

The main idea of this programming assignment is that the bloom filter will hold strings. When we add a string, its reference to indices will be set to 1. If we add it again, it will increase to 2. Since different strings can map to the same index, returning the indices for an string will yield different values (2, 3, 7). We will average them to get an estimate.

Problem:

BloomFilter Class

table \leftarrow empty list of numbers

createTable algorithm:

Creates new BloomFilter

Input: *size* - the size of the table

Output: none

table \leftarrow new list of bits of size *size*

for each element in table

 element \leftarrow 0

end for

Runtime is $O(n)$ where n is the size of the table

Add/Increment algorithm:

Adds string to BloomFilter

Input: *o*

Output: none

firstIndex \leftarrow h1(*o*)

secondIndex \leftarrow h2(*o*)

thirdIndex \leftarrow h3(*o*)

table[firstIndex] \leftarrow table[firstIndex] + 1

table[secondIndex] \leftarrow table[secondIndex] + 1

table[thirdIndex] \leftarrow 1 + table[thirdIndex]

Runtime is $O(n)$ where n is the length of the string. This is due to the call of $h3$, which has the same runtime

h1 algorithm:

Hashes string

Input: string

Output: an index in table

return (hash value of string) % size of table

Constant runtime

h2 algorithm:

Hashes string

Input: string

Output: an index in table

temp \leftarrow string

temp \leftarrow temp but in reverse order

return (hash value of temp) % size of table

Constant runtime

h3 algorithm:

Hashes string

Input: string

Output: an index in table

temp \leftarrow string but in reverse order

Combination \leftarrow empty list

for i in temp.length*2

if $i \% 2 == 0$

combination \leftarrow ith letter of string

else

combination \leftarrow ith letter of temp

end if

end for

return (hash value of combination) % size of table

Runtime is $O(n)$ where n is the length of the string

Count Algorithm:

Return frequency of string

Input: o - string

Output: frequency of string

firstIndex $\leftarrow h1(o)$

secondIndex $\leftarrow h2(o)$

thirdIndex $\leftarrow h3(o)$

If firstIndex or secondIndex or thirdIndex $\equiv 0$

return 0

End if

Return average of firstIndex, secondIndex, thirdIndex

Constant runtime

New algorithm:

Clear table

Input: none

Output: none

for each element in table

element $\leftarrow 0$

end for

Runtime is $O(n)$ where n is the size of the table

Example:

BloomFilter filter = new BloomFilter(10);

filter.add("192.12.235.36") - adds string

return filter.count("192.12.235.36") - returns the average of 1, 1 and 1 = 1.

filter.new() - clears table and sets everything to 0