq N} \omega(n) = \log \log N + O(1).$$

Proof:

$$\sum_{n \leq N} \omega(n) = \sum_{n \leq N} \sum_{p|n} 1 = \sum_{p \leq N} \sum_{\substack{n \leq N \\ n=pr}} 1 = \sum_{p \leq N} \sum_{r \leq \frac{N}{p}} 1 = \sum_{p \leq N} \left[\frac{N}{p} \right]$$

AVERAGE OF $\omega(n)$

$$\begin{aligned} &= \sum_{p \leq N} \left(\frac{N}{p} + O(1) \right) = N \sum_{p \leq N} \frac{1}{p} + O(N) \\ &= N \sum_{p \leq N} \frac{1}{p} + O(N) = N(\log \log N + O(1)) + O(N) \\ &= N \log \log N + O(N). \end{aligned}$$

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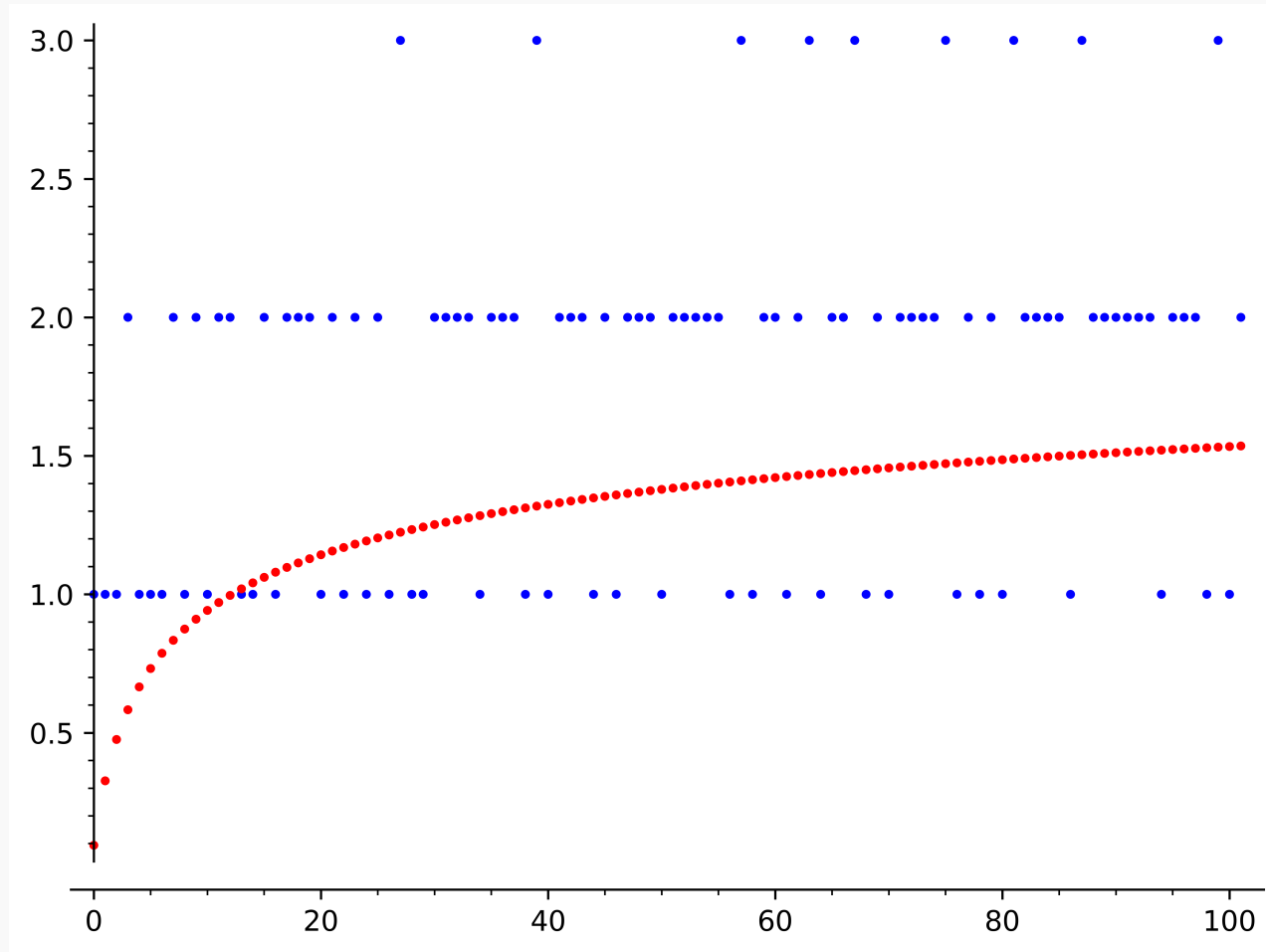
Therefore,

$$E(N) = \frac{1}{N} \sum_{n \leq N} \omega(n) = \log \log N + O(1).$$

□

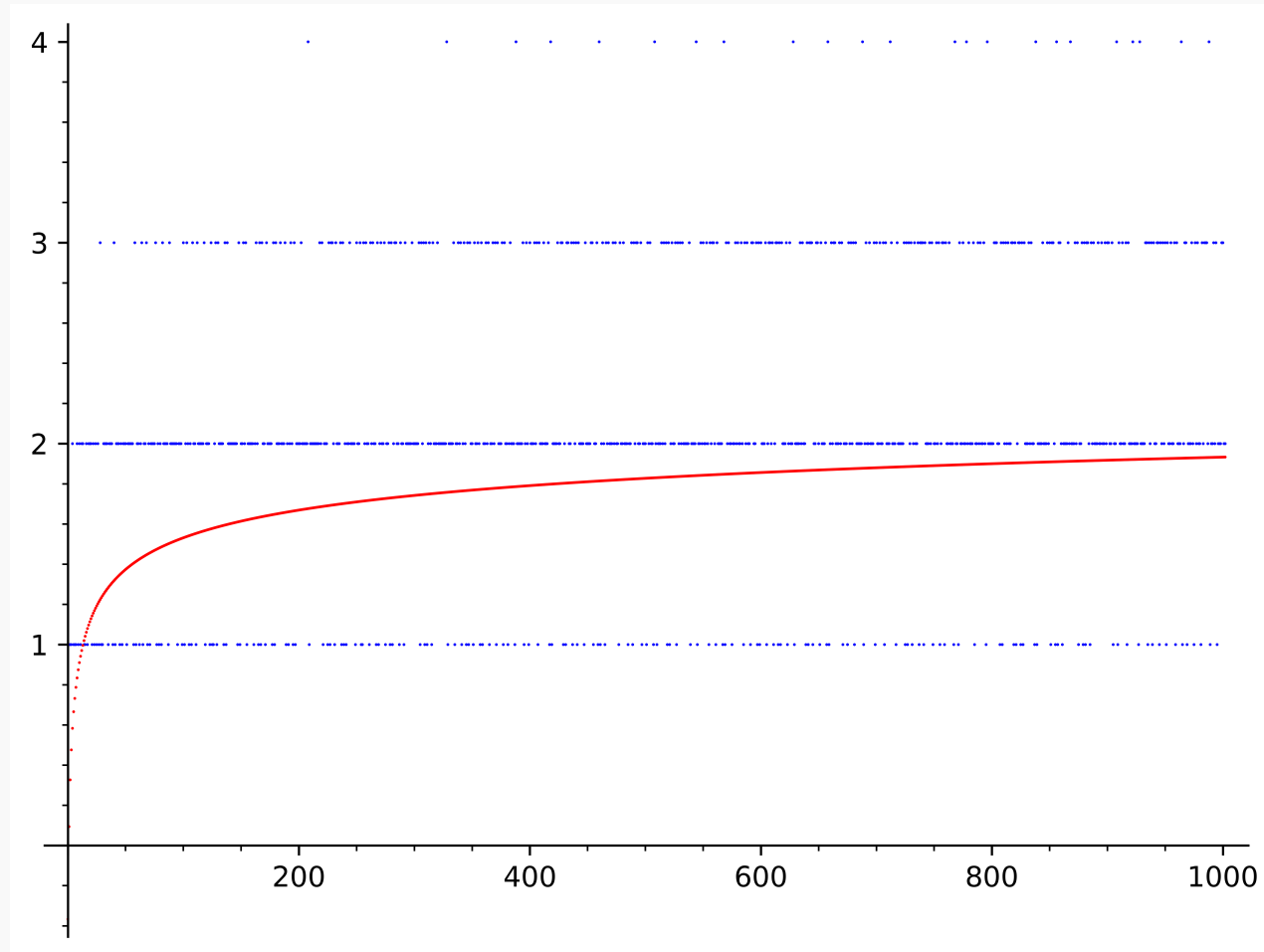
Deviation from Average

VISUAL: DEVIATION FROM AVERAGE



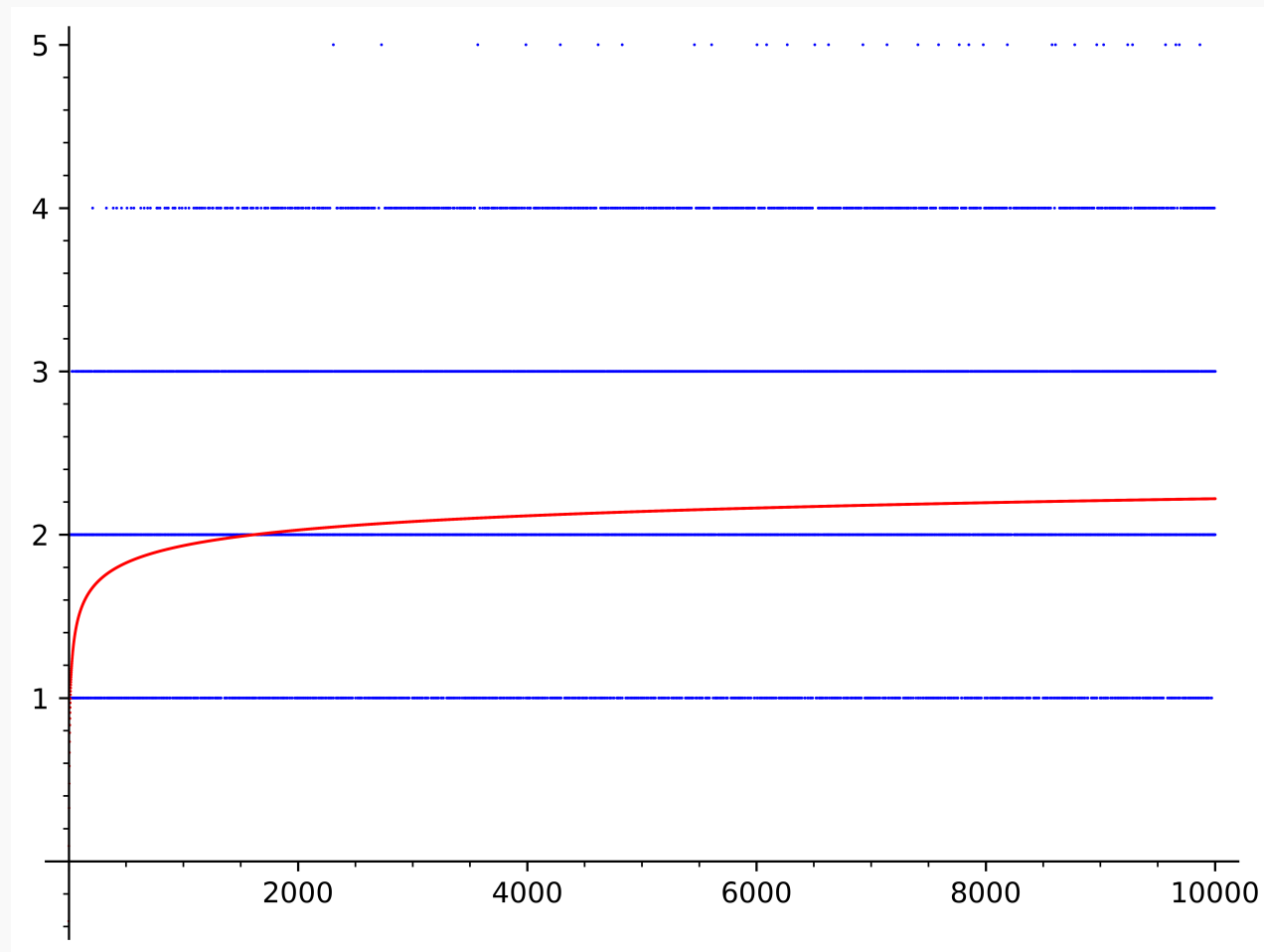
BLUE: $\omega(n)$; RED: $\log \log(n)$

VISUAL: DEVIATION FROM AVERAGE



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One apparent candidate could be an estimation of *variance*:

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In 1943, *Pál Turán* showed:

$$V(N) = O(\log \log N).$$

SIGNIFICANCE OF TURÁN'S ESTIMATE

- Shows variance is of same order as the average.
- Motivates to calculate higher moments and ...

PROOF OF TURÁN'S ESTIMATE

$$\begin{aligned}& \frac{1}{N} \sum_{1 < n \leq N} (\omega(n) - \log \log(N))^2 \\&= \frac{1}{N} \underbrace{\sum_{1 < n \leq N} \omega(n)^2}_{S(N)} - 2 \log \log(N) \underbrace{\frac{1}{N} \sum_{1 < n \leq N} \omega(n)}_{\log \log(N) + O(1)} + (\log \log(N))^2 \underbrace{\frac{1}{N} \sum_{n \leq N} 1}_{N + O(1)} \\&= \frac{1}{N} \sum_{1 < n \leq N} \omega(n)^2 \\&\quad - 2(\log \log(N))^2 + O(\log \log(N)) \\&\quad + (\log \log(N))^2 + O\left(\frac{(\log \log(N))^2}{N}\right) \\&= \frac{1}{N} \sum_{1 < n \leq N} \omega(n)^2 - (\log \log N)^2 + O(\log \log N)\end{aligned}$$

PROOF OF TURÁN'S ESTIMATE

$$S(N) = \sum_{1 < n \leq N} \omega(n)^2 = \sum_{1 < n \leq N} \sum_{p|n} 1 \sum_{q|n} 1 = \sum_{1 < n \leq N} \sum_{p|n} \sum_{q|n} 1,$$

which can be broken into following two sums:

$$S_1(N) = \sum_{1 < n \leq N} \sum_{p|n} 1 \text{ and } S_2(N) = \sum_{1 < n \leq N} \sum_{\substack{pq|n \\ p \neq q}} 1.$$

$$\begin{aligned} S_1(N) &= \sum_{1 < n \leq N} \sum_{p|n} 1 = \sum_{1 < n \leq N} \omega(n) \\ &= N \log \log(N) + O(N). \end{aligned}$$

PROOF OF TURÁN'S ESTIMATE

$$\begin{aligned} S_2(N) &= \sum_{\substack{pq \leq N \\ p \neq q}} \sum_{r: n=r} 1 \\ &= \sum_{\substack{pq \leq N \\ p \neq q}} \left[\frac{N}{pq} \right] \\ &= N \sum_{\substack{pq \leq N \\ p \neq q}} \frac{1}{pq} + O \left(\sum_{\substack{pq \leq N \\ p \neq q}} 1 \right). \end{aligned}$$

PROOF OF TURÁN'S ESTIMATE

$$\begin{aligned} S_3(N) &= \sum_{\substack{pq \leq N \\ p \neq q}} 1 \\ &= \sum_{p \leq \sqrt{N}} \sum_{q \leq \frac{N}{p}} 1 \ll \sum_{p \leq \sqrt{N}} \left[\frac{N}{p} \right] \leq \sum_{p \leq \sqrt{N}} \frac{N}{p} \ll N \log \log(N), \end{aligned}$$

Then

$$S_2(N) = N \sum_{\substack{pq \leq N \\ p \neq q}} \frac{1}{pq} + O(N \log \log(N)),$$

PROOF OF TURÁN'S ESTIMATE

$$S_4(x) = N \sum_{\substack{pq \leq N \\ p \neq q}} \frac{1}{pq} = N \sum_{pq \leq N} \frac{1}{pq} - N \underbrace{\sum_{p^2 \leq N} \frac{1}{p^2}}_{O(1)} = N \sum_{pq \leq N} \frac{1}{pq} + O(1).$$

Observe that,

$$\left(\sum_{p \leq \sqrt{N}} \frac{1}{p} \right)^2 \leq \sum_{pq \leq N} \frac{1}{pq} \leq \left(\sum_{p \leq N} \frac{1}{p} \right)^2,$$

which implies

$$\left(\log \log(\sqrt{N}) + R_1(N) \right)^2 \leq \sum_{pq \leq N} \frac{1}{pq} \leq \left(\log \log(N) + R_2(N) \right)^2,$$

where $R_1(N), R_2(N)$ both are $O(1)$. The above implies

$$\sum_{pq \leq N} \frac{1}{pq} = (\log \log(N))^2 + O(\log \log(N)).$$

Therefore,

$$S_2(N) = N(\log \log(N))^2 + O(N \log \log(N)),$$

which plugging back yields,

$$S(N) = N(\log \log(N))^2 + O(N \log \log(N)).$$

If we estimate the total pointwise deviation:

$$\tilde{V}(N) = \frac{1}{N} \sum_{1 < n \leq N} (\omega(n) - \log \log(n))^2,$$

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- Upto $N^{\frac{1}{e}}$ there are $o(N)$ number of numbers & $\omega(n) \ll \log N$.

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Because:

- Upto $N^{\frac{1}{e}}$ there are $o(N)$ number of numbers & $\omega(n) \ll \log N$.
- For $N^{\frac{1}{e}} \leq n \leq N$,

$$\log \log(N) - 1 \leq \log \log(n) \leq \log \log(N)$$

A positive real-valued function g is called *normal order* of the arithmetic function f , if for every $\varepsilon > 0$,

$$\#\left\{n \leq N : \left|\frac{f(n)}{g(n)} - 1\right| \geq \varepsilon\right\} = o(N) \text{ as } N \rightarrow \infty.$$

We also say

$$f(n) \sim g(n) \quad \text{a.e.}$$

to convey the same notion.

HARDY-RAMANUJAN THEOREM (1917)

Theorem 0.3: For every $0 < \delta < \frac{1}{2}$,

$$\underbrace{\# \left\{ 1 < n \leq N : \left| \frac{\omega(n)}{\log \log(n)} - 1 \right| \geq \frac{1}{(\log \log(n))^{\frac{1}{2}-\delta}} \right\}}_{=: B_\delta(N)} = o(N).$$

HARDY-RAMANUJAN THEOREM (1917)

Theorem 0.4: For every $0 < \delta < \frac{1}{2}$,

$$\underbrace{\# \left\{ 1 < n \leq N : \left| \frac{\omega(n)}{\log \log(n)} - 1 \right| \geq \frac{1}{(\log \log(n))^{\frac{1}{2}-\delta}} \right\}}_{=: B_\delta(N)} = o(N).$$

- RHS becomes small for sufficiently large n .

HARDY-RAMANUJAN THEOREM (1917)

Theorem 0.5: For every $0 < \delta < \frac{1}{2}$,

$$\underbrace{\# \left\{ 1 < n \leq N : \left| \frac{\omega(n)}{\log \log(n)} - 1 \right| \geq \frac{1}{(\log \log(n))^{\frac{1}{2}-\delta}} \right\}}_{=: B_\delta(N)} = o(N).$$

- RHS becomes small for sufficiently large n .

Corollary 0.5.1: The normal order of $\omega(n)$ is $\log \log(n)$.

WHY IT'S IMPORTANT?

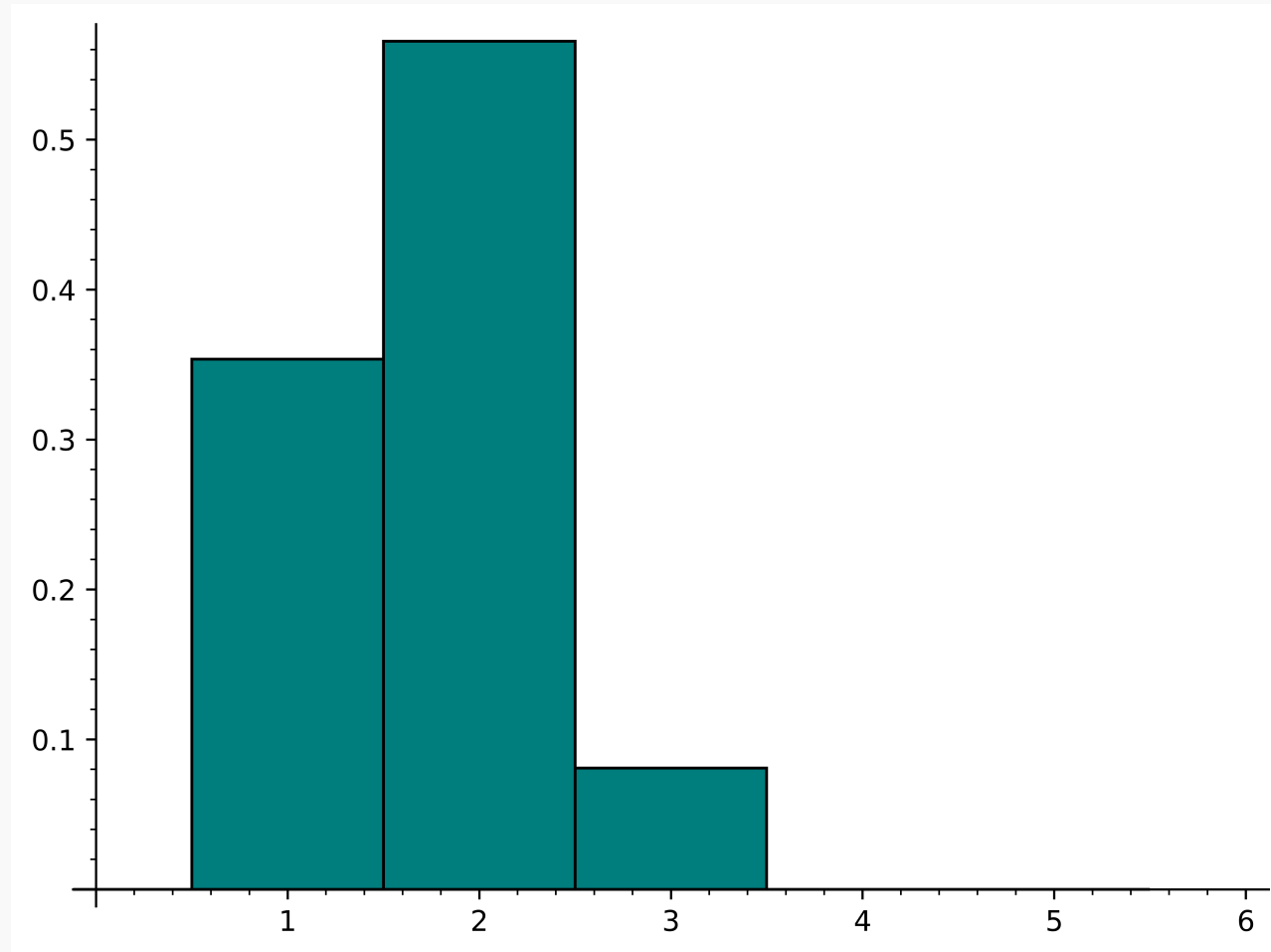
- Looks like Chebyshev's inequality in probability. Opens up the question: whether results of probability can be used to solve number theoretic problems.
- Whether this can be generalized for other additive functions.
- Computationally, the approximation is good enough in many situations.

Suppose to the contrary that for some $\delta > 0$ there is an $\alpha > 0$, such that for each $N \in \mathbb{N}$, $\#B_\delta(N) \geq \alpha N$. For any large N , then:

$$\begin{aligned}\tilde{V}(N) &= \frac{1}{N} \sum_{1 < n \leq N} (\omega(n) - \log \log(n))^2 \\ &\geq \frac{1}{N} \sum_{N^{\frac{1}{e}} < n \leq N} (\omega(n) - \log \log(n))^2 \\ &\geq \frac{1}{N} \alpha N (\log \log(N) - 1)^{1+2\delta} \\ &= \alpha' (\log \log(N))^{1+2\delta} \text{ for some } \alpha' > 0,\end{aligned}$$

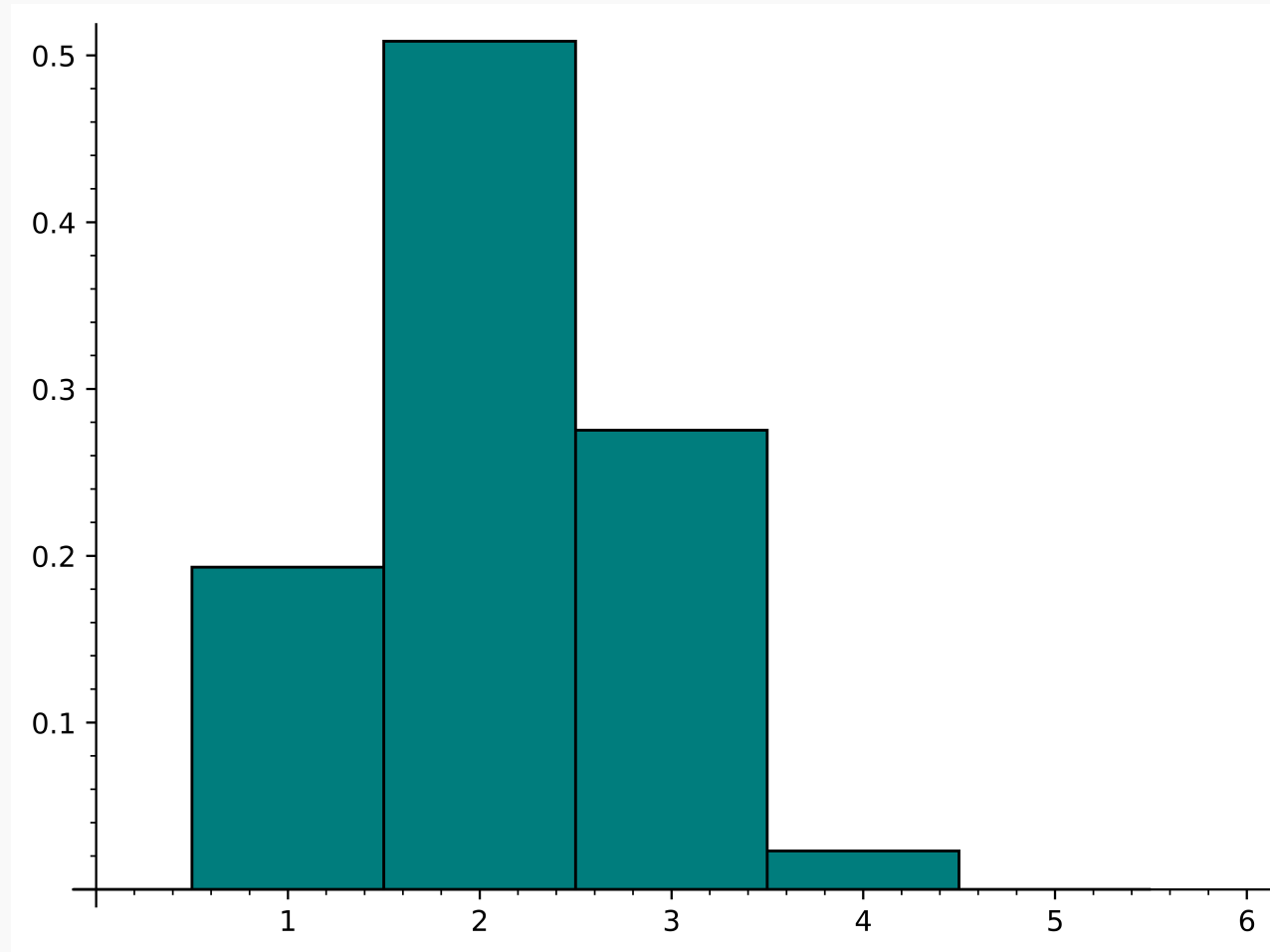
which is a contradiction since $\tilde{V}(N) = O(\log \log N)$.

FURTHER AMBITION



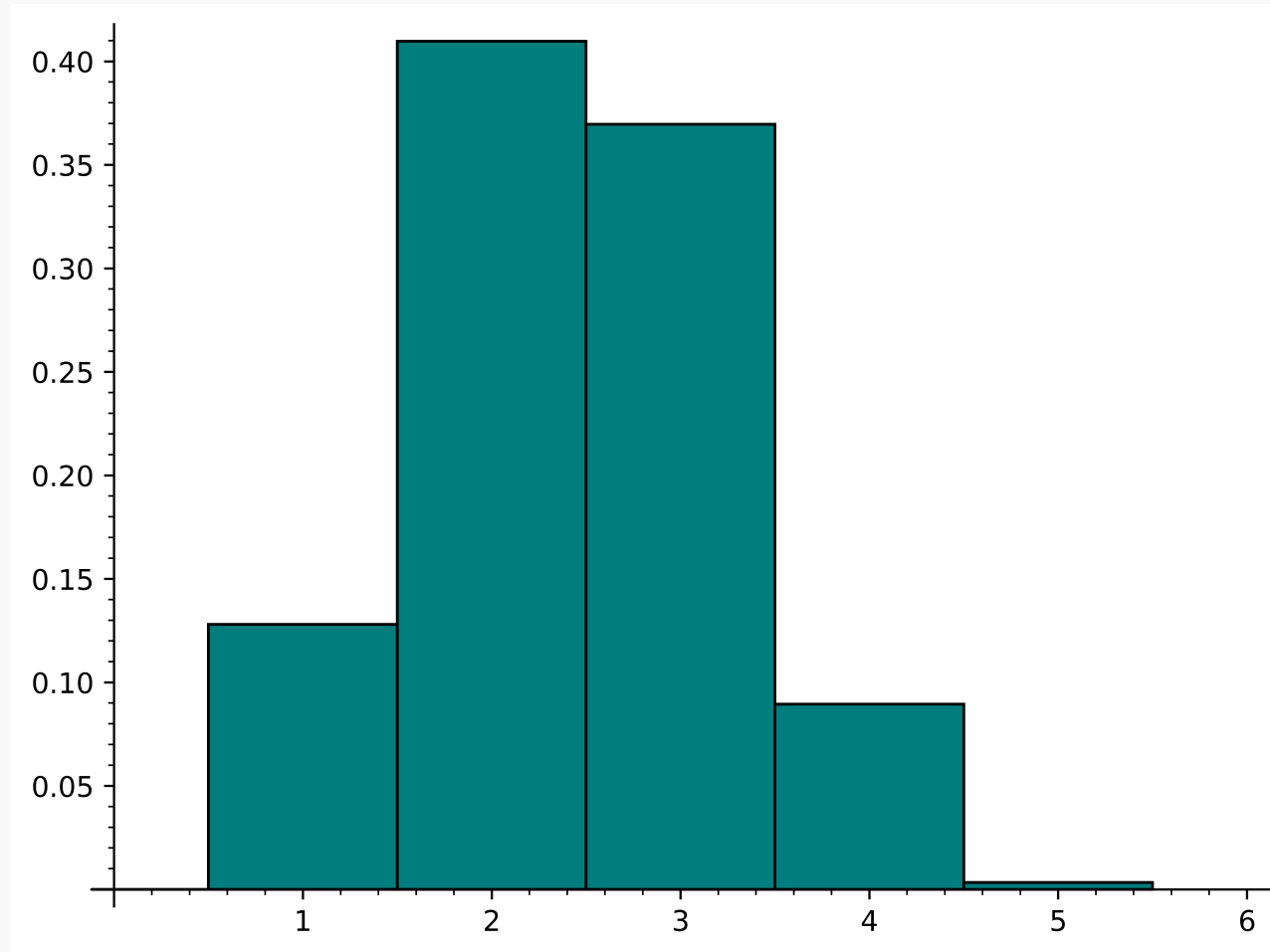
$x : \omega$ value; y : density with $N = 10^2$

FURTHER AMBITION



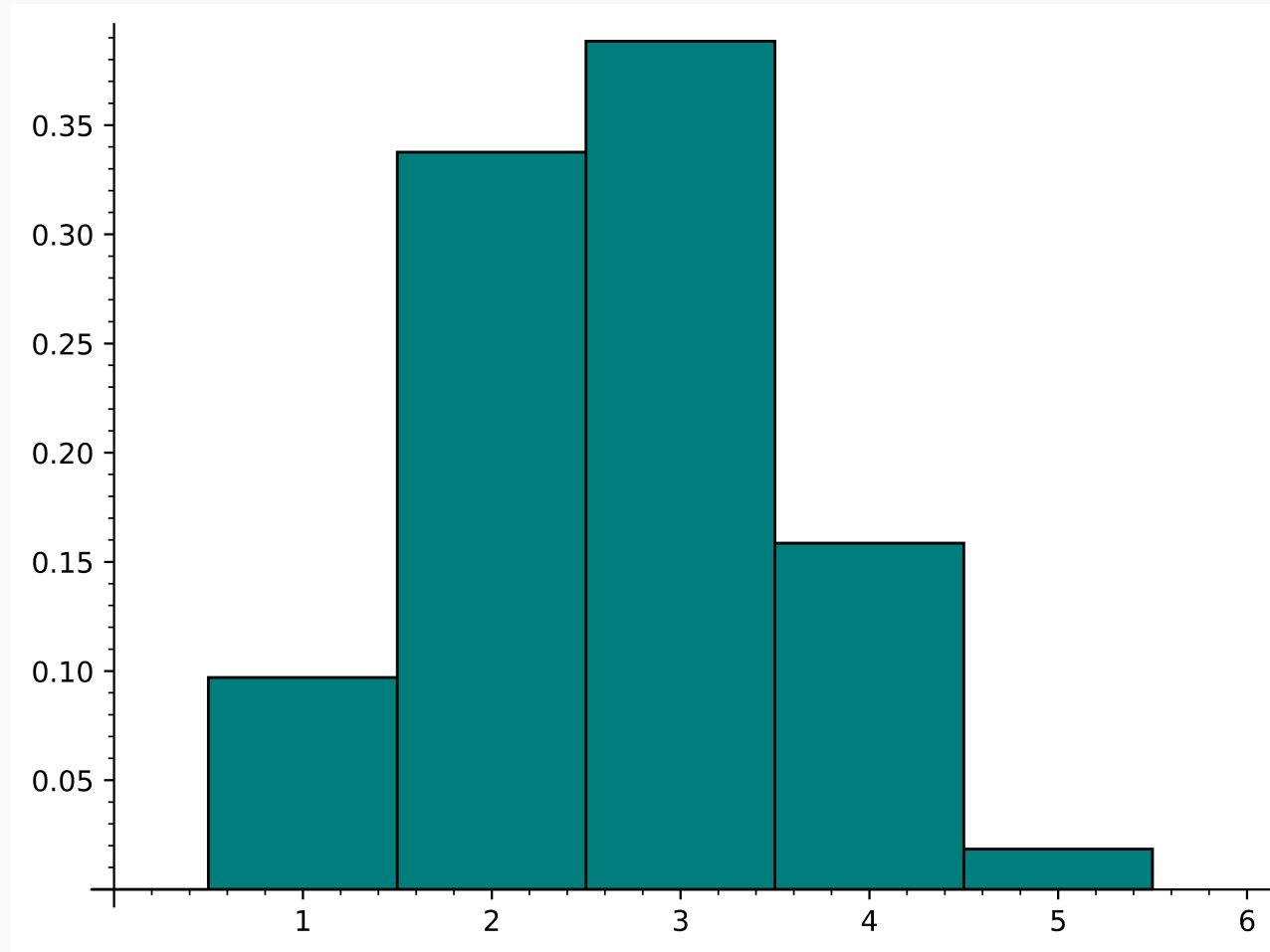
$x : \omega$ value; y : density with $N = 10^3$

FURTHER AMBITION



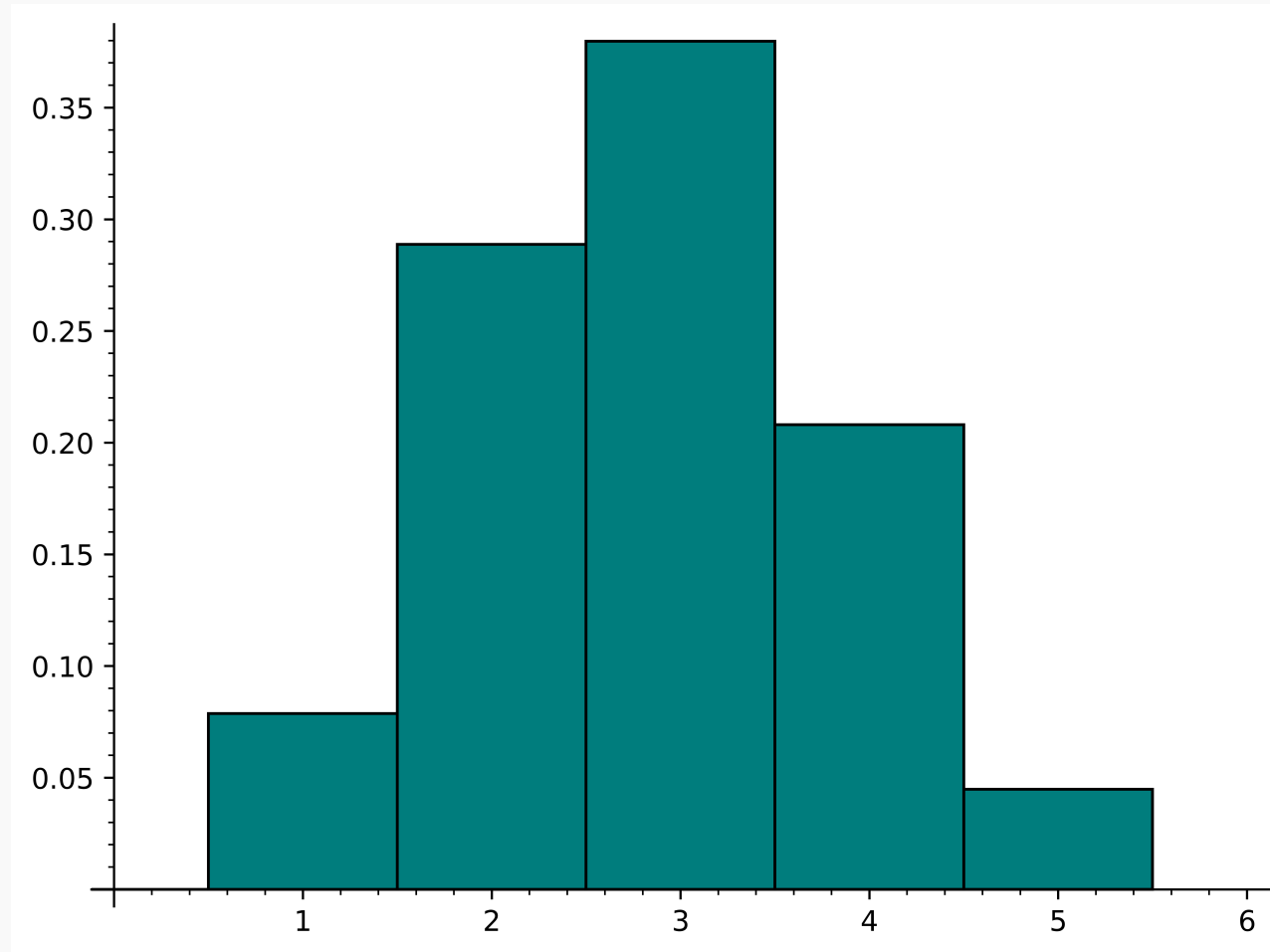
$x : \omega$ value; y : density with $N = 10^4$

FURTHER AMBITION



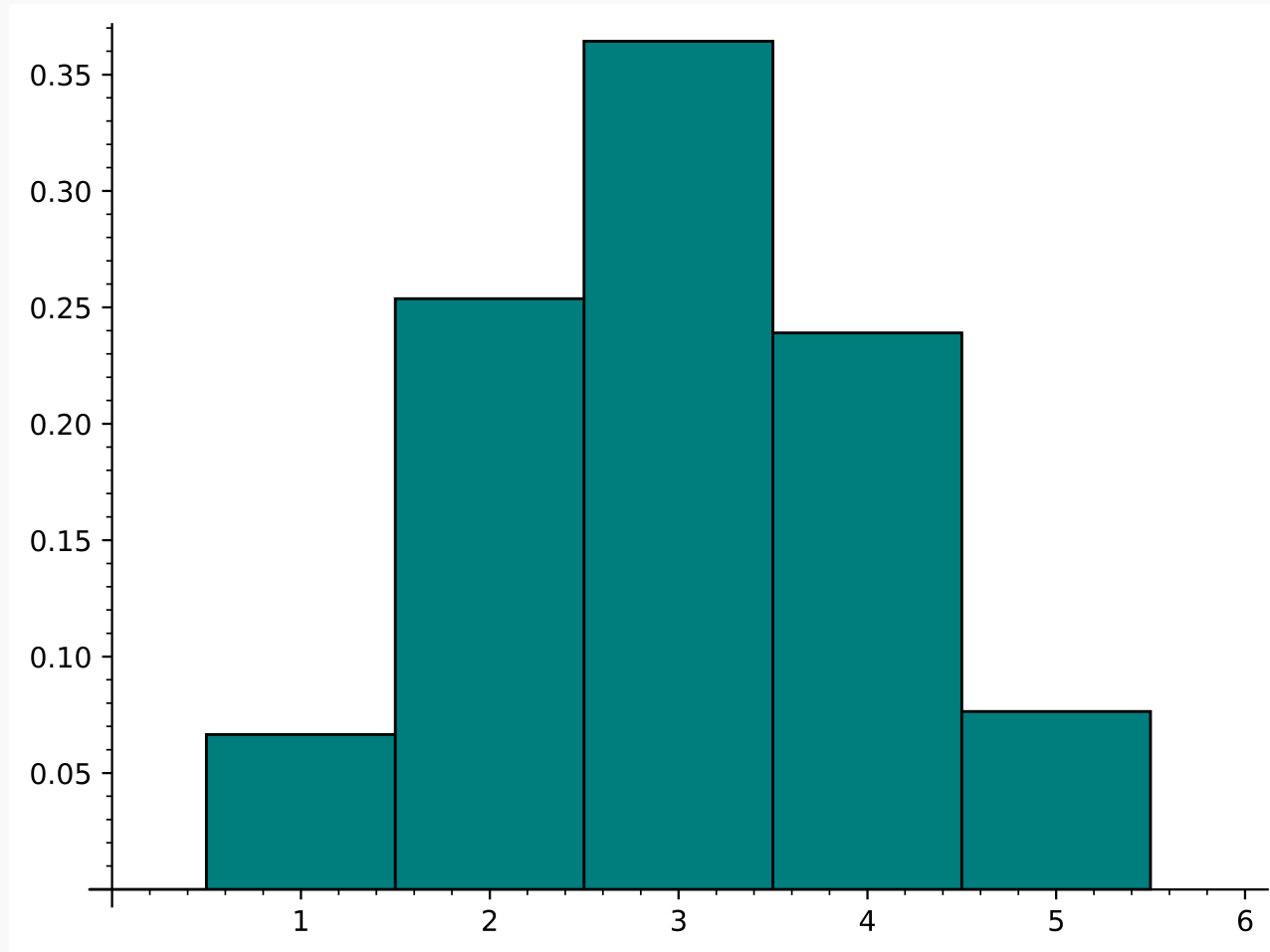
$x : \omega$ value; y : density with $N = 10^5$

FURTHER AMBITION



$x : \omega$ value; y : density with $N = 10^6$

FURTHER AMBITION



$x : \omega$ value; y : density with $N = 10^7$

'NORMAL' ORDER MAKES SENSE!

The elegant **Erdős-Kac Theorem**:

$$\lim_{N \rightarrow \infty} \frac{1}{N} \# \left\{ n \leq N : \frac{\omega(n) - \log \log n}{\sqrt{\log \log n}} \leq x \right\} = \int_{-\infty}^x e^{-\frac{t^2}{2}} dt,$$

which is called the Central Limit Theorem of probabilistic number theory.

THANK YOU!

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