${\bf IUT\ Shongshoptok, Islamic\ University\ of\ Technology}$

Contents		6	Stri 6.1	ing Aho Corasick	1. 1.
1	All headers	1	6.2 6.3	Prefix Function	10
2	Data Structure	1	6.4	double hash	1
	2.1 BIT-2D		Ear	nations and Formulas	18
	2.2 BIT		7.1	Catalan Numbers	1
	2.3 DSU With Rollbacks + Queue Undo 7	Trick 1	7.2	Stirling Numbers First Kind	18
	2.4 MO with Update	2	7.3	Stirling Numbers Second Kind	18
	2.5 SegmentTree	2	$7.4 \\ 7.5$	Other Combinatorial Identities	18 18
3	Geometry	2	7.6	GCD and LCM	1
	3.1 Circular	2			
	3.2 Convex	3			
	3.3 Enclosing Circle	5			
	3.4 Half Planar				
	3.5 Intersecting Segments	6			
	3.6 Linear	6			
	3.7 Point Rotation Trick	7			
	3.8 Point				
	3.9 Polygon	8			
4	Graphs	9			
	4.1 Bridge Tree				
	4.2 Centroid Decomposition				
	4.3 Dinic				
	4.4 HLD				
	4.5 LCA In $O(1)$				
	4.6 SCC				
	4.7 StoerWanger				
	4.8 Tree Algo	11			
5	Math	12			
	5.1 Adaptive Simpsons				
	5.2 Berlekamp Massey	12			
	5.3 Combi				
	5.4 FFT				
	5.5 Fractional Binary Search				
	5.6 Gaussian Elimination				
	5.7 Linear Sieve				
	5.8 Pollard Rho				
	5.9 Subset-Convolution	15			

VIM Init

Sublime Build

"working_dir" : "\$file_path"

```
filetype on
filetype plugin on
filetype plugin indent on
syntax on
" Some useful settings
set smartindent
set expandtab
                    "tab to spaces
                    "the width of a tab
set tabstop=4
set shiftwidth=4
                    "the width for indent
" set foldenable
set foldmethod=indent "folding by indent
                    "ignore the case when search texts
set ignorecase
set incsearch
set relativenumber "line number
                   "hilight the line of the cursor
set cursorline
set nowrap
                   "no line wrapping
set number
:autocmd BufNewFile *.cpp Or ~/.temp.temp
map<F5> :w <CR> : ! g++ -std=c++17 -DSPTK % && ./a.out <CR>
```

1 All headers

```
//#pragma GCC optimize("Ofast")
//#pragma GCC optimization ("03")
//#pragma comment(linker, "/stack:200000000")
//#pragma GCC optimize("unroll-loops")
//#pragma GCC target("sse,sse2,sse3,sse3,sse4,popcnt,abm,mmx,
    avx.tune=native")
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
   //find_by_order(k) --> returns iterator to the kth largest
        element counting from 0
   //order of kev(val) --> returns the number of items in a
        set that are strictly smaller than our item
template <typename DT>
using ordered_set = tree <DT, null_type, less<DT>, rb_tree_tag,
    tree_order_statistics_node_update>;
// debug template
void __print(int x) {cerr << x;}</pre>
void print(long long x) {cerr << x:}</pre>
void __print(unsigned long long x) {cerr << x;}</pre>
void __print(double x) {cerr << x;}</pre>
```

```
void __print(long double x) {cerr << x;}</pre>
void __print(char x) {cerr << '\'' << x << '\'';}</pre>
void __print(const string &x) {cerr << '\"' << x << '\"';}</pre>
void print(bool x) {cerr << (x ? "true" : "false"):}</pre>
template<typename T, typename V>
void __print(const pair<T, V> &x) {cerr << ','; __print(x.first</pre>
    ); cerr << ", "; __print(x.second); cerr << '}';}
template<tvpename T>
void __print(const T &x) {int f = 0; cerr << "{"; for (auto &i:</pre>
     x) cerr << (f++ ? ", " : ""), __print(i); cerr << "}";}</pre>
void _print() {cerr << "]\n";}</pre>
template <typename T, typename... V>
void _print(T t, V... v) {__print(t); if (sizeof...(v)) cerr <</pre>
      ". ": print(v...):}
#ifdef SPTK
#define debug(x...) cerr << "[" << #x << "] = ["; _print(x)</pre>
#define debug(x...)
#endif
#define fastio
                        ios base::svnc with stdio(0):cin.tie(0):
#define Make(x.p)
                       (x | (1 << p))
#define DeMake(x,p)
                      (x & ~(1<<p))
#define Check(x,p)
                       (x & (1<<p))
template<size_t N>
bitset<N> rotl( std::bitset<N> const& bits, unsigned count ) {
   count %= N; // Limit count to range [0,N)
   return bits << count | bits >> (N - count):
```

2 Data Structure

2.1 BIT-2D

```
#include "bits/stdc++.h"
using namespace std;
const int N = 1008;
int bit[N][N], n, m;
int a[N][N], q;
void update(int x, int y, int val) {
   for (; x < N; x += -x & x)
       for (int j = y; j < N; j += -j & j) bit[x][j] += val;</pre>
int get(int x, int y) {
   int ans = 0:
   for (; x; x -= x & -x)
       for (int j = y; j; j -= j \& -j) ans += bit[x][j];
   return ans:
int get(int x1, int v1, int x2, int v2) {
   return get(x2, y2) - get(x1 - 1, y2) - get(x2, y1 - 1) +
        get(x1 - 1, y1 - 1);
```

```
2.2 BIT
```

```
#include "bits/stdc++.h"
using namespace std;
struct BIT {
   int n;
   vector<int> bit:
   BIT(int n) {
       this \rightarrow n = n:
       bit.resize(n);
   void update(int x, int delta) {
       for (; x <= n; x += x & -x) bit[x] += delta;</pre>
   int query(int x) {
       int sum = 0:
       for (; x > 0; x -= x & -x) sum += bit[x];
   }
};
int main() {}
```

2.3 DSU With Rollbacks + Queue Undo Trick

```
struct Rollback_DSU {
  int n:
  vector <int> par, sz;
   vector <pair <int, int>> op;
  Rollback_DSU(int n) : par(n), sz(n, 1) {
      iota(par.begin(), par.end(), 0);
      op.reserve(n);
  }
  int Anc(int node) {
      for(; node != par[node]; node = par[node]);
       return node;
  void Unite(int x, int y) {
      if(sz[x = Anc(x)] < sz[y = Anc(y)])
          swap(x, y);
      op.emplace_back(x, y);
      par[y] = x;
       sz[x] += sz[v]:
  void Undo(size_t t) {
      for(; op.size() > t; op.pop_back()) {
          par[op.back().second] = op.back().second;
          sz[op.back().first] -= sz[op.back().second];
```

2.4 MO with Update

```
const int N = 1e5 + 5, sz = 2700, bs = 25;
int arr[N], freq[2 * N], cnt[2*N], id[N], ans[N];
struct querv{
    int 1, r, t, L, R;
    query(int l = 1, int r = 0, int t = 1, int id = -1): l(1),
         r(r), t(t), L(1 / sz), R(r / sz) {}
    bool operator < (const query &rhs) const {</pre>
       return (L < rhs.L) or (L == rhs.L and R < rhs.R) or (L
            == rhs.L and R == rhs.R and t < rhs.t);
   }
} Q[N];
struct update{
    int idx. val. last:
int qi = 0, ui = 0;
int 1 = 1, r = 0, t = 0;
void add(int idx) {
    --cnt[freg[arr[idx]]]:
    freq[arr[idx]]++;
    cnt[freq[arr[idx]]]++;
void remove(int idx){
    --cnt[freq[arr[idx]]];
    freq[arr[idx]]--;
    cnt[freq[arr[idx]]]++;
void apply(int t) {
    const bool f = 1 <= Up[t].idx and Up[t].idx <= r;</pre>
    if(f) remove(Up[t].idx);
    arr[Up[t].idx] = Up[t].val;
    if(f) add(Up[t].idx);
void undo(int t) {
    const bool f = 1 <= Up[t].idx and Up[t].idx <= r;</pre>
    if(f) remove(Up[t].idx);
    arr[Up[t].idx] = Up[t].last;
    if(f) add(Up[t].idx);
int mex(){
    for(int i = 1; i <= N; i++)</pre>
       if(!cnt[i])
           return i:
    assert(0);
int main() {
    int n, q;
    cin >> n >> q;
    int counter = 0;
    map <int, int> M;
    for(int i = 1; i <= n; i++){</pre>
       cin >> arr[i];
       if(!M[arr[i]])
           M[arr[i]] = ++counter;
       arr[i] = M[arr[i]];
    }
```

```
iota(id. id + N. 0):
while(q--){
    int tp, x, y;
    cin >> tp >> x >> y;
    if(tp == 1) Q[++qi] = query(x, y, ui);
    else {
       if(!M[y]) M[y] = ++counter;
       v = M[v];
       Up[++ui] = \{x, y, arr[x]\};
       arr[x] = y;
   }
t = ui;
cnt[0] = 3 * n:
sort(id + 1, id + qi + 1, [\&](int x, int v) {return Q[x] < }
    Q[v];});
for(int i = 1; i <= qi; i++) {</pre>
   int x = id[i]:
    while(Q[x].t > t) apply(++t);
    while(Q[x].t < t) undo(t--);</pre>
    while (Q[x].1 < 1) add(--1);
   while (0[x].r > r) add (++r):
    while(Q[x].1 > 1) remove(1++);
    while(Q[x].r < r) remove(r--);</pre>
    ans[x] = mex():
for(int i = 1; i <= qi; i++)</pre>
    cout << ans[i] << '\n';</pre>
```

2.5 SegmentTree

```
#include <bits/stdc++.h>
using namespace std;
typedef long long LL;
const int N = 1 << 18:
LL tree[N * 4], lazy[N * 4], a[N];
LL combine(LL u, LL v){
   return u + v;
void build(int u, int st, int en) {
   lazy[u] = 0;
   if (st == en) {
       tree[u] = a[st]; return;
   int mid = (st + en) / 2;
   build(u * 2, st, mid);
   build(u * 2 + 1, mid + 1, en);
   tree[u] = combine(tree[u * 2], tree[u * 2 + 1]);
//for max, min query just add lazy * 1;
void propagate(int u,int st, int en) {
```

```
int mid = (st+en)/2:
   lazv[u * 2] += lazv[u];
   tree[u * 2] += lazy[u]*(mid-st+1);
   lazv[u * 2 + 1] += lazv[u]:
   tree[u * 2 + 1] += lazv[u]*(en-mid);
   lazv[u] = 0:
void update(int u, int st, int en, int i, int i, int val) {
   if (en < i || st> j) return;
   if (st >= i && en <= j) {</pre>
       // for max, min query just add lazy * 1;
       lazy[u] += val; tree[u] += val*(en-st+1); return;
   propagate(u.st.en):
   int mid = (st + en) / 2;
   update(u * 2, st, mid, i, j, val);
   update(u * 2 + 1, mid + 1, en, i, j, val);
   tree[u] = combine(tree[u * 2], tree[u * 2 + 1]);
LL query(int u, int st, int en, int i, int j) {
   if (st > j || en < i) return 0; // return appropriate
        identity value
   if (st >= i && en <= j) return tree[u];</pre>
   propagate(u,st,en);
   int mid = (st + en) / 2;
   return combine(query(u * 2, st, mid, i, j), query(u * 2 +
        1, mid + 1, en, i, j));
```

3 Geometry

3.1 Circular

```
// Extremely inaccurate for finding near touches
// compute intersection of line l with circle c
// The intersections are given in order of the ray (l.a. l.b)
vector<Point> circleLineIntersection(Circle c, Line 1) {
   static_assert(is_same<Tf, Ti>::value);
   vector<Point> ret:
   Point b = 1.b - 1.a, a = 1.a - c.o;
   Tf A = dot(b, b), B = dot(a, b);
   Tf C = dot(a, a) - c.r * c.r, D = B*B - A*C;
   if (D < -EPS) return ret:</pre>
   ret.push_back(1.a + b * (-B - sqrt(D + EPS)) / A);
   if (D > EPS)
       ret.push_back(1.a + b * (-B + sqrt(D)) / A);
   return ret;
// signed area of intersection of circle(c.o, c.r) &&
// triangle(c.o, s.a, s.b) [cross(a-o, b-o)/2]
Tf circleTriangleIntersectionArea(Circle c, Segment s) {
   using Linear::distancePointSegment;
```

```
Tf OA = length(c.o - s.a):
        Tf OB = length(c.o - s.b);
        // sector
        if(dcmp(distancePointSegment(c.o, s) - c.r) >= 0)
                return angleBetween(s.a-c.o, s.b-c.o) * (c.r * c.r) /
                           2.0:
        // triangle
        if(dcmp(OA - c.r) \le 0 \&\& dcmp(OB - c.r) \le 0)
                return cross(c.o - s.b, s.a - s.b) / 2.0;
        // three part: (A, a) (a, b) (b, B)
        vector<Point> Sect = circleLineIntersection(c, s);
        return circleTriangleIntersectionArea(c, Segment(s,a, Sect
                  [0])
                + circleTriangleIntersectionArea(c, Segment(Sect[0],
                           Sect[1]))
                + circleTriangleIntersectionArea(c, Segment(Sect[1], s.b
}
// area of intersecion of circle(c.o, c.r) && simple polyson(p
 // Tested : https://codeforces.com/gym/100204/problem/F -
           Little Mammoth
Tf circlePolyIntersectionArea(Circle c, Polygon p) {
        Tf res = 0:
        int n = p.size();
        for(int i = 0; i < n; ++i)</pre>
                res += circleTriangleIntersectionArea(c, Segment(p[i], p vector<Point> pointCircleTangents(Point p, Circle c) {
                           [(i + 1) \% n]);
        return abs(res):
// locates circle c2 relative to c1
// interior
                                                (d < R - r)
// interior tangents (d = R - r)
                                                                                  ----> -1
// concentric (d = 0)
 // secants
                                              (R - r < d < R + r) \longrightarrow 0
// exterior tangents (d = R + r)
                                                (d > R + r)
// exterior
int circleCirclePosition(Circle c1, Circle c2) {
        Tf d = length(c1.o - c2.o);
        int in = dcmp(d - abs(c1.r - c2.r)), ex = dcmp(d - (c1.r + abs(c1.r + abs
        return in < 0 ? -2 : in == 0 ? -1 : ex == 0 ? 1 : ex > 0 ?
                  2:0:
// compute the intersection points between two circles c1 && c2
 vector<Point> circleCircleIntersection(Circle c1, Circle c2) {
        static_assert(is_same<Tf, Ti>::value);
        vector<Point> ret:
        Tf d = length(c1.o - c2.o);
        if(dcmp(d) == 0) return ret;
        if(dcmp(c1.r + c2.r - d) < 0) return ret;</pre>
```

```
if(dcmp(abs(c1.r - c2.r) - d) > 0) return ret;
   Point v = c2.0 - c1.0;
   Tf co = (c1.r * c1.r + sqLength(v) - c2.r * c2.r) / (2 * c1)
        .r * length(v));
   Tf si = sqrt(abs(1.0 - co * co));
   Point p1 = scale(rotatePrecise(v, co, -si), c1.r) + c1.o;
   Point p2 = scale(rotatePrecise(v, co, si), c1.r) + c1.o;
   ret.push_back(p1);
   if(p1 != p2) ret.push_back(p2);
   return ret;
// intersection area between two circles c1. c2
Tf circleCircleIntersectionArea(Circle c1, Circle c2) {
   Point AB = c2.0 - c1.0;
   Tf d = length(AB);
   if(d \ge c1.r + c2.r) return 0;
   if(d + c1.r <= c2.r) return PI * c1.r * c1.r;</pre>
   if(d + c2.r <= c1.r) return PI * c2.r * c2.r;</pre>
   Tf alpha1 = acos((c1.r * c1.r + d * d - c2.r * c2.r) / (2.0)
         * c1.r * d));
   Tf alpha2 = acos((c2.r * c2.r + d * d - c1.r * c1.r) / (2.0)
         * c2.r * d));
   return c1.sector(2 * alpha1) + c2.sector(2 * alpha2);
// returns tangents from a point p to circle c
   static_assert(is_same<Tf, Ti>::value);
   vector<Point> ret;
   Point u = c.o - p;
   Tf d = length(u);
   if(d < c.r);
   else if(dcmp(d - c.r) == 0) {
       ret = { rotate(u, PI / 2) }:
   else {
       Tf ang = asin(c.r / d);
       ret = { rotate(u, -ang), rotate(u, ang) };
   return ret;
// returns the points on tangents that touches the circle
vector<Point> pointCircleTangencyPoints(Point p, Circle c) {
   static_assert(is_same<Tf, Ti>::value);
   Point u = p - c.o;
   Tf d = length(u);
   if(d < c.r) return {};</pre>
   else if(dcmp(d - c.r) == 0) return {c.o + u};
   else {
       Tf ang = acos(c.r / d);
```

u = u / length(u) * c.r;

```
return { c.o + rotate(u, -ang), c.o + rotate(u, ang) };
// for two circles c1 && c2, returns two list of points a && b
// such that a[i] is on c1 && b[i] is c2 && for every i
// Line(a[i], b[i]) is a tangent to both circles
// CAUTION: a[i] = b[i] in case they touch | -1 for c1 = c2
int circleCircleTangencyPoints(Circle c1, Circle c2, vector<</pre>
    Point> &a, vector<Point> &b) {
   a.clear(), b.clear();
   int cnt = 0:
   if(dcmp(c1.r - c2.r) < 0) {
       swap(c1, c2); swap(a, b);
   Tf d2 = sqLength(c1.o - c2.o);
   Tf rdif = c1.r - c2.r, rsum = c1.r + c2.r;
   if(dcmp(d2 - rdif * rdif) < 0) return 0;</pre>
   if(dcmp(d2) == 0 && dcmp(c1.r - c2.r) == 0) return -1;
   Tf base = angle(c2.o - c1.o);
   if(dcmp(d2 - rdif * rdif) == 0) {
       a.push_back(c1.point(base));
       b.push_back(c2.point(base));
       cnt++:
       return cnt;
   Tf ang = acos((c1.r - c2.r) / sqrt(d2));
   a.push_back(c1.point(base + ang));
   b.push_back(c2.point(base + ang));
   a.push_back(c1.point(base - ang));
   b.push_back(c2.point(base - ang));
   if(dcmp(d2 - rsum * rsum) == 0) {
       a.push_back(c1.point(base));
       b.push_back(c2.point(PI + base));
   else if(dcmp(d2 - rsum * rsum) > 0) {
       Tf ang = acos((c1.r + c2.r) / sqrt(d2));
       a.push_back(c1.point(base + ang));
       b.push_back(c2.point(PI + base + ang));
       cnt++:
       a.push_back(c1.point(base - ang));
       b.push_back(c2.point(PI + base - ang));
       cnt++:
   return cnt;
```

3.2 Convex

```
///minkowski sum of two polygons in O(n)
Polygon minkowskiSum(Polygon A, Polygon B){
   int n = A.size(), m = B.size();
```

```
rotate(A.begin(), min_element(A.begin(), A.end()), A.end())
   rotate(B.begin(), min_element(B.begin(), B.end()), B.end())
   A.push_back(A[0]); B.push_back(B[0]);
   for(int i = 0; i < n; i++) A[i] = A[i+1] - A[i];</pre>
   for(int i = 0; i < m; i++) B[i] = B[i+1] - B[i];</pre>
   Polygon C(n+m+1);
   C[0] = A.back() + B.back();
   merge(A.begin(), A.end()-1, B.begin(), B.end()-1, C.begin()
        +1, polarComp(Point(0, 0), Point(0, -1)));
   for(int i = 1; i < C.size(); i++) C[i] = C[i] + C[i-1];</pre>
   C.pop back():
   return C;
/// finds the rectangle with minimum area enclosing a convex
    polygon and
/// the rectangle with minimum perimeter enclosing a convex
/// Tested on https://open.kattis.com/problems/fenceortho
pair< Tf, Tf >rotatingCalipersBoundingBox(const Polygon &p) {
   using Linear::distancePointLine;
   static_assert(is_same<Tf, Ti>::value);
   int n = p.size();
   int 1 = 1, r = 1, j = 1;
   Tf area = 1e100;
   Tf perimeter = 1e100;
   for(int i = 0: i < n: i++) {</pre>
       Point v = (p[(i+1)\%n] - p[i]) / length(p[(i+1)\%n] - p[i])
           1):
       while (dcmp(dot(v, p[r/n] - p[i]) - dot(v, p[(r+1)/n] - p
            [i])) < 0) r++;
       while(j < r \mid | dcmp(cross(v, p[j/n] - p[i]) - cross(v, p
            [(j+1)\%n] - p[i])) < 0) j++;
       while(l < j \mid | dcmp(dot(v, p[1%n] - p[i]) - dot(v, p[(1
            +1)%n] - p[i])) > 0) 1++;
       Tf w = dot(v, p[r\n] - p[i]) - dot(v, p[1\n] - p[i]);
       Tf h = distancePointLine(p[j/n], Line(p[i], p[(i+1)/n]))
       area = min(area, w * h);
       perimeter = min(perimeter, 2 * w + 2 * h);
   return make_pair(area, perimeter);
// returns the left side of polygon u after cutting it by ray a int extremePoint(const Polygon &poly, Point u) {
Polygon cutPolygon(Polygon u, Point a, Point b) {
   using Linear::lineLineIntersection;
   using Linear::onSegment;
   Polygon ret;
   int n = u.size();
   for(int i = 0: i < n: i++) {</pre>
       Point c = u[i], d = u[(i + 1) \% n];
```

```
if(dcmp(cross(b-a, c-a)) >= 0) ret.push_back(c);
       if(dcmp(cross(b-a, d-c)) != 0) {
           Point t;
           lineLineIntersection(a, b - a, c, d - c, t):
           if(onSegment(t, Segment(c, d))) ret.push_back(t);
    return ret;
// returns true if point p is in or on triangle abc
bool pointInTriangle(Point a, Point b, Point c, Point p) {
   return dcmp(cross(b - a, p - a)) >= 0
       && dcmp(cross(c - b, p - b)) >= 0
       && dcmp(cross(a - c, p - c)) >= 0:
// Tested : https://www.spoj.com/problems/INOROUT
// pt must be in ccw order with no three collinear points
// returns inside = -1, on = 0, outside = 1
int pointInConvexPolygon(const Polygon &pt, Point p) {
   int n = pt.size():
   assert(n >= 3):
   int lo = 1, hi = n - 1:
   while(hi - lo > 1) {
       int mid = (lo + hi) / 2;
       if(dcmp(cross(pt[mid] - pt[0], p - pt[0])) > 0) lo = mid |// of p that touch or intersect line 1.
       else hi = mid:
   bool in = pointInTriangle(pt[0], pt[lo], pt[hi], p);
   if(!in) return 1;
    if(dcmp(cross(pt[lo] - pt[lo - 1], p - pt[lo - 1])) == 0)
   if(dcmp(cross(pt[hi] - pt[lo], p - pt[lo])) == 0) return 0;
   if(dcmp(cross(pt[hi] - pt[(hi + 1) % n], p - pt[(hi + 1) %
        n])) == 0) return 0:
   return -1;
// Extreme Point for a direction is the farthest point in that
// also https://codeforces.com/blog/entry/48868
// u is the direction for extremeness
// weakly tested on https://open.kattis.com/problems/fenceortho
   int n = (int) poly.size();
   int a = 0, b = n;
   while(b - a > 1) {
       int c = (a + b) / 2:
       if(dcmp(dot(poly[c] - poly[(c + 1) % n], u)) >= 0 &&
            dcmp(dot(poly[c] - poly[(c - 1 + n) \% n], u)) >= 0)
           return c;
```

```
bool a_up = dcmp(dot(poly[(a + 1) \% n] - poly[a], u)) >=
       bool c_up = dcmp(dot(poly[(c + 1) \% n] - poly[c], u)) >=
       bool a_above_c = dcmp(dot(poly[a] - poly[c], u)) > 0;
       if(a_up \&\& !c_up) b = c;
       else if(!a_up && c_up) a = c;
       else if(a_up && c_up) {
          if(a_above_c) b = c;
          else a = c;
      }
       else {
          if(!a above c) b = c:
          else a = c;
   }
   if(dcmp(dot(poly[a] - poly[(a + 1) % n], u)) > 0 && dcmp(
       dot(poly[a] - poly[(a - 1 + n) \% n], u)) > 0)
       return a:
   return b % n:
// For a convex polygon p and a line 1, returns a list of
// the i'th segment is considered (p[i], p[(i + 1) modulo |p|])
// #1 If a segment is collinear with the line, only that is
// #2 Else if 1 goes through i'th point, the i'th segment is
// Complexity: O(lg |p|)
vector<int> lineConvexPolyIntersection(const Polygon &p, Line 1
   assert((int) p.size() >= 3);
   assert(1.a != 1.b):
   int n = p.size();
   vector<int> ret;
   Point v = 1.b - 1.a;
   int lf = extremePoint(p, rotate90(v));
   int rt = extremePoint(p, rotate90(v) * Ti(-1));
   int olf = orient(l.a, l.b, p[lf]);
   int ort = orient(l.a, l.b, p[rt]);
   if(!olf || !ort) {
       int idx = (!olf ? lf : rt);
       if(orient(l.a, l.b, p[(idx - 1 + n) \% n]) == 0)
          ret.push_back((idx - 1 + n) \% n);
       else ret.push_back(idx);
       return ret;
   if(olf == ort) return ret;
   for(int i=0; i<2; ++i) {</pre>
```

```
int lo = i ? rt : lf:
       int hi = i ? lf : rt;
       int olo = i ? ort : olf;
       while(true) {
          int gap = (hi - lo + n) \% n;
          if(gap < 2) break;</pre>
          int mid = (lo + gap / 2) % n;
          int omid = orient(l.a, l.b, p[mid]);
          if(!omid) {
              lo = mid:
              break;
          if(omid == olo) lo = mid:
          else hi = mid;
       ret.push_back(lo);
   }
   return ret;
// Tested : https://toph.co/p/cover-the-points
// Calculate [ACW, CW] tangent pair from an external point
constexpr int CW = -1, ACW = 1;
bool isGood(Point u, Point v, Point Q, int dir) { return orient
    (Q, u, v) != -dir; }
Point better(Point u, Point v, Point Q, int dir) { return
    orient(Q, u, v) == dir ? u : v; }
Point pointPolyTangent(const Polygon &pt, Point Q, int dir, int
     lo, int hi) {
   while(hi - lo > 1) {
       int mid = (lo + hi) / 2:
       bool pvs = isGood(pt[mid], pt[mid - 1], Q, dir);
       bool nxt = isGood(pt[mid], pt[mid + 1], Q, dir);
       if(pvs && nxt) return pt[mid];
       if(!(pvs || nxt)) {
          Point p1 = pointPolyTangent(pt, Q, dir, mid + 1, hi)
          Point p2 = pointPolyTangent(pt, Q, dir, lo, mid - 1)
          return better(p1, p2, Q, dir);
       if(!pvs) {
                                                         hi =
          if(orient(Q, pt[mid], pt[lo]) == dir)
               mid - 1;
          else if(better(pt[lo], pt[hi], Q, dir) == pt[lo]) hi
                = mid - 1;
          else lo = mid + 1;
      }
       if(!nxt) {
          if(orient(Q, pt[mid], pt[lo]) == dir)
               mid + 1:
          else if(better(pt[lo], pt[hi], Q, dir) == pt[lo]) hi
                = mid - 1:
          else lo = mid + 1;
```

```
}
   Point ret = pt[lo]:
   for(int i = lo + 1; i <= hi; i++) ret = better(ret, pt[i],</pre>
        Q. dir):
   return ret:
// [ACW, CW] Tangent
    ) {
   int n = pt.size();
   Point acw_tan = pointPolyTangent(pt, Q, ACW, 0, n - 1);
   Point cw_tan = pointPolyTangent(pt, Q, CW, 0, n - 1);
   return make_pair(acw_tan, cw_tan);
```

// returns false if points are collinear, true otherwise

Enclosing Circle

```
// circle p touch each arm of triangle abo
bool inscribedCircle(Point a, Point b, Point c, Circle &p) {
   using Linear::distancePointLine;
   static_assert(is_same<Tf, Ti>::value);
   if(orient(a, b, c) == 0) return false;
   Tf u = length(b - c):
   Tf v = length(c - a);
   Tf w = length(a - b);
   p.o = (a * u + b * v + c * w) / (u + v + w);
   p.r = distancePointLine(p.o, Line(a, b));
   return true:
// set of points A(x, y) such that PA : QA = rp : rq
Circle apolloniusCircle(Point P, Point Q, Tf rp, Tf rq) {
   static_assert(is_same<Tf, Ti>::value);
   rq *= rq; rp *= rp;
   Tf a = rq - rp;
   assert(dcmp(a)):
   Tf g = (rq * P.x - rp * Q.x)/a;
   Tf h = (rq * P.y - rp * Q.y)/a;
   Tf c = (rq * P.x * P.x - rp * Q.x * Q.x + rq * P.y * P.y -
        rp * Q.v * Q.v)/a;
   Point o(g, h);
   Tf R = sqrt(g * g + h * h - c);
   return Circle(o, R);
// returns false if points are collinear, true otherwise
// circle p goes through points a, b && c
bool circumscribedCircle(Point a, Point b, Point c, Circle &p)
   using Linear::lineLineIntersection;
   if(orient(a, b, c) == 0) return false;
   Point d = (a + b) / 2, e = (a + c) / 2;
   Point vd = rotate90(b - a), ve = rotate90(a - c):
   bool f = lineLineIntersection(d, vd, e, ve, p.o);
   if(f) p.r = length(a - p.o);
```

```
return f:
                                                                 // Following three methods implement Welzl's algorithm for
                                                                // the smallest enclosing circle problem: Given a set of
                                                                 // points, find out the minimal circle that covers them all.
                                                                 // boundary(p) determines (if possible) a circle that goes
                                                                 // through the points in p. Ideally |p| \le 3.
                                                                 // welzl() is a recursive helper function doing the most part
pair<Point, Point> pointPolyTangents(const Polygon &pt, Point Q // of Welzl's algorithm. Call minidisk with the set of points
                                                                 // Randomized Complexity: O(CN) with C~10 (practically lesser)
                                                                Circle boundary(const vector<Point> &p) {
                                                                    Circle ret:
                                                                    int sz = p.size():
                                                                    if(sz == 0)
                                                                                      ret.r = 0;
                                                                    else if(sz == 1) ret.o = p[0], ret.r = 0;
                                                                    else if(sz == 2) ret.o = (p[0] + p[1]) / 2, ret.r = length
                                                                         (p[0] - p[1]) / 2;
                                                                    else if(!circumscribedCircle(p[0], p[1], p[2], ret)) ret.r
                                                                         = 0:
                                                                    return ret:
                                                                Circle welzl(const vector<Point> &p, int fr, vector<Point> &b)
                                                                    if(fr >= (int) p.size() || b.size() == 3) return boundary(b
                                                                         );
                                                                    Circle c = welzl(p, fr + 1, b);
                                                                    if(!c.contains(p[fr])) {
                                                                        b.push_back(p[fr]);
                                                                        c = welzl(p, fr + 1, b);
                                                                       b.pop_back();
                                                                    }
                                                                    return c;
                                                                Circle minidisk(vector<Point> p) {
                                                                    random_shuffle(p.begin(), p.end());
                                                                    vector<Point> a:
                                                                    return welzl(p, 0, q);
```

3.4 Half Planar

```
using Linear::lineLineIntersection;
struct DirLine {
   Point p, v;
   Tf ang:
   DirLine() {}
   /// Directed line containing point P in the direction v
   DirLine(Point p, Point v) : p(p), v(v) { ang = atan2(v.y, v
   /// Directed Line for ax+by+c >=0
   DirLine(Tf a, Tf b, Tf c) {
       assert(dcmp(a) || dcmp(b));
       p = dcmp(a) ? Point(-c/a, 0) : Point(0,-c/b);
      v = Point(b, -a);
       ang = atan2(v.y, v.x);
```

```
bool operator<(const DirLine& u) const { return ang < u.ang</pre>
    bool onLeft(Point x) const { return dcmp(cross(v, x-p)) >=
};
// Returns the region bounded by the left side of some directed
      lines
// MAY CONTAIN DUPLICATE POINTS
// OUTPUT IS UNDEFINED if intersection is unbounded
// Complexity: O(n \log n) for sorting, O(n) afterwards
Polygon halfPlaneIntersection(vector<DirLine> li) {
    int n = li.size(), first = 0, last = 0;
    sort(li.begin(), li.end());
    vector<Point> p(n);
    vector<DirLine> q(n);
    q[0] = li[0];
    for(int i = 1: i < n: i++) {</pre>
       while(first < last && !li[i].onLeft(p[last - 1])) last</pre>
       while(first < last && !li[i].onLeft(p[first])) first++;</pre>
        q[++last] = li[i];
       if(dcmp(cross(q[last].v, q[last-1].v)) == 0) {
           if(q[last].onLeft(li[i].p)) q[last] = li[i];
       if(first < last)</pre>
           lineLineIntersection(q[last - 1].p, q[last - 1].v, q
                [last].p. q[last].v. p[last - 1]);
   }
    while(first < last && !q[first].onLeft(p[last - 1])) last</pre>
    if(last - first <= 1) return {}:</pre>
    lineLineIntersection(q[last].p, q[last].v, q[first].p, q[
        first].v, p[last]);
    return Polygon(p.begin()+first, p.begin()+last+1);
// O(n^2 lg n) implementation of Voronoi Diagram bounded by INF
// returns region, where regions[i] = set of points for which
// point is site[i]. This region is a polygon.
const Tf INF = 1e10:
vector<Polygon> voronoi(const vector<Point> &site, Tf bsq) {
    int n = site.size();
    vector<Polygon> region(n);
    Point A(-bsq, -bsq), B(bsq, -bsq), C(bsq, bsq), D(-bsq, bsq
        );
    for(int i = 0: i < n: ++i) {</pre>
       vector<DirLine> li(n - 1):
       for(int j = 0, k = 0; j < n; ++j) {
           if(i == i) continue;
```

3.5 Intersecting Segments

```
// Given a list of segments v, finds a pair (i, j)
// st v[i], v[i] intersects. If none, returns {-1, -1}
// Tested Timus 1469, CF 1359F
struct Event {
   Tf x:
   int tp, id;
   bool operator < (const Event &p) const {</pre>
       if(dcmp(x - p.x)) return x < p.x;</pre>
       return tp > p.tp;
   }
pair<int, int> anyIntersection(const vector<Segment> &v) {
   using Linear::segmentsIntersect;
   static assert(is same<Tf. Ti>::value):
   vector<Event> ev:
   for(int i=0: i<v.size(): i++) {</pre>
       ev.push_back({min(v[i].a.x, v[i].b.x), +1, i});
       ev.push_back({max(v[i].a.x, v[i].b.x), -1, i});
   sort(ev.begin(), ev.end());
   auto comp = [&v] (int i, int j) {
       Segment p = v[i], q = v[j];
       Tf x = max(min(p.a.x, p.b.x), min(q.a.x, q.b.x));
       auto yvalSegment = [&x](const Line &s) {
           if(dcmp(s.a.x - s.b.x) == 0) return s.a.y;
           return s.a.v + (s.b.v - s.a.v) * (x - s.a.x) / (s.b.
               x - s.a.x;
       };
       return dcmp(yvalSegment(p) - yvalSegment(q)) < 0;</pre>
   multiset<int, decltype(comp)> st(comp);
   typedef decltype(st)::iterator iter;
   auto prev = [&st](iter it) {
       return it == st.begin() ? st.end() : --it;
   auto next = [&st](iter it) {
       return it == st.end() ? st.end() : ++it:
   vector<iter> pos(v.size());
```

```
for(auto &cur : ev) {
   int id = cur.id;
   if(cur.tp == 1) {
       iter nxt = st.lower bound(id):
       iter pre = prev(nxt);
       if(pre != st.end() && segmentsIntersect(v[*pre], v[
            idl)) return {*pre, id};
       if(nxt != st.end() && segmentsIntersect(v[*nxt], v[
            idl)) return {*nxt, id}:
       pos[id] = st.insert(nxt, id);
   }
    else {
       iter nxt = next(pos[id]);
       iter pre = prev(pos[id]);
       if(pre != st.end() && nxt != st.end() &&
            segmentsIntersect(v[*pre], v[*nxt]))
           return {*pre, *nxt};
       st.erase(pos[id]);
}
return {-1, -1};
```

3.6 Linear

```
// returns true if point p is on segment s
bool onSegment(Point p, Segment s) {
   return dcmp(cross(s.a - p, s.b - p)) == 0 && dcmp(dot(s.a -
        p, s.b - p)) <= 0;
// returns true if segment p && q touch or intersect
bool segmentsIntersect(Segment p, Segment q) {
   if(onSegment(p.a, q) || onSegment(p.b, q)) return true;
   if(onSegment(q.a, p) || onSegment(q.b, p)) return true;
   Ti c1 = cross(p.b - p.a. g.a - p.a):
   Ti c2 = cross(p.b - p.a, q.b - p.a);
   Ti c3 = cross(q.b - q.a, p.a - q.a);
   Ti c4 = cross(q.b - q.a, p.b - q.a);
   return dcmp(c1) * dcmp(c2) < 0 && dcmp(c3) * dcmp(c4) < 0;
bool linesParallel(Line p, Line q) {
   return dcmp(cross(p.b - p.a, q.b - q.a)) == 0;
// lines are represented as a ray from a point: (point, vector)
// returns false if two lines (p, v) && (q, w) are parallel or
// true otherwise, intersection point is stored at o via
bool lineLineIntersection(Point p, Point v, Point q, Point w,
    Point& o) {
   static_assert(is_same<Tf, Ti>::value);
   if(dcmp(cross(v, w)) == 0) return false:
   Point u = p - q;
   o = p + v * (cross(w,u)/cross(v,w));
```

```
return true:
// returns false if two lines p && g are parallel or collinear
// true otherwise, intersection point is stored at o via
     reference
bool lineLineIntersection(Line p, Line q, Point& o) {
   return lineLineIntersection(p.a, p.b - p.a, q.a, q.b - q.a,
         0):
// returns the distance from point a to line 1
Tf distancePointLine(Point p, Line 1) {
   return abs(cross(l.b - l.a, p - l.a) / length(l.b - l.a));
// returns the shortest distance from point a to segment s
Tf distancePointSegment(Point p, Segment s) {
   if(s.a == s.b) return length(p - s.a);
   Point v1 = s.b - s.a, v2 = p - s.a, v3 = p - s.b;
   if(dcmp(dot(v1, v2)) < 0)     return length(v2);</pre>
   else if(dcmp(dot(v1, v3)) > 0) return length(v3);
   else return abs(cross(v1, v2) / length(v1));
// returns the shortest distance from segment p to segment q
Tf distanceSegmentSegment(Segment p, Segment q) {
   if(segmentsIntersect(p, q)) return 0;
   Tf ans = distancePointSegment(p.a, q);
   ans = min(ans, distancePointSegment(p.b, q));
   ans = min(ans, distancePointSegment(q.a, p));
   ans = min(ans, distancePointSegment(q.b, p));
   return ans:
// returns the projection of point p on line 1
Point projectPointLine(Point p, Line 1) {
   static_assert(is_same<Tf, Ti>::value);
   Point v = 1.b - 1.a:
   return 1.a + v * ((Tf) dot(v, p - 1.a) / dot(v, v));
```

3.7 Point Rotation Trick

```
/// you may define the processor function in this namespace
/// instead of passing as an argument; testing shows function
/// defined using lambda and passed as argument performs better
/// tested on:
/// constant width strip - https://codeforces.com/gym/100016/
    problem/I
/// constant area triangle - https://codeforces.com/contest
    /1019/problem/D
/// smallest area quadrilateral - https://codingcompetitions.
    withgoogle.com/codejamio/round/00000000019ff03
    /0000000001b5e89
/// disjoint triangles count - https://codeforces.com/contest
    /1025/problem/F
```

```
/// smallest and largest triangle - http://serjudging.vanb.org
typedef pair< int , int >PII;
void performTrick(vector< Point >pts. const function<void(const</pre>
     vector< Point >&, int)> &processor) {
   int n = pts.size();
   sort(pts.begin(), pts.end());
   vector< int >position(n);
   vector< PII >segments:
   segments.reserve((n*(n-1))/2);
   for (int i = 0; i < n; i++) {</pre>
       position[i] = i;
       for (int j = i+1; j < n; j++) {</pre>
          segments.emplace_back(i, j);
   assert(segments.capacity() == segments.size());
   sort(segments.begin(), segments.end(), [&](PII p, PII q) {
       Ti prod = cross(pts[p.second]-pts[p.first], pts[q.second
           l-pts[a.first]):
       if (prod != 0) return prod > 0;
       return p < a:
   }):
   for (PII seg : segments) {
       int i = position[seg.first];
       assert(position[seg.second] == i+1);
       processor(pts, i);
       swap(pts[i], pts[i+1]);
       swap(position[seg.first], position[seg.second]);
```

3.8 Point

```
typedef double Tf;
typedef double Ti;
                          /// use long long for exactness
const Tf PI = acos(-1), EPS = 1e-9:
int dcmp(Tf x) { return abs(x) < EPS ? 0 : (x<0 ? -1 : 1);}
struct Point {
   Ti x, y;
   Point(Ti x = 0, Ti y = 0) : x(x), y(y) {}
   Point operator + (const Point& u) const { return Point(x +
        u.x, y + u.y); }
   Point operator - (const Point& u) const { return Point(x -
        u.x, y - u.y); }
   Point operator * (const long long u) const { return Point(x
         * u, v * u); }
   Point operator * (const Tf u) const { return Point(x * u, y
   Point operator / (const Tf u) const { return Point(x / u, y
   bool operator == (const Point& u) const { return dcmp(x - u
        (x) == 0 && dcmp(v - u.v) == 0; 
   bool operator != (const Point& u) const { return !(*this ==
         u); }
```

```
bool operator < (const Point& u) const { return dcmp(x - u.
        x) < 0 || (dcmp(x - u.x) == 0 && dcmp(y - u.y) < 0); }
   friend istream &operator >> (istream &is, Point &p) {
        return is >> p.x >> p.v: }
   friend ostream &operator << (ostream &os, const Point &p) {</pre>
         return os << p.x << " " << p.y; }</pre>
Ti dot(Point a, Point b) { return a.x * b.x + a.y * b.y; }
Ti cross(Point a, Point b) { return a.x * b.y - a.y * b.x; }
Tf length(Point a) { return sqrt(dot(a, a)); }
Ti sqLength(Point a) { return dot(a, a); }
Tf distance(Point a, Point b) {return length(a-b);}
Tf angle(Point u) { return atan2(u.y, u.x); }
// returns angle between oa, ob in (-PI, PI]
Tf angleBetween(Point a, Point b) {
   Tf ans = angle(b) - angle(a);
   return ans <= -PI ? ans + 2*PI : (ans > PI ? ans - 2*PI :
// Rotate a ccw by rad radians
Point rotate(Point a, Tf rad) {
   static_assert(is_same<Tf, Ti>::value);
   return Point(a.x * cos(rad) - a.y * sin(rad), a.x * sin(rad
        ) + a.y * cos(rad));
// rotate a ccw by angle th with cos(th) = co && sin(th) = si
Point rotatePrecise(Point a. Tf co. Tf si) {
   static_assert(is_same<Tf, Ti>::value);
   return Point(a.x * co - a.y * si, a.y * co + a.x * si);
Point rotate90(Point a) { return Point(-a.y, a.x); }
// scales vector a by s such that length of a becomes s
Point scale(Point a, Tf s) {
   static_assert(is_same<Tf, Ti>::value);
   return a / length(a) * s;
// returns an unit vector perpendicular to vector a
Point normal(Point a) {
   static assert(is same<Tf. Ti>::value):
   Tf 1 = length(a):
   return Point(-a.v / 1, a.x / 1);
// returns 1 if c is left of ab, 0 if on ab && -1 if right of
int orient(Point a, Point b, Point c) {
   return dcmp(cross(b - a, c - a));
///Use as sort(v.begin(), v.end(), polarComp(0, dir))
///Polar comparator around O starting at direction dir
```

```
struct polarComp {
   Point O, dir;
   polarComp(Point 0 = Point(0, 0), Point dir = Point(1, 0))
       : 0(0), dir(dir) {}
   bool half(Point p) {
       return dcmp(cross(dir, p)) < 0 || (dcmp(cross(dir, p))</pre>
            == 0 && dcmp(dot(dir, p)) > 0);
   }
   bool operator()(Point p, Point q) {
       return make_tuple(half(p), 0) < make_tuple(half(q),</pre>
            cross(p, q));
   }
};
struct Segment {
   Point a. b:
   Segment(Point aa, Point bb) : a(aa), b(bb) {}
};
typedef Segment Line;
struct Circle {
   Point o;
   Tf r:
   Circle(Point o = Point(0, 0), Tf r = 0) : o(o), r(r) {}
   // returns true if point p is in || on the circle
   bool contains(Point p) {
       return dcmp(sqLength(p - o) - r * r) <= 0;</pre>
   // returns a point on the circle rad radians away from +X
        CCW
   Point point(Tf rad) {
       static_assert(is_same<Tf, Ti>::value);
       return Point(o.x + cos(rad) * r, o.y + sin(rad) * r);
   // area of a circular sector with central angle rad
   Tf area(Tf rad = PI + PI) { return rad * r * r / 2: }
   // area of the circular sector cut by a chord with central
   Tf sector(Tf alpha) { return r * r * 0.5 * (alpha - sin(
        alpha)); }
};
```

3.9 Polygon

```
typedef vector<Point> Polygon;
// removes redundant colinear points
// polygon can't be all colinear points
Polygon RemoveCollinear(const Polygon& poly) {
   Polygon ret;
   int n = poly.size();
   for(int i = 0; i < n; i++) {
      Point a = poly[i];
      Point b = poly[(i + 1) % n];
      Point c = poly[(i + 2) % n];</pre>
```

```
if(dcmp(cross(b-a, c-b)) != 0 && (ret.empty() || b !=
            ret.back()))
           ret.push_back(b);
   return ret;
// returns the signed area of polygon p of n vertices
Tf signedPolygonArea(const Polygon &p) {
   Tf ret = 0:
   for(int i = 0; i < (int) p.size() - 1; i++)</pre>
       ret += cross(p[i]-p[0], p[i+1]-p[0]);
   return ret / 2;
// given a polygon p of n vertices, generates the convex hull
// Tested on https://acm.timus.ru/problem.aspx?space=1&num=1185
// Caution: when all points are colinear AND removeRedundant ==
// output will be contain duplicate points (from upper hull) at
Polygon convexHull(Polygon p, bool removeRedundant) {
   int check = removeRedundant ? 0 : -1;
   sort(p.begin(), p.end());
   p.erase(unique(p.begin(), p.end()), p.end());
   int n = p.size();
   Polygon ch(n+n);
   int m = 0;  // preparing lower hull
   for(int i = 0: i < n: i++) {</pre>
       while(m > 1 && dcmp(cross(ch[m - 1] - ch[m - 2], p[i] -
            ch[m-1])) \le check) m--:
       ch[m++] = p[i];
   int k = m; // preparing upper hull
   for(int i = n - 2; i \ge 0; i--) {
       while(m > k && dcmp(cross(ch[m - 1] - ch[m - 2], p[i] -
            ch[m - 2]) <= check) m--:
       ch[m++] = p[i];
   if(n > 1) m--;
   ch.resize(m);
   return ch:
// Tested : https://www.spoj.com/problems/INOROUT
// returns inside = -1, on = 0, outside = 1
int pointInPolygon(const Polygon &p, Point o) {
   using Linear::onSegment;
   int wn = 0, n = p.size();
   for(int i = 0; i < n; i++) {</pre>
       int j = (i + 1) \% n;
       if(onSegment(o, Segment(p[i], p[j])) || o == p[i])
       int k = dcmp(cross(p[j] - p[i], o - p[i]));
       int d1 = dcmp(p[i].y - o.y);
       int d2 = dcmp(p[j].y - o.y);
```

```
if(k > 0 \&\& d1 \le 0 \&\& d2 > 0) wn++:
       if(k < 0 \&\& d2 <= 0 \&\& d1 > 0) wn--;
   }
   return wn ? -1 : 1:
// Tested: Timus 1955, CF 598F
// Given a simple polygon p, and a line l, returns (x, y)
// x = longest segment of 1 in p, y = total length of 1 in p.
pair<Tf, Tf> linePolygonIntersection(Line 1, const Polygon &p)
   using Linear::lineLineIntersection;
   int n = p.size();
   vector<pair<Tf, int>> ev;
   for(int i=0: i<n: ++i) {</pre>
       Point a = p[i], b = p[(i+1)\%n], z = p[(i-1+n)\%n];
       int ora = orient(1.a, 1.b, a), orb = orient(1.a, 1.b, b)
            , orz = orient(l.a, l.b, z);
       if(!ora) {
          Tf d = dot(a - 1.a, 1.b - 1.a):
          if(orz && orb) {
              if(orz != orb) ev.emplace back(d, 0);
              //else // Point Touch
          else if(orz) ev.emplace_back(d, orz);
          else if(orb) ev.emplace_back(d, orb);
       else if(ora == -orb) {
          Point ins;
          lineLineIntersection(l, Line(a, b), ins);
          ev.emplace back(dot(ins - 1.a, 1.b - 1.a), 0):
   sort(ev.begin(), ev.end());
   Tf ans = 0, len = 0, last = 0, tot = 0:
   bool active = false;
   int sign = 0;
   for(auto &qq : ev) {
       int tp = qq.second;
       Tf d = qq.first; /// current Segment is (last, d)
                         /// On Border
       if(sign) {
          len += d-last; tot += d-last;
          ans = max(ans. len):
          if(tp != sign) active = !active;
          sign = 0;
      }
       else {
          if(active) { ///Strictly Inside
              len += d-last; tot += d-last;
              ans = max(ans, len);
          if(tp == 0) active = !active;
          else sign = tp;
       }
       last = d;
       if(!active) len = 0:
```

```
ans /= length(l.b-l.a);
tot /= length(l.b-l.a);
return {ans, tot};
}
```

4 Graphs

4.1 Bridge Tree

```
class Bridge_Tree {
   int n;
   vector <vector <int>> components;
   vector <int> depth, low;
   stack <int> st:
   void find_bridges(int node, Graph &G, int par = -1, int d =
       low[node] = depth[node] = d;
       st.push(node);
       for (int e : G[node]) if(par != e) {
           if(depth[e] == -1) {
              find_bridges(e, G, node, d + 1);
              if (low[e] > depth[node]){
                  bridges.emplace back(node, e):
                  components.push_back({});
                  for(int x = -1; x! = e; x = st.top(), st.pop()
                      components.back().push_back(st.top());
              }
           }
           low[node] = min(low[node], low[e]);
       if(par == -1){
           components.push_back({});
           while(!st.empty())
              components.back().push_back(st.top()), st.pop();
   }
public:
   vector <int> id:
   vector <edge> bridges;
   Graph tree;
   void create_tree() {
       for(auto &comp : components){
           int idx = tree.add_node();
           for(auto &e: comp)
              id[e] = idx;
       for(auto &[1,r]: bridges)
           tree.add_edge(id[1], id[r]);
   Bridge_Tree(Graph &G): n(G.n) {
       depth.assign(n,-1), id.assign(n, -1), low.resize(n);
       for(int i = 0; i < n; i++)</pre>
           if(depth[i] == -1)
              find_bridges(i, G);
};
```

4.2 Centroid Decomposition

```
class Centroid_Decomposition {
   vector <bool> blocked:
   vector <int> CompSize:
   int CompDFS(tree &T, int node, int par) {
       CompSize[node] = 1;
       for(int &e: T[node]) if(e != par and !blocked[e])
          CompSize[node] += CompDFS(T, e, node);
       return CompSize[node];
   int FindCentroid(tree &T, int node, int par, int sz) {
       for(int &e: T[node]) if(e != par and !blocked[e]) if(
            CompSize[e] > sz / 2)
          return FindCentroid(T, e, node, sz);
   pair <int,int> GetCentroid(tree &T, int entry) {
       int sz = CompDFS(T, entry, entry);
       return {FindCentroid(T, entry, entry, sz), sz};
   c_vector <LL> left[2], right[2];
   int GMin. GMax:
   void dfs(tree &T, int node, int par, int Min, int Max, int
       if(blocked[node])
          return;
       right [Max < sum or Min > sum] [sum] ++;
       Max = max(Max, sum), Min = min(Min, sum);
       GMin = min(GMin, sum), GMax = max(GMax, sum);
       for(int i = 0; i < T[node].size(); i++) if(T[node][i] !=</pre>
             par) {
          dfs(T, T[node][i], node, Min, Max, sum + T.col[node
               l[i]):
   LL solve(tree &T, int c, int sz) {
       LL ans = 0;
       left[0].clear(-sz, sz), left[1].clear(-sz, sz);
       for(int i = 0; i < T[c].size(); i++) {</pre>
          GMin = 1, GMax = -1;
          dfs(T, T[c][i], c, GMin, GMax, T.col[c][i]);
          ans += right[0][0] + left[1][0] * right[1][0];
          for(int j:\{0, 1\}) for(int k = GMin; k \le GMax; k++)
              ans += right[j][k] * (left[0][-k] + (j == 0) *
                   left[1][-k]):
          for(int j:\{0, 1\}) for(int k = GMin; k \le GMax; k++)
              left[j][k] += right[j][k];
              right[j][k] = 0;
          }
      }
       return ans;
   }
public:
   LL operator () (tree &T, int entry) {
```

4.3 Dinic

```
#include <bits/stdc++.h>
using namespace std;
struct FlowEdge {
   int v. u:
   long long cap, flow = 0;
   FlowEdge(int v, int u, long long cap) : v(v), u(u), cap(cap
        ) {}
};
struct Dinic {
   const long long flow_inf = 1e18;
   vector<FlowEdge> edg;
   vector<vector<int>> adj;
   int n. m = 0, s. t:
   vector<int> lvl, ptr;
   queue<int> q;
   // number of nodes, source, sink
   Dinic(int n, int s, int t) : n(n), s(s), t(t) {
       adj.resize(n + 100), lvl.resize(n + 100), ptr.resize(n +
             100):
   }
   //directed edge from v to u (not u to v)
   void add_edge(int v, int u, long long cap) {
       edg.emplace_back(v, u, cap);
       edg.emplace_back(u, v, 0);
       adj[v].push_back(m++);
       adj[u].push_back(m++);
   bool bfs() {
       while (!q.empty()) {
          int v = q.front(); q.pop();
          for (int id : adj[v]) {
              if (edg[id].cap - edg[id].flow < 1) continue;</pre>
              if (lvl[edg[id].u] != -1) continue;
              lvl[edg[id].u] = lvl[v] + 1, q.push(edg[id].u);
```

```
return lvl[t] != -1:
   }
   long long dfs(int v, long long pushed) {
       if (pushed == 0) return 0:
       if (v == t) return pushed;
       for (int& cid = ptr[v]; cid < (int)adj[v].size(); cid++)</pre>
           int id = adi[v][cid]. u = edg[id].u;
           if (lvl[v] + 1 != lvl[u] || edg[id].cap - edg[id].
                flow < 1) continue;</pre>
           long long tr = dfs(u, min(pushed, edg[id].cap - edg[
                id].flow));
           if (tr == 0) continue:
           edg[id].flow += tr. edg[id ^ 1].flow -= tr:
           return tr;
       return 0;
   long long flow() {
       long long f = 0:
       while (true) {
           fill(lvl.begin(), lvl.end(), -1);
           lvl[s] = 0, q.push(s);
           if (!bfs()) break;
           fill(ptr.begin(), ptr.end(), 0);
           while (long long pushed = dfs(s, flow_inf)) f +=
                pushed;
       }
       return f:
};
 **directed edge from v to u (not u to v)
 ** to find matching edges go over all edges where s!=u && s !=
     v && t!=u && t!=v and find flow of 1
```

4.4 HLD

```
const int N = 1e6+7;
template <typename DT>
struct Segtree {
    vector<DT> tree, prob, a;
    Segtree(int n) {
        tree.resize(n * 4);
        prob.resize(n), a.resize(n);
    }
    void build(int u, int l, int r) {
        if (1 == r) {
            tree[u] = a[1];
            return;
        }
        int mid = (1 + r) / 2;
        build(u * 2, l, mid);
        build(u * 2 + 1, mid + 1, r);
        tree[u] = (tree[u * 2] + tree[u * 2 + 1]);
```

```
void propagate(int u) {
       prob[u * 2] += prob[u], tree[u * 2] += prob[u];
       prob[u * 2 + 1] += prob[u], tree[u * 2 + 1] += prob[u];
   void update(int u, int l, int r, int i, int j, int val) {
       if (r < i || 1 > j) return;
       if (1 >= i && r <= i) {
          tree[u] = val;
          return;
       int mid = (1 + r) / 2;
       update(u * 2, 1, mid, i, j, val);
       update(u * 2 + 1, mid + 1, r, i, i, val):
       tree[u] = (tree[u * 2] + tree[u * 2 + 1]);
   DT query(int u, int 1, int r, int i, int j) {
       if (1 > j || r < i) return 0;
       if (1 >= i && r <= j) return tree[u];</pre>
       int mid = (1 + r) / 2;
       return (query(u * 2, 1, mid, i, j) + query(u * 2 + 1,
           mid + 1, r, i, i)):
Segtree<int>tree(N);
vector<int> adj[N];
int depth[N], par[N], pos[N];
int head[N], heavy[N], cnt;
int dfs(int u, int p) {
   int SZ = 1, mxsz = 0, heavyc;
   depth[u] = depth[p] + 1;
   for (auto v : adj[u]) {
       if (v == p) continue:
       par[v] = u;
       int subsz = dfs(v, u);
       if (subsz > mxsz) heavv[u] = v. mxsz = subsz:
   return SZ;
void decompose(int u. int h) {
   head[u] = h, pos[u] = ++cnt;
   if(heavy[u]!=-1) decompose(heavy[u], h);
   for(int v : adj[u]) {
       if(v==par[u]) continue;
       if(v!=heavy[u]) decompose(v, v);
int query(int a, int b) {
   int ret = 0:
   for(;head[a]!=head[b]; b=par[head[b]]){
      if(depth[head[a]]>depth[head[b]]) swap(a,b);
      ret += tree.query(1, 0, cnt, pos[head[b]], pos[b]);
```

```
if(depth[a]>depth[b]) swap(a,b);
ret += tree.query(1,0,cnt,pos[a],pos[b]);
return ret;
}
```

4.5 LCA In O(1)

```
* LCA in O(1)
 * depth calculates weighted distance
 * level calculates distance by number of edges
 * Preprocessing in NlongN
#include <bits/stdc++.h>
using namespace std;
typedef long long LL;
typedef pair<int, int> PII;
const int N = 1e6 + 7;
const int L = 21;
namespace LCA {
LL depth[N];
int level[N];
int st[N], en[N], LOG[N], par[N];
int a[N], id[N], table[L][N];
vector<PII> adj[N];
int n, root, Time, cur;
void init(int nodes, int root_) {
   n = nodes, root = root_, LOG[0] = LOG[1] = 0;
   for (int i = 2; i <= n; i++) LOG[i] = LOG[i >> 1] + 1;
   for (int i = 0: i <= n: i++) adi[i].clear():</pre>
void addEdge(int u, int v, int w) {
   adj[u].push_back(PII(v, w));
   adj[v].push_back(PII(u, w));
int lca(int u, int v) {
   if (en[u] > en[v]) swap(u, v):
   if (st[v] <= st[u] && en[u] <= en[v]) return v;</pre>
   int 1 = LOG[id[v] - id[u] + 1];
   int p1 = id[u], p2 = id[v] - (1 << 1) + 1;
   int d1 = level[table[1][p1]], d2 = level[table[1][p2]];
   if (d1 < d2) return par[table[1][p1]];</pre>
   else return par[table[1][p2]]:
```

```
LL dist(int u. int v) {
   int 1 = lca(u, v);
   return (depth[u] + depth[v] - (depth[1] * 2));
/* Euler tour */
void dfs(int u, int p) {
   st[u] = ++Time, par[u] = p;
   for (auto [v, w] : adj[u]) {
       if (v == p) continue;
       depth[v] = depth[u] + w;
       level[v] = level[u] + 1;
       dfs(v. u):
   }
   en[u] = ++Time:
   a[++cur] = u, id[u] = cur;
/* RMQ */
void pre() {
   cur = Time = 0, dfs(root, root);
   for (int i = 1; i <= n; i++) table[0][i] = a[i];</pre>
   for (int 1 = 0; 1 < L - 1; 1++) {
       for (int i = 1; i <= n; i++) {</pre>
           table[1 + 1][i] = table[1][i];
           bool C1 = (1 << 1) + i <= n;
           bool C2 = level[table[1][i + (1 << 1)]] < level[</pre>
                table[1][i]];
           if (C1 && C2) table[1 + 1][i] = table[1][i + (1 << 1</pre>
/* namespace LCA */
//tested on kattis-greatestpair
using namespace LCA;
int main() {
   ios::sync_with_stdio(0);
   cin.tie(0);
```

4.6 SCC

```
typedef long long LL;
const LL N = 1e6 + 7;
bool vis[N];
vector<int> adj[N], adjr[N];
```

```
vector<int> order, component:
// tp = 0 ,finding topo order, tp = 1 , reverse edge traversal
void dfs(int u, int tp = 0) {
   vis[u] = true:
   if (tp) component.push_back(u);
   auto& ad = (tp ? adjr : adj);
   for (int v : ad[u])
       if (!vis[v]) dfs(v, tp):
   if (!tp) order.push_back(u);
int main() {
   for (int i = 1; i <= n; i++) {</pre>
       if (!vis[i]) dfs(i):
   memset(vis, 0, sizeof vis);
   reverse(order.begin(), order.end());
   for (int i : order) {
       if (!vis[i]) {
           // one component is found
           dfs(i, 1), component.clear();
   }
```

4.7 StoerWanger

```
/* for finding the min cut of a graph without specifing the
     source and the sink.
  all the edges are directed and no need to make any edge
       bidirectional.
const int N = 1407;
// O(n^3) but faster, 1 indexed
mt19937 rnd(chrono::steady_clock::now().time_since_epoch().
    count()):
struct StoerWagner {
   int n. idx[N]:
   LL G[N][N], dis[N];
   bool vis[N];
   const LL inf = 1e18;
   StoerWagner() {}
   StoerWagner(int _n) {
       n = _n;
       memset(G, 0, sizeof G);
   void add_edge(int u, int v, LL w) { // undirected edge,
        multiple edges are merged into one edge
       if (u != v) G[u][v] += w, G[v][u] += w;
   LL solve() {
       LL ans = inf:
       for (int i = 0; i < n; ++i) idx[i] = i + 1;
       shuffle(idx, idx + n, rnd);
```

```
while (n > 1) {
          int t = 1, s = 0;
          for (int i = 1; i < n; ++i) {</pre>
              dis[idx[i]] = G[idx[0]][idx[i]]:
              if (dis[idx[i]] > dis[idx[t]]) t = i;
          memset(vis, 0, sizeof vis);
          vis[idx[0]] = true:
          for (int i = 1; i < n; ++i) {</pre>
              if (i == n - 1) {
                  if (ans > dis[idx[t]])
                      ans = dis[idx[t]]: // idx[s] - idx[t] is
                           in two halves of the mincut
                  if (ans == 0) return 0;
                  for (int j = 0; j < n; ++j) {
                      G[idx[s]][idx[j]] += G[idx[j]][idx[t]];
                      G[idx[j]][idx[s]] += G[idx[j]][idx[t]];
                  }
                  idx[t] = idx[--n];
              vis[idx[t]] = true;
              s = t, t = -1;
              for (int j = 1; j < n; ++j) {
                  if (!vis[idx[j]]) {
                      dis[idx[j]] += G[idx[s]][idx[j]];
                      if (t == -1 || dis[idx[t]] < dis[idx[j]])</pre>
                            t = j;
              }
       return ans;
};
```

4.8 Tree Algo

```
struct tree {
   int n;
   vector <vector <int> > adj;
   inline vector<iint>& operator[](int u) {
      return adj[u];
   }
   tree(int n = 0) : n(n), adj(n) {}
   void addEdge(int u, int v) {
      adj[u].push_back(v);
      adj[v].push_back(u);
   }
};
struct lca_table {
   tree &T;
   int n, LOG = 20;
   vector <vector <int>> anc;
   vector <int> level;
```

```
void setupLifting(int node, int par) {
       for (int v : T[node]) if (v != par) {
               anc[v][0] = node. level[v] = level[node] + 1:
              for (int k = 1; k < LOG; k++)
                  anc[v][k] = anc[anc[v][k - 1]][k - 1];
              setupLifting(v, node);
   lca_table(tree &T, int root = 0): T(T), n(T.n) {
       LOG = 33 - __builtin_clz(n);
       anc.assign(n, vector <int> (LOG, root));
       level.resize(n);
       setupLifting(root, root);
   int lca(int u, int v) {
       if (level[u] > level[v])
           swap(u, v);
       for (int k = LOG - 1; ~k; k--)
           if (level[u] + (1 << k) <= level[v])</pre>
              v = anc[v][k]:
       if (u == v)
           return u:
       for (int k = LOG - 1; "k; k--)
           if (anc[u][k] != anc[v][k])
              u = anc[u][k], v = anc[v][k];
       return anc[u][0];
   int getAncestor(int node, int ht) {
       for (int k = 0; k < LOG; k++)
           if (ht & (1 << k))
              node = anc[node][k];
       return node:
   int distance(int u, int v) {
       int g = lca(u, v);
       return level[u] + level[v] - 2 * level[g];
}:
struct euler_tour {
   int time = 0;
   tree &T;
   int n;
   vector <int> start, finish, level, par;
   euler_tour(tree &T, int root = 0) : T(T), n(T.n), start(n),
         finish(n), level(n), par(n) {
       time = 0:
       call(root);
   void call(int node, int p = -1) {
       if (p != -1) level[node] = level[p] + 1;
       start[node] = time++:
       for (int e : T[node]) if (e != p)
              call(e, node);
       par[node] = p:
       finish[node] = time++;
   bool isAncestor(int node, int par) {
```

```
return start[par] <= start[node] and finish[par] >=
            finish[node];
   int subtreeSize(int node) {
       return finish[node] - start[node] + 1 >> 1;
   }
tree virtual_tree(vector <int> &nodes, lca_table &table,
     euler tour &tour) {
   sort(nodes.begin(), nodes.end(), [&](int x, int y) {
       return tour.start[x] < tour.start[y];</pre>
   }):
   int n = nodes.size();
   for (int i = 0: i + 1 < n: i++)
       nodes.push back(table.lca(nodes[i], nodes[i + 1]));
   sort(nodes.begin(), nodes.end());
   nodes.erase(unique(nodes.begin(), nodes.end()), nodes.end()
   sort(nodes.begin(), nodes.end(), [&](int x, int y) {
       return tour.start[x] < tour.start[v]:</pre>
   });
   n = nodes.size();
   stack <int> st:
   st.push(0);
   tree ans(n):
   for (int i = 1; i < n; i++) {</pre>
       while (!tour.isAncestor(nodes[i], nodes[st.top()])) st.
       ans.addEdge(st.top(), i);
       st.push(i);
   return ans;
set <int> getCenters(tree &T) {
   int n = T.n;
   vector <int> deg(n), q;
   set <int> s;
   for (int i = 0: i < n: i++) {</pre>
       deg[i] = T[i].size():
       if (deg[i] == 1)
           q.push_back(i);
       s.insert(i);
   for (vector <int> t ; s.size() > 2; q = t) {
       for (auto x : q) {
           for (auto e : T[x])
              if (--deg[e] == 1)
                  t.push_back(e);
           s.erase(x):
       }
   return s;
bool check(tree &T) {
   for (int i = 0: i < T.n: i++)</pre>
       if (T[i].size() > 2) return false;
   return true:
```

5 Math

5.1 Adaptive Simpsons

```
For finding the length of an arc in a range
            L = integrate(ds) from start to end of range
            where ds = sqrt(1+(d/dy(x))^2)dy
const double SIMPSON TERMINAL EPS = 1e-12:
/// Function whose integration is to be calculated
double F(double x):
double simpson(double minx, double maxx)
             return (\max - \min x) / 6 * (F(\min x) + 4 * F((\min x + \max x)) / 6 * (F(\min x) + 4 * F((\min x + \max x)) / 6 * (F(\min x) + 4 * F((\min x) + \max x)) / 6 * (F(\min x) + 4 * F((\min x) + \max x)) / 6 * (F(\min x) + 4 * F((\min x) + \max x)) / 6 * (F(\min x) + 4 * F((\min x) + \max x)) / 6 * (F(\min x) + 4 * F((\min x) + \max x)) / 6 * (F(\min x) + 4 * F((\min x) + \max x)) / 6 * (F(\min x) + \max x) / 6 * (F(\min x) + \max x
                                    2.) + F(maxx):
double adaptive_simpson(double minx, double maxx, double c,
                 double EPS)
                   if(maxx - minx < SIMPSON TERMINAL EPS) return 0:</pre>
             double midx = (minx + maxx) / 2;
            double a = simpson(minx, midx):
            double b = simpson(midx, maxx);
             if(fabs(a + b - c) < 15 * EPS) return a + b + (a + b - c) /
                                    15.0:
            return adaptive_simpson(minx, midx, a, EPS / 2.) +
                                adaptive_simpson(midx, maxx, b, EPS / 2.);
double adaptive_simpson(double minx, double maxx, double EPS)
            return adaptive_simpson(minx, maxx, simpson(minx, maxx, i),
                                    EPS):
```

5.2 Berlekamp Massey

```
struct berlekamp_massey { // for linear recursion
typedef long long LL;
static const int SZ = 2e5 + 5:
static const int MOD = 1e9 + 7; /// mod must be a prime
LL m, a[SZ], h[SZ], t_{SZ}, s[SZ], t[SZ];
// bigmod goes here
inline vector <LL> BM( vector <LL> &x ) {
  LL lf . ld:
  vector <LL> ls . cur:
  for ( int i = 0; i < int(x.size()); ++i ) {</pre>
    for ( int j = 0; j < int(cur.size()); ++j ) t = (t + x[i -</pre>
          i - 1] * cur[i]) % MOD;
    if ((t - x[i]) \% MOD == 0) continue;
    if (!cur.size()) {
      cur.resize( i + 1 );
      lf = i: ld = (t - x[i]) % MOD:
      continue;
```

```
LL k = -(x[i] - t) * bigmod(ld . MOD - 2 . MOD) % MOD:
     vector <LL> c(i - lf - 1);
     c.push_back( k );
     for ( int j = 0; j < int(ls.size()); ++j ) c.push_back(-ls</pre>
          [i] * k % MOD):
     if ( c.size() < cur.size() ) c.resize( cur.size() );</pre>
     for ( int j = 0; j < int(cur.size()); ++j ) c[j] = (c[j] +</pre>
           cur[i]) % MOD;
     if (i - lf + (int)ls.size() >= (int)cur.size() ) ls = cur.
           lf = i, ld = (t - x[i]) \% MOD;
     cur = c;
    for ( int i = 0; i < int(cur.size()); ++i ) cur[i] = (cur[i])</pre>
        1 % MOD + MOD) % MOD:
   return cur:
  inline void mull( LL *p , LL *q ) {
   for ( int i = 0; i < m + m; ++i ) t_{-}[i] = 0;
   for ( int i = 0; i < m; ++i ) if ( p[i] )</pre>
       for ( int j = 0; j < m; ++j ) t_{i} = (t_{i} + j) + p
            [i] * q[i]) % MOD;
   for ( int i = m + m - 1; i \ge m; --i ) if ( t [i] )
       for ( int j = m - 1; (i - j - 1) = (t_[i - j])
             -1] + t_{[i]} * h[i]) % MOD;
   for ( int i = 0; i < m; ++i ) p[i] = t_[i];</pre>
  inline LL calc( LL K ) {
   for ( int i = m; ~i; --i ) s[i] = t[i] = 0;
   s[0] = 1; if (m!=1) t[1] = 1; else t[0] = h[0];
   while ( K ) {
     if (K & 1 ) mull(s, t);
     mull( t , t ); K >>= 1;
   LL su = 0;
   for ( int i = 0; i < m; ++i ) su = (su + s[i] * a[i]) % MOD
   return (su % MOD + MOD) % MOD;
  /// already calculated upto k . now calculate upto n.
  inline vector <LL> process( vector <LL> &x , int n , int k )
      {
   auto re = BM( x );
   x.resize(n+1);
   for ( int i = k + 1: i <= n: i++ ) {
     for ( int j = 0; j < re.size(); j++ ) {</pre>
       x[i] += 1LL * x[i - j - 1] % MOD * re[j] % MOD; x[i] %=
            MOD:
     }
   return x;
  inline LL work( vector <LL> &x , LL n ) {
   if ( n < int(x.size()) ) return x[n] % MOD:</pre>
   vector <LL> v = BM( x ); m = v.size(); if ( !m ) return 0;
   for ( int i = 0; i < m; ++i ) h[i] = v[i], a[i] = x[i];
   return calc( n ) % MOD;
} rec;
```

```
vector <LL> v:
void solve() {
 int n;
 cin >> n:
 cout << rec.work(v, n - 1) << endl;</pre>
5.3 Combi
```

```
const int N = 2e5 + 5:
const int mod = 1e9 + 7:
array < int, N + 1 > fact, inv, inv_fact, Drng;
void init() {
  fact[0] = inv_fact[0] = 1;
   for (int i = 1: i <= N: i++) {
       inv[i] = i == 1 ? 1 : (LL) inv[i - mod % i] * (mod / i +
             1) % mod;
       fact[i] = (LL) fact[i - 1] * i % mod:
       inv fact[i] = (LL) inv_fact[i - 1] * inv[i] % mod;
   Drng[0] = 1, Drng[1] = 0;
   for(int i = 2; i <= N; i++)</pre>
       Drng[i] = (LL) (i - 1) * (Drng[i - 1] + Drng[i - 2]) %
int C(int n. int r) {
   if (fact[0] != 1) init():
   return (r < 0 \text{ or } r > n) ? 0 : (LL) fact[n] * inv_fact[r] %
        mod * inv fact[n - r] % mod:
int D(int n) {
   return n < 0 ? 0 : Drng[n]:
```

5.4 FFT

```
using CD = complex<double>;
typedef long long LL;
const double PI = acos(-1.0L):
int N;
vector<int> perm;
vector<CD> wp[2];
void precalculate(int n) {
   assert((n & (n - 1)) == 0), N = n;
   perm = vector<int>(N, 0);
   for (int k = 1: k < N: k <<= 1) {
       for (int i = 0: i < k: i++) {</pre>
           perm[i] <<= 1;
           perm[i + k] = 1 + perm[i]:
   wp[0] = wp[1] = vector < CD > (N);
   for (int i = 0; i < N; i++) {</pre>
       wp[0][i] = CD(cos(2 * PI * i / N), sin(2 * PI * i / N));
       Wp[1][i] = CD(cos(2 * PI * i / N), -sin(2 * PI * i / N))
```

```
void fft(vector<CD> &v, bool invert = false) {
   if (v.size() != perm.size()) precalculate(v.size());
   for (int i = 0: i < N: i++)</pre>
       if (i < perm[i]) swap(v[i], v[perm[i]]);</pre>
   for (int len = 2; len <= N; len *= 2) {</pre>
       for (int i = 0, d = N / len: i < N: i += len) {</pre>
           for (int j = 0, idx = 0; j < len / 2; j++, idx += d)
               CD x = v[i + j];
              CD y = wp[invert][idx] * v[i + j + len / 2];
              v[i + j] = x + y;
              v[i + j + len / 2] = x - y;
       }
   }
   if (invert) {
       for (int i = 0; i < N; i++) v[i] /= N;
void pairfft(vector<CD> &a, vector<CD> &b, bool invert = false)
   int N = a.size():
   vector<CD> p(N);
   for (int i = 0; i < N; i++) p[i] = a[i] + b[i] * CD(0, 1);</pre>
   fft(p, invert);
   p.push_back(p[0]);
   for (int i = 0; i < N; i++) {</pre>
       if (invert) {
           a[i] = CD(p[i].real(), 0);
           b[i] = CD(p[i].imag(), 0):
       } else {
           a[i] = (p[i] + conj(p[N - i])) * CD(0.5, 0);
           b[i] = (p[i] - conj(p[N - i])) * CD(0, -0.5);
   }
vector<LL> multiply(const vector<LL> &a. const vector<LL> &b) {
   int. n = 1:
   while (n < a.size() + b.size()) n <<= 1;</pre>
   vector<CD> fa(a.begin(), a.end()), fb(b.begin(), b.end());
   fa.resize(n);
   fb.resize(n):
            fft(fa): fft(fb):
   pairfft(fa, fb);
   for (int i = 0; i < n; i++) fa[i] = fa[i] * fb[i];
   fft(fa, true):
   vector<LL> ans(n);
   for (int i = 0: i < n: i++) ans[i] = round(fa[i].real()):</pre>
   return ans;
const int M = 1e9 + 7, B = sqrt(M) + 1;
vector<LL> anyMod(const vector<LL> &a, const vector<LL> &b) {
   int. n = 1:
   while (n < a.size() + b.size()) n <<= 1;</pre>
   vector<CD> al(n), ar(n), bl(n), br(n);
   for (int i = 0: i < a.size(): i++)</pre>
       al[i] = a[i] % M / B, ar[i] = a[i] % M % B;
```

```
for (int i = 0: i < b.size(): i++)</pre>
   bl[i] = b[i] % M / B, br[i] = b[i] % M % B;
pairfft(al, ar);
pairfft(bl, br);
         fft(al); fft(ar); fft(bl); fft(br);
for (int i = 0; i < n; i++) {</pre>
   CD 11 = (al[i] * bl[i]), lr = (al[i] * br[i]);
   CD rl = (ar[i] * bl[i]), rr = (ar[i] * br[i]);
   al[i] = 11:
   ar[i] = lr;
   bl[i] = rl;
   br[i] = rr;
pairfft(al. ar. true):
pairfft(bl. br. true):
         fft(al, true); fft(ar, true); fft(bl, true); fft(
    br. true):
vector<LL> ans(n):
for (int i = 0; i < n; i++) {</pre>
   LL right = round(br[i].real()), left = round(al[i].real
        ()):
   LL mid = round(round(bl[i].real()) + round(ar[i].real())
   ans[i] = ((left % M) * B * B + (mid % M) * B + right) %
return ans;
```

Fractional Binary Search

```
Given a function f and n, finds the smallest fraction p / q in
    [0, 1] or [0,n]
such that f(p / q) is true, and p, q <= n.
Time: O(\log(n))
struct frac { long long p, q; };
bool f(frac x) {
return 6 + 8 * x.p >= 17 * x.q + 12;
frac fracBS(long long n) {
 bool dir = 1, A = 1, B = 1;
 frac lo\{0, 1\}, hi\{1, 0\}; // Set hi to 1/0 to search within
      [0, n] and \{1, 1\} to search within [0, 1]
 if (f(lo)) return lo:
 assert(f(hi)): //checking if any solution exists or not
 while (A || B) {
   long long adv = 0, step = 1; // move hi if dir, else lo
   for (int si = 0; step; (step *= 2) >>= si) {
     frac mid{lo.p * adv + hi.p, lo.q * adv + hi.q};
     if (abs(mid.p) > n || mid.q > n || dir == !f(mid)) {
       adv -= step: si = 2:
     }
   hi.p += lo.p * adv;
```

```
hi.q += lo.q * adv:
 dir = !dir;
 swap(lo, hi);
 A = B: B = !!adv:
return dir ? hi : lo;
```

Gaussian Elimination

```
double gaussian elimination(int row, int col) {
   int basis[30]:
   for (int j = 0; j < row; j++) {</pre>
       MAT[j][j + col] = 1;
   memset(basis, -1, sizeof basis);
   double det = 1:
   for (int i = 0; i < col; i++) {</pre>
       for (int p = 0; p < row ; p++) {</pre>
          for (int q = 0; q < col; q++)</pre>
              cout << MAT[p][q] << ' ';
          cout << '\n':
      }
       int x = -1;
       for (int k = 0: k < row: k++) {
          if (abs(MAT[k][i]) > eps and basis[k] == -1) {
              x = k, det *= MAT[k][i], basis[x] = i:
              break:
       }
       if (x < 0) continue;
       for (int j = 0; j < col; j++)
          if (j != i) for (int k = 0; k < row; k++) if (k != x | 5.8 Pollard Rho
                     MAT[k][j] = (MAT[k][i] * MAT[x][j]) /
                          MAT[x][i]:
       for (int k = 0; k < col; k++) if (k != i)
              MAT[x][k] /= MAT[x][i]:
       for (int j = 0; j < row; j++)
          MAT[j][i] = (j == i);
   for (int i = 0; i < row ; i++) {</pre>
       for (int j = 0; j < col; j++)</pre>
          cout << MAT[i][j] << ' ';
       cout << '\n';
   for (int k = 0; k < row; k++)
       if (basis[k] == -1)
          return 0;
   return det;
```

Linear Sieve

```
const int maxn = 1e7;
vector <int> primes:
int spf[maxn+5], phi[maxn+5], NOD[maxn+5], cnt[maxn+5], POW[
    maxn+51:
bool prime[maxn+5];
```

```
int SOD[maxn+5]:
void init(){
   fill(prime+2, prime+maxn+1, 1);
   SOD[1] = NOD[1] = phi[1] = spf[1] = 1;
   for(LL i=2;i<=maxn;i++){</pre>
       if(prime[i]) {
          primes.push_back(i), spf[i] = i;
          phi[i] = i-1;
          NOD[i] = 2, cnt[i] = 1:
          SOD[i] = i+1, POW[i] = i;
       for(auto p:primes){
          if(p*i>maxn or p > spf[i]) break;
          prime[p*i] = false, spf[p*i] = p;
          if(i\%p == 0){
              phi[p*i]=p*phi[i];
              NOD[p*i]=NOD[i]/(cnt[i]+1)*(cnt[i]+2), cnt[p*i]=
              SOD[p*i]=SOD[i]/SOD[POW[i]]*(SOD[POW[i]]+p*POW[i
                   ]),POW[p*i]=p*POW[i];
              break;
          } else {
              phi[p*i]=phi[p]*phi[i];
              NOD[p*i]=NOD[p]*NOD[i], cnt[p*i]=1;
              SOD[p*i]=SOD[p]*SOD[i], POW[p*i]=p;
       }
   }
```

```
ULL mul(ULL a, ULL b, ULL mod) {
   LL ans = a * b - mod * (ULL) (1.L / mod * a * b);
   return ans + mod * (ans < 0) - mod * (ans >= (LL) mod);
ULL bigmod(ULL num,ULL pow,ULL mod){
   ULL ans = 1:
   for( ; pow > 0; pow >>= 1, num = mul(num, num, mod))
       if(pow&1) ans = mul(ans,num,mod);
   return ans:
bool is_prime(ULL n){
   if(n < 2 or n % 6 % 4 != 1)
       return (n|1) == 3;
   ULL a[] = {2, 325, 9375, 28178, 450775, 9780504,
        1795265022}:
   ULL s = __builtin_ctzll(n-1), d = n >> s;
   for(ULL x: a){
       ULL p = bigmod(x \% n, d, n), i = s;
       for (; p != 1 \text{ and } p != n-1 \text{ and } x \% n \text{ and } i--; p = mul(p,
             p, n));
       if(p != n-1 and i != s)
           return false:
   return true;
```

```
ULL get factor(ULL n) {
   auto f = [&](LL x) { return mul(x, x, n) + 1; };
   ULL x = 0, y = 0, t = 0, prod = 2, i = 2, q;
   for( : t++ \frac{40}{40} or gcd(prod, n) == 1: x = f(x), y = f(f(y))
       (x == y) ? x = i++, y = f(x) : 0;
       prod = (q = mul(prod, max(x,y) - min(x,y), n)) ? q :
   return gcd(prod, n);
map <ULL, int> factorize(ULL n){
   map <ULL, int> res;
   if(n < 2) return res:</pre>
   ULL small primes [] = \{2, 3, 5, 7, 11, 13, 17, 19, 23, 29, ...\}
        31, 37, 41, 43, 47, 53, 59, 61, 67, 71, 73, 79, 83, 89,
         97 }:
   for (ULL p: small_primes)
       for( ; n % p == 0; n /= p, res[p]++);
   auto _factor = [&](ULL n, auto &_factor) {
       if(n == 1) return;
       if(is_prime(n))
           res[n]++;
       else {
           ULL x = get_factor(n);
           _factor(x, _factor);
           _factor(n / x, _factor);
   };
   _factor(n, _factor);
   return res;
```

5.9 Subset-Convolution

```
#include <bits/stdc++.h>
#define monke_flip ios_base::sync_with_stdio(false); cin.tie(
    NULL);
#define random_monke chrono::system_clock::now().
    time_since_epoch().count()
#ifdef LEL
#include <dbg.h>
#else
#define dbg(...) {/* temon kichu na; */}
#endif

using namespace std;
using LL = long long;
using ULL = unsigned long long;
mt19937_64 rng(random_monke);
const int MONKE = 0;

/*
```

```
inline int sgn(int mask) {
   return 1 - 2 * (__builtin_popcount(mask) & 1);
} // returns 1 if set cardinality is even, -1 otherwise
template <typename T, int b> struct Subset {
   static const int N = 1 << b:
   array <T, N> F;
   void Zeta() { // SOS
       for(int i = 0; i < b; i++)</pre>
           for(int mask = 0; mask < N; mask++)</pre>
               if(mask & 1 << i)
                   F[mask] += F[mask ^ 1 << i]:</pre>
   void OddEven() {
       for(int mask = 0: mask < N: mask++)</pre>
           F[mask] *= sgn(mask);
   void MobiusOld() {
       OddEven():
       Zeta():
       OddEven();
   void Mobius(){
       for(int i = 0; i < b; i++)</pre>
           for(int mask = 0; mask < N; mask++)</pre>
               if(mask & 1 << i)</pre>
                   F[mask] -= F[mask ^ 1 << i]:</pre>
   void operator *= (Subset &R) {
       auto &G = R.F:
       array < array < int, N>, b> Fh = {0}, Gh = {0}, H = {0};
       for(int mask = 0: mask < N: mask++)</pre>
           Fh[__builtin_popcount(mask)][mask] = F[mask], Gh[
                __builtin_popcount(mask)] [mask] = G[mask];
       for(int i = 0: i < b: i++)
           for(int j = 0; j < b; j++)
               for(int mask = 0; mask < N; mask++)</pre>
                   if((mask & (1 << j)) != 0)</pre>
                       Fh[i][mask] += Fh[i][mask ^ (1 << i)]. Gh
                            [i] [mask] += Gh[i] [mask ^ (1 << j)];
       for(int mask = 0; mask < N; mask++)</pre>
           for(int i = 0; i < b; i++)</pre>
               for(int i = 0: i <= i: i++)
                   H[i][mask] += Fh[i][mask] * Gh[i - i][mask];
       for(int i = 0; i < b; i++)</pre>
           for(int j = 0; j < b; j++)
               for(int mask = 0; mask < N; mask++)</pre>
                   if((mask & (1 << j)) != 0)</pre>
                       H[i][mask] -= H[i][mask ^ (1 << j)];</pre>
       for(int mask = 0; mask < N; mask++)</pre>
```

```
F[mask] = H[__builtin_popcount(mask)][mask];
....}...
Subset operator * (Subset &R) {
    Subset ans = *this;
    return ans;
};
int main()
{
    monke_flip
    Subset <int, 3> S, T;
    S.F = {1, 3, 4, 9, 3, 7, 9, 8};
    T = S;
    S.Zeta();
    S.Mobius();
    dbg(S.F == T.F);
    return MONKE;
}
```

5.10 Xor Basis

```
struct YorBasis {
   static const int sz = 64;
   array <ULL, sz> base = {0}, back;
   array <int, sz> pos;
   void insert(ULL x, int p) {
      ULL cur = 0:
      for(int i = sz - 1; i; i--) if (x >> i & 1) {
          if(!base[i]) {
              base[i] = x, back[i] = cur, pos[i] = p;
          } else x ^= base[i], cur |= 1ULL << i;</pre>
      }
   }
   pair <ULL. vector <int>> construct(ULL mask) {
      ULL ok = 0, x = mask;
      for(int i = sz - 1; ~i; i--)
          if(mask >> i & 1 and base[i])
              mask ^= base[i], ok |= 1ULL << i;
       vector <int> ans:
       for(int i = 0; i < sz; i++) if(ok >> i & 1) {
          ans.push_back(pos[i]);
          ok ^= back[i]:
       return {x ^ mask, ans};
```

|6 String

6.1 Aho Corasick

```
const int sg = 26, N = 1e3 + 9;
struct aho_corasick {
    struct node{
        node *link, *out, *par;
        bool leaf;
        LL val;
```

```
int cnt. last. len:
       char p_ch;
       array <node*, sg> to;
       node(node* par = NULL, char p_ch = '$', int len = 0):
       par(par), p_ch(p_ch), len(len) {
           val = leaf = cnt = last = 0;
           link = out = NULL:
       }
   };
   vector <node> trie;
   node *root;
   aho_corasick(){
       trie.reserve(N), trie.emplace_back();
       root = &trie[0]:
       root-> link = root -> out = root;
   inline int f(char c){
       return c - 'a':
   inline node* add_node(node* par = NULL, char p_ch = '$',
        int len = 0){
       trie.emplace_back(par, p_ch, len);
       return &trie.back();
   void add_str(const string& s, LL val = 1){
       node* now = root:
       for(char c: s){
           int i = f(c):
           if(!now->to[i])
              now->to[i] = add_node(now, c, now->len + 1);
           now = now \rightarrow to[i]:
       now -> leaf = true. now -> val++:
   void push_links(){
       queue <node*> q;
       for(q.push(root); q.empty(); q.pop()){
           node *cur = a.front(). *link = cur -> link:
           cur -> out = link -> leaf ? link : link-> out:
           int idx = 0:
           for(auto &next: cur -> to) {
               if(next != NULL){
                  next -> link = cur != root ? link -> to [idx
                       ++] : root;
                  q.push(next);
               else next = link -> to[idx++];
       cur -> val += link -> val;
   }
};
```

6.2 Prefix Function

```
vector<int> prefix_function(string s) {
  int n = (int)s.length();
```

```
vector<int> pi(n);
for (int i = 1; i < n; i++) {
   int j = pi[i - 1];
   while (j > 0 && s[i] != s[j]) j = pi[j - 1];
   if (s[i] == s[j]) j++;
   pi[i] = j;
}
return pi;
}
```

6.3 Z Algo

```
vector<int> calcz(string s) {
   int n = s.size():
   vector<int> z(n);
   int 1. r: 1 = r = 0:
   for (int i = 1; i < n; i++) {</pre>
       if (i > r) {
          1 = r = i:
           while (r < n \&\& s[r] == s[r - 1]) r++:
          z[i] = r - 1; r--;
       } else {
           int k = i - 1;
           if (z[k] < r - i + 1) z[i] = z[k];
              while (r < n \&\& s[r] == s[r - 1]) r++:
              z[i] = r - 1; r--;
       }
   return z;
```

6.4 double hash

ans = (ans * ans) % mod;

```
ostream& operator << (ostream& os, PLL hash) {
    return os << "(" << hash.ff << ", " << hash.ss << ")";
}

PLL operator + (PLL a, LL x) {return PLL(a.ff + x, a.ss + x);}

PLL operator - (PLL a, LL x) {return PLL(a.ff - x, a.ss - x);}

PLL operator * (PLL a, LL x) {return PLL(a.ff * x, a.ss * x);}

PLL operator + (PLL a, PLL x) {return PLL(a.ff * x, a.ss * x);}

PLL operator + (PLL a, PLL x) {return PLL(a.ff + x.ff, a.ss + x .ss);}

PLL operator - (PLL a, PLL x) {return PLL(a.ff - x.ff, a.ss - x .ss);}

PLL operator * (PLL a, PLL x) {return PLL(a.ff * x.ff, a.ss * x .ss);}

PLL operator % (PLL a, PLL m) {return PLL(a.ff % m.ff, a.ss % m .ss);}

PLL base(1949313259, 1997293877);

PLL power (PLL a, LL p) {
    if (!p) return PLL(1, 1);
    PLL ans = power(a, p / 2);
</pre>
```

```
if (p \% 2) ans = (ans * a) \% mod:
 return ans;
PLL inverse(PLL a) {
 return power(a, (mod.ff - 1) * (mod.ss - 1) - 1);
PLL inv_base = inverse(base);
PLL val;
vector<PLL> P;
void hash_init(int n) {
 P.resize(n + 1):
 P[0] = PLL(1, 1):
 for (int i = 1; i <= n; i++) P[i] = (P[i - 1] * base) % mod;
///appends c to string
PLL append(PLL cur, char c) {
 return (cur * base + c) % mod;
///prepends c to string with size k
PLL prepend(PLL cur, int k, char c) {
 return (P[k] * c + cur) % mod;
///replaces the i-th (0-indexed) character from right from a to
PLL replace(PLL cur, int i, char a, char b) {
 cur = (cur + P[i] * (b - a)) \% mod;
 return (cur + mod) % mod:
///Erases c from the back of the string
PLL pop_back(PLL hash, char c) {
 return (((hash - c) * inv base) % mod + mod) % mod:
///Erases c from front of the string with size len
PLL pop_front(PLL hash, int len, char c) {
 return ((hash - P[len - 1] * c) % mod + mod) % mod;
///concatenates two strings where length of the right is k
PLL concat(PLL left, PLL right, int k) {
 return (left * P[k] + right) % mod;
///Calculates hash of string with size len repeated cnt times
///This is O(log n). For O(1), pre-calculate inverses
PLL repeat(PLL hash, int len, LL cnt) {
 PLL mul = (P[len * cnt] - 1) * inverse(P[len] - 1);
 mul = (mul % mod + mod) % mod:
 PLL ret = (hash * mul) % mod:
 if (P[len].ff == 1) ret.ff = hash.ff * cnt;
 if (P[len].ss == 1) ret.ss = hash.ss * cnt:
 return ret;
```

```
LL get(PLL hash) {
 return ( (hash.ff << 32) ^ hash.ss );</pre>
struct hashlist {
 int len;
 vector<PLL> H, R;
 hashlist() {}
 hashlist(string &s) {
   len = (int)s.size();
   hash_init(len);
   H.resize(len + 1, PLL(0, 0)), R.resize(len + 2, PLL(0, 0));
   for (int i = 1; i <= len; i++) H[i] = append(H[i - 1], s[i</pre>
   for (int i = len; i >= 1; i--) R[i] = append(R[i + 1], s[i
        - 1]);
  /// 1-indexed
  inline PLL range_hash(int 1, int r) {
   int len = r - 1 + 1;
   return ((H[r] - H[1 - 1] * P[len]) % mod + mod) % mod;
  inline PLL reverse_hash(int 1, int r) {
   int len = r - 1 + 1;
   return ((R[1] - R[r + 1] * P[len]) % mod + mod) % mod;
  inline PLL concat_range_hash(int 11, int r1, int 12, int r2)
   int len_2 = r2 - 12 + 1;
   return concat(range_hash(11, r1), range_hash(12, r2), len_2
        );
 inline PLL concat_reverse_hash(int 11, int r1, int 12, int r2
      ) {
   int len_1 = r1 - l1 + 1;
   return concat(reverse_hash(12, r2), reverse_hash(11, r1),
        len_1);
};
```

Equations and Formulas

Catalan Numbers

$$C_n = \frac{1}{n+1} {2n \choose n}$$
 $C_0 = 1, C_1 = 1$ and $C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}$

The number of ways to completely parenthesize n+1 factors. The number of triangulations of a convex polygon with n+2sides (i.e. the number of partitions of polygon into disjoint triangles by using the diagonals).

The number of ways to connect the 2n points on a circle to form n disjoint i.e. non-intersecting chords.

The number of rooted full binary trees with n+1 leaves (vertices are not numbered). A rooted binary tree is full if every vertex has either two children or no children.

Number of permutations of $1, \dots, n$ that avoid the pattern 123 **7.4** Other Combinatorial Identities (or any of the other patterns of length 3); that is, the number of permutations with no three-term increasing sub-sequence. For n = 3, these permutations are 132, 213, 231, 312 and 321.

Stirling Numbers First Kind

The Stirling numbers of the first kind count permutations according to their number of cycles (counting fixed points as $n, r \in N, n > r, \sum_{i=1}^{n} {i \choose r} = {n+1 \choose r+1}$ cycles of length one).

S(n,k) counts the number of permutations of n elements with If $P(n) = \sum_{k=0}^{n} \binom{n}{k} \cdot Q(k)$, then,

$$S(n,k) = (n-1) \cdot S(n-1,k) + S(n-1,k-1), \ where, \ S(0,0) = (n-1) \cdot S(n-1,k) + S(n-1,k-1), \ where, \ S(0,0) = (n-1) \cdot S(n-1,k) + S(n-1,k-1), \ where \in S(0,0)$$

$$1, S(n,0) = S(0,n) = 0 \sum_{k=0}^{n} S(n,k) = n!$$

The unsigned Stirling numbers may also be defined algebraically, as the coefficient of the rising factorial:

$$x^{\bar{n}} = x(x+1)...(x+n-1) = \sum_{k=0}^{n} S(n,k)x^{k}$$

Lets [n, k] be the stirling number of the first kind, then

$${n \brack n-k} = \sum_{0 \le i_1 < i_2 < i_k < n} i_1 i_2 i_k.$$

Stirling Numbers Second Kind

Stirling number of the second kind is the number of ways to partition a set of n objects into k non-empty subsets.

$$S(n,k) = k \cdot S(n-1,k) + S(n-1,k-1), \text{ where } S(0,0) = \text{ given by the formula } g_k = \frac{k(3k-1)}{2} \text{ for } k = 1,-1,2,-2,3,\cdots$$
 and are called (generalized) pentagonal numbers. It is useful to find the partition number in $O(n\sqrt{n})$
$$F(n) = \sum_{i=1}^n \sum_{j=1}^n \operatorname{lcm}(i,j) = \sum_{l=1}^n \left(\frac{\left(1 + \lfloor \frac{n}{l} \rfloor\right) \left(\lfloor \frac{n}{l} \rfloor\right)}{2}\right)^2 \sum_{d|l} \mu(d) l d$$
 each color is used at least once.

ber of ways to partition a set of n objects into k subsets, with at and bt such that $aat \equiv 1 \mod b$ and $bbt \equiv 1 \mod a$. Then each subset containing at least r elements. It is denoted by the number of representations of a positive integers (n) as a $S_r(n,k)$ and obeys the recurrence relation. $S_r(n+1,k) =$ non negative linear combination of a and b is

$$kS_r(n,k) + \binom{n}{r-1}S_r(n-r+1,k-1)$$

Denote the n objects to partition by the integers $1, 2, \ldots, n$. Define the reduced Stirling numbers of the second kind, denoted $S^d(n,k)$, to be the number of ways to partition the integers $1, 2, \ldots, n$ into k nonempty subsets such that all elements in each subset have pairwise distance at least d. That is, for any integers i and j in a given subset, it is required that if a_1 and a_2 are relatively prime, then $gcd(a_1 \cdot a_2, b) =$ $|i-j| \geq d$. It has been shown that these numbers satisfy, $S^{d}(n,k) = S(n-d+1,k-d+1), n \ge k \ge d$

$$\binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}$$

$$\sum_{i=0}^{k} \binom{n+i}{i} = \sum_{i=0}^{k} \binom{n+i}{n} = \binom{n+k+1}{k}$$

$$n, r \in N, n > r, \sum_{i=r}^{n} \binom{i}{r} = \binom{n+1}{r+1}$$

$$\text{If } P(n) = \sum_{k=0}^{n} \binom{n}{k} \cdot Q(k), \text{ then,}$$

$$Q(n) = \sum_{k=0}^{n} (-1)^{n-k} \binom{n}{k} \cdot P(k)$$

If
$$P(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot Q(k)$$
, then,
$$Q(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot P(k)$$

Different Math Formulas

$$(1-x)(1-x^2)(1-x^3)\cdots = 1-x-x^2+x^5+x^7-x^{12}-x^{15}+x^{22}+x^{26}-\cdots$$

The exponents $1, 2, 5, 7, 12, \cdots$ on the right hand side are $\sum_{i=1}^{n} \sum_{j=1}^{n} i \cdot j[\gcd(i, j) = 1] = \sum_{i=1}^{n} \phi(i)i^2$ wen by the formula $a_k = \frac{k(3k-1)}{n}$ for k = 1, 2, 3, 2 $S(n,k) = k \cdot S(n-1,k) + S(n-1,k-1), \text{ where } S(0,0) = \text{ given by the formula } g_k = \frac{k(3k-1)}{2} \text{ for } k = 1,-1,2,-2,3,\cdots$

An r-associated Stirling number of the second kind is the num-Let a and b be coprime positive integers, and find integers

$$\frac{n}{ab}$$
 - $\left\{\frac{b'n}{a}\right\}$ - $\left\{\frac{a'n}{b}\right\}$ + 1

7.6 GCD and LCM

if m is any integer, then $gcd(a + m \cdot b, b) = gcd(a, b)$

The gcd is a multiplicative function in the following sense: $\gcd(a_1,b)\cdot\gcd(a_2,b).$

 $\gcd(a, \operatorname{lcm}(b, c)) = \operatorname{lcm}(\gcd(a, b), \gcd(a, c)).$

 $\operatorname{lcm}(a, \gcd(b, c)) = \gcd(\operatorname{lcm}(a, b), \operatorname{lcm}(a, c)).$

For non-negative integers a and b, where a and b are not both zero, $gcd(n^a - 1, n^b - 1) = n^{gcd(a,b)} - 1$

$$\gcd(a,b) = \sum_{k|a \text{ and } k|b} \phi(k)$$

$$\sum_{i=1}^{n} [\gcd(i,n) = k] = \phi\left(\frac{n}{k}\right)$$

$$\sum_{k=1}^{n} \gcd(k,n) = \sum_{d|n} d \cdot \phi\left(\frac{n}{d}\right)$$

$$\sum_{k=1}^{n} x^{\gcd(k,n)} = \sum_{d|n} x^d \cdot \phi\left(\frac{n}{d}\right)$$

$$\sum_{k=1}^{n} \frac{1}{\gcd(k,n)} = \sum_{d|n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{1}{n} \sum_{d|n} d \cdot \phi(d)$$

$$\sum_{k=1}^{n} \frac{k}{\gcd(k,n)} = \frac{n}{2} \cdot \sum_{d|n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{n}{2} \cdot \frac{1}{n} \cdot \sum_{d|n} d \cdot \phi(d)$$

$$\sum_{k=1}^{n} \frac{n}{\gcd(k,n)} = 2 * \sum_{k=1}^{n} \frac{k}{\gcd(k,n)} - 1, \text{ for } n > 1$$

$$\sum_{i=1}^{n} \sum_{j=1}^{n} [\gcd(i,j) = 1] = \sum_{d=1}^{n} \mu(d) \lfloor \frac{n}{d} \rfloor^2$$

$$\sum_{i=1}^{n} \sum_{j=1}^{n} \gcd(i,j) = \sum_{d=1}^{n} \phi(d) \lfloor \frac{n}{d} \rfloor^2$$

$$\sum_{i=1}^{n} \sum_{j=1}^{n} \gcd(i,j) = 1 = \sum_{d=1}^{n} \phi(i)i^2$$