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```
Sublime Build
   "cmd" : ["g++ -std=c++14 -DSONIC $file_name -o
       $file_base_name && timeout 4s ./$file_base_name<</pre>
       inputf.in>outputf.in"],
   "selector" : "source.cpp",
   "file_regex": "^(..[^:]*):([0-9]+):?([0-9]+)?:? (.*)$
   "shell": true,
   "working_dir" : "$file_path"
vimrc
set mouse=a
 set termguicolors
 filetype plugin indent on
 syntax on
" Some useful settings
 set smartindent expandtab ignorecase smartcase
      incsearch relativenumber nowrap autoread splitright
       splitbelow
 set tabstop=4
                      "the width of a tab
                      "the width for indent
 set shiftwidth=4
 colorscheme torte
"auto pair curlybraces
 inoremap {<ENTER> {}<LEFT><CR><ESC><S-o>
" mapping jj to esc
 inoremap jj <ESC>
 "compile and run using file input put
 autocmd filetype cpp map <F5> :wa<CR>:!clear && g++ % - using namespace __gnu_pbds;
      D LOCAL -std=c++17 -Wall -Wextra -Wconversion -
      Wshadow -Wfloat-equal -o ~/Codes/prog && (timeout 5
      ~/Codes/prog < ~/Codes/in) > ~/Codes/out<CR>
 "copy to input file
 map <F4> :!xclip -o -sel clip > ~/Codes/in <CR><CR>
 map <F6> :vsplit ~/Codes/in<CR>:split ~/Codes/out<CR><C using ordered_set = tree <DT, null_type, less<DT>,
      -w>=20<C-w><<C-w><C-h>
 " Leader key
 let mapleader=',,'
 " Copy template
 noremap <Leader>t :!cp ~/Codes/temp.cpp %<CR><CR>
 :autocmd BufNewFile *.cpp Or ~/Codes/temp.cpp
```

```
"note if vim-features +clipboard is not found, it will
  "for fast check :echo has('clipboard) = 0 if clipboard
      features not present,
  "need vim-gtk / vim-gtk3 package for this
 set clipboard=unnamedplus
Stress-tester
#!/bin/bash
# Call as stresstester GENERATOR SOL1 SOL2 ITERATIONS [-- void debug_out(string s) {
    countl
for i in $(seq 1 "$4"); do
   [[ $* == *--count* ]] && echo "Attempt $i/$4"
   $1 > in.txt
   $2 < in.txt > out1.txt
   $3 < in.txt > out2.txt
   diff -y out1.txt out2.txt > diff.txt
   if [ $? -ne 0 ] ; then
       echo "Differing Testcase Found:"; cat in.txt
       echo -e "\nOutputs:"; cat diff.txt
       break
   fi
done
    All Macros
```

```
//#pragma GCC optimize("Ofast")
//#pragma GCC optimization ("03")
//#pragma comment(linker, "/stack:20000000")
//#pragma GCC optimize("unroll-loops")
//#pragma GCC target("sse,sse2,sse3,ssse3,sse4,popcnt,abm|_2 |_DP
    ,mmx,avx,tune=native")
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
   //find_by_order(k) --> returns iterator to the kth
        largest element counting from 0
   //order_of_key(val) --> returns the number of items
        in a set that are strictly smaller than our item
template <typename DT>
    rb_tree_tag,tree_order_statistics_node_update>;
/*--- DEBUG TEMPLATE STARTS HERE ---*/
#ifdef SFT
void show(int x) {cerr << x;}</pre>
void show(long long x) {cerr << x;}</pre>
void show(double x) {cerr << x;}</pre>
void show(char x) {cerr << '\'' << x << '\'';}</pre>
void show(const string &x) {cerr << '\"' << x << '\"';}</pre>
```

```
void show(bool x) {cerr << (x ? "true" : "false"):}</pre>
template<typename T, typename V>
void show(pair<T, V> x) { cerr << '\f'; show(x.first);</pre>
    cerr << ", "; show(x.second); cerr << '}'; }</pre>
template<typename T>
void show(T x) {int f = 0; cerr << "{"; for (auto &i: x)</pre>
    cerr << (f++ ? ", " : ""), show(i); cerr << "}";}</pre>
   cerr << '\n';
template <typename T, typename... V>
void debug_out(string s, T t, V... v) {
   s.erase(remove(s.begin(), s.end(), ''), s.end());
   cerr << "
                   "; // 8 spaces
   cerr << s.substr(0, s.find(','));</pre>
   s = s.substr(s.find(',') + 1);
   cerr << " = ":
   show(t):
   cerr << endl;
   if(sizeof...(v)) debug_out(s, v...);
#define debug(x...) cerr << "LINE: " << __LINE__ << endl;</pre>
     debug_out(#x, x); cerr << endl;</pre>
#else
#define debug(x...)
#endif
```

#### 2.1 1D-1D

```
/// Author: anachor
#include <bits/stdc++.h>
using namespace std;
/// Solves dp[i] = min(dp[j] + cost(j+1, i)) given that
    cost() is QF
long long solve1D(int n, long long cost(int, int)) {
 vector<long long> dp(n + 1), opt(n + 1);
 deque<pair<int, int>> dq;
 dq.push_back({0, 1});
 dp[0] = 0;
 for (int i = 1; i <= n; i++) {</pre>
   opt[i] = dq.front().first;
   dp[i] = dp[opt[i]] + cost(opt[i] + 1, i);
   if (i == n) break:
```

```
dq[0].second++;
   if (dq.size() > 1 && dq[0].second == dq[1].second) dq
        .pop_front();
    int en = n;
    while (dq.size()) {
     int o = dq.back().first, st = dq.back().second;
     if (dp[o] + cost(o + 1, st) >= dp[i] + cost(i + 1,
       dq.pop_back();
     else {
       int lo = st, hi = en;
       while (lo < hi) {</pre>
         int mid = (lo + hi + 1) / 2;
         if (dp[o] + cost(o + 1, mid) < dp[i] + cost(i +</pre>
           lo = mid:
         else
           hi = mid - 1:
       if (lo < n) dq.push_back({i, lo + 1});</pre>
       break;
     }
     en = st - 1;
   if (dq.empty()) dq.push_back({i, i + 1});
 return dp[n];
/// Solves https://open.kattis.com/problems/
    coveredwalkwav
const int N = 1e6 + 7;
long long x[N];
int c;
long long cost(int 1, int r) { return (x[r] - x[1]) * (x[
    r] - x[1]) + c; }
int main() {
  ios::sync_with_stdio(false);
  cin.tie(0);
  int n:
  cin >> n >> c;
  for (int i = 1; i <= n; i++) cin >> x[i];
  cout << solve1D(n, cost) << endl;</pre>
```

# 2.2 Convex Hull Trick

```
struct line {
```

```
ll m. c:
 line() {}
 line(ll m, ll c) : m(m), c(c) {}
struct convex_hull_trick {
 vector<line> lines;
 int ptr = 0;
 convex_hull_trick() {}
  bool bad(line a, line b, line c) {
   return 1.0 * (c.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c)
         a.c) * (a.m - c.m);
 void add(line L) {
   int sz = lines.size();
   while (sz >= 2 && bad(lines[sz - 2], lines[sz - 1], L
     lines.pop_back();
     sz--;
   lines.pb(L);
 ll get(int idx, int x) { return (1ll * lines[idx].m * x
       + lines[idx].c); }
 11 query(int x) {
   if (lines.empty()) return 0;
   if (ptr >= lines.size()) ptr = lines.size() - 1;
   while (ptr < lines.size() - 1 && get(ptr, x) > get(
        ptr + 1, x)) ptr++;
   return get(ptr, x);
11 sum[MAX];
11 dp[MAX];
int arr[MAX];
int main() {
 fastio;
 int t;
  cin >> t:
 while (t--) {
   int n, a, b, c;
   cin >> n >> a >> b >> c;
   for (int i = 1; i <= n; i++) cin >> sum[i];
   for (int i = 1; i <= n; i++) dp[i] = 0, sum[i] += sum
        [i - 1];
   convex_hull_trick cht;
   cht.add(line(0, 0));
   for (int pos = 1; pos <= n; pos++) {</pre>
     dp[pos] = cht.query(sum[pos]) - 111 * a * sqr(sum[
         pos]) - c;
```

# 2.3 Divide and Conquer dp

```
const int K = 805, N = 4005;
LL dp[2][N], _cost[N][N];
// 1-indexed for convenience
LL cost(int 1, int r) {
 return _cost[r][r] - _cost[l - 1][r] - _cost[r][l - 1]
     + _cost[1 - 1][1 - 1] >> 1;
void compute(int cnt, int 1, int r, int optl, int optr) {
 if (1 > r) return:
 int mid = 1 + r >> 1;
 LL best = INT_MAX;
 int opt = -1;
 for (int i = optl; i <= min(mid, optr); i++) {</pre>
   LL cur = dp[cnt ^1][i - 1] + cost(i, mid);
   if (cur < best) best = cur, opt = i;</pre>
 dp[cnt][mid] = best;
 compute(cnt, 1, mid - 1, optl, opt);
 compute(cnt, mid + 1, r, opt, optr);
LL dnc_dp(int k, int n) {
 fill(dp[0] + 1, dp[0] + n + 1, INT_MAX);
 for (int cnt = 1; cnt <= k; cnt++) {</pre>
   compute(cnt & 1, 1, n, 1, n);
 return dp[k & 1][n];
```

#### 2.4 Knuth optimization

```
const int N = 1005;
LL dp[N][N], a[N];
int opt[N][N];
LL cost(int i, int j) { return a[j + 1] - a[i]; }
LL knuth_optimization(int n) {
  for (int i = 0; i < n; i++) {
    dp[i][i] = 0;
    opt[i][i] = i;
  }
  for (int i = n - 2; i >= 0; i--) {
    for (int j = i + 1; j < n; j++) {</pre>
```

#### 2.5 Li Chao Tree

```
struct line {
 LL m, c;
 line(LL m = 0, LL c = 0) : m(m), c(c) {}
LL calc(line L, LL x) { return 1LL * L.m * x + L.c; }
struct node {
 LL m, c;
 line L:
  node *lft, *rt;
 node(LL m = 0, LL c = 0, node *lft = NULL, node *rt =
      NULL)
     : L(line(m, c)), lft(lft), rt(rt) {}
};
struct LiChao {
 node *root;
  LiChao() { root = new node(); }
  void update(node *now, int L, int R, line newline) {
   int mid = L + (R - L) / 2:
   line lo = now->L, hi = newline;
   if (calc(lo, L) > calc(hi, L)) swap(lo, hi);
   if (calc(lo, R) <= calc(hi, R)) {</pre>
     now->L = hi;
     return;
   if (calc(lo, mid) < calc(hi, mid)) {</pre>
     now->L = hi;
     if (now->rt == NULL) now->rt = new node();
     update(now->rt, mid + 1, R, lo);
   } else {
     now->L = lo;
     if (now->lft == NULL) now->lft = new node():
     update(now->lft, L, mid, hi);
```

```
LL query(node *now, int L, int R, LL x) {
  if (now == NULL) return -inf;
  int mid = L + (R - L) / 2;
  if (x <= mid)
    return max(calc(now->L, x), query(now->lft, L, mid, x));
  else
    return max(calc(now->L, x), query(now->rt, mid + 1, R, x));
};
```

#### Data Structure

# 3.1 Segment Tree $\frac{\text{const int N} = 1000006}{\text{const}}$

```
using DT = LL;
using LT = LL;
constexpr DT I = 0;
constexpr LT None = 0;
DT val[4 * N];
LT lazy[4 * N];
int L, R;
void pull(int s, int e, int node) {
 val[node] = val[node << 1] + val[node << 1 | 1];</pre>
void apply(const LT &U, int s, int e, int node) {
 val[node] += (e - s + 1) * U;
 lazv[node] += U;
void reset(int node) { lazy[node] = None; }
DT merge(const DT &a, const DT &b) { return a + b; }
DT get(int s, int e, int node) { return val[node]; }
void push(int s, int e, int node) {
 if (s == e) return;
 apply(lazy[node], s, s + e \gg 1, node \ll 1);
 apply(lazy[node], s + e + 2 >> 1, e, node << 1 | 1);
 reset(node);
void build(int s, int e, vector<DT> &v, int node = 1) {
 int m = s + e >> 1;
 if (s == e) {
   val[node] = v[s];
   return:
 build(s, m, v, node * 2);
 build(m + 1, e, v, node * 2 + 1);
 pull(s, e, node);
```

```
void update(int S, int E, LT uval, int s = L, int e = R,
    int node = 1) {
 if (S > E) return:
 if (S == s \text{ and } E == e) {
   apply(uval, s, e, node);
   return;
 }
 push(s, e, node);
 int m = s + e >> 1:
 update(S, min(m, E), uval, s, m, node * 2);
 update(max(S, m + 1), E, uval, m + 1, e, node * 2 + 1);
 pull(s, e, node);
DT query(int S, int E, int s = L, int e = R, int node =
    1) {
 if (S > E) return I;
 if (s == S and e == E) return get(s, e, node);
 push(s, e, node);
 int m = s + e >> 1:
 DT L = query(S, min(m, E), s, m, node * 2);
 DT R = query(max(S, m + 1), E, m + 1, e, node * 2 + 1);
 return merge(L, R);
void init(int _L, int _R, vector<DT> &v) {
 L = L, R = R:
 build(L, R, v);
```

#### 3.2 Persistent Segment Tree

```
struct Node {
 Node *1, *r;
 int sum:
 Node(int val) : l(nullptr), r(nullptr), sum(val) {}
 Node(Node* 1, Node* r) : 1(1), r(r), sum(0) {
   if (1) sum += 1->sum;
   if (r) sum += r->sum;
 }
};
int a[MAXN];
Node* root[MAXN];
Node* Build(int bg, int ed) {
 if (bg == ed) return new Node(a[bg]);
 int mid = (bg + ed) / 2;
 return new Node(Build(bg, mid), Build(mid + 1, ed));
int Query(Node* v, int bg, int ed, int 1, int r) {
```

```
if (1 > ed || r < bg) return 0;</pre>
  if (1 <= bg && ed <= r) return v->sum;
  int mid = (bg + ed) / 2;
  return Query(v->1, bg, mid, 1, r) + Query(v->r, mid +
      1, ed, l, r);
Node* Update(Node* v, int bg, int ed, int pos, int
    new val) {
  if (bg == ed) return new Node(v->sum + new_val);
  int mid = (bg + ed) / 2;
  if (pos <= mid)</pre>
   return new Node(Update(v->1, bg, mid, pos, new_val),
        v->r);
  else
   return new Node(v->1, Update(v->r, mid + 1, ed, pos,
        new_val));
```

# 3.3 Implicit Segment Tree

```
struct node {
  int val:
  node *lft, *rt;
  node() {}
  node(int val = 0) : val(val), lft(NULL), rt(NULL) {}
};
struct implicit_segtree {
  node *root;
  implicit_segtree() {}
  implicit_segtree(int n) {
   root = new node(n);
  void update(node *now, int L, int R, int idx, int val)
   if (L == R) {
     now -> val += val;
     return;
   int mid = L + (R - L) / 2:
   if (now->lft == NULL) now->lft = new node(mid - L +
        1);
   if (now->rt == NULL) now->rt = new node(R - mid);
    if (idx <= mid) update(now->lft, L, mid, idx, val);
    else update(now->rt, mid + 1, R, idx, val);
   now->val = (now->lft)->val + (now->rt)->val;
  int query(node *now, int L, int R, int k) {
   if (L == R) return L;
```

```
int mid = L + (R - L) / 2:
   if (now->lft == NULL) now->lft = new node(mid - L +
   if (now->rt == NULL) now->rt = new node(R - mid):
   if (k <= (now->lft)->val) return query(now->lft, L,
        mid, k);
   else return query(now->rt, mid + 1, R, k - (now->lft)
        ->val);
 }
};
```

### 3.4 HashTable

```
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
const int RANDOM = chrono::high_resolution_clock::now().
    time_since_epoch().count();
unsigned hash_f(unsigned x) {
 x = ((x >> 16) ^x) * 0x45d9f3b;
 x = ((x >> 16) ^x) * 0x45d9f3b;
 return x = (x >> 16) \hat{x};
unsigned hash_combine(unsigned a, unsigned b) { return a
    * 31 + b: }
struct chash {
 int operator()(int x) const { return hash_f(x); }
typedef gp_hash_table<int, int, chash> gp;
gp table;
```

#### 3.5 DSU With Rollbacks

```
struct Rollback_DSU {
 int n:
 vector<int> par, sz;
 vector<pair<int, int>> op;
 Rollback_DSU(int n) : par(n), sz(n, 1) {
   iota(par.begin(), par.end(), 0);
   op.reserve(n);
 int Anc(int node) {
   for (; node != par[node]; node = par[node])
     ; // no path compression
   return node;
 }
 void Unite(int x, int y) {
   if (sz[x = Anc(x)] < sz[y = Anc(y)]) swap(x, y);
   op.emplace_back(x, y);
   par[v] = x;
   sz[x] += sz[y];
```

```
void Undo(int t) {
  for (; op.size() > t; op.pop_back()) {
    par[op.back().second] = op.back().second;
    sz[op.back().first] -= sz[op.back().second];
}
```

```
3.6 Binary Trie
const int N = 1e7 + 5, b = 30:
int tc = 1;
struct node {
 int vis = 0;
 int to[2] = \{0, 0\};
 int val[2] = \{0, 0\};
 void update() {
   to[0] = to[1] = 0;
   val[0] = val[1] = 0;
   vis = tc;
 }
T[N + 2];
node *root = T;
int ptr = 0;
node *nxt(node *cur, int x) {
 if (cur->to[x] == 0) cur->to[x] = ++ptr;
 assert(ptr < N);</pre>
 int idx = cur->to[x];
 if (T[idx].vis < tc) T[idx].update();</pre>
 return T + idx;
int query(int j, int aj) {
 int ans = 0, jaj = j ^ aj;
 node *cur = root:
 for (int k = b - 1; "k; k--) {
   maximize(ans, nxt(cur, (jaj >> k & 1) ^ 1)->val[1 ^ (
        aj >> k & 1)]);
   cur = nxt(cur, (jaj >> k & 1));
 return ans;
void insert(int j, int aj, int val) {
 int jaj = j ^ aj;
 node *cur = root;
 for (int k = b - 1; "k; k--) {
   cur = nxt(cur, (jaj >> k & 1));
   maximize(cur->val[j >> k & 1], val);
 }
void clear() {
```

tc++:

ptr = 0;

root->update();

# 3.8 Divide And Conquer Query Offline

```
namespace up {
int 1[N], r[N], u[N], v[N], tm;
void push(int _1, int _r, int _u, int _v) {
 1[tm] = _1, r[tm] = _r, u[tm] = _u, v[tm] = _v;
 tm++;
} // namespace up
namespace que {
int node[N], tm;
LL ans[N];
void push(int _node) { node[++tm] = _node; }
} // namespace que
namespace edge_set {
void push(int i) { dsu ::merge(up ::u[i], up ::v[i]); }
void pop(int t) { dsu ::rollback(t); }
int time() { return dsu ::op.size(); }
LL query(int u) { return a[dsu ::root(u)]; }
} // namespace edge_set
namespace dncq {
vector<int> tree[4 * N];
void update(int idx, int 1 = 0, int r = que ::tm, int
    node = 1) {
 int ul = up ::1[idx], ur = up ::r[idx];
 if (r 
 if (ul <= l and r <= ur) {</pre>
```

```
tree[node].push_back(idx);
   return;
 }
 int m = 1 + r >> 1:
 update(idx, 1, m, node << 1);
 update(idx, m + 1, r, node \langle \langle 1 | 1 \rangle \rangle;
void dfs(int 1 = 0, int r = que ::tm, int node = 1) {
 int cur = edge_set ::time();
 for (int e : tree[node]) edge_set ::push(e);
 if (1 == r) {
   que ::ans[1] = edge_set ::query(que ::node[1]);
 } else {
   int m = 1 + r >> 1;
   dfs(1, m, node << 1);
   dfs(m + 1, r, node << 1 | 1);
  edge_set ::pop(cur);
} // namespace dncq
void push_updates() {
 for (int i = 0; i < up ::tm; i++) dncq ::update(i);</pre>
```

# 3.9 Merge Sort Tree

```
vector<LL> Tree[4 * MAXN];
LL arr[MAXN];
vector<LL> merge(vector<LL> v1, vector<LL> v2) {
 LL i = 0, j = 0;
  vector<LL> ret;
  while (i < v1.size() || j < v2.size()) {</pre>
   if (i == v1.size()) {
     ret.push_back(v2[j]);
     j++;
   } else if (j == v2.size()) {
     ret.push_back(v1[i]);
     i++;
   } else {
     if (v1[i] < v2[i]) {</pre>
       ret.push_back(v1[i]);
       i++:
     } else {
       ret.push_back(v2[j]);
       j++;
     }
   }
 }
```

```
return ret:
void Build(LL node, LL bg, LL ed) {
 if (bg == ed) {
   Tree[node].push_back(arr[bg]);
   return;
 }
 LL leftNode = 2 * node, rightNode = 2 * node + 1;
 LL mid = (bg + ed) / 2;
 Build(leftNode, bg, mid);
 Build(rightNode, mid + 1, ed);
 Tree[node] = merge(Tree[leftNode], Tree[rightNode]);
LL query(LL node, LL bg, LL ed, LL 1, LL r, LL k) {
 if (ed < 1 || bg > r) return 0;
 if (1 <= bg && ed <= r)</pre>
   return upper_bound(Tree[node].begin(), Tree[node].end
       (), k) -
         Tree[node].begin();
 LL leftNode = 2 * node, rightNode = 2 * node + 1;
 LL mid = (bg + ed) / 2;
 return query(leftNode, bg, mid, l, r, k) +
       query(rightNode, mid + 1, ed, 1, r, k);
```

#### 3.10 MO with Update

```
const int N = 1e5 + 5, sz = 2700, bs = 25;
int arr[N], freq[2 * N], cnt[2 * N], id[N], ans[N];
struct query {
 int 1, r, t, L, R;
 query(int l = 1, int r = 0, int t = 1, int id = -1)
     : 1(1), r(r), t(t), L(1 / sz), R(r / sz) {}
 bool operator<(const query &rhs) const {</pre>
   return (L < rhs.L) or (L == rhs.L and R < rhs.R) or
          (L == rhs.L and R == rhs.R and t < rhs.t):
 }
} Q[N];
struct update {
 int idx, val, last;
} Up[N];
int gi = 0, ui = 0;
int 1 = 1, r = 0, t = 0;
```

```
void add(int idx) {
  --cnt[freq[arr[idx]]];
  freq[arr[idx]]++;
  cnt[freq[arr[idx]]]++;
void remove(int idx) {
  --cnt[freq[arr[idx]]];
  freq[arr[idx]]--;
  cnt[freq[arr[idx]]]++;
void apply(int t) {
  const bool f = 1 <= Up[t].idx and Up[t].idx <= r;</pre>
  if (f) remove(Up[t].idx);
  arr[Up[t].idx] = Up[t].val;
 if (f) add(Up[t].idx);
void undo(int t) {
  const bool f = 1 <= Up[t].idx and Up[t].idx <= r;</pre>
  if (f) remove(Up[t].idx);
  arr[Up[t].idx] = Up[t].last;
  if (f) add(Up[t].idx);
int mex() {
 for (int i = 1; i <= N; i++)</pre>
   if (!cnt[i]) return i;
  assert(0);
int main() {
 int n, q;
  cin >> n >> q;
  int counter = 0;
  map<int, int> M;
  for (int i = 1; i <= n; i++) {</pre>
   cin >> arr[i];
   if (!M[arr[i]]) M[arr[i]] = ++counter;
   arr[i] = M[arr[i]];
 iota(id, id + N, 0);
  while (q--) {
   int tp, x, y;
   cin >> tp >> x >> y;
   if (tp == 1)
     Q[++qi] = query(x, y, ui);
     if (!M[y]) M[y] = ++counter;
     y = M[y];
     Up[++ui] = \{x, y, arr[x]\};
     arr[x] = y;
```

```
}
t = ui;
cnt[0] = 3 * n;
sort(id + 1, id + qi + 1, [&](int x, int y) { return Q[
    x] < Q[y]; });
for (int i = 1; i <= qi; i++) {
    int x = id[i];
    while (Q[x].t > t) apply(++t);
    while (Q[x].t < t) undo(t--);
    while (Q[x].t < 1) add(--1);
    while (Q[x].r > r) add(++r);
    while (Q[x].t > 1) remove(1++);
    while (Q[x].r < r) remove(r--);
    ans[x] = mex();
}
for (int i = 1; i <= qi; i++) cout << ans[i] << '\n';
}
</pre>
```

### 3.11 SparseTable (Rectangle Query)

```
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 505;
const int LOGN = 9;
// O(n^2 (logn)^2
// Supports Rectangular Query
int A[MAXN][MAXN];
int M[MAXN][MAXN][LOGN][LOGN];
void Build2DSparse(int N) {
 for (int i = 1; i <= N; i++) {</pre>
   for (int j = 1; j <= N; j++) {
     M[i][j][0][0] = A[i][j];
   for (int q = 1; (1 << q) <= N; q++) {
     int add = 1 << (q - 1);</pre>
     for (int j = 1; j + add <= N; j++) {
       M[i][j][0][q] = max(M[i][j][0][q - 1], M[i][j +
            add][0][q - 1]);
     }
   }
 }
 for (int p = 1; (1 << p) <= N; p++) {</pre>
   int add = 1 << (p - 1);
   for (int i = 1; i + add <= N; i++) {</pre>
     for (int q = 0; (1 << q) <= N; q++) {
       for (int j = 1; j <= N; j++) {</pre>
```

# 4 Geometry

#### 4.1 Point

```
typedef double Tf;
typedef double Ti; /// use long long for exactness
const Tf PI = acos(-1), EPS = 1e-9;
int dcmp(Tf x) \{ return abs(x) < EPS ? 0 : (x < 0 ? -1 : 
    1); }
struct Point {
 Ti x, y;
 Point(Ti x = 0, Ti y = 0) : x(x), y(y) {}
 Point operator+(const Point& u) const { return Point(x
     + u.x, y + u.y); }
 Point operator-(const Point& u) const { return Point(x
     -u.x, y - u.y); }
 Point operator*(const LL u) const { return Point(x * u,
      y * u); }
 Point operator*(const Tf u) const { return Point(x * u,
      v * u); }
 Point operator/(const Tf u) const { return Point(x / u,
      y / u); }
 bool operator==(const Point& u) const {
   return dcmp(x - u.x) == 0 && dcmp(y - u.y) == 0;
```

```
bool operator!=(const Point& u) const { return !(*this
      == u): }
  bool operator<(const Point& u) const {</pre>
   return dcmp(x - u.x) < 0 \mid \mid (dcmp(x - u.x) == 0 &&
        dcmp(y - u.y) < 0);
 }
};
Ti dot(Point a, Point b) { return a.x * b.x + a.y * b.y;
Ti cross(Point a, Point b) { return a.x * b.y - a.y * b.x
Tf length(Point a) { return sqrt(dot(a, a)); }
Ti sqLength(Point a) { return dot(a, a); }
Tf distance(Point a, Point b) { return length(a - b); }
Tf angle(Point u) { return atan2(u.y, u.x); }
// returns angle between oa, ob in (-PI, PI]
Tf angleBetween(Point a, Point b) {
 Tf ans = angle(b) - angle(a);
 return ans <= -PI ? ans + 2 * PI : (ans > PI ? ans - 2
      * PI : ans);
// Rotate a ccw by rad radians, Tf Ti same
Point rotate(Point a, Tf rad) {
 return Point(a.x * cos(rad) - a.y * sin(rad),
             a.x * sin(rad) + a.y * cos(rad));
// rotate a ccw by angle th with cos(th) = co && sin(th)
    = si, tf ti same
Point rotatePrecise(Point a, Tf co, Tf si) {
 return Point(a.x * co - a.y * si, a.y * co + a.x * si);
Point rotate90(Point a) { return Point(-a.y, a.x); }
// scales vector a by s such that length of a becomes s,
    Tf Ti same
Point scale(Point a, Tf s) { return a / length(a) * s; }
// returns an unit vector perpendicular to vector a, Tf
    Ti same
Point normal(Point a) {
 Tf 1 = length(a);
 return Point(-a.y / 1, a.x / 1);
// returns 1 if c is left of ab, 0 if on ab && -1 if
    right of ab
int orient(Point a, Point b, Point c) { return dcmp(cross
    (b - a, c - a)): 
/// Use as sort(v.begin(), v.end(), polarComp(0, dir))
/// Polar comparator around O starting at direction dir
struct polarComp {
 Point O, dir;
```

```
polarComp(Point 0 = Point(0, 0), Point dir = Point(1,
      0)) : O(O), dir(dir) {}
 bool half(Point p) {
   return dcmp(cross(dir, p)) < 0 ||</pre>
          (dcmp(cross(dir, p)) == 0 && dcmp(dot(dir, p))
 bool operator()(Point p, Point q) {
   return make_tuple(half(p), 0) < make_tuple(half(q),</pre>
       cross(p, q));
struct Segment {
 Point a, b;
 Segment(Point aa, Point bb) : a(aa), b(bb) {}
typedef Segment Line;
struct Circle {
 Point o:
 Tf r:
 Circle(Point o = Point(0, 0), Tf r = 0) : o(o), r(r) {}
 // returns true if point p is in || on the circle
 bool contains(Point p) { return dcmp(sqLength(p - o) -
      r * r) <= 0: 
 // returns a point on the circle rad radians away from
 Point point(Tf rad) {
   static_assert(is_same<Tf, Ti>::value);
 }
 // area of a circular sector with central angle rad
 Tf area(Tf rad = PI + PI) { return rad * r * r / 2; }
 // area of the circular sector cut by a chord with
      central angle alpha
 Tf sector(Tf alpha) { return r * r * 0.5 * (alpha - sin // returns the distance from point a to line 1
      (alpha)); }
4.2 Linear
```

```
// **** LINE LINE INTERSECTION START ****
// returns true if point p is on segment s
bool onSegment(Point p, Segment s) {
 return dcmp(cross(s.a - p, s.b - p)) == 0 && dcmp(dot(s
      .a - p, s.b - p)) <= 0;
// returns true if segment p && q touch or intersect
bool segmentsIntersect(Segment p, Segment q) {
 if (onSegment(p.a, q) || onSegment(p.b, q)) return true
```

```
if (onSegment(q.a, p) || onSegment(q.b, p)) return true
                                                        Ti c1 = cross(p.b - p.a, q.a - p.a);
                                                        Ti c2 = cross(p.b - p.a, q.b - p.a);
                                                        Ti c3 = cross(q.b - q.a, p.a - q.a);
                                                        Ti c4 = cross(q.b - q.a, p.b - q.a);
                                                        return dcmp(c1) * dcmp(c2) < 0 && dcmp(c3) * dcmp(c4) <
                                                       bool linesParallel(Line p, Line q) {
                                                        return dcmp(cross(p.b - p.a, q.b - q.a)) == 0;
                                                       // lines are represented as a ray from a point: (point,
                                                       // returns false if two lines (p, v) && (q, w) are
                                                           parallel or collinear
                                                       // true otherwise, intersection point is stored at o via
                                                           reference, Tf Ti Same
                                                       bool lineLineIntersection(Point p, Point v, Point q,
                                                           Point w, Point& o) {
                                                        if (dcmp(cross(v, w)) == 0) return false;
                                                        Point u = p - q;
                                                        o = p + v * (cross(w, u) / cross(v, w));
                                                        return true:
                                                       // returns false if two lines p && q are parallel or
return Point(o.x + cos(rad) * r, o.y + sin(rad) * r); // true otherwise, intersection point is stored at o via
                                                           reference
                                                       bool lineLineIntersection(Line p, Line q, Point& o) {
                                                        return lineLineIntersection(p.a, p.b - p.a, q.a, q.b -
                                                             q.a, o);
                                                       // **** LINE LINE INTERSECTION FINISH ****
                                                       Tf distancePointLine(Point p, Line 1) {
                                                        return abs(cross(1.b - 1.a, p - 1.a) / length(1.b - 1.a
                                                       // returns the shortest distance from point a to segment
                                                       Tf distancePointSegment(Point p, Segment s) {
                                                        if (s.a == s.b) return length(p - s.a);
                                                        Point v1 = s.b - s.a, v2 = p - s.a, v3 = p - s.b;
                                                        if (dcmp(dot(v1, v2)) < 0)
                                                          return length(v2);
                                                        else if (dcmp(dot(v1, v3)) > 0)
                                                          return length(v3);
                                                        else
```

```
return abs(cross(v1, v2) / length(v1));
}
// returns the shortest distance from segment p to
    segment q

Tf distanceSegmentSegment(Segment p, Segment q) {
    if (segmentsIntersect(p, q)) return 0;
    Tf ans = distancePointSegment(p.a, q);
    ans = min(ans, distancePointSegment(p.b, q));
    ans = min(ans, distancePointSegment(q.a, p));
    ans = min(ans, distancePointSegment(q.b, p));
    return ans;
}
// returns the projection of point p on line 1, Tf Ti
    Same
Point projectPointLine(Point p, Line 1) {
    Point v = l.b - l.a;
    return l.a + v * ((Tf)dot(v, p - l.a) / dot(v, v));
}
```

#### 4.3 Circular

```
// Extremely inaccurate for finding near touches
// compute intersection of line 1 with circle c
// The intersections are given in order of the ray (l.a,
    1.b), Tf Ti same
vector<Point> circleLineIntersection(Circle c, Line 1) {
  vector<Point> ret;
 Point b = 1.b - 1.a, a = 1.a - c.o;
 Tf A = dot(b, b), B = dot(a, b);
 Tf C = dot(a, a) - c.r * c.r, D = B * B - A * C;
 if (D < -EPS) return ret;</pre>
 ret.push_back(1.a + b * (-B - sqrt(D + EPS)) / A);
  if (D > EPS) ret.push_back(1.a + b * (-B + sqrt(D)) / A
      );
 return ret;
// signed area of intersection of circle(c.o, c.r) &&
// triangle(c.o, s.a, s.b) [cross(a-o, b-o)/2]
Tf circleTriangleIntersectionArea(Circle c, Segment s) {
  using Linear::distancePointSegment;
 Tf OA = length(c.o - s.a);
 Tf OB = length(c.o - s.b);
 // sector
  if (dcmp(distancePointSegment(c.o, s) - c.r) >= 0)
   return angleBetween(s.a - c.o, s.b - c.o) * (c.r * c.
       r) / 2.0;
  // triangle
  if (dcmp(OA - c.r) \le 0 \&\& dcmp(OB - c.r) \le 0)
   return cross(c.o - s.b, s.a - s.b) / 2.0;
  // three part: (A, a) (a, b) (b, B)
 vector<Point> Sect = circleLineIntersection(c, s);
```

```
return circleTriangleIntersectionArea(c, Segment(s.a,
      Sect[0])) +
        circleTriangleIntersectionArea(c, Segment(Sect
            [0]. Sect[1])) +
        circleTriangleIntersectionArea(c, Segment(Sect
            [1], s.b));
// area of intersecion of circle(c.o, c.r) && simple
    polyson(p[])
Tf circlePolyIntersectionArea(Circle c, Polygon p) {
 Tf res = 0:
 int n = p.size();
 for (int i = 0; i < n; ++i)</pre>
   res += circleTriangleIntersectionArea(c, Segment(p[i
       ], p[(i + 1) % n]));
 return abs(res);
// locates circle c2 relative to c1
// interior
                      (d < R - r)
// interior tangents (d = R - r)
// concentric
                   (d = 0)
// secants
                     (R - r < d < R + r) \longrightarrow 0
// exterior tangents (d = R + r)
// exterior
                      (d > R + r)
                                        ----> 2
int circleCirclePosition(Circle c1, Circle c2) {
 Tf d = length(c1.o - c2.o);
 int in = dcmp(d - abs(c1.r - c2.r)), ex = dcmp(d - (c1.r))
      r + c2.r):
 return in < 0 ? -2 : in == 0 ? -1 : ex == 0 ? 1 : ex >
      0 ? 2 : 0:
^{\prime}/ compute the intersection points between two circles c1
     && c2, Tf Ti same
vector<Point> circleCircleIntersection(Circle c1, Circle
    c2) {
 vector<Point> ret;
 Tf d = length(c1.o - c2.o);
 if (dcmp(d) == 0) return ret;
 if (dcmp(c1.r + c2.r - d) < 0) return ret;
 if (dcmp(abs(c1.r - c2.r) - d) > 0) return ret;
 Point v = c2.0 - c1.0;
 Tf co = (c1.r * c1.r + sqLength(v) - c2.r * c2.r) / (2)
      * c1.r * length(v));
 Tf si = sqrt(abs(1.0 - co * co));
 Point p1 = scale(rotatePrecise(v, co, -si), c1.r) + c1.
 Point p2 = scale(rotatePrecise(v, co, si), c1.r) + c1.o
```

```
ret.push_back(p1);
 if (p1 != p2) ret.push_back(p2);
 return ret;
// intersection area between two circles c1, c2
Tf circleCircleIntersectionArea(Circle c1, Circle c2) {
 Point AB = c2.0 - c1.0:
 Tf d = length(AB);
 if (d \ge c1.r + c2.r) return 0;
 if (d + c1.r <= c2.r) return PI * c1.r * c1.r;</pre>
 if (d + c2.r <= c1.r) return PI * c2.r * c2.r;</pre>
 Tf alpha1 = acos((c1.r * c1.r + d * d - c2.r * c2.r) /
     (2.0 * c1.r * d));
 Tf alpha2 = acos((c2.r * c2.r + d * d - c1.r * c1.r) /
     (2.0 * c2.r * d));
 return c1.sector(2 * alpha1) + c2.sector(2 * alpha2);
// returns tangents from a point p to circle c, Tf Ti
vector<Point> pointCircleTangents(Point p, Circle c) {
 vector<Point> ret;
 Point u = c.o - p;
 Tf d = length(u);
 if (d < c.r)
 else if (dcmp(d - c.r) == 0) {
   ret = {rotate(u, PI / 2)};
 } else {
   Tf ang = asin(c.r / d);
   ret = {rotate(u, -ang), rotate(u, ang)};
 return ret;
// returns the points on tangents that touches the circle
    , Tf Ti Same
vector<Point> pointCircleTangencyPoints(Point p, Circle c
   ) {
 Point u = p - c.o;
 Tf d = length(u);
 if (d < c.r)
   return {};
 else if (dcmp(d - c.r) == 0)
   return {c.o + u};
 else {
   Tf ang = acos(c.r / d);
   u = u / length(u) * c.r;
   return {c.o + rotate(u, -ang), c.o + rotate(u, ang)};
```

```
// for two circles c1 && c2, returns two list of points a }
// such that a[i] is on c1 && b[i] is c2 && for every i
// Line(a[i], b[i]) is a tangent to both circles
// CAUTION: a[i] = b[i] in case they touch | -1 for c1 =
int circleCircleTangencyPoints(Circle c1, Circle c2,
    vector<Point> &a,
                            vector<Point> &b) {
  a.clear(), b.clear();
  int cnt = 0;
  if (dcmp(c1.r - c2.r) < 0) {
   swap(c1, c2);
   swap(a, b);
 Tf d2 = sqLength(c1.o - c2.o);
 Tf rdif = c1.r - c2.r, rsum = c1.r + c2.r;
  if (dcmp(d2 - rdif * rdif) < 0) return 0;</pre>
  if (dcmp(d2) == 0 \&\& dcmp(c1.r - c2.r) == 0) return -1;
 Tf base = angle(c2.o - c1.o);
  if (dcmp(d2 - rdif * rdif) == 0) {
   a.push_back(c1.point(base));
   b.push_back(c2.point(base));
   cnt++:
   return cnt;
 Tf ang = acos((c1.r - c2.r) / sqrt(d2));
  a.push_back(c1.point(base + ang));
  b.push_back(c2.point(base + ang));
  a.push_back(c1.point(base - ang));
  b.push_back(c2.point(base - ang));
  cnt++;
  if (dcmp(d2 - rsum * rsum) == 0) {
   a.push_back(c1.point(base));
   b.push_back(c2.point(PI + base));
   cnt++:
  } else if (dcmp(d2 - rsum * rsum) > 0) {
   Tf ang = acos((c1.r + c2.r) / sqrt(d2));
   a.push_back(c1.point(base + ang));
   b.push_back(c2.point(PI + base + ang));
   a.push_back(c1.point(base - ang));
   b.push_back(c2.point(PI + base - ang));
   cnt++;
 return cnt;
```

#### 4.4 Convex

```
/// minkowski sum of two polygons in O(n)
Polygon minkowskiSum(Polygon A, Polygon B) {
 int n = A.size(), m = B.size();
 rotate(A.begin(), min_element(A.begin(), A.end()), A.
      end()):
 rotate(B.begin(), min_element(B.begin(), B.end()), B.
      end()):
 A.push_back(A[0]);
  B.push_back(B[0]);
  for (int i = 0; i < n; i++) A[i] = A[i + 1] - A[i];
 for (int i = 0; i < m; i++) B[i] = B[i + 1] - B[i];
 Polygon C(n + m + 1);
  C[0] = A.back() + B.back();
  merge(A.begin(), A.end() - 1, B.begin(), B.end() - 1, C
      .begin() + 1,
       polarComp(Point(0, 0), Point(0, -1)));
 for (int i = 1; i < C.size(); i++) C[i] = C[i] + C[i -</pre>
 C.pop_back();
 return C;
// finds the rectangle with minimum area enclosing a
    convex polygon and
// the rectangle with minimum perimeter enclosing a
    convex polygon
// Tf Ti Same
pair<Tf, Tf> rotatingCalipersBoundingBox(const Polygon &p
 using Linear::distancePointLine;
 int n = p.size();
 int 1 = 1, r = 1, j = 1;
 Tf area = 1e100;
 Tf perimeter = 1e100;
 for (int i = 0; i < n; i++) {</pre>
   Point v = (p[(i + 1) \% n] - p[i]) / length(p[(i + 1)
        % n] - p[i]);
   while (dcmp(dot(v, p[r \% n] - p[i]) - dot(v, p[(r +
        1) \% n] - p[i])) < 0)
     r++:
   while (j < r \mid | dcmp(cross(v, p[j % n] - p[i]) -
                       cross(v, p[(j + 1) % n] - p[i])) <
                            0)
     j++;
   while (1 < j ||
```

```
dcmp(dot(v, p[1 \% n] - p[i]) - dot(v, p[(1 + 1)
               n - p[i] > 0
     1++:
   Tf w = dot(v, p[r \% n] - p[i]) - dot(v, p[1 \% n] - p[
   Tf h = distancePointLine(p[j % n], Line(p[i], p[(i +
       1) % n]));
   area = min(area, w * h);
   perimeter = min(perimeter, 2 * w + 2 * h);
 }
 return make_pair(area, perimeter);
// returns the left side of polygon u after cutting it by
     ray a->b
Polygon cutPolygon(Polygon u, Point a, Point b) {
 using Linear::lineLineIntersection;
 using Linear::onSegment;
 Polygon ret;
 int n = u.size();
 for (int i = 0; i < n; i++) {</pre>
   Point c = u[i], d = u[(i + 1) \% n];
   if (dcmp(cross(b - a, c - a)) >= 0) ret.push_back(c);
   if (dcmp(cross(b - a, d - c)) != 0) {
     Point t:
     lineLineIntersection(a, b - a, c, d - c, t);
     if (onSegment(t, Segment(c, d))) ret.push_back(t);
 }
 return ret:
// returns true if point p is in or on triangle abc
bool pointInTriangle(Point a, Point b, Point c, Point p)
 return dcmp(cross(b - a, p - a)) >= 0 && dcmp(cross(c -
      b, p - b)) >= 0 &&
       dcmp(cross(a - c, p - c)) >= 0;
// pt must be in ccw order with no three collinear points
// returns inside = -1, on = 0, outside = 1
int pointInConvexPolygon(const Polygon &pt, Point p) {
 int n = pt.size();
 assert(n >= 3);
 int lo = 1, hi = n - 1;
 while (hi - lo > 1) {
   int mid = (lo + hi) / 2;
   if (dcmp(cross(pt[mid] - pt[0], p - pt[0])) > 0)
     lo = mid;
   else
```

```
hi = mid:
 bool in = pointInTriangle(pt[0], pt[lo], pt[hi], p);
 if (!in) return 1;
 if (dcmp(cross(pt[lo] - pt[lo - 1], p - pt[lo - 1])) == |}
       0) return 0;
 if (dcmp(cross(pt[hi] - pt[lo], p - pt[lo])) == 0)
      return 0;
 if (dcmp(cross(pt[hi] - pt[(hi + 1) % n], p - pt[(hi +
      1) % n)) == 0)
   return 0:
 return -1;
// Extreme Point for a direction is the farthest point in
     that direction
// u is the direction for extremeness
int extremePoint(const Polygon &poly, Point u) {
 int n = (int)poly.size();
 int a = 0, b = n;
 while (b - a > 1) {
   int c = (a + b) / 2;
   if (dcmp(dot(poly[c] - poly[(c + 1) % n], u)) >= 0 &&
       dcmp(dot(poly[c] - poly[(c - 1 + n) \% n], u)) >=
           0) {
     return c;
   }
   bool a_up = dcmp(dot(poly[(a + 1) % n] - poly[a], u))
   bool c_{up} = dcmp(dot(poly[(c + 1) % n] - poly[c], u))
   bool a_above_c = dcmp(dot(poly[a] - poly[c], u)) > 0;
   if (a_up && !c_up)
     b = c;
   else if (!a_up && c_up)
     a = c;
   else if (a_up && c_up) {
     if (a_above_c)
      b = c;
     else
       a = c;
   } else {
     if (!a_above_c)
      b = c;
     else
       a = c;
```

```
}
 if (dcmp(dot(poly[a] - poly[(a + 1) % n], u)) > 0 &&
     dcmp(dot(poly[a] - poly[(a - 1 + n) % n], u)) > 0)
   return a:
 return b % n;
// For a convex polygon p and a line 1, returns a list of
// of p that touch or intersect line 1.
// the i'th segment is considered (p[i], p[(i + 1) modulo
// #1 If a segment is collinear with the line, only that
    is returned
// #2 Else if 1 goes through i'th point, the i'th segment }
// Complexity: O(lg |p|)
vector<int> lineConvexPolyIntersection(const Polygon &p,
    Line 1) {
 assert((int)p.size() >= 3);
 assert(1.a != 1.b);
 int n = p.size();
 vector<int> ret;
 Point v = 1.b - 1.a;
 int lf = extremePoint(p, rotate90(v));
 int rt = extremePoint(p, rotate90(v) * Ti(-1));
 int olf = orient(l.a, l.b, p[lf]);
 int ort = orient(l.a, l.b, p[rt]);
 if (!olf || !ort) {
   int idx = (!olf ? lf : rt);
   if (orient(1.a, 1.b, p[(idx - 1 + n) \% n]) == 0)
     ret.push_back((idx - 1 + n) \% n);
   else
     ret.push_back(idx);
   return ret;
 if (olf == ort) return ret;
 for (int i = 0; i < 2; ++i) {</pre>
   int lo = i ? rt : lf;
   int hi = i ? lf : rt;
   int olo = i ? ort : olf;
   while (true) {
     int gap = (hi - lo + n) \% n;
     if (gap < 2) break;</pre>
```

```
int mid = (lo + gap / 2) \% n;
     int omid = orient(l.a, l.b, p[mid]);
     if (!omid) {
       lo = mid:
       break;
     if (omid == olo)
       lo = mid;
     else
       hi = mid;
   ret.push_back(lo);
 return ret;
// Calculate [ACW, CW] tangent pair from an external
    point
constexpr int CW = -1, ACW = 1;
bool isGood(Point u, Point v, Point Q, int dir) {
 return orient(Q, u, v) != -dir;
Point better(Point u, Point v, Point Q, int dir) {
 return orient(Q, u, v) == dir ? u : v;
Point pointPolyTangent(const Polygon &pt, Point Q, int
    dir, int lo, int hi) {
 while (hi - lo > 1) {
   int mid = (lo + hi) / 2;
   bool pvs = isGood(pt[mid], pt[mid - 1], Q, dir);
   bool nxt = isGood(pt[mid], pt[mid + 1], Q, dir);
   if (pvs && nxt) return pt[mid];
   if (!(pvs || nxt)) {
     Point p1 = pointPolyTangent(pt, Q, dir, mid + 1, hi
     Point p2 = pointPolyTangent(pt, Q, dir, lo, mid -
     return better(p1, p2, Q, dir);
   }
   if (!pvs) {
     if (orient(Q, pt[mid], pt[lo]) == dir)
       hi = mid - 1:
     else if (better(pt[lo], pt[hi], Q, dir) == pt[lo])
       hi = mid - 1;
     else
       lo = mid + 1;
   }
   if (!nxt) {
     if (orient(Q, pt[mid], pt[lo]) == dir)
```

```
lo = mid + 1:
     else if (better(pt[lo], pt[hi], Q, dir) == pt[lo])
       hi = mid - 1;
     else
       lo = mid + 1;
   }
 }
 Point ret = pt[lo];
 for (int i = lo + 1; i <= hi; i++) ret = better(ret, pt</pre>
      [i], Q, dir);
 return ret;
// [ACW, CW] Tangent
pair<Point, Point> pointPolyTangents(const Polygon &pt,
    Point Q) {
 int n = pt.size();
  Point acw_tan = pointPolyTangent(pt, Q, ACW, 0, n - 1);
 Point cw_tan = pointPolyTangent(pt, Q, CW, 0, n - 1);
 return make_pair(acw_tan, cw_tan);
```

### 4.5 Polygon

```
typedef vector<Point> Polygon;
// removes redundant colinear points
// polygon can't be all colinear points
Polygon RemoveCollinear(const Polygon &poly) {
 Polygon ret;
 int n = poly.size();
 for (int i = 0; i < n; i++) {</pre>
   Point a = polv[i];
   Point b = poly[(i + 1) \% n];
   Point c = poly[(i + 2) \% n];
   if (dcmp(cross(b - a, c - b)) != 0 && (ret.empty() ||
         b != ret.back()))
     ret.push_back(b);
 }
 return ret;
// returns the signed area of polygon p of n vertices
Tf signedPolygonArea(const Polygon &p) {
 Tf ret = 0;
 for (int i = 0; i < (int)p.size() - 1; i++)</pre>
   ret += cross(p[i] - p[0], p[i + 1] - p[0]);
 return ret / 2:
// given a polygon p of n vertices, generates the convex
    hull in in CCW
// Tested on https://acm.timus.ru/problem.aspx?space=1&
    num=1185
```

```
// Caution: when all points are colinear AND
    removeRedundant == false
// output will be contain duplicate points (from upper
    hull) at back
Polygon convexHull(Polygon p, bool removeRedundant) {
 int check = removeRedundant ? 0 : -1;
 sort(p.begin(), p.end());
 p.erase(unique(p.begin(), p.end()), p.end());
 int n = p.size();
 Polygon ch(n + n);
 int m = 0; // preparing lower hull
 for (int i = 0; i < n; i++) {</pre>
   while (m > 1 &&
          dcmp(cross(ch[m-1]-ch[m-2], p[i]-ch[m
               1])) <= check)
     m--:
   ch[m++] = p[i];
 int k = m; // preparing upper hull
 for (int i = n - 2; i \ge 0; i--) {
   while (m > k &&
          dcmp(cross(ch[m-1]-ch[m-2], p[i]-ch[m
               2])) <= check)
     m--:
   ch[m++] = p[i];
 if (n > 1) m--:
 ch.resize(m);
 return ch:
// returns inside = -1, on = 0, outside = 1
int pointInPolygon(const Polygon &p, Point o) {
 using Linear::onSegment;
 int wn = 0, n = p.size();
 for (int i = 0; i < n; i++) {</pre>
   int j = (i + 1) \% n;
   if (onSegment(o, Segment(p[i], p[j])) || o == p[i])
   int k = dcmp(cross(p[j] - p[i], o - p[i]));
   int d1 = dcmp(p[i].y - o.y);
   int d2 = dcmp(p[j].y - o.y);
   if (k > 0 \&\& d1 \le 0 \&\& d2 > 0) wn++;
   if (k < 0 \&\& d2 \le 0 \&\& d1 > 0) wn--;
 return wn ? -1 : 1:
// Given a simple polygon p, and a line l, returns (x, y)
// x = longest segment of 1 in p, y = total length of 1
    in p.
```

```
pair<Tf, Tf> linePolygonIntersection(Line 1, const
    Polygon &p) {
 using Linear::lineLineIntersection;
 int n = p.size();
 vector<pair<Tf, int>> ev;
 for (int i = 0; i < n; ++i) {</pre>
   Point a = p[i], b = p[(i + 1) \% n], z = p[(i - 1 + n)
        % n];
   int ora = orient(l.a, l.b, a), orb = orient(l.a, l.b,
       orz = orient(l.a, l.b, z);
   if (!ora) {
     Tf d = dot(a - 1.a, 1.b - 1.a);
     if (orz && orb) {
       if (orz != orb) ev.emplace_back(d, 0);
       // else // Point Touch
     } else if (orz)
       ev.emplace_back(d, orz);
     else if (orb)
       ev.emplace_back(d, orb);
   } else if (ora == -orb) {
     Point ins:
     lineLineIntersection(1, Line(a, b), ins);
     ev.emplace_back(dot(ins - 1.a, 1.b - 1.a), 0);
 }
 sort(ev.begin(), ev.end());
 Tf ans = 0, len = 0, last = 0, tot = 0;
 bool active = false:
 int sign = 0;
 for (auto &qq : ev) {
   int tp = qq.second;
   Tf d = qq.first; /// current Segment is (last, d)
   if (sign) {
                   /// On Border
     len += d - last;
     tot += d - last;
     ans = max(ans, len);
     if (tp != sign) active = !active;
     sign = 0;
   } else {
     if (active) { /// Strictly Inside
       len += d - last;
       tot += d - last;
       ans = max(ans, len);
     if (tp == 0)
       active = !active;
       sign = tp;
```

```
}
  last = d;
  if (!active) len = 0;
}
ans /= length(l.b - l.a);
tot /= length(l.b - l.a);
return {ans, tot};
}
```

# 5 Graph

# 5.1 Graph Template

```
struct edge {
 int u, v;
  edge(int u = 0, int v = 0) : u(u), v(v) {}
 int to(int node) { return u ^ v ^ node; }
struct graph {
 int n;
 vector<vector<int>> adj;
  vector<edge> edges;
  graph(int n = 0) : n(n), adj(n) {}
  void addEdge(int u, int v, bool dir = 1) {
   adj[u].push_back(edges.size());
   if (dir) adj[v].push_back(edges.size());
   edges.emplace_back(u, v);
  int addNode() {
   adj.emplace_back();
   return n++;
  edge &operator()(int idx) { return edges[idx]; }
 vector<int> &operator[](int u) { return adj[u]; }
```

# 5.2 LCA, ETT, VT

```
struct lca_table {
  tree &T;
  int n, LOG = 20;
  vector<vector<int>> anc;
  vector<iint> level;

void setupLifting(int node, int par) {
  for (int v : T[node])
    if (v != par) {
     anc[v][0] = node, level[v] = level[node] + 1;
     for (int k = 1; k < LOG; k++) anc[v][k] = anc[anc
        [v][k - 1]][k - 1];
     setupLifting(v, node);
  }
}</pre>
```

```
lca_table(tree &T, int root = 0) : T(T), n(T.n) {
   LOG = 33 - \_builtin\_clz(n);
   anc.assign(n, vector<int>(LOG, root));
   level.resize(n):
   setupLifting(root, root);
 int lca(int u, int v) {
   if (level[u] > level[v]) swap(u, v);
   for (int k = LOG - 1; ~k; k--)
     if (level[u] + (1 << k) <= level[v]) v = anc[v][k];</pre>
   if (u == v) return u:
   for (int k = LOG - 1; ~k; k--)
     if (anc[u][k] != anc[v][k]) u = anc[u][k], v = anc[
         vl[k]:
   return anc[u][0];
 int getAncestor(int node, int ht) {
   for (int k = 0; k < LOG; k++)
     if (ht & (1 << k)) node = anc[node][k]:</pre>
   return node:
 int distance(int u, int v) {
   int g = lca(u, v);
   return level[u] + level[v] - 2 * level[g];
struct euler_tour {
 int time = 0:
 tree &T;
 int n:
 vector<int> start, finish, level, par;
 euler_tour(tree &T, int root = 0)
     : T(T), n(T.n), start(n), finish(n), level(n), par(
         n) {
   time = 0:
   call(root);
 void call(int node, int p = -1) {
   if (p != -1) level[node] = level[p] + 1;
   start[node] = time++:
   for (int e : T[node])
     if (e != p) call(e, node);
   par[node] = p;
   finish[node] = time++;
 bool isAncestor(int node, int par) {
   return start[par] <= start[node] and finish[par] >=
        finish[node];
 }
```

```
int subtreeSize(int node) { return finish[node] - start
      [node] + 1 >> 1; }
tree virtual tree(vector<int> &nodes. lca table &table.
    euler_tour &tour) {
 sort(nodes.begin(), nodes.end(),
      [&](int x, int y) { return tour.start[x] < tour.</pre>
          start[v]; });
 int n = nodes.size():
 for (int i = 0; i + 1 < n; i++)</pre>
   nodes.push_back(table.lca(nodes[i], nodes[i + 1]));
 sort(nodes.begin(), nodes.end());
 nodes.erase(unique(nodes.begin(), nodes.end()), nodes.
      end()):
 sort(nodes.begin(), nodes.end(),
      [&] (int x, int y) { return tour.start[x] < tour.
          start[v]: }):
 n = nodes.size();
 stack<int> st:
 st.push(0);
 tree ans(n);
 for (int i = 1; i < n; i++) {
   while (!tour.isAncestor(nodes[i], nodes[st.top()]))
       st.pop();
   ans.addEdge(st.top(), i);
   st.push(i);
 }
 return ans;
set<int> getCenters(tree &T) {
 int n = T.n;
 vector<int> deg(n), q;
 set<int> s:
 for (int i = 0; i < n; i++) {</pre>
   deg[i] = T[i].size();
   if (deg[i] == 1) q.push_back(i);
   s.insert(i);
 for (vector<int> t; s.size() > 2; q = t) {
   for (auto x : q) {
     for (auto e : T[x])
       if (--deg[e] == 1) t.push_back(e);
     s.erase(x):
 }
 return s:
```

#### 5.3 SCC

typedef long long LL;

```
const LL N = 1e6 + 7;
bool vis[N];
vector<int> adj[N], adjr[N];
vector<int> order, component;
// tp = 0 ,finding topo order, tp = 1 , reverse edge
    traversal
void dfs(int u, int tp = 0) {
 vis[u] = true;
  if (tp) component.push_back(u);
  auto& ad = (tp ? adjr : adj);
 for (int v : ad[u])
   if (!vis[v]) dfs(v, tp);
 if (!tp) order.push_back(u);
int main() {
 for (int i = 1; i <= n; i++) {
   if (!vis[i]) dfs(i);
 memset(vis, 0, sizeof vis);
 reverse(order.begin(), order.end());
 for (int i : order) {
   if (!vis[i]) {
     // one component is found
     dfs(i, 1), component.clear();
   }
 }
```

```
5.4 Euler Tour on Edge
// for simplicity, G[idx] contains the adjacency list of
// while G(e) is a reference to the e-th edge.
const int N = 2e5 + 5;
int in[N], out[N], fwd[N], bck[N];
int t = 0;
void dfs(graph &G, int node, int par) {
 out[node] = t;
 for (int e : G[node]) {
   int v = G(e).to(node);
   if (v == par) continue;
   fwd[e] = t++;
   dfs(G, v, node);
   bck[e] = t++;
 in[node] = t - 1;
void init(graph &G, int node) {
 t = 0;
```

```
dfs(G, node, node);
|5.5 \quad LCA \text{ In O}(1)|
/* LCA in O(1)
* depth calculates weighted distance
* level calculates distance by number of edges
* Preprocessing in NlongN */
LL depth[N];
int level[N];
int st[N], en[N], LOG[N], par[N];
int a[N], id[N], table[L][N];
vector<PII> adj[N];
int n, root, Time, cur;
void init(int nodes, int root_) {
 n = nodes, root = root_, LOG[0] = LOG[1] = 0;
 for (int i = 2; i <= n; i++) LOG[i] = LOG[i >> 1] + 1;
 for (int i = 0; i <= n; i++) adj[i].clear();</pre>
void addEdge(int u, int v, int w) {
 adj[u].push_back(PII(v, w));
 adj[v].push_back(PII(u, w));
int lca(int u, int v) {
 if (en[u] > en[v]) swap(u, v);
 if (st[v] <= st[u] && en[u] <= en[v]) return v;</pre>
 int 1 = LOG[id[v] - id[u] + 1];
 int p1 = id[u], p2 = id[v] - (1 << 1) + 1;
 int d1 = level[table[1][p1]], d2 = level[table[1][p2]]; struct Segtree {
 if (d1 < d2)
   return par[table[1][p1]];
   return par[table[1][p2]];
LL dist(int u, int v) {
 int 1 = lca(u, v);
 return (depth[u] + depth[v] - (depth[1] * 2));
/* Euler tour */
void dfs(int u, int p) {
 st[u] = ++Time, par[u] = p;
```

```
for (auto [v, w] : adj[u]) {
   if (v == p) continue;
   depth[v] = depth[u] + w;
   level[v] = level[u] + 1;
   dfs(v, u);
 }
 en[u] = ++Time:
 a[++cur] = u, id[u] = cur;
/* RMQ */
void pre() {
 cur = Time = 0, dfs(root, root);
 for (int i = 1; i <= n; i++) table[0][i] = a[i];
 for (int 1 = 0; 1 < L - 1; 1++) {
   for (int i = 1; i <= n; i++) {
     table[1 + 1][i] = table[1][i];
     bool C1 = (1 << 1) + i <= n;
     bool C2 = level[table[1][i + (1 << 1)]] < level[</pre>
         table[1][i]];
     if (C1 && C2) table[1 + 1][i] = table[1][i + (1 <<</pre>
         1)];
   }
}
```

#### 5.6 HLD

```
const int N = 1e6 + 7;
template <typename DT>
 // write lazy segtree here
Segtree<int> tree(N);
vector<int> adj[N];
int depth[N], par[N], pos[N];
int head[N], heavy[N], cnt;
int dfs(int u, int p) {
 int SZ = 1, mxsz = 0, heavyc;
 depth[u] = depth[p] + 1;
 for (auto v : adj[u]) {
   if (v == p) continue;
   par[v] = u;
   int subsz = dfs(v, u);
```

```
if (subsz > mxsz) heavy[u] = v, mxsz = subsz;
   SZ += subsz;
 return SZ:
void decompose(int u, int h) {
 head[u] = h, pos[u] = ++cnt;
 if (heavy[u] != -1) decompose(heavy[u], h);
 for (int v : adj[u]) {
   if (v == par[u]) continue;
   if (v != heavy[u]) decompose(v, v);
int query(int a, int b) {
 int ret = 0;
 for (; head[a] != head[b]; b = par[head[b]]) {
   if (depth[head[a]] > depth[head[b]]) swap(a, b);
   ret += tree.query(1, 0, cnt, pos[head[b]], pos[b]);
  if (depth[a] > depth[b]) swap(a, b);
 ret += tree.query(1, 0, cnt, pos[a], pos[b]);
 return ret;
```

#### 5.7 Centroid Decomposition

```
class Centroid_Decomposition {
 vector<bool> blocked;
 vector<int> CompSize;
 int CompDFS(tree &T, int node, int par) {
   CompSize[node] = 1;
   for (int &e : T[node])
     if (e != par and !blocked[e]) CompSize[node] +=
         CompDFS(T, e, node);
   return CompSize[node];
 int FindCentroid(tree &T, int node, int par, int sz) {
   for (int &e : T[node])
     if (e != par and !blocked[e])
       if (CompSize[e] > sz / 2) return FindCentroid(T,
           e, node, sz);
   return node;
 pair<int, int> GetCentroid(tree &T, int entry) {
   int sz = CompDFS(T, entry, entry);
   return {FindCentroid(T, entry, entry, sz), sz};
 c_vector<LL> left[2], right[2];
 int GMin, GMax;
```

```
void dfs(tree &T, int node, int par, int Min, int Max,
      int sum) {
   if (blocked[node]) return;
   right[Max < sum or Min > sum][sum]++;
   Max = max(Max, sum), Min = min(Min, sum);
   GMin = min(GMin, sum), GMax = max(GMax, sum);
   for (int i = 0; i < T[node].size(); i++)</pre>
     if (T[node][i] != par) {
       dfs(T, T[node][i], node, Min, Max, sum + T.col[
           node][i]);
     }
 }
 LL solve(tree &T, int c, int sz) {
   LL ans = 0;
   left[0].clear(-sz, sz), left[1].clear(-sz, sz);
   for (int i = 0; i < T[c].size(); i++) {</pre>
     GMin = 1, GMax = -1;
     dfs(T, T[c][i], c, GMin, GMax, T.col[c][i]);
     ans += right[0][0] + left[1][0] * right[1][0];
     for (int j : {0, 1})
       for (int k = GMin; k <= GMax; k++) {</pre>
         ans += right[j][k] * (left[0][-k] + (j == 0) *
             left[1][-k]);
       }
     for (int j : {0, 1})
       for (int k = GMin; k <= GMax; k++) {</pre>
         left[j][k] += right[j][k];
         right[j][k] = 0;
       }
   }
   return ans;
 public:
 LL operator()(tree &T, int entry) {
   blocked.resize(T.n);
   CompSize.resize(T.n);
   for (int i : {0, 1})
     left[i].resize(2 * T.n + 5), right[i].resize(2 * T.
         n + 5):
   auto [c, sz] = GetCentroid(T, entry);
   LL ans = solve(T, c, sz);
   blocked[c] = true:
   for (int e : T[c])
     if (!blocked[e]) ans += (*this)(T, e);
   return ans:
 }
5.8 Dinic Max Flow
```

```
// add_edge(new src, u, sum(in_demand[u]))
// add_edge(u, new sink, sum(out_demand[u]))
// add_edge(old sink, old src, inf)
// if (sum of lower bound == flow) then demand satisfied
// flow in every edge i = demand[i] + e.flow
using Ti = long long;
const Ti INF = 1LL << 60;</pre>
struct edge {
 int v, u;
 Ti cap, flow = 0;
 edge(int v, int u, Ti cap) : v(v), u(u), cap(cap) {}
const int N = 1e5 + 50;
vector<edge> edges;
vector<int> adj[N];
int m = 0, n;
int level[N], ptr[N];
queue<int> q;
bool bfs(int s, int t) {
 for (q.push(s), level[s] = 0; !q.empty(); q.pop()) {
   for (int id : adj[q.front()]) {
     auto &ed = edges[id];
     if (ed.cap - ed.flow > 0 and level[ed.u] == -1)
       level[ed.u] = level[ed.v] + 1, q.push(ed.u);
 }
 return level[t] != -1;
Ti dfs(int v, Ti pushed, int t) {
 if (pushed == 0) return 0;
 if (v == t) return pushed;
 for (int &cid = ptr[v]; cid < adj[v].size(); cid++) {</pre>
   int id = adj[v][cid];
   auto &ed = edges[id];
   if (level[v] + 1 != level[ed.u] || ed.cap - ed.flow <</pre>
         1) continue;
   Ti tr = dfs(ed.u, min(pushed, ed.cap - ed.flow), t);
   if (tr == 0) continue;
   ed.flow += tr;
   edges[id ^ 1].flow -= tr;
   return tr;
 }
 return 0:
void init(int nodes) {
 m = 0, n = nodes;
```

/// flow with demand(lower bound) only for DAG

// create new src and sink

```
for (int i = 0; i < n; i++) level[i] = -1, ptr[i] = 0,
      adj[i].clear();
void addEdge(int v, int u, Ti cap) {
  edges.emplace_back(v, u, cap), adj[v].push_back(m++);
  edges.emplace_back(u, v, 0), adj[u].push_back(m++);
Ti maxFlow(int s, int t) {
 Ti f = 0:
  for (auto &ed : edges) ed.flow = 0;
  for (; bfs(s, t); memset(level, -1, n * 4)) {
   for (memset(ptr, 0, n * 4); Ti pushed = dfs(s, INF, t
        ); f += pushed)
     ;
 }
 return f;
```

### 5.9 Min Cost Max Flow

```
mt19937 rnd(chrono::steady_clock::now().time_since_epoch
    ().count());
const LL inf = 1e9;
struct edge {
  int v, rev;
  LL cap, cost, flow;
  edge() {}
  edge(int v, int rev, LL cap, LL cost)
     : v(v), rev(rev), cap(cap), cost(cost), flow(0) {}
};
struct mcmf {
  int src, sink, n;
  vector<int> par, idx, Q;
  vector<bool> inq;
  vector<LL> dis;
  vector<vector<edge>> g;
  mcmf() {}
  mcmf(int src, int sink, int n)
     : src(src),
       sink(sink),
       n(n),
       par(n),
       idx(n),
       inq(n),
       dis(n),
       g(n),
       \mathbb{Q}(10000005) {} // use \mathbb{Q}(n) if not using random
  void add_edge(int u, int v, LL cap, LL cost, bool
      directed = true) {
    edge _u = edge(v, g[v].size(), cap, cost);
    edge _v = edge(u, g[u].size(), 0, -cost);
```

```
g[u].pb(_u);
 g[v].pb(_v);
 if (!directed) add_edge(v, u, cap, cost, true);
bool spfa() {
 for (int i = 0; i < n; i++) {</pre>
   dis[i] = inf, inq[i] = false;
 int f = 0, 1 = 0;
 dis[src] = 0, par[src] = -1, Q[1++] = src, inq[src] =
       true:
 while (f < 1) {
   int u = Q[f++];
   for (int i = 0; i < g[u].size(); i++) {</pre>
     edge &e = g[u][i];
     if (e.cap <= e.flow) continue;</pre>
     if (dis[e.v] > dis[u] + e.cost) {
       dis[e.v] = dis[u] + e.cost;
       par[e.v] = u, idx[e.v] = i;
       if (!inq[e.v]) inq[e.v] = true, Q[1++] = e.v;
       // if (!ing[e.v]) {
       // inq[e.v] = true;
       // if (f \&\& rnd() \& 7) Q[--f] = e.v;
       // else Q[1++] = e.v;
       // }
     }
   inq[u] = false;
 return (dis[sink] != inf);
pair<LL, LL> solve() {
 LL mincost = 0, maxflow = 0;
 while (spfa()) {
   LL bottleneck = inf;
   for (int u = par[sink], v = idx[sink]; u != -1; v =
         idx[u], u = par[u]) {
     edge &e = g[u][v];
     bottleneck = min(bottleneck, e.cap - e.flow);
   for (int u = par[sink], v = idx[sink]; u != -1; v =
         idx[u], u = par[u]) {
     edge &e = g[u][v];
     e.flow += bottleneck;
     g[e.v][e.rev].flow -= bottleneck;
   mincost += bottleneck * dis[sink], maxflow +=
        bottleneck;
 }
 return make_pair(mincost, maxflow);
```

```
// want to minimize cost and don't care about flow
// add edge from sink to dummy sink (cap = inf, cost = 0)
// add edge from source to sink (cap = inf, cost = 0)
// run mcmf, cost returned is the minimum cost
```

```
5.10 Block Cut Tree
vector<vector<int> > components;
vector<int> cutpoints, start, low;
vector<bool> is_cutpoint;
stack<int> st;
void find_cutpoints(int node, graph &G, int par = -1, int
 low[node] = start[node] = d++;
 st.push(node);
 int cnt = 0;
 for (int e : G[node])
   if (int to = G(e).to(node); to != par) {
     if (start[to] == -1) {
       find_cutpoints(to, G, node, d + 1);
       cnt++:
       if (low[to] >= start[node]) {
         is_cutpoint[node] = par != -1 or cnt > 1;
         components.push_back({node}); // starting a new
             block with the point
         while (st.top() != node)
          components.back().push_back(st.top()), st.pop
               ();
     low[node] = min(low[node], low[to]);
   }
graph tree;
vector<int> id;
void init(graph &G) {
 int n = G.n;
 start.assign(n, -1), low.resize(n), is_cutpoint.resize(
     n), id.assign(n, -1);
 find_cutpoints(0, G);
 for (int u = 0; u < n; ++u)
   if (is_cutpoint[u]) id[u] = tree.addNode();
 for (auto &comp : components) {
   int node = tree.addNode();
   for (int u : comp)
     if (!is_cutpoint[u])
       id[u] = node;
       tree.addEdge(node, id[u]);
```

```
if (id[0] == -1) // corner - 1
   id[0] = tree.addNode();
}
5.11 Bridge Tree
```

vector<vector<int>> components;

```
vector<int> depth, low;
stack<int> st:
vector<int> id;
vector<edge> bridges;
graph tree;
void find_bridges(int node, graph &G, int par = -1, int d
 low[node] = depth[node] = d;
  st.push(node);
  for (int id : G[node]) {
   int to = G(id).to(node);
   if (par != to) {
     if (depth[to] == -1) {
       find_bridges(to, G, node, d + 1);
       if (low[to] > depth[node]) {
         bridges.emplace_back(node, to);
         components.push_back({});
         for (int x = -1; x != to; x = st.top(), st.pop()
           components.back().push_back(st.top());
       }
     low[node] = min(low[node], low[to]);
 }
 if (par == -1) {
   components.push_back({});
   while (!st.empty()) components.back().push_back(st.
       top()), st.pop();
 }
graph &create_tree() {
 for (auto &comp : components) {
   int idx = tree.addNode();
   for (auto &e : comp) id[e] = idx;
 for (auto &[1, r] : bridges) tree.addEdge(id[1], id[r])
 return tree;
void init(graph &G) {
 int n = G.n;
 depth.assign(n, -1), id.assign(n, -1), low.resize(n); | vector<int> primes;
```

```
for (int i = 0: i < n: i++)
 if (depth[i] == -1) find_bridges(i, G);
```

# 5.12 Tree Isomorphism

```
mp["01"] = 1;
ind = 1;
int dfs(int u, int p) {
 int cnt = 0:
 vector<int> vs;
 for (auto v : g1[u]) {
   if (v != p) {
     int got = dfs(v, u);
     vs.pb(got);
     cnt++;
   }
 }
 if (!cnt) return 1;
  sort(vs.begin(), vs.end());
  string s = "0";
 for (auto i : vs) s += to_string(i);
 vs.clear();
 s.pb('1');
 if (mp.find(s) == mp.end()) mp[s] = ++ind;
 int ret = mp[s];
 return ret;
```

#### Math

#### 6.1 Combi

const int N = 1e7;

```
array<int, N + 1> fact, inv, inv_fact;
void init() {
 fact[0] = inv_fact[0] = 1;
 for (int i = 1; i <= N; i++) {</pre>
   inv[i] = i == 1 ? 1 : (LL)inv[i - mod % i] * (mod / i | LL mul(LL a, LL b, LL mod) {
         + 1) % mod;
   fact[i] = (LL)fact[i - 1] * i % mod;
   inv fact[i] = (LL)inv fact[i - 1] * inv[i] % mod:
LL C(int n, int r) {
 return (r < 0 \text{ or } r > n) ? 0 : (LL)fact[n] * inv_fact[r]
       % mod * inv_fact[n - r] % mod;
6.2 Linear Sieve
```

```
int spf[N + 5], phi[N + 5], NOD[N + 5], cnt[N + 5], POW[N]
     + 5];
bool prime[N + 5];
int SOD[N + 5]:
void init() {
 fill(prime + 2, prime + N + 1, 1);
 SOD[1] = NOD[1] = phi[1] = spf[1] = 1;
 for (LL i = 2; i <= N; i++) {</pre>
   if (prime[i]) {
     primes.push_back(i), spf[i] = i;
     phi[i] = i - 1;
     NOD[i] = 2, cnt[i] = 1;
     SOD[i] = i + 1, POW[i] = i;
   for (auto p : primes) {
     if (p * i > N \text{ or } p > spf[i]) break;
     prime[p * i] = false, spf[p * i] = p;
     if (i % p == 0) {
       phi[p * i] = p * phi[i];
       NOD[p * i] = NOD[i] / (cnt[i] + 1) * (cnt[i] + 2)
              cnt[p * i] = cnt[i] + 1;
       SOD[p * i] = SOD[i] / SOD[POW[i]] * (SOD[POW[i]]
           + p * POW[i]),
              POW[p * i] = p * POW[i];
       break:
     } else {
       phi[p * i] = phi[p] * phi[i];
       NOD[p * i] = NOD[p] * NOD[i], cnt[p * i] = 1;
       SOD[p * i] = SOD[p] * SOD[i], POW[p * i] = p;
     }
   }
 }
```

#### 6.3 Pollard Rho

```
return (__int128)a * b % mod;
 // LL ans = a * b - mod * (LL) (1.L / mod * a * b);
 // return ans + mod * (ans < 0) - mod * (ans >= (LL)
     mod);
LL bigmod(LL num, LL pow, LL mod) {
 LL ans = 1;
 for (; pow > 0; pow >>= 1, num = mul(num, num, mod))
  if (pow & 1) ans = mul(ans, num, mod);
 return ans:
bool is_prime(LL n) {
 if (n < 2 or n % 6 % 4 != 1) return (n | 1) == 3;
```

```
LL a[] = \{2, 325, 9375, 28178, 450775, 9780504,
      1795265022};
  LL s = \__builtin\_ctzll(n - 1), d = n >> s;
  for (LL x : a) {
   LL p = bigmod(x \% n, d, n), i = s;
   for (; p != 1 and p != n - 1 and x % n and i--; p =
        mul(p, p, n))
   if (p != n - 1 and i != s) return false;
  return true;
LL get_factor(LL n) {
  auto f = [\&](LL x) \{ return mul(x, x, n) + 1; \};
  LL x = 0, y = 0, t = 0, prod = 2, i = 2, q;
  for (; t++ % 40 or gcd(prod, n) == 1; x = f(x), y = f(f)
      (y))) {
   (x == y) ? x = i++, y = f(x) : 0;
   prod = (q = mul(prod, max(x, y) - min(x, y), n)) ? q
  return gcd(prod, n);
map<LL, int> factorize(LL n) {
  map<LL, int> res;
  if (n < 2) return res;
  LL small_primes[] = {2, 3, 5, 7, 11, 13, 17, 19, 23,
      29, 31, 37, 41,
                     43, 47, 53, 59, 61, 67, 71, 73, 79,
                         83, 89, 97};
  for (LL p : small_primes)
   for (; n % p == 0; n /= p, res[p]++)
  auto _factor = [&](LL n, auto &_factor) {
   if (n == 1) return;
   if (is_prime(n))
     res[n]++:
   else {
     LL x = get_factor(n);
     _factor(x, _factor);
     _factor(n / x, _factor);
  _factor(n, _factor);
  return res;
```

#### 6.4 Chinese Remainder Theorem

```
\ensuremath{//} given a, b will find solutions for
```

```
// ax + by = 1
tuple<LL, LL, LL> EGCD(LL a, LL b) {
   if (b == 0)
     return {1, 0, a};
   else {
     auto [x, y, g] = EGCD(b, a % b);
     return {y, x - a / b * y, g};
   }
}
// given modulo equations, will apply CRT
PLL CRT(vector<PLL> &v) {
   LL V = 0, M = 1;
   for (auto &[v, m] : v) { // value % mod
     auto [x, y, g] = EGCD(M, m);
     if ((v - V) % g != 0) return {-1, 0};
     V += x * (v - V) / g % (m / g) * M, M *= m / g;
     V = (V % M + M) % M;
}
return make_pair(V, M);
}
```

### 6.5 Mobius Function

```
const int N = 1e6 + 5;
int mob[N];
void mobius() {
  memset(mob, -1, sizeof mob);
  mob[1] = 1;
  for (int i = 2; i < N; i++)
    if (mob[i]) {
      for (int j = i + i; j < N; j += i) mob[j] -= mob[i ];
      }
}</pre>
```

#### 6.6 FFT

```
using CD = complex<double>;
typedef long long LL;
const double PI = acos(-1.0L);
int N;
vector<int> perm;
vector<CD> wp[2];
void precalculate(int n) {
  assert((n & (n - 1)) == 0), N = n;
  perm = vector<int>(N, 0);
  for (int k = 1; k < N; k <<= 1) {
    for (int i = 0; i < k; i++) {
       perm[i] <<= 1;
       perm[i + k] = 1 + perm[i];
}</pre>
```

```
wp[0] = wp[1] = vector < CD > (N);
 for (int i = 0; i < N; i++) {</pre>
   wp[0][i] = CD(cos(2 * PI * i / N), sin(2 * PI * i / N)
   wp[1][i] = CD(cos(2 * PI * i / N), -sin(2 * PI * i / N))
       N)):
 }
void fft(vector<CD> &v, bool invert = false) {
 if (v.size() != perm.size()) precalculate(v.size());
 for (int i = 0; i < N; i++)</pre>
   if (i < perm[i]) swap(v[i], v[perm[i]]);</pre>
 for (int len = 2; len <= N; len *= 2) {</pre>
   for (int i = 0, d = N / len; i < N; i += len) {</pre>
     for (int j = 0, idx = 0; j < len / 2; j++, idx += d
         ) {
       CD x = v[i + j];
       CD y = wp[invert][idx] * v[i + j + len / 2];
       v[i + j] = x + y;
       v[i + j + len / 2] = x - y;
   }
 }
 if (invert) {
   for (int i = 0; i < N; i++) v[i] /= N;</pre>
void pairfft(vector<CD> &a, vector<CD> &b, bool invert =
    false) {
 int N = a.size();
 vector<CD> p(N);
 for (int i = 0; i < N; i++) p[i] = a[i] + b[i] * CD(0,
     1);
 fft(p, invert);
 p.push_back(p[0]);
 for (int i = 0; i < N; i++) {</pre>
   if (invert) {
     a[i] = CD(p[i].real(), 0);
     b[i] = CD(p[i].imag(), 0);
   } else {
     a[i] = (p[i] + conj(p[N - i])) * CD(0.5, 0);
     b[i] = (p[i] - conj(p[N - i])) * CD(0, -0.5);
   }
 }
vector<LL> multiply(const vector<LL> &a, const vector<LL>
     &b) {
 int n = 1;
 while (n < a.size() + b.size()) n <<= 1;</pre>
```

```
vector CD fa(a.begin(), a.end()), fb(b.begin(), b.end 6.7 NTT
      ());
 fa.resize(n);
 fb.resize(n):
          fft(fa); fft(fb);
 pairfft(fa, fb);
 for (int i = 0; i < n; i++) fa[i] = fa[i] * fb[i];</pre>
 fft(fa, true);
 vector<LL> ans(n):
 for (int i = 0; i < n; i++) ans[i] = round(fa[i].real()</pre>
     );
 return ans;
const int M = 1e9 + 7, B = sqrt(M) + 1;
vector<LL> anyMod(const vector<LL> &a, const vector<LL> &
 int n = 1:
 while (n < a.size() + b.size()) n <<= 1;</pre>
 vector<CD> al(n), ar(n), bl(n), br(n);
 for (int i = 0; i < a.size(); i++) al[i] = a[i] % M / B</pre>
      , ar[i] = a[i] % M % B;
 for (int i = 0; i < b.size(); i++) bl[i] = b[i] % M / B</pre>
      , br[i] = b[i] % M % B;
 pairfft(al, ar);
 pairfft(bl, br);
          fft(al); fft(ar); fft(bl); fft(br);
 for (int i = 0; i < n; i++) {
   CD ll = (al[i] * bl[i]), lr = (al[i] * br[i]);
   CD rl = (ar[i] * bl[i]), rr = (ar[i] * br[i]);
   al[i] = 11:
   ar[i] = lr;
   bl[i] = rl;
   br[i] = rr;
 pairfft(al, ar, true);
 pairfft(bl, br, true);
           fft(al, true); fft(ar, true); fft(bl, true);
      fft(br, true);
 vector<LL> ans(n);
 for (int i = 0; i < n; i++) {
   LL right = round(br[i].real()), left = round(al[i].
       real());
   LL mid = round(round(bl[i].real()) + round(ar[i].real
   ans[i] = ((left \% M) * B * B + (mid \% M) * B + right)
        % M;
 }
 return ans;
```

```
const LL N = 1 << 18;</pre>
const LL MOD = 786433;
vector<LL> P[N]:
LL rev[N], w[N | 1], a[N], b[N], inv_n, g;
LL Pow(LL b, LL p) {
 LL ret = 1;
 while (p) {
   if (p & 1) ret = (ret * b) % MOD;
   b = (b * b) \% MOD;
   p >>= 1;
 return ret;
LL primitive_root(LL p) {
 vector<LL> factor;
 LL phi = p - 1, n = phi;
 for (LL i = 2; i * i <= n; i++) {
   if (n % i) continue;
   factor.emplace_back(i);
   while (n \% i == 0) n /= i;
 if (n > 1) factor.emplace_back(n);
 for (LL res = 2; res <= p; res++) {</pre>
   bool ok = true:
   for (LL i = 0; i < factor.size() && ok; i++)</pre>
     ok &= Pow(res, phi / factor[i]) != 1;
   if (ok) return res;
 }
 return -1;
void prepare(LL n) {
 LL sz = abs(31 - __builtin_clz(n));
 LL r = Pow(g, (MOD - 1) / n);
 inv_n = Pow(n, MOD - 2);
 w[0] = w[n] = 1;
 for (LL i = 1; i < n; i++) w[i] = (w[i - 1] * r) % MOD; const LL MOD = 30011;
 for (LL i = 1: i < n: i++)
   rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << (sz - 1)); | LL BigMod(LL b, LL p) {
void NTT(LL *a, LL n, LL dir = 0) {
 for (LL i = 1; i < n - 1; i++)
   if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
 for (LL m = 2; m <= n; m <<= 1) {
   for (LL i = 0; i < n; i += m) {
     for (LL j = 0; j < (m >> 1); j++) {
       LL &u = a[i + j], &v = a[i + j + (m >> 1)];
       LL t = v * w[dir ? n - n / m * j : n / m * j] %
```

MOD;

```
v = u - t < 0 ? u - t + MOD : u - t;
       u = u + t >= MOD ? u + t - MOD : u + t;
   }
 }
 if (dir)
   for (LL i = 0; i < n; i++) a[i] = (inv_n * a[i]) %</pre>
        MOD;
vector<LL> mul(vector<LL> p, vector<LL> q) {
 LL n = p.size(), m = q.size();
 LL t = n + m - 1, sz = 1;
 while (sz < t) sz <<= 1;</pre>
 prepare(sz);
 for (LL i = 0; i < n; i++) a[i] = p[i];
 for (LL i = 0; i < m; i++) b[i] = q[i];
 for (LL i = n; i < sz; i++) a[i] = 0;</pre>
 for (LL i = m; i < sz; i++) b[i] = 0;</pre>
 NTT(a, sz);
 NTT(b, sz);
 for (LL i = 0; i < sz; i++) a[i] = (a[i] * b[i]) % MOD;</pre>
 NTT(a, sz, 1);
 vector<LL> c(a, a + sz);
 while (c.size() && c.back() == 0) c.pop_back();
 return c;
```

```
6.8 WalshHadamard
#include <bits/stdc++.h>
using namespace std;
typedef long long LL;
#define bitwiseXOR 1
// #define bitwiseAND 2
// #define bitwiseOR 3
 LL ret = 1;
 while (p > 0) {
   if (p % 2 == 1) {
    ret = (ret * b) \% MOD;
   }
   p = p / 2;
   b = (b * b) \% MOD:
 return ret % MOD;
```

```
void FWHT(vector<LL>& p, bool inverse) {
 LL n = p.size();
 assert((n & (n - 1)) == 0):
 for (LL len = 1; 2 * len <= n; len <<= 1) {
   for (LL i = 0: i < n: i += len + len) {
    for (LL j = 0; j < len; j++) {</pre>
      LL u = p[i + j];
      LL v = p[i + len + j];
#ifdef bitwiseXOR
       p[i + j] = (u + v) \% MOD;
      p[i + len + j] = (u - v + MOD) \% MOD;
#endif // bitwiseXOR
#ifdef bitwiseAND
       if (!inverse) {
        p[i + j] = v \% MOD;
        p[i + len + j] = (u + v) \% MOD;
      } else {
        p[i + j] = (-u + v) \% MOD;
        p[i + len + j] = u \% MOD;
#endif // bitwiseAND
#ifdef bitwiseOR
       if (!inverse) {
        p[i + j] = u + v;
        p[i + len + j] = u;
      } else {
        p[i + j] = v;
        p[i + len + j] = u - v;
#endif // bitwiseOR
    }
#ifdef bitwiseXOR
 if (inverse) {
   LL val = BigMod(n, MOD - 2); // Option 2: Exclude
   for (LL i = 0; i < n; i++) {</pre>
    // assert(p[i]%n==0); //Option 2: Include
     p[i] = (p[i] * val) % MOD; // Option 2: p[i]/=n;
   }
 }
#endif // bitwiseXOR
```

#### 6.9 Adaptive Simpsons

```
For finding the length of an arc in a range
   L = integrate(ds) from start to end of range
   where ds = sqrt(1+(d/dy(x))^2)dy
const double SIMPSON TERMINAL EPS = 1e-12:
/// Function whose integration is to be calculated
double F(double x);
double simpson(double minx, double maxx) {
 return (\max - \min x) / 6 * (F(\min x) + 4 * F((\min x + \max x)))
      maxx) / 2.) + F(maxx));
double adaptive_simpson(double minx, double maxx, double
    c, double EPS) {
 // if(maxx - minx < SIMPSON_TERMINAL_EPS) return 0;</pre>
 double midx = (minx + maxx) / 2:
 double a = simpson(minx, midx);
 double b = simpson(midx, maxx);
 if (fabs(a + b - c) < 15 * EPS) return a + b + (a + b - c)
       c) / 15.0;
 return adaptive_simpson(minx, midx, a, EPS / 2.) +
        adaptive_simpson(midx, maxx, b, EPS / 2.);
double adaptive_simpson(double minx, double maxx, double
 return adaptive_simpson(minx, maxx, simpson(minx, maxx,
       i), EPS);
```

#### 6.10 Berlekamp Massey

```
struct berlekamp_massey { // for linear recursion
 typedef long long LL;
 static const int SZ = 2e5 + 5;
 static const int MOD = 1e9 + 7; /// mod must be a prime
 LL m , a[SZ] , h[SZ] , t_{SZ} , s[SZ] , t[SZ];
 // bigmod goes here
 inline vector <LL> BM( vector <LL> &x ) {
  LL lf , ld;
   vector <LL> ls , cur;
   for ( int i = 0; i < int(x.size()); ++i ) {</pre>
    LL t = 0:
    for ( int j = 0; j < int(cur.size()); ++j ) t = (t</pre>
         + x[i - j - 1] * cur[j]) % MOD;
     if ((t - x[i]) \% MOD == 0) continue;
     if (!cur.size()) {
       cur.resize( i + 1 );
```

```
lf = i; ld = (t - x[i]) % MOD;
     continue;
   }
   LL k = -(x[i] - t) * bigmod(ld, MOD - 2, MOD) %
         MOD:
   vector <LL> c(i - lf - 1);
   c.push_back( k );
   for ( int j = 0; j < int(ls.size()); ++j ) c.</pre>
        push_back(-ls[j] * k % MOD);
   if ( c.size() < cur.size() ) c.resize( cur.size() )</pre>
   for ( int j = 0; j < int(cur.size()); ++j ) c[j] =</pre>
        (c[i] + cur[i]) % MOD;
   if (i - lf + (int)ls.size() >= (int)cur.size() ) ls
         = cur, lf = i, ld = (t - x[i]) % MOD;
  for ( int i = 0; i < int(cur.size()); ++i ) cur[i] =</pre>
      (cur[i] % MOD + MOD) % MOD;
 return cur:
inline void mull( LL *p , LL *q ) {
 for ( int i = 0; i < m + m; ++i ) t_{-}[i] = 0;
 for ( int i = 0; i < m; ++i ) if ( p[i] )</pre>
     for ( int j = 0; j < m; ++j ) t_{i} = (t_{i} + t_{i})
          j] + p[i] * q[j]) % MOD;
 for ( int i = m + m - 1; i >= m; --i ) if ( t_[i] )
     for (int j = m - 1; \tilde{j}; --j) t_{i} = 1
          t_{i} = i - 1 + t_{i} * h[i] * h[i] % MOD;
 for ( int i = 0; i < m; ++i ) p[i] = t_[i];</pre>
inline LL calc( LL K ) {
 for ( int i = m; ~i; --i ) s[i] = t[i] = 0;
 s[0] = 1; if ( m != 1 ) t[1] = 1; else t[0] = h[0];
 while (K) {
   if (K&1) mull(s,t);
   mull( t , t ); K >>= 1;
 LL su = 0:
 for ( int i = 0; i < m; ++i ) su = (su + s[i] * a[i])
       % MOD:
 return (su % MOD + MOD) % MOD;
}
/// already calculated upto k , now calculate upto n.
inline vector <LL> process( vector <LL> &x , int n ,
    int k ) {
  auto re = BM( x );
 x.resize(n+1);
 for ( int i = k + 1; i <= n; i++ ) {
   for ( int j = 0; j < re.size(); j++ ) {</pre>
```

```
x[i] += 1LL * x[i - j - 1] % MOD * re[j] % MOD; x
            [i] %= MOD;
     }
   }
   return x;
  inline LL work( vector <LL> &x , LL n ) {
   if ( n < int(x.size()) ) return x[n] % MOD;</pre>
   vector \langle LL \rangle v = BM(x); m = v.size(); if (!m)
        return 0;
   for ( int i = 0; i < m; ++i ) h[i] = v[i], a[i] = x[i</pre>
        ];
   return calc( n ) % MOD;
  }
} rec;
vector <LL> v;
void solve() {
 int n;
  cin >> n:
  cout << rec.work(v, n - 1) << endl;</pre>
```

### 6.11 Fractional Binary Search

```
/**
Given a function f and n, finds the smallest fraction p /
     q in [0, 1] or [0,n]
such that f(p / q) is true, and p, q <= n.
Time: O(log(n))
struct frac { long long p, q; };
bool f(frac x) {
return 6 + 8 * x.p >= 17 * x.q + 12;
frac fracBS(long long n) {
 bool dir = 1, A = 1, B = 1;
 frac lo{0, 1}, hi{1, 0}; // Set hi to 1/0 to search
      within [0, n] and \{1, 1\} to search within [0, 1]
  if (f(lo)) return lo;
 assert(f(hi)); //checking if any solution exists or not
  while (A || B) {
   long long adv = 0, step = 1; // move hi if dir, else
   for (int si = 0; step; (step *= 2) >>= si) {
     adv += step;
     frac mid{lo.p * adv + hi.p, lo.q * adv + hi.q};
     if (abs(mid.p) > n \mid\mid mid.q > n \mid\mid dir == !f(mid))
         {
       adv -= step; si = 2;
     }
   }
```

```
hi.p += lo.p * adv;
   hi.q += lo.q * adv;
   dir = !dir;
   swap(lo, hi);
   A = B; B = !!adv;
 return dir ? hi : lo;
6.12 Lagrange
// p is a polynomial with n points.
// p(0), p(1), p(2), \dots p(n-1) are given.
// Find p(x).
LL Lagrange(vector<LL> &p, LL x) {
 LL n = p.size(), L, i, ret;
 if (x < n) return p[x];</pre>
 L = 1:
 for (i = 1; i < n; i++) {
  L = (L * (x - i)) \% MOD:
   L = (L * bigmod(MOD - i, MOD - 2)) % MOD;
 ret = (L * p[0]) % MOD;
 for (i = 1; i < n; i++) {</pre>
   L = (L * (x - i + 1)) \% MOD;
   L = (L * bigmod(x - i, MOD - 2)) \% MOD;
   L = (L * bigmod(i, MOD - 2)) % MOD;
   L = (L * (MOD + i - n)) % MOD;
   ret = (ret + L * p[i]) % MOD;
 return ret;
```

# 6.13 Shanks' Baby Step, Giant Step

```
// Finds a^x = b (mod p)

LL bigmod(LL b, LL p, LL m) {}

LL babyStepGiantStep(LL a, LL b, LL p) {
    LL i, j, c, sq = sqrt(p);
    map<LL, LL> babyTable;

for (j = 0, c = 1; j <= sq; j++, c = (c * a) % p)
        babyTable[c] = j;

LL giant = bigmod(a, sq * (p - 2), p);

for (i = 0, c = 1; i <= sq; i++, c = (c * giant) % p) {
    if (babyTable.find((c * b) % p) != babyTable.end())
        return i * sq + babyTable[(c * b) % p];
}</pre>
```

```
return -1;
```

### 6.14 Xor Basis

```
struct XorBasis {
 static const int sz = 64;
 array<ULL, sz> base = {0}, back;
 array<int, sz> pos;
 void insert(ULL x, int p) {
  ULL cur = 0;
  for (int i = sz - 1; ~i; i--)
    if (x >> i & 1) {
      if (!base[i]) {
        base[i] = x, back[i] = cur, pos[i] = p;
      } else x ^= base[i], cur |= 1ULL << i;</pre>
    }
}
 pair<ULL, vector<int>> construct(ULL mask) {
  ULL ok = 0, x = mask;
  for (int i = sz - 1; ~i; i--)
    if (mask >> i & 1 and base[i]) mask ^= base[i], ok
         |= 1ULL << i;
  vector<int> ans;
  for (int i = 0; i < sz; i++)</pre>
    if (ok >> i & 1) {
      ans.push_back(pos[i]);
      ok ^= back[i];
  return {x ^ mask, ans};
}
```

# 7 String

#### 7.1 Aho Corasick

```
const int sg = 26, N = 1e3 + 9;
struct aho_corasick {
   struct node {
     node *link, *out, *par;
     bool leaf;
   LL val;
   int cnt, last, len;
   char p_ch;
   array<node*, sg> to;
   node(node* par = NULL, char p_ch = '$', int len = 0)
     : par(par), p_ch(p_ch), len(len) {
     val = leaf = cnt = last = 0;
     link = out = NULL;
```

```
}
  };
  vector<node> trie;
  node* root:
  aho_corasick() {
   trie.reserve(N), trie.emplace_back();
   root = &trie[0];
   root->link = root->out = root;
  inline int f(char c) { return c - 'a'; }
  inline node* add_node(node* par = NULL, char p_ch = '$' PLL inv_base = inverse(base);
      , int len = 0) {
   trie.emplace_back(par, p_ch, len);
   return &trie.back();
  void add_str(const string& s, LL val = 1) {
   node* now = root:
   for (char c : s) {
     int i = f(c):
     if (!now->to[i]) now->to[i] = add_node(now, c, now
          ->len + 1);
     now = now - > to[i]:
   }
   now->leaf = true, now->val++;
  void push_links() {
   queue<node*> q;
   for (q.push(root); q.empty(); q.pop()) {
     node *cur = q.front(), *link = cur->link;
     cur->out = link->leaf ? link : link->out:
     int idx = 0;
     for (auto& next : cur->to) {
       if (next != NULL) {
         next->link = cur != root ? link->to[idx++] :
             root;
         q.push(next);
       } else
         next = link->to[idx++];
     }
   }
   cur->val += link->val;
 }
};
```

### 7.2 Double hash

```
// define +, -, * for (PLL, LL) and (PLL, PLL), % for (
    PLL. PLL):
PLL base(1949313259, 1997293877);
PLL mod(2091573227, 2117566807);
```

```
PLL power(PLL a, LL p) {
 PLL ans = PLL(1, 1);
 for(; p; p >>= 1, a = a * a % mod) {
     if(p & 1) ans = ans * a % mod;
 return ans;
PLL inverse(PLL a) { return power(a, (mod.ff - 1) * (mod.
    ss - 1) - 1); }
PLL val:
vector<PLL> P;
void hash_init(int n) {
 P.resize(n + 1);
 P[0] = PLL(1, 1):
 for (int i = 1; i \le n; i++) P[i] = (P[i-1] * base) %
PLL append(PLL cur, char c) { return (cur * base + c) %
/// prepends c to string with size k
PLL prepend(PLL cur, int k, char c) { return (P[k] * c +
    cur) % mod: }
/// replaces the i-th (0-indexed) character from right
    from a to b:
PLL replace(PLL cur, int i, char a, char b) {
 cur = (cur + P[i] * (b - a)) \% mod;
  return (cur + mod) % mod;
/// Erases c from the back of the string
PLL pop_back(PLL hash, char c) {
 return (((hash - c) * inv_base) % mod + mod) % mod;
/// Erases c from front of the string with size len
PLL pop_front(PLL hash, int len, char c) {
 return ((hash - P[len - 1] * c) % mod + mod) % mod;
/// concatenates two strings where length of the right is 7.3 Manacher's
PLL concat(PLL left, PLL right, int k) { return (left * P
    [k] + right) % mod; }
/// Calculates hash of string with size len repeated cnt
/// This is O(\log n). For O(1), pre-calculate inverses
PLL repeat(PLL hash, int len, LL cnt) {
 PLL mul = (P[len * cnt] - 1) * inverse(P[len] - 1);
 mul = (mul % mod + mod) % mod;
 PLL ret = (hash * mul) % mod;
```

```
if (P[len].ff == 1) ret.ff = hash.ff * cnt:
 if (P[len].ss == 1) ret.ss = hash.ss * cnt;
 return ret:
LL get(PLL hash) { return ((hash.ff << 32) ^ hash.ss); }
struct hashlist {
 int len:
 vector<PLL> H, R;
 hashlist() {}
 hashlist(string& s) {
   len = (int)s.size();
   hash_init(len);
   H.resize(len + 1, PLL(0, 0)), R.resize(len + 2, PLL
       (0, 0));
   for (int i = 1; i <= len; i++) H[i] = append(H[i -</pre>
       1], s[i - 1]);
   for (int i = len; i >= 1; i--) R[i] = append(R[i +
       1], s[i - 1]);
 /// 1-indexed
 PLL range_hash(int 1, int r) {
   return ((H[r] - H[l - 1] * P[r - l + 1]) % mod + mod)
        % mod;
 PLL reverse hash(int 1, int r) {
   return ((R[1] - R[r + 1] * P[r - 1 + 1]) \% mod + mod)
        % mod;
 }
 PLL concat_range_hash(int 11, int r1, int 12, int r2) {
   return concat(range_hash(11, r1), range_hash(12, r2),
        r2 - 12 + 1);
 PLL concat_reverse_hash(int 11, int r1, int 12, int r2)
   return concat(reverse_hash(12, r2), reverse_hash(11,
       r1), r1 - 11 + 1);
 }
```

```
vector<int> d1(n);
// d[i] = number of palindromes taking s[i] as center
for (int i = 0, l = 0, r = -1; i < n; i++) {
 int k = (i > r) ? 1 : min(d1[1 + r - i], r - i + 1);
 while (0 \le i - k \&\& i + k \le n \&\& s[i - k] == s[i + k])
       k++;
 d1[i] = k--:
 if (i + k > r) l = i - k, r = i + k;
vector<int> d2(n);
```

```
// d[i] = number of palindromes taking s[i-1] and s[i] as }
     center
for (int i = 0, l = 0, r = -1; i < n; i++) {
 int k = (i > r) ? 0 : min(d2[1 + r - i + 1], r - i + 1) //find occurrences of t in s where '?'s are automatically
      [i + k]) k++:
 d2[i] = k--;
 if (i + k > r) l = i - k - 1, r = i + k;
```

#### 7.4 Palindromic Tree

```
struct state {
 int len, link;
 map<char, int> next;
state st[MAX];
int id, last;
string s;
11 ans[MAX];
void init() {
 for (int i = 0; i <= id; i++) {</pre>
   st[i].len = 0; st[i].link = 0;
   st[i].next.clear(); ans[i] = 0;
  st[1].len = -1; st[1].link = 1;
 st[2].len = 0; st[2].link = 1;
 id = 2; last = 2;
void extend(int pos) {
 while (s[pos - st[last].len - 1] != s[pos]) last = st[
      last].link:
 int newlink = st[last].link;
  char c = s[pos];
  while (s[pos - st[newlink].len - 1] != s[pos]) newlink
      = st[newlink].link;
  if (!st[last].next.count(c)) {
   st[last].next[c] = ++id;
   st[id].len = st[last].len + 2;
   st[id].link = (st[id].len == 1 ? 2 : st[newlink].next
        [c]);
   ans[id] += ans[st[id].link];
   if (st[id].len > 2) {
     int 1 = st[id].len / 2 + (st[id].len % 2 ? 1 : 0);
     if (h.range_hash(pos - st[id].len + 1, pos - st[id
         ].len + 1) == h.reverse_hash(pos - st[id].len +
          1, pos - st[id].len + 1)) ans[id]++;
   }
 last = st[last].next[c];
```

# 7.5 String Match FFT

```
matched with any character
while (0 <= i - k - 1 && i + k < n && s[i - k - 1] == s \left| \frac{1}{res[i + m - 1]} \right| = sum_j = 0 to m - 1_{s[i + j] * t[j] * (s]
                                                                [i + i] - t[i]
                                                           vector<int> string_matching(string &s, string &t) {
                                                             int n = s.size(), m = t.size();
                                                             vector<int> s1(n), s2(n), s3(n);
                                                             for(int i = 0; i < n; i++) s1[i] = s[i] == '?' ? 0 : s[
                                                                 i] - 'a' + 1; //assign any non zero number for non
                                                             for(int i = 0; i < n; i++) s2[i] = s1[i] * s1[i];</pre>
                                                             for(int i = 0; i < n; i++) s3[i] = s1[i] * s2[i];
                                                             vector<int> t1(m), t2(m), t3(m);
                                                             for(int i = 0: i < m: i++) t1[i] = t[i] == '?' ? 0 : t[
                                                                 i] - 'a' + 1;
                                                            for(int i = 0; i < m; i++) t2[i] = t1[i] * t1[i];</pre>
                                                             for(int i = 0; i < m; i++) t3[i] = t1[i] * t2[i];</pre>
                                                             reverse(t1.begin(), t1.end());
                                                             reverse(t2.begin(), t2.end());
                                                             reverse(t3.begin(), t3.end());
                                                             vector<int> s1t3 = multiply(s1, t3);
                                                             vector<int> s2t2 = multiply(s2, t2);
                                                             vector<int> s3t1 = multiply(s3, t1);
                                                             vector<int> res(n);
                                                             for(int i = 0; i < n; i++) res[i] = s1t3[i] - s2t2[i] *</pre>
                                                                  2 + s3t1[i];
                                                             vector<int> oc:
                                                             for(int i = m - 1; i < n; i++) if(res[i] == 0) oc.
                                                                 push_back(i - m + 1);
                                                             return oc;
```

# 7.6 Suffix Array

```
vector<VI> c:
VI sort_cyclic_shifts(const string &s) {
 int n = s.size();
 const int alphabet = 256;
 VI p(n), cnt(alphabet, 0);
  c.clear():
  c.emplace_back();
  c[0].resize(n);
 for (int i = 0; i < n; i++) cnt[s[i]]++;</pre>
 for (int i = 1; i < alphabet; i++) cnt[i] += cnt[i -</pre>
 for (int i = 0; i < n; i++) p[--cnt[s[i]]] = i;</pre>
```

```
c[0][p[0]] = 0;
 int classes = 1;
 for (int i = 1; i < n; i++) {</pre>
   if (s[p[i]] != s[p[i - 1]]) classes++;
   c[0][p[i]] = classes - 1;
 VI pn(n), cn(n);
 cnt.resize(n);
 for (int h = 0; (1 << h) < n; h++) {
   for (int i = 0; i < n; i++) {</pre>
     pn[i] = p[i] - (1 << h);
     if (pn[i] < 0) pn[i] += n;</pre>
   fill(cnt.begin(), cnt.end(), 0);
   /// radix sort
   for (int i = 0; i < n; i++) cnt[c[h][pn[i]]]++;</pre>
   for (int i = 1; i < classes; i++) cnt[i] += cnt[i -</pre>
   for (int i = n - 1; i >= 0; i--) p[--cnt[c[h][pn[i
       ]]]] = pn[i];
   cn[p[0]] = 0;
   classes = 1;
   for (int i = 1; i < n; i++) {</pre>
     PII cur = \{c[h][p[i]], c[h][(p[i] + (1 << h)) \% n\}
         1}:
     PII prev = \{c[h][p[i-1]], c[h][(p[i-1] + (1 <<
         h)) % n]}:
     if (cur != prev) ++classes;
     cn[p[i]] = classes - 1;
   c.push_back(cn);
 return p;
VI suffix_array_construction(string s) {
 s += "!";
 VI sorted_shifts = sort_cyclic_shifts(s);
 sorted_shifts.erase(sorted_shifts.begin());
 return sorted_shifts;
/// LCP between the ith and jth (i != j) suffix of the
    STRING
int suffixLCP(int i, int j) {
 assert(i != j);
 int log_n = c.size() - 1;
```

```
int ans = 0;
 for (int k = log_n; k \ge 0; k--) {
   if (c[k][i] == c[k][j]) {
     ans += 1 << k;
     i += 1 << k;
     j += 1 << k;
 }
 return ans;
VI lcp_construction(const string &s, const VI &sa) {
 int n = s.size();
 VI rank(n, 0);
 VI lcp(n - 1, 0);
 for (int i = 0; i < n; i++) rank[sa[i]] = i;</pre>
 for (int i = 0, k = 0; i < n; i++, k -= (k != 0)) {
   if (rank[i] == n - 1) {
     k = 0:
     continue;
   int j = sa[rank[i] + 1];
   while (i + k < n \&\& j + k < n \&\& s[i + k] == s[j + k]
       ]) k++;
   lcp[rank[i]] = k;
 return lcp;
```

#### 7.7 Suffix Automata

```
const int MAXN = 1e5 + 7, ALPHA = 26;
int len[2 * MAXN], link[2 * MAXN], nxt[2 * MAXN][ALPHA];
int sz;
int last;

void sa_init() {
   memset(nxt, -1, sizeof nxt);
   len[0] = 0, link[0] = -1, sz = 1, last = 0;
}

void add(char ch) {
   int c = ch - 'a';
   int cur = sz++;
   len[cur] = len[last] + 1;

int u = last;
   while (u != -1 && nxt[u][c] == -1) {
```

```
nxt[u][c] = cur:
   u = link[u];
  if (u == -1) link[cur] = 0;
  else {
   int v = nxt[u][c];
   if (len[v] == len[u] + 1) link[cur] = v;
     int clone = sz++;
     len[clone] = 1 + len[u], link[clone] = link[v];
     for (int i = 0; i < ALPHA; i++) nxt[clone][i] = nxt</pre>
          [v][i];
     while (u != -1 && nxt[u][c] == v) {
       nxt[u][c] = clone:
       u = link[u];
     link[v] = link[cur] = clone;
   }
 }
 last = cur;
vector<int> edge[2 * MAXN];
/// Optional, Call after adding all characters
void makeEdge() {
 for (int i = 0; i < sz; i++) {</pre>
   edge[i].clear();
   for (int j = 0; j < ALPHA; j++)
     if (nxt[i][j] != -1) edge[i].push_back(j);
 }
// Given a string S, you have to answer some queries:
// If all distinct substrings of string S were sorted
// lexicographically, which one will be the K-th smallest
long long dp[2 * MAXN];
bool vis[2 * MAXN];
void dfs(int u) {
 if (vis[u]) return;
 vis[u] = 1;
 dp[u] = 1;
  for (int i : edge[u]) {
   if (nxt[u][i] == -1) continue;
   dfs(nxt[u][i]);
   dp[u] += dp[nxt[u][i]];
```

```
}

void go(int u, long long rem, string &s) {
    if (rem == 1) return;
    long long sum = 1;
    for (int i : edge[u]) {
        if (nxt[u][i] == -1) continue;
        if (sum + dp[nxt[u][i]] < rem) {
            sum += dp[nxt[u][i]];
        } else {
            s += ('a' + i);
            go(nxt[u][i], rem - sum, s);
            return;
        }
    }
}
</pre>
```

#### 7.8 Z Algo

```
vector<int> calcz(string s) {
 int n = s.size();
 vector<int> z(n):
 int 1 = 0, r = 0:
 for (int i = 1; i < n; i++) {
  if (i > r) {
    1 = r = i:
     while (r < n \&\& s[r] == s[r - 1]) r++;
    z[i] = r - 1, r--;
  } else {
     int k = i - 1:
     if (z[k] < r - i + 1) z[i] = z[k];
     else {
      1 = i:
      while (r < n \&\& s[r] == s[r - 1]) r++;
      z[i] = r - 1, r--;
     }
  }
 return z;
```

# Equations and Formulas

### Catalan Numbers

$$C_n = \frac{1}{n+1} {2n \choose n}$$
  $C_0 = 1, C_1 = 1$  and  $C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}$ 

The number of ways to completely parenthesize n+1 factors. The number of triangulations of a convex polygon with n+2sides (i.e. the number of partitions of polygon into disjoint tegers  $1, 2, \ldots, n$  into k nonempty subsets such that all ele- $\sum [\gcd(i, n) = k] = \phi(\frac{n}{i})$ triangles by using the diagonals).

form n disjoint i.e. non-intersecting chords.

The number of rooted full binary trees with n+1 leaves (ver- $S^d(n,k) = S(n-d+1,k-d+1), n \ge k \ge d$ tices are not numbered). A rooted binary tree is full if every 8.4 Other Combinatorial Identities vertex has either two children or no children.

vertex has either two children or no children. Number of permutations of  $1, \ldots, n$  that avoid the pattern 123  $\binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}$ (or any of the other patterns of length 3); that is, the number of permutations with no three-term increasing sub-sequence. For n = 3, these permutations are 132, 213, 231, 312 and 321

# 8.2 Stirling Numbers First Kind

The Stirling numbers of the first kind count permutations according to their number of cycles (counting fixed points as cycles of length one).

S(n,k) counts the number of permutations of n elements with k disjoint cycles.

$$S(n,k) = (n-1) \cdot S(n-1,k) + S(n-1,k-1), where, S(0,0) =$$

$$1, S(n,0) = S(0,n) = 0 \sum_{k=0}^{n} S(n,k) = n!$$

The unsigned Stirling numbers may also be defined algebraically, as the coefficient of the rising factorial:

$$x^{\bar{n}} = x(x+1)...(x+n-1) = \sum_{k=0}^{n} S(n,k)x^{k}$$

Lets [n, k] be the stirling number of the first kind, then

$$[n - k] = \sum_{0 \le i_1 < i_2 < i_k < n} i_1 i_2 \dots i_k.$$

# 8.3 Stirling Numbers Second Kind

Stirling number of the second kind is the number of wavs to partition a set of n objects into k non-empty subsets.

$$S(n,k) = k \cdot S(n-1,k) + S(n-1,k-1)$$
, where  $S(0,0) = 1$ ,  $S(n,0) = S(0,n) = 0$   $S(n,2) = 2^{n-1} - 1$   $S(n,k) \cdot k! = \text{number of ways to color } n \text{ nodes using colors from } 1 \text{ to } k \text{ such that if } m \text{ is any integer, then } \gcd(a+m\cdot b,b) = \gcd(a,b)$  The gcd is a multiplicative function in the follow

An r-associated Stirling number of the second kind is the num- if  $a_1$  and  $a_2$  are relatively prime, then  $gcd(a_1 \cdot a_2, b) =$ ber of ways to partition a set of n objects into k subsets, with  $\gcd(a_1,b) \cdot \gcd(a_2,b)$ .

each subset containing at least r elements. It is denoted by  $\gcd(a, \operatorname{lcm}(b, c)) = \operatorname{lcm}(\gcd(a, b), \gcd(a, c))$ .  $S_r(n,k)$  and obeys the recurrence relation.  $S_r(n+1,k) = |\operatorname{lcm}(a,\operatorname{gcd}(b,c))| = \operatorname{gcd}(\operatorname{lcm}(a,b),\operatorname{lcm}(a,c)).$ 

$$kS_r(n,k) + \binom{n}{r-1}S_r(n-r+1,k-1)$$

Denote the n objects to partition by the integers  $1, 2, \dots, n$ .  $\gcd(a, b) = \sum_{a} \phi(k)$ Define the reduced Stirling numbers of the second kind, denoted  $S^d(n,k)$ , to be the number of ways to partition the inments in each subset have pairwise distance at least d. That  $\overline{i=1}$ The number of ways to connect the 2n points on a circle to is, for any integers i and j in a given subset, it is required that  $|i-j| \geq d$ . It has been shown that these numbers satisfy,

$$\binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}$$

$$\sum_{i=0}^{k} \binom{n+i}{i} = \sum_{i=0}^{k} \binom{n+i}{n} = \binom{n+k+1}{k}$$

$$n, r \in N, n > r, \sum_{i=r}^{n} \binom{i}{r} = \binom{n+1}{r+1}$$
If  $R(x) = \sum_{i=0}^{n} \binom{n}{i}$   $C(k)$ , there

If 
$$P(n) = \sum_{k=0}^{n} {n \choose k} \cdot Q(k)$$
, then,

$$Q(n) = \sum_{k=0}^{n} (-1)^{n-k} \binom{n}{k} \cdot P(k)$$

If 
$$P(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot Q(k)$$
, then,

$$Q(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot P(k)$$

### 8.5 Different Math Formulas

Picks Theorem: A = i + b/2 - 1

**Deragements:**  $d(i) = (i-1) \times (d(i-1) + d(i-2))$ 

$$\frac{n}{ab}$$
 -  $\left\{\frac{b\prime n}{a}\right\}$  -  $\left\{\frac{a\prime n}{b}\right\}$  +

The gcd is a multiplicative function in the following sense:

For non-negative integers a and b, where a and b are not both zero,  $gcd(n^a - 1, n^b - 1) = n^{gcd(a,b)} - 1$  $\sum_{k=1}^{n} \gcd(k, n) = \sum_{d \mid n} d \cdot \phi\left(\frac{n}{d}\right)$  $\sum_{k=1}^{n} x^{\gcd(k,n)} = \sum_{l=1}^{n} x^{d} \cdot \phi\left(\frac{n}{d}\right)$  $\sum_{k=1}^{n} \frac{1}{\gcd(k,n)} = \sum_{\text{all}} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{1}{n} \sum_{\text{all}} d \cdot \phi(d)$  $\sum_{k=1}^{n} \frac{k}{\gcd(k,n)} = \frac{n}{2} \cdot \sum_{d|n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{n}{2} \cdot \frac{1}{n} \cdot \sum_{d|n} d \cdot \phi(d)$  $\sum_{k=1}^{n} \frac{n}{\gcd(k,n)} = 2 * \sum_{k=1}^{n} \frac{k}{\gcd(k,n)} - 1, \text{ for } n > 1$  $\left| \sum_{i=1}^{n} \sum_{j=1}^{n} [\gcd(i,j) = 1] = \sum_{d=1}^{n} \mu(d) \left\lfloor \frac{n}{d} \right\rfloor^{2} \right|$ 

 $F(n) = \sum_{i=1}^{n} \sum_{i=1}^{n} \operatorname{lcm}(i,j) = \sum_{l=1}^{n} \left( \frac{\left(1 + \lfloor \frac{n}{l} \rfloor\right) \left(\lfloor \frac{n}{l} \rfloor\right)}{2} \right)^{2} \sum_{l \mid l} \mu(d) l d$ 

 $\left|\sum_{i=1}^{n}\sum_{j=1}^{n}\gcd(i,j)=\sum_{d=1}^{n}\phi(d)\left\lfloor\frac{n}{d}\right\rfloor^{2}\right|$ 

 $\left| \sum_{i=1}^{n} \sum_{j=1}^{n} i \cdot j [\gcd(i, j) = 1] \right| = \sum_{i=1}^{n} \phi(i) i^{2}$