# $IUT\_Shongshop tok, Islamic\ University\ of\ Technology$

Contents				5.7	Dinic Max Flow	14
1	All Macros	1		5.8	Min Cost Max Flow	15
1	All Macros	1		5.9	Block Cut Tree	15
2	DP	1			Bridge Tree	16
_	2.1 1D-1D	1		5.11	Tree Isomorphism	16
	2 Convex Hull Trick		6	Mat	Math	
	2.3 Divide and Conquer dp	2	_	6.1	Combi	16 16
	2.4 Dynamic CHT	2		6.2	Linear Sieve	16
	2.5 FFT Online	3		6.3	Pollard Rho	17
	2.6 Knuth optimization	3		6.4	Chinese Remainder Theorem	17
	2.7 Li Chao Tree	3		6.5	Mobius Function	17
	2.1 Li Chao i i c	3		6.6	FFT	17
3	Data Structure	3		6.7	NTT	18
	3.1 Segment Tree	3		6.8	WalshHadamard	19
	3.2 Spare Table	4		6.9	Berlekamp Massey	19
	3.3 Persistent Segment Tree	$\overline{4}$			Lagrange	20
	3.4 SegTree Beats	5			Shanks' Baby Step, Giant Step	20
	3.5 HashTable	5			Xor Basis	20
	3.6 DSU With Rollbacks	6				
	3.7 Binary Trie	6	7	Stri	ng	20
	3.8 BIT-2D	6		7.1	Aho Corasick	20
	3.9 Divide And Conquer Query Offline	6		7.2	Double hash	21
	3.10 MO with Update	6		7.3	Manacher's	21
	3.11 SparseTable (Rectangle Query)	7		7.4	Suffix Array	21
	5.11 Sparserable (Rectangle Query)	•		7.5	Z Algo	22
4	Geometry			<b>T</b>		0.0
	4.1 Point	7	8	_	ations and Formulas	23
	4.2 Linear	8		8.1	Catalan Numbers	23
	4.3 Polygon	8		8.2 8.3	Stirling Numbers First Kind	$\frac{23}{23}$
	4.4 Convex	9		8.4	Stirling Numbers Second Kind Other Combinatorial Identities	23 23
	4.5 Circular	11		8.5	Different Math Formulas	23 23
	4.6 Half Plane	12		8.6		$\frac{23}{23}$
				0.0	GCD and LCM	20
<b>5</b>	Graph	<b>12</b>				
	5.1 Graph Template	12				
	5.2 Lifting, LCA, HLD	12				
	5.3 SCC	13				
	5.4 Centroid Decompose	13				
	5.5 Euler Tour on Edge	14				
	5.6 Virtual Tree	14				

```
Sublime Build
```

### vimrc

```
set mouse=a
set termguicolors
filetype plugin indent on
syntax on
set smartindent expandtab ignorecase smartcase incsearch
    relativenumber nowrap autoread splitright splitbelow
set tabstop=4
                    "the width of a tab
                    "the width for indent
set shiftwidth=4
colorscheme torte
inoremap {<ENTER> {}<LEFT><CR><ESC><S-o>
inoremap jj <ESC>
autocmd filetype cpp map <F5> :wa<CR>:!clear && g++ % -
    DLOCAL -std=c++20 -Wall -Wextra -Wconversion -
    Wshadow -Wfloat-equal -o ~/Codes/prog && (timeout 5
    ~/Codes/prog < ~/Codes/in) > ~/Codes/out<CR>
map <F4> :!xclip -o -sel c > ~/Codes/in <CR><CR>
map <F3> :!xclip -sel c % <CR><CR>
map <F6> :vsplit ~/Codes/in<CR>:split ~/Codes/out<CR><C-w
    >=20<C-w><<C-w><C-h>
:autocmd BufNewFile *.cpp Or ~/Codes/temp.cpp
set clipboard=unnamedplus
```

#### Stress-tester

```
#!/bin/bash
# Call as stresstester ITERATIONS [--count]
g++ gen.cpp -o gen
g++ sol.cpp -o sol
g++ brute.cpp -o brute

for i in $(seq 1 "$1") ; do
    echo "Attempt $i/$1"
```

```
./gen > in.txt
./sol < in.txt > out1.txt
./brute < in.txt > out2.txt
diff -y out1.txt out2.txt > diff.txt
if [ $? -ne 0 ] ; then
        echo "Differing Testcase Found:"; cat in.txt
        echo -e "\nOutputs:"; cat diff.txt
        break
fi
done
```

# 1 All Macros

/\*--- DEBUG TEMPLATE STARTS HERE ---\*/

void show(long long x) {cerr << x;}</pre>

void show(int x) {cerr << x;}</pre>

```
void show(double x) {cerr << x;}</pre>
void show(char x) {cerr << '\'' << x << '\'';}</pre>
void show(const string &x) {cerr << '\"' << x << '\"';}</pre>
void show(bool x) {cerr << (x ? "true" : "false");}</pre>
template<typename T, typename V>
void show(pair<T, V> x) { cerr << '\footnote{'}; show(x.first);</pre>
    cerr << ", "; show(x.second); cerr << '}'; }</pre>
template<typename T>
void show(T x) {int f = 0; cerr << "{"; for (auto &i: x)</pre>
    cerr << (f++ ? ", " : ""), show(i); cerr << "}";}</pre>
void debug_out(string s) {
s.clear();
 cerr << s << '\n';
template <typename T, typename... V>
void debug_out(string s, T t, V... v) {
 s.erase(remove(s.begin(), s.end(), ''), s.end());
                 "; // 8 spaces
 cerr << "
 cerr << s.substr(0, s.find(','));</pre>
 s = s.substr(s.find(',') + 1);
 cerr << " = ";
 show(t):
 cerr << endl;
 if(sizeof...(v)) debug_out(s, v...);
#define dbg(x...) cerr << "LINE: " << __LINE__ << endl;</pre>
    debug_out(#x, x); cerr << endl;</pre>
/*--- DEBUG TEMPLATE ENDS HERE ---*/
//#pragma GCC optimize("Ofast")
//#pragma GCC optimization ("03")
//#pragma comment(linker, "/stack:200000000")
//#pragma GCC optimize("unroll-loops")
```

```
//#pragma GCC target("sse,sse2,sse3,sse3,sse4,popcnt,abm
    ,mmx,avx,tune=native")
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
//find by order(k) --> returns iterator to the kth
     largest element counting from 0
//order_of_key(val) --> returns the number of items in a
      set that are strictly smaller than our item
template <typename DT>
using ordered_set = tree <DT, null_type, less<DT>,
    rb_tree_tag,tree_order_statistics_node_update>;
mt19937 rnd(chrono::steady_clock::now().time_since_epoch
    ().count());
#ifdef LOCAL
#include "dbg.h"
#else
#define dbg(x...)
#endif
int main() {
cin.tie(0) -> sync_with_stdio(0);
```

# 2 DP

#### $2.1 ext{ } 1D-1D$

```
#include <bits/stdc++.h>
using namespace std;

/// Solves dp[i] = min(dp[j] + cost(j+1, i)) given that
    cost() is QF

long long solve1D(int n, long long cost(int, int)) {
  vector<long long> dp(n + 1), opt(n + 1);
  deque<pair<int, int>> dq;
  dq.push_back({0, 1});
  dp[0] = 0;

for (int i = 1; i <= n; i++) {
    opt[i] = dq.front().first;
    dp[i] = dp[opt[i]] + cost(opt[i] + 1, i);
    if (i == n) break;

dq[0].second++;</pre>
```

```
if (dq.size() > 1 \&\& dq[0].second == dq[1].second) dq.
      pop_front();
  int en = n;
  while (dq.size()) {
  int o = dq.back().first, st = dq.back().second;
  if (dp[o] + cost(o + 1, st) >= dp[i] + cost(i + 1, st) struct CHT {
   dq.pop_back();
  else {
   int lo = st, hi = en;
   while (lo < hi) {
    int mid = (lo + hi + 1) / 2:
    if (dp[o] + cost(o + 1, mid) < dp[i] + cost(i + 1,</pre>
         mid))
     lo = mid:
    else
     hi = mid - 1;
   if (lo < n) dq.push_back({i, lo + 1});</pre>
   break;
  en = st - 1;
 if (dq.empty()) dq.push_back({i, i + 1});
return dp[n];
/// Solves https://open.kattis.com/problems/
    coveredwalkway
const int N = 1e6 + 7;
long long x[N];
int c:
long long cost(int 1, int r) { return (x[r] - x[1]) * (x[ //(slope dec+query max), (slope inc+query min) -> x
    r] - x[1]) + c; }
int main() {
ios::sync_with_stdio(false);
cin.tie(0);
int n;
cin >> n >> c;
for (int i = 1; i \le n; i++) cin >> x[i];
cout << solve1D(n, cost) << endl;</pre>
```

### 2.2 Convex Hull Trick

```
#include <bits/stdc++.h>
using namespace std;
```

```
using LL = long long;
const int N = 3e5 + 9;
const int M = 1e9 + 7;
vector<LL> m, b;
int ptr = 0;
bool bad(int 11, int 12, int 13) {
 return 1.0 * (b[13] - b[11]) * (m[11] - m[12]) <= 1.0 * \
       (b[12] - b[11]) * (m[11] - m[13]); //(slope dec+
      query min), (slope inc+query max)
 return 1.0 * (b[13] - b[11]) * (m[11] - m[12]) > 1.0 *
      (b[12] - b[11]) * (m[11] - m[13]); //(slope dec+
      query max), (slope inc+query min)
}
 void add(LL _m, LL _b) {
 m.push_back(_m);
 b.push_back(_b);
  int s = m.size();
  while (s >= 3 \&\& bad(s - 3, s - 2, s - 1)) {
  m.erase(m.end() - 2);
  b.erase(b.end() - 2);
LL f(int i, LL x) { return m[i] * x + b[i]; }
//(slope dec+query min), (slope inc+query max) -> x
     increasing
     decreasing
 LL query(LL x) {
 if (ptr >= m.size()) ptr = m.size() - 1;
  while (ptr < m.size() - 1 && f(ptr + 1, x) < f(ptr, x))
       ptr++;
 return f(ptr, x);
LL bs(int 1, int r, LL x) {
 int mid = (1 + r) / 2:
 if (mid + 1 < m.size() && f(mid + 1, x) < f(mid, x))</pre>
      return bs(mid + 1, r, x); // > for max
 if (mid - 1 \ge 0 \&\& f(mid - 1, x) < f(mid, x)) return
      bs(1, mid - 1, x); // > for max
```

```
return f(mid, x);
```

### 2.3 Divide and Conquer dp

```
const int K = 805, N = 4005;
LL dp[2][N], cost[N][N];
// 1-indexed for convenience
LL cost(int 1, int r) {
return _cost[r][r] - _cost[l - 1][r] - _cost[r][l - 1] +
      _cost[1 - 1][1 - 1] >> 1;
void compute(int cnt, int 1, int r, int opt1, int optr) {
if (1 > r) return;
int mid = 1 + r \gg 1;
LL best = INT MAX:
int opt = -1;
for (int i = optl; i <= min(mid, optr); i++) {</pre>
 LL cur = dp[cnt ^1][i - 1] + cost(i, mid);
 if (cur < best) best = cur, opt = i;</pre>
dp[cnt][mid] = best;
compute(cnt, 1, mid - 1, optl, opt);
compute(cnt, mid + 1, r, opt, optr);
LL dnc_dp(int k, int n) {
fill(dp[0] + 1, dp[0] + n + 1, INT_MAX);
for (int cnt = 1; cnt <= k; cnt++) {</pre>
 compute(cnt & 1, 1, n, 1, n);
return dp[k & 1][n];
```

#### 2.4 Dynamic CHT

```
typedef long long LL;
const LL IS_QUERY = -(1LL << 62);</pre>
struct line {
LL m, b;
mutable function <const line*()> succ;
 bool operator < (const line &rhs) const {</pre>
 if (rhs.b != IS QUERY) return m < rhs.m;</pre>
 const line *s = succ();
 if (!s) return 0;
 LL x = rhs.m:
 return b - s -> b < (s -> m - m) * x;
}
```

```
struct HullDynamic : public multiset <line> {
bool bad (iterator y) {
  auto z = next(y);
  if (y == begin()) {
  if (z == end()) return 0;
  return y -> m == z -> m && y -> b <= z -> b;
  auto x = prev(y);
  if (z == end()) return y \rightarrow m == x \rightarrow m && y \rightarrow b <= x
       -> b;
 return 1.0 * (x \rightarrow b - y \rightarrow b) * (z \rightarrow m - y \rightarrow m) >=
      1.0 * (y \rightarrow b - z \rightarrow b) * (y \rightarrow m - x \rightarrow m);
 void insert_line (LL m, LL b) {
  auto y = insert({m, b});
  y \rightarrow succ = [=] \{return next(y) == end() ? 0 : \&*next(y) \}
  if (bad(y)) {erase(y); return;}
  while (next(y) != end() && bad(next(y))) erase(next(y))
  while (y != begin() && bad(prev(y))) erase(prev(y));
LL eval (LL x) {
  auto 1 = *lower_bound((line) {x, IS_QUERY});
 return 1.m * x + 1.b:
};
```

#### 2.5 FFT Online

```
void fftOnline(vector <LL> &a, vector <LL> b) {
  int n = a.size();
  auto call = [&](int l, int r, auto &call){
    if(l >= r) return;
    int mid = l + r >> 1;
    call(l, mid, call);

  vector <LL> _a(a.begin() + l, a.begin() + mid + 1);
  vector <LL> _b(b.begin(), b.begin() + (r - l + 1));
  auto c = fft :: anyMod(_a, _b);

  for(int i = mid + 1; i <= r; i++) {
    a[i] += c[i - l];
    a[i] -= (a[i] >= mod) * mod;
  }
  call(mid + 1, r, call);
};
  call(0, n - 1, call);
```

# 2.6 Knuth optimization

```
const int N = 1005;
LL dp[N][N], a[N];
int opt[N][N];
LL cost(int i, int j) { return a[j + 1] - a[i]; }
LL knuth_optimization(int n) {
for (int i = 0; i < n; i++) {</pre>
 dp[i][i] = 0;
 opt[i][i] = i;
for (int i = n - 2; i \ge 0; i--) {
 for (int j = i + 1; j < n; j++) {
  LL mn = LLONG MAX;
  LL c = cost(i, j);
  for (int k = opt[i][j-1]; k \le min(j-1, opt[i+1])
       1][j]); k++) {
   if (mn > dp[i][k] + dp[k + 1][j] + c) {
    mn = dp[i][k] + dp[k + 1][j] + c;
    opt[i][j] = k;
   }
  dp[i][j] = mn;
return dp[0][n - 1];
```

#### 2.7 Li Chao Tree

```
struct line {
LL m. c:
line(LL m = 0, LL c = 0) : m(m), c(c) {}
LL calc(line L, LL x) { return 1LL * L.m * x + L.c; }
struct node {
LL m, c;
line L;
node *lft, *rt;
node(LL m = 0, LL c = 0, node *lft = NULL, node *rt =
     NULL)
  : L(line(m, c)), lft(lft), rt(rt) {}
struct LiChao {
node *root:
LiChao() { root = new node(); }
void update(node *now, int L, int R, line newline) {
 int mid = L + (R - L) / 2;
 line lo = now->L, hi = newline;
 if (calc(lo, L) > calc(hi, L)) swap(lo, hi);
```

```
if (calc(lo, R) <= calc(hi, R)) {</pre>
 now->L = hi;
 return:
if (calc(lo, mid) < calc(hi, mid)) {</pre>
 now->L = hi;
 if (now->rt == NULL) now->rt = new node();
 update(now->rt, mid + 1, R, lo);
} else {
 now->L = lo:
 if (now->lft == NULL) now->lft = new node();
 update(now->lft, L, mid, hi);
LL query(node *now, int L, int R, LL x) {
if (now == NULL) return -inf;
int mid = L + (R - L) / 2;
if (x \le mid)
 return max(calc(now->L, x), query(now->lft, L, mid, x)
else
 return max(calc(now->L, x), query(now->rt, mid + 1, R,
```

### 3 Data Structure

### 3.1 Segment Tree

```
template <typename VT>
struct segtree {
using DT = typename VT::DT;
using LT = typename VT::LT;
int L, R;
vector <VT> tr;
segtree(int n): L(0), R(n-1), tr(n << 2) {}
segtree(int 1, int r): L(1), R(r), tr((r - 1 + 1) << 2)
    {}
void propagate(int 1, int r, int u) {
 if(l == r) return;
 VT :: apply(tr[u << 1], tr[u].lz, 1, (1 + r) >> 1);
 VT :: apply(tr[u << 1 | 1], tr[u].lz, (1 + r + 2) >> 1,
      r):
tr[u].lz = VT :: None;
void build(int 1, int r, vector <DT> &v, int u = 1 ) {
if(1 == r) {
```

```
tr[u].val = v[1];
  return;
 int m = (1 + r) >> 1, lft = u << 1, ryt = u << 1 | 1;
 build(1, m, v, lft);
 build(m + 1, r, v, ryt);
 tr[u].val = VT :: merge(tr[lft].val, tr[ryt].val, 1, r)
void update(int ql,int qr, LT up, int l, int r, int u =
    1) {
 if(ql > qr) return;
  if(ql == l and qr == r) {
  VT :: apply(tr[u], up, 1, r);
  return;
 }
 propagate(1, r, u);
 int m = (1 + r) >> 1, lft = u << 1, ryt = u << 1 | 1;
  update(ql, min(m, qr), up, l, m, lft);
  update(max(ql, m + 1), qr, up, m + 1, r, ryt);
 tr[u].val = VT :: merge(tr[lft].val, tr[ryt].val, 1, r)
}
DT query(int ql, int qr, int l, int r, int u = 1) {
 if(ql > qr) return VT::I;
 if(1 == ql and r == qr)
  return tr[u].val;
  propagate(1, r, u);
 int m = (1 + r) >> 1, lft = u << 1, ryt = u << 1 | 1;
 DT ansl = query(ql, min(m, qr), l, m, lft);
 DT ansr = query(max(ql, m + 1), qr, m + 1, r, ryt);
 return tr[u].merge(ansl, ansr, l, r);
void build(vector <DT> &v) { build(L, R, v); }
void update(int ql, int qr, LT U) { update(ql, qr, U, L,
DT query(int ql, int qr) { return query(ql, qr, L, R); } }.
};
struct add_sum {
using DT = LL;
using LT = LL;
DT val:
LT lz;
static constexpr DT I = 0;
```

```
static constexpr LT None = 0;
add_sum(DT _val = I, LT _lz = None): val(_val), lz(_lz)
    {}
static void apply(add_sum &u, const LT &up, int 1, int r|}
u.val += (r - l + 1) * up;
u.lz += up;
static DT merge(const DT &a, const DT &b, int 1, int r)
return a + b;
```

# 3.2 Spare Table

```
template <typename T> struct sparse_table {
vector <vector<T>> tbl;
function < T(T, T) > f:
sparse_table(const vector <T> &v, function <T(T, T)> _f,
      T _id) : f(_f), id(_id) {
 int n = (int) v.size(), b = __lg(n);
 tbl.assign(b + 1, v);
 for(int k = 1; k <= b; k++) {</pre>
  for(int i = 0; i + (1 << k) <= n; i++) {</pre>
   tbl[k][i] = f(tbl[k-1][i], tbl[k-1][i+(1 << (k)
       - 1))]);
  }
 }
T query(int 1, int r) {
 if(1 > r) return id;
 int pow = __lg(r - l + 1);
 return f(tbl[pow][1], tbl[pow][r - (1 << pow) + 1]);</pre>
```

# 3.3 Persistent Segment Tree

```
struct Node {
int 1 = 0, r = 0, val = 0;
tr[20 * N];
int ptr = 0;
int build(int st, int en) {
int u = ++ptr;
if (st == en) return u;
int mid = (st + en) / 2;
```

```
auto& [1, r, val] = tr[u];
1 = build(st, mid);
r = build(mid + 1, en);
val = tr[1].val + tr[r].val;
return u:
int update(int pre, int st, int en, int idx, int v) {
int u = ++ptr;
tr[u] = tr[pre];
if (st == en) {
 tr[u].val += v;
 return u:
int mid = (st + en) / 2;
auto& [1, r, val] = tr[u];
if (idx <= mid) {</pre>
 r = tr[pre].r;
 1 = update(tr[pre].1, st, mid, idx, v);
} else {
 1 = tr[pre].1;
 r = update(tr[pre].r, mid + 1, en, idx, v);
tr[u].val = tr[1].val + tr[r].val;
return u:
// finding the kth elelment in a range
int query(int left, int right, int st, int en, int k) {
if (st == en) return st;
int cnt = tr[tr[right].1].val - tr[tr[left].1].val;
int mid = (st + en) / 2;
if (cnt >= k) return query(tr[left].1, tr[right].1, st,
     mid. k):
else return query(tr[left].r, tr[right].r, mid + 1, en,
     k - cnt);
int V[N], root[N], a[N];
int main() {
map<int, int> mp; int n, q;
cin >> n >> q;
for (int i = 1; i <= n; i++) cin >> a[i], mp[a[i]];
for (auto x : mp) mp[x.first] = ++c, V[c] = x.first;
root[0] = build(1, n);
for (int i = 1; i <= n; i++) {</pre>
 root[i] = update(root[i - 1], 1, n, mp[a[i]], 1);
while (q--) {
 int 1, r, k; cin >> 1 >> r >> k; 1++, k++;
```

```
cout << V[query(root[1 - 1], root[r], 1, n, k)] << '\n'
;
}
return 0;
}</pre>
```

# 3.4 SegTree Beats

```
const int N = 2e5 + 5:
LL mx[4 * N], mn[4 * N], smx[4 * N], smn[4 * N], sum[4 * N]
    N], add[4 * N];
int mxcnt[4 * N]. mncnt[4 * N]:
int L, R;
void applyMax(int u, LL x) {
sum[u] += mncnt[u] * (x - mn[u]):
if (mx[u] == mn[u]) mx[u] = x;
if (smx[u] == mn[u]) smx[u] = x;
mn[u] = x:
void applyMin(int u, LL x) {
sum[u] -= mxcnt[u] * (mx[u] - x):
if (mn[u] == mx[u]) mn[u] = x;
if (smn[u] == mx[u]) smn[u] = x;
mx[u] = x;
void applyAdd(int u, LL x, int tl, int tr) {
sum[u] += (tr - tl + 1) * x:
add[u] += x;
mx[u] += x, mn[u] += x;
if (smx[u] != -INF) smx[u] += x;
if (smn[u] != INF) smn[u] += x:
void push(int u, int tl, int tr) {
int lft = u << 1, ryt = lft | 1, mid = tl + tr >> 1;
if (add[u] != 0) {
 applyAdd(lft, add[u], tl, mid);
 applyAdd(ryt, add[u], mid + 1, tr);
 add[u] = 0;
if (mx[u] < mx[lft]) applyMin(lft, mx[u]);</pre>
if (mx[u] < mx[ryt]) applyMin(ryt, mx[u]);</pre>
if (mn[u] > mn[lft]) applyMax(lft, mn[u]);
if (mn[u] > mn[ryt]) applyMax(ryt, mn[u]);
void merge(int u) {
int lft = u << 1, ryt = lft | 1;</pre>
sum[u] = sum[lft] + sum[ryt];
```

```
mx[u] = max(mx[lft], mx[rvt]);
smx[u] = max(smx[lft], smx[ryt]);
if (mx[lft] != mx[ryt]) smx[u] = max(smx[u], min(mx[lft
     ], mx[rvt]));
mxcnt[u] = (mx[u] == mx[lft]) * mxcnt[lft] + (mx[u] ==
     mx[ryt]) * mxcnt[ryt];
mn[u] = min(mn[lft], mn[ryt]);
smn[u] = min(smn[lft], smn[ryt]);
 if (mn[lft] != mn[ryt]) smn[u] = min(smn[u], max(mn[lft
     ], mn[rvt]));
mncnt[u] = (mn[u] == mn[lft]) * mncnt[lft] + (mn[u] ==
     mn[rvt]) * mncnt[rvt];
void minimize(int 1, int r, LL x, int tl = L, int tr = R,
     int u = 1) {
if (1 > tr or tl > r or mx[u] <= x) return;</pre>
if (1 \le t1 \text{ and } tr \le r \text{ and } smx[u] \le x) 
 applyMin(u, x);
 return:
push(u, tl, tr);
int mid = tl + tr >> 1, lft = u << 1, ryt = lft | 1;</pre>
minimize(l, r, x, tl, mid, lft);
minimize(l, r, x, mid + 1, tr, ryt);
merge(u):
void maximize(int 1, int r, LL x, int t1 = L, int tr = R.
     int u = 1) {
if (1 > tr or tl > r or mn[u] >= x) return;
if (1 \le t1 \text{ and } tr \le r \text{ and } smn[u] > x) {
 applyMax(u, x);
 return:
push(u, tl, tr);
int mid = tl + tr >> 1, lft = u << 1, ryt = lft | 1;</pre>
maximize(l, r, x, tl, mid, lft);
maximize(1, r, x, mid + 1, tr, ryt);
merge(u);
void increase(int 1, int r, LL x, int tl = L, int tr = R.
     int u = 1) {
if (1 > tr or tl > r) return;
if (1 <= t1 and tr <= r) {</pre>
 applyAdd(u, x, tl, tr);
 return:
push(u, tl, tr);
int mid = tl + tr >> 1, lft = u << 1, ryt = lft | 1;</pre>
increase(l, r, x, tl, mid, lft);
```

```
increase(l, r, x, mid + 1, tr, ryt);
merge(u);
LL getSum(int 1, int r, int tl = L, int tr = R, int u =
   1) {
if (1 > tr or tl > r) return 0;
if (1 <= t1 and tr <= r) return sum[u]:</pre>
push(u, tl, tr);
int mid = tl + tr >> 1, lft = u << 1, ryt = lft | 1;</pre>
return getSum(1, r, tl, mid, lft) + getSum(1, r, mid +
    1, tr, rvt);
void build(LL a\Pi, int tl = L, int tr = R, int u = 1) {
if (t1 == tr) {
 sum[u] = mn[u] = mx[u] = a[t1];
 mxcnt[u] = mncnt[u] = 1;
 smx[u] = -INF;
 smn[u] = INF:
 return;
int mid = tl + tr >> 1, lft = u << 1, ryt = lft | 1;</pre>
build(a, tl, mid, lft);
build(a, mid + 1, tr, ryt);
merge(u);
void init(LL a[], int _L, int _R) {
L = L, R = R;
build(a):
```

#### 3.5 HashTable

```
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;

const int RANDOM = chrono::high_resolution_clock::now().
    time_since_epoch().count();
unsigned hash_f(unsigned x) {
    x = ((x >> 16) ^ x) * 0x45d9f3b;
    x = ((x >> 16) ^ x) * 0x45d9f3b;
    return x = (x >> 16) ^ x;
}

unsigned hash_combine(unsigned a, unsigned b) { return a
        * 31 + b; }
struct chash {
    int operator()(int x) const { return hash_f(x); }
};
typedef gp_hash_table<int, int, chash> gp;
gp table;
```

### 3.6 DSU With Rollbacks

```
struct Rollback DSU {
int n;
vector<int> par, sz;
vector<pair<int, int>> op;
Rollback_DSU(int n) : par(n), sz(n, 1) {
 iota(par.begin(), par.end(), 0);
 op.reserve(n):
int Anc(int node) {
 for (; node != par[node]; node = par[node])
  ; // no path compression
 return node;
void Unite(int x, int y) {
 if (sz[x = Anc(x)] < sz[y = Anc(y)]) swap(x, y);
  op.emplace_back(x, y);
 par[y] = x;
 sz[x] += sz[y];
void Undo(int t) {
 for (; op.size() > t; op.pop_back()) {
  par[op.back().second] = op.back().second;
  sz[op.back().first] -= sz[op.back().second];
};
```

# 3.7 Binary Trie

```
const int N = 1e7 + 5, b = 30;
int tc = 1:
struct node {
 int vis = 0:
 int to [2] = \{0, 0\};
 int val[2] = {0, 0};
 void update() {
  to[0] = to[1] = 0;
 val[0] = val[1] = 0;
  vis = tc:
T[N + 2];
node *root = T;
int ptr = 0;
node *nxt(node *cur, int x) {
if (cur->to[x] == 0) cur->to[x] = ++ptr;
 assert(ptr < N);</pre>
 int idx = cur->to[x];
 if (T[idx].vis < tc) T[idx].update();</pre>
 return T + idx;
```

```
int query(int j, int aj) {
int ans = 0, jaj = j ^ aj;
node *cur = root;
for (int k = b - 1; ~k; k--) {
 maximize(ans, nxt(cur, (jaj >> k & 1) ^ 1)->val[1 ^ (aj |void push(int _node) { node[++tm] = _node; }
      >> k & 1)]):
 cur = nxt(cur, (jaj >> k & 1));
return ans;
void insert(int j, int aj, int val) {
int jaj = j ^ aj;
node *cur = root:
for (int k = b - 1; ~k; k--) {
 cur = nxt(cur, (jaj >> k & 1));
 maximize(cur->val[j >> k & 1], val);
void clear() {
tc++;
ptr = 0;
root->update();
```

### 3.8 BIT-2D

```
const int N = 1008;
int bit[N][N], n, m;
int a[N][N], q;
void update(int x, int y, int val) {
  for (; x < N; x += -x & x)
    for (int j = y; j < N; j += -j & j) bit[x][j] += val;
}
int get(int x, int y) {
  int ans = 0;
  for (; x; x -= x & -x)
    for (int j = y; j; j -= j & -j) ans += bit[x][j];
  return ans;
}
int get(int x1, int y1, int x2, int y2) {
  return get(x2, y2) - get(x1 - 1, y2) - get(x2, y1 - 1) +
        get(x1 - 1, y1 - 1);
}</pre>
```

```
} // namespace up
namespace que {
int node[N], tm;
LL ans[N]:
} // namespace que
namespace edge_set {
void push(int i) { dsu ::merge(up ::u[i], up ::v[i]); }
void pop(int t) { dsu ::rollback(t); }
int time() { return dsu ::op.size(); }
LL query(int u) { return a[dsu ::root(u)]; }
} // namespace edge_set
namespace dncq {
vector<int> tree[4 * N];
void update(int idx, int 1 = 0, int r = que ::tm, int
    node = 1) {
int ul = up ::1[idx], ur = up ::r[idx];
if (r 
if (ul \le l and r \le ur) {
 tree[node].push_back(idx);
 return:
}
int m = 1 + r >> 1;
update(idx, 1, m, node << 1);
update(idx, m + 1, r, node << 1 | 1);
void dfs(int 1 = 0, int r = que ::tm, int node = 1) {
int cur = edge_set ::time();
for (int e : tree[node]) edge set ::push(e);
if (1 == r) {
 que ::ans[1] = edge set ::query(que ::node[1]);
} else {
 int m = 1 + r >> 1:
 dfs(1, m, node << 1);
 dfs(m + 1, r, node << 1 | 1);
edge_set ::pop(cur);
} // namespace dncg
void push_updates() {
for (int i = 0; i < up ::tm; i++) dncq ::update(i);</pre>
```

# 3.9 Divide And Conquer Query Offline

```
namespace up {
int 1[N], r[N], u[N], v[N], tm;
void push(int _1, int _r, int _u, int _v) {
 1[tm] = _1, r[tm] = _r, u[tm] = _u, v[tm] = _v;
 tm++;
```

# 3.10 MO with Update

```
const int N = 1e5 + 5, sz = 2700, bs = 25;
int arr[N], freq[2 * N], cnt[2 * N], id[N], ans[N];
struct query {
  int l, r, t, L, R;
  query(int l = 1, int r = 0, int t = 1, int id = -1)
```

```
: l(1), r(r), t(t), L(1 / sz), R(r / sz) {}
  bool operator<(const query &rhs) const {</pre>
    return (L < rhs.L) or (L == rhs.L and R < rhs.R) or</pre>
          (L == rhs.L and R == rhs.R and t < rhs.t);
  }
} Q[N];
struct update {
  int idx, val, last;
; [N]qU {
int qi = 0, ui = 0;
int 1 = 1, r = 0, t = 0;
void add(int idx) {
  --cnt[freq[arr[idx]]];
  freq[arr[idx]]++;
  cnt[freq[arr[idx]]]++;
void remove(int idx) {
  --cnt[freq[arr[idx]]];
  freq[arr[idx]]--;
  cnt[freq[arr[idx]]]++;
void apply(int t) {
  const bool f = 1 <= Up[t].idx and Up[t].idx <= r;</pre>
  if (f) remove(Up[t].idx);
  arr[Up[t].idx] = Up[t].val;
  if (f) add(Up[t].idx);
void undo(int t) {
  const bool f = 1 <= Up[t].idx and Up[t].idx <= r;</pre>
  if (f) remove(Up[t].idx);
  arr[Up[t].idx] = Up[t].last;
  if (f) add(Up[t].idx);
}
int mex() {
  for (int i = 1; i <= N; i++)</pre>
    if (!cnt[i]) return i;
  assert(0);
int main() {
  sort(id + 1, id + qi + 1, [&](int x, int y) { return Q[
      x] < Q[y]; \});
  for (int i = 1; i <= qi; i++) {</pre>
    int x = id[i]:
    while (Q[x].t > t) apply(++t);
    while (Q[x].t < t) undo(t--);
    while (Q[x].1 < 1) add(--1);
    while (Q[x].r > r) add(++r);
    while (Q[x].1 > 1) remove(1++);
    while (Q[x].r < r) remove(r--);
```

```
ans[x] = mex();
}
```

# 3.11 SparseTable (Rectangle Query)

```
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 505;
const int LOGN = 9;
// O(n^2 (logn)^2
// Supports Rectangular Query
int A[MAXN][MAXN];
int M[MAXN] [MAXN] [LOGN] [LOGN];
void Build2DSparse(int N) {
for (int i = 1; i <= N; i++) {</pre>
 for (int j = 1; j <= N; j++) {</pre>
  M[i][j][0][0] = A[i][j];
 }
  for (int q = 1; (1 << q) <= N; q++) {
  int add = 1 << (q - 1);
  for (int j = 1; j + add <= N; j++) {</pre>
   M[i][j][0][q] = max(M[i][j][0][q - 1], M[i][j + add
       ][0][q - 1]);
  }
 }
 for (int p = 1; (1 << p) <= N; p++) {
 int add = 1 << (p - 1);
 for (int i = 1; i + add <= N; i++) {</pre>
  for (int q = 0; (1 << q) <= N; q++) {</pre>
   for (int j = 1; j <= N; j++) {
    M[i][j][p][q] = max(M[i][j][p - 1][q], M[i + add][j]
         ][p - 1][q]);
   }
// returns max of all A[i][j], where x1<=i<=x2 and y1<=j</pre>
    <=y2
int Query(int x1, int y1, int x2, int y2) {
int kX = log2(x2 - x1 + 1);
int kY = log2(y2 - y1 + 1);
 int addX = 1 << kX;
int addY = 1 << kY;
```

# 4 Geometry

#### 4.1 Point

```
typedef double Tf;
typedef Tf Ti; /// use long long for exactness
const Tf PI = acos(-1), EPS = 1e-9;
int dcmp(Tf x) \{ return abs(x) < EPS ? 0 : (x<0 ? -1 : 1) \}
    ;}
struct Pt {
Ti x, y;
Pt(Ti x = 0, Ti y = 0) : x(x), y(y) {}
 Pt operator + (const Pt& u) const { return Pt(x + u.x, y
      + u.y); }
 Pt operator - (const Pt& u) const { return Pt(x - u.x, y
      - u.y); }
 Pt operator * (const long long u) const { return Pt(x *
     u, y * u); }
 Pt operator * (const Tf u) const { return Pt(x * u, y *
     u): }
 Pt operator / (const Tf u) const { return Pt(x / u, y /
     u); }
 bool operator == (const Pt& u) const { return dcmp(x - u
     .x) == 0 && dcmp(y - u.y) == 0; }
bool operator != (const Pt& u) const { return !(*this ==
      u); }
 bool operator < (const Pt& u) const { return dcmp(x - u.
     x) < 0 \mid | (dcmp(x - u.x) == 0 && dcmp(y - u.y) < 0)
     ; }
 friend istream &operator >> (istream &is, Pt &p) {
     return is >> p.x >> p.y; }
 friend ostream &operator << (ostream &os, const Pt &p) {</pre>
      return os << p.x << " " << p.y; }
using vec_p = vector<Pt>;
Ti dot(Pt a, Pt b) { return a.x * b.x + a.y * b.y; }
Ti crs(Pt a, Pt b) { return a.x * b.y - a.y * b.x; }
Tf len(Pt a) { return sqrt(dot(a, a)); }
Ti sqlen(Pt a) { return dot(a, a); }
Tf dis(Pt a, Pt b) {return len(a-b);}
```

```
Tf angle(Pt u) { return atan2(u.v, u.x); }
// returns angle between oa, ob in (-PI, PI]
Tf angleBetween(Pt a, Pt b) {
double ans = angle(b) - angle(a);
return ans <= -PI ? ans + 2*PI : (ans > PI ? ans - 2*PI
     : ans):
// Rotate a ccw by rad radians
Pt rotate(Pt a, Tf rad) {
return Pt(a.x * cos(rad) - a.y * sin(rad), a.x * sin(rad)
     ) + a.v * cos(rad);
// rotate a ccw by angle th with cos(th) = co && sin(th) |};
Pt rotatePrecise(Pt a, Tf co, Tf si) {
return Pt(a.x * co - a.y * si, a.y * co + a.x * si);
Pt rotate90(Pt a) { return Pt(-a.y, a.x); }
// scales vector a by s such that len of a becomes s
Pt scale(Pt a. Tf s) {
return a / len(a) * s:
// returns an unit vector perpendicular to vector a
Pt normal(Pt a) {
Tf l = len(a):
return Pt(-a.y / 1, a.x / 1);
int ornt(Pt a, Pt b, Pt c) {
return dcmp(crs(b - a, c - a));
bool half(Pt p){      // returns true for pt above x axis
    or on negative x axis
return p.y > 0 || (p.y == 0 \&\& p.x < 0);
bool polarComp(Pt p, Pt q){ //to be used in sort()
    function
return make_tuple(half(p), 0) < make_tuple(half(q), crs(</pre>
     p, q));
struct Seg {
Pt a. b:
Seg(Pt aa, Pt bb) : a(aa), b(bb) {}
typedef Seg Line;
struct Crc {
 Pt o:
 Tf r:
 Crc(Pt \ o = Pt(0, \ 0), \ Tf \ r = 0) : o(o), \ r(r) \ \{\}
 // returns true if pt p is in || on the crc
 bool contains(Pt p) {
```

```
return dcmp(sqlen(p - o) - r * r) <= 0;</pre>
// returns a pt on the crc rad radians away from +X CCW }
Pt pt(Tf rad) {
return Pt(o.x + cos(rad) * r, o.y + sin(rad) * r);
// area of a circular sector with central angle rad
Tf area(Tf rad = PI + PI) { return rad * r * r / 2; }
// area of the circular sector cut by a chord with
    central angle alpha
Tf sector(Tf alpha) { return r * r * 0.5 * (alpha - sin | Pt v1 = s.b - s.a, v2 = p - s.a, v3 = p - s.b;
    (alpha)); }
```

#### 4.2 Linear

```
// returns true if pt p is on segs s
bool onSeg(Pt p, Seg s) {
return dcmp(crs(s.a - p, s.b - p)) == 0 && dcmp(dot(s.a
     - p, s.b - p)) <= 0;
// returns true if segs p && q touch or intersect
bool segssIntersect(Seg p, Seg q) {
if(onSeg(p.a, q) || onSeg(p.b, q)) return true;
if(onSeg(q.a, p) || onSeg(q.b, p)) return true;
Ti c1 = crs(p.b - p.a, q.a - p.a);
Ti c2 = crs(p.b - p.a, q.b - p.a);
Ti c3 = crs(q.b - q.a, p.a - q.a);
Ti c4 = crs(q.b - q.a, p.b - q.a);
return dcmp(c1) * dcmp(c2) < 0 && dcmp(c3) * dcmp(c4) <
bool linesParallel(Line p, Line q) {
return dcmp(crs(p.b - p.a, q.b - q.a)) == 0;
// lines are represented as a ray from a pt: (pt, vector) for(int i = 0; i < (int) p.size() - 1; i++)
// returns false if two lines (p, v) && (q, w) are
    parallel or collinear
// true otherwise, intersection pt is stored at o via
    reference
bool lineLineIntscn(Pt p, Pt v, Pt q, Pt w, Pt& o) {
if(dcmp(crs(v, w)) == 0) return false;
Pt u = p - q;
o = p + v * (crs(w,u)/crs(v,w));
return true;
// returns false if two lines p && q are parallel or
    collinear
// true otherwise, intersection pt is stored at o via
    reference
bool lineLineIntscn(Line p, Line q, Pt& o) {
```

```
return lineLineIntscn(p.a, p.b - p.a, q.a, q.b - q.a, o)
// returns the dis from pt a to line 1
Tf disPtLine(Pt p, Line 1) {
return abs(crs(1.b - 1.a, p - 1.a) / len(1.b - 1.a));
// returns the shortest dis from pt a to segs s
Tf disPtSeg(Pt p, Seg s) {
if(s.a == s.b) return len(p - s.a);
if(dcmp(dot(v1, v2)) < 0) return len(v2);</pre>
else if(dcmp(dot(v1, v3)) > 0) return len(v3);
else return abs(crs(v1, v2) / len(v1));
// returns the shortest dis from segs p to segs q
Tf disSegSeg(Seg p, Seg q) {
if(segssIntersect(p, q)) return 0;
Tf ans = disPtSeg(p.a, q);
ans = min(ans, disPtSeg(p.b, q));
ans = min(ans, disPtSeg(q.a, p));
ans = min(ans, disPtSeg(q.b, p));
return ans:
// returns the projection of pt p on line l
Pt projectPtLine(Pt p, Line 1) {
Pt v = 1.b - 1.a;
return 1.a + v * ((Tf) dot(v, p - 1.a) / dot(v, v));
```

#### 4.3 Polygon

```
using Poly = vector<Pt>;
Tf signedPolyArea(Poly p) {
If ret = 0:
 ret += crs(p[i]-p[0], p[i+1]-p[0]);
return ret / 2:
// given a polygon p of n vertices, generates the convex
// in CCW && returns the number of vertices in the convex
int convexHull(Poly p, Poly &ch) {
sort(p.begin(), p.end());
int n = p.size();
ch.resize(n + n);
int m = 0; // preparing lower hull
for(int i = 0; i < n; i++) {</pre>
 while(m > 1 && dcmp(crs(ch[m - 1] - ch[m - 2], p[i] -
      ch[m - 1])) \le 0) m--;
```

```
ch[m++] = p[i];
int k = m; // preparing upper hull
for(int i = n - 2; i \ge 0; i--) {
 while (m > k \&\& dcmp(crs(ch[m - 1] - ch[m - 2], p[i] -
      ch[m - 2])) <= 0) m--;
 ch[m++] = p[i]:
if(n > 1) m--;
ch.resize(m):
return m;
// for a pt o and polygon p returns:
// -1 if o is strictly inside p
// 0 if o is on a segs of p
// 1 if o is strictly outside p
// computes via winding numbers
int ptInPoly(Pt o, Poly p) {
using Linear::onSeg;
int wn = 0, n = p.size();
for(int i = 0; i < n; i++) {</pre>
 int j = (i + 1) \% n;
 if(onSeg(o, Seg(p[i], p[j])) || o == p[i]) return 0;
 int k = dcmp(crs(p[j] - p[i], o - p[i]));
 int d1 = dcmp(p[i].y - o.y);
 int d2 = dcmp(p[j].y - o.y);
 if(k > 0 \&\& d1 \le 0 \&\& d2 > 0) wn++;
 if(k < 0 \&\& d2 <= 0 \&\& d1 > 0) wn--:
return wn ? -1 : 1;
// returns the longest line segs of 1 that is inside or
// simply polygon p. O(n lg n). TESTED: TIMUS 1955
Tf longestSegInPoly(Line 1, const Poly &p) {
using Linear::lineLineIntscn;
int n = p.size();
vector<pair<Tf, int>> ev;
for(int i=0; i<n; ++i) {</pre>
 Pt a = p[i], b = p[(i + 1) \% n], z = p[(i - 1 + n) \% n]
      ];
  int ora = ornt(1.a, 1.b, a), orb = ornt(1.a, 1.b, b),
      orz = ornt(1.a, 1.b, z);
  if(!ora) {
  Tf d = dot(a - 1.a, 1.b - 1.a);
  if(orz && orb) {
   if(orz != orb) ev.emplace_back(d, 0);
  else if(orz) ev.emplace_back(d, orz);
  else if(orb) ev.emplace_back(d, orb);
```

```
else if(ora == -orb) {
 Pt ins:
 lineLineIntscn(l, Line(a, b), ins);
 ev.emplace_back(dot(ins - 1.a, 1.b - 1.a), 0);
}
sort(ev.begin(), ev.end());
Tf ret = 0, cur = 0, pre = 0;
bool active = false;
int sign = 0;
for(auto &qq : ev) {
int tp = qq.second;
Tf d = qq.first;
 if(sign) {
 cur += d - pre;
 ret = max(ret, cur);
 if(tp != sign) active = !active;
 sign = 0;
 else {
 if(active) cur += d - pre, ret = max(ret, cur);
 if(tp == 0) active = !active;
 else sign = tp;
pre = d;
if(!active) cur = 0;
ret /= len(l.b - l.a);
return ret;
```

#### 4.4 Convex

```
/// minkowski sum of two polygons in O(n)
Poly minkowskiSum(Poly A, Poly B) {
int n = A.size(), m = B.size();
rotate(A.begin(), min_element(A.begin(), A.end()), A.end
     ());
rotate(B.begin(), min_element(B.begin(), B.end()), B.end
     ()):
A.push_back(A[0]);
B.push_back(B[0]);
for (int i = 0; i < n; i++) A[i] = A[i + 1] - A[i];
for (int i = 0; i < m; i++) B[i] = B[i + 1] - B[i];
Poly C(n + m + 1);
C[0] = A.back() + B.back();
merge(A.begin(), A.end() - 1, B.begin(), B.end() - 1, C.
     begin() + 1, polarComp(Pt(0, 0), Pt(0, -1));
for (int i = 1; i < C.size(); i++) C[i] = C[i] + C[i -</pre>
```

```
C.pop_back();
 return C;
void rotatingCalipersGetRectangle(Pt* p, int n, Tf& area,
            Tf& perimeter) {
 using Linear::disPtLine;
 p[n] = p[0];
 int l = 1, r = 1, j = 1;
 area = perimeter = 1e100;
  for(int i = 0; i < n; i++) {</pre>
   Pt v = (p[i + 1] - p[i]) / len(p[i + 1] - p[i]);
    while (dcmp(dot(v, p[r % n] - p[i]) - dot(v, p[(r + 1) % ot(v, p
                n] - p[i])) < 0) r++;
    while(j < r \mid | dcmp(crs(v, p[j % n] - p[i]) - crs(v, p
              [(j + 1) \% n] - p[i])) < 0) j++;
    while(1 < j \mid | dcmp(dot(v, p[1 % n] - p[i]) - dot(v, p
              [(1 + 1) \% n] - p[i])) > 0) 1++;
    Tf w = dot(v, p[r \% n] - p[i]) - dot(v, p[1 \% n] - p[i])
    Tf h = disPtLine(p[j % n], Line(p[i], p[i + 1]));
    area = min(area, w * h):
   perimeter = min(perimeter, 2 * w + 2 * h);
// returns the left side of polygon u after cutting it by
            rav a->b
Poly cutPoly(Poly u, Pt a, Pt b) {
 using Linear::lineLineIntscn, Linear::onSeg;
 Poly ret;
 int n = u.size();
  for(int i = 0; i < n; i++) {</pre>
   Pt c = u[i], d = u[(i + 1) \% n];
    if(dcmp(crs(b-a, c-a)) >= 0) ret.push_back(c);
    if(dcmp(crs(b-a, d-c)) != 0) {
     lineLineIntscn(a, b - a, c, d - c, t);
      if(onSeg(t, Seg(c, d))) ret.push_back(t);
 return ret;
// returns true if pt p is in or on tri abc
bool ptInTri(Pt a, Pt b, Pt c, Pt p) {
 return dcmp(crs(b - a, p - a)) >= 0 \&\& dcmp(crs(c - b, p))
              -b)) >= 0 && dcmp(crs(a - c, p - c)) >= 0;
 // pt must be in ccw order with no three collinear pts
 // returns inside = -1, on = 0, outside = 1
int ptInConvexPoly(const Poly &pt, Pt p) {
 int n = pt.size();
```

```
assert(n >= 3);
int lo = 1, hi = n - 1;
while(hi - lo > 1) {
 int mid = (lo + hi) / 2;
 if(dcmp(crs(pt[mid] - pt[0], p - pt[0])) > 0) lo = mid;
 else hi = mid;
}
bool in = ptInTri(pt[0], pt[lo], pt[hi], p);
if(!in) return 1;
if(dcmp(crs(pt[lo] - pt[lo - 1], p - pt[lo - 1])) == 0)
     return 0;
if(dcmp(crs(pt[hi] - pt[lo], p - pt[lo])) == 0) return
if(dcmp(crs(pt[hi] - pt[(hi + 1) % n], p - pt[(hi + 1) % |// any of them (only one, not all) should be returned (
      n])) == 0) return 0;
return -1;
}
// Extreme Pt for a direction is the farthest pt in that
// poly is a convex polygon, sorted in CCW, doesn't
    contain redundant pts
// u is the direction for extremeness
int extremePt(const Poly &poly, Pt u = Pt(0, 1)) {
int n = (int) poly.size();
int a = 0, b = n;
while(b - a > 1) {
 int c = (a + b) / 2;
 if(dcmp(dot(poly[c] - poly[(c + 1) % n], u)) >= 0 &&
      dcmp(dot(poly[c] - poly[(c - 1 + n) % n], u)) >= 0)
      {
  return c;
  bool a_up = dcmp(dot(poly[(a + 1) % n] - poly[a], u))
      >= 0:
  bool c_{up} = dcmp(dot(poly[(c + 1) % n] - poly[c], u))
  bool a_above_c = dcmp(dot(poly[a] - poly[c], u)) > 0;
  if (a up && !c up) b = c;
  else if(!a_up && c_up) a = c;
  else if(a up && c up) {
  if(a_above_c) b = c;
  else a = c;
  else {
  if(!a_above_c) b = c;
  else a = c;
 }
if(dcmp(dot(poly[a] - poly[(a + 1) % n], u)) > 0 && dcmp
     (dot(poly[a] - poly[(a - 1 + n) % n], u)) > 0)
```

```
return a;
return b % n;
// For a convex polygon p and a line 1, returns a list of
// of p that are touch or intersect line 1.
// the i'th segs is considered (p[i], p[(i + 1) modulo |p return ret;
// #1 If a segs is collinear with the line, only that is
// #2 Else if 1 goes through i'th pt, the i'th segs is
// If there are 2 or more such collinear segss for #1,
    not tested)
// Complexity: O(lg |p|)
vector<int> lineConvexPolyIntscn(const Poly &p, Line 1) {
assert((int) p.size() >= 3);
assert(1.a != 1.b);
int n = p.size();
vector<int> ret;
Pt v = 1.b - 1.a;
int lf = extremePt(p, rotate90(v));
int rt = extremePt(p, rotate90(v) * Ti(-1));
 int olf = ornt(l.a, l.b, p[lf]);
 int ort = ornt(l.a, l.b, p[rt]);
 if(!olf || !ort) {
 int idx = (!olf ? lf : rt);
 if(ornt(1.a, 1.b, p[(idx - 1 + n) \% n]) == 0)
  ret.push_back((idx - 1 + n) \% n);
  else ret.push back(idx):
 return ret;
if(olf == ort) return ret;
for(int i=0; i<2; ++i) {</pre>
 int lo = i ? rt : lf:
 int hi = i ? lf : rt;
 int olo = i ? ort : olf;
  while(true) {
  int gap = (hi - lo + n) \% n;
  if(gap < 2) break;</pre>
  int mid = (lo + gap / 2) % n;
  int omid = ornt(l.a, l.b, p[mid]);
  if(!omid) {
   lo = mid:
   break;
```

```
if(omid == olo) lo = mid;
  else hi = mid;
 ret.push_back(lo);
// Calculate [ACW, CW] tangent pair from an external pt
constexpr int CW = -1, ACW = 1;
bool isGood(Pt u, Pt v, Pt Q, int dir) { return ornt(Q, u
    , v) != -dir; }
Pt better(Pt u, Pt v, Pt Q, int dir) { return ornt(Q, u,
    v) == dir ? u : v; }
Pt ptPolyTng(const Poly &pt, Pt Q, int dir, int lo, int
    hi) {
while(hi - lo > 1) {
  int mid = (lo + hi) / 2;
 bool pvs = isGood(pt[mid], pt[mid - 1], Q, dir);
  bool nxt = isGood(pt[mid], pt[mid + 1], Q, dir);
  if(pvs && nxt) return pt[mid];
  if(!(pvs || nxt)) {
  Pt p1 = ptPolyTng(pt, Q, dir, mid + 1, hi);
  Pt p2 = ptPolyTng(pt, Q, dir, lo, mid - 1);
  return better(p1, p2, Q, dir);
  if(!pvs) {
  if(ornt(Q, pt[mid], pt[lo]) == dir)
                                          hi = mid - 1:
  else if(better(pt[lo], pt[hi], Q, dir) == pt[lo]) hi =
        mid - 1;
  else lo = mid + 1;
 }
  if(!nxt) {
  if(ornt(Q, pt[mid], pt[lo]) == dir)
                                          lo = mid + 1:
  else if(better(pt[lo], pt[hi], Q, dir) == pt[lo]) hi =
        mid - 1:
  else lo = mid + 1;
}
Pt ret = pt[lo];
 for(int i = lo + 1; i <= hi; i++) ret = better(ret, pt[i</pre>
     ], Q, dir);
return ret;
// [ACW, CW] Tng
pair<Pt, Pt> ptPolyTngs(const Poly &pt, Pt Q) {
 int n = pt.size();
Pt acw_tan = ptPolyTng(pt, Q, ACW, 0, n - 1);
 Pt cw_tan = ptPolyTng(pt, Q, CW, 0, n - 1);
```

return make\_pair(acw\_tan, cw\_tan);

```
4.5 Circular
// Extremely inaccurate for finding near touches
// compute intersection of line 1 with crc c
// The intersections are given in order of the ray (l.a,
vec_p crcLineIntscn(Crc c, Line 1) {
static_assert(is_same<Tf, Ti>::value);
vec_p ret;
Pt b = 1.b - 1.a, a = 1.a - c.o;
Tf A = dot(b, b), B = dot(a, b);
Tf C = dot(a, a) - c.r * c.r, D = B*B - A*C;
if (D < -EPS) return ret:</pre>
ret.push_back(l.a + b * (-B - sqrt(D + EPS)) / A);
if (D > EPS)
 ret.push back(l.a + b * (-B + sqrt(D)) / A);
return ret:
// signed area of intersection of crc(c.o, c.r) &&
// tri(c.o, s.a, s.b) [crs(a-o, b-o)/2]
Tf crcTriIntscnArea(Crc c, Seg s) {
using Linear::disPtSeg;
Tf OA = len(c.o - s.a);
Tf OB = len(c.o - s.b);
// sector
if(dcmp(disPtSeg(c.o, s) - c.r) >= 0)
 return angleBetween(s.a-c.o, s.b-c.o) * (c.r * c.r) /
      2.0;
// tri
if(dcmp(OA - c.r) \le 0 && dcmp(OB - c.r) \le 0)
 return crs(c.o - s.b, s.a - s.b) / 2.0;
// three part: (A, a) (a, b) (b, B)
vec_p Sect = crcLineIntscn(c, s);
return crcTriIntscnArea(c, Seg(s.a, Sect[0])) +
     crcTriIntscnArea(c, Seg(Sect[0], Sect[1])) +
     crcTriIntscnArea(c, Seg(Sect[1], s.b));
// area of intersecion of crc(c.o, c.r) && simple polyson
    (p[])
// Tested : https://codeforces.com/gym/100204/problem/F
     Little Mammoth
Tf crcPolyIntscnArea(Crc c, Poly p) {
Tf res = 0:
int n = p.size();
for(int i = 0; i < n; ++i)</pre>
 res += crcTriIntscnArea(c, Seg(p[i], p[(i + 1) % n]));
return abs(res):
```

```
// locates crc c2 relative to c1
// interior (d < R - r) \longrightarrow -2
// interior tangents (d = R - r) \longrightarrow -1
// concentric (d = 0)
// secants
               (R - r < d < R + r) \longrightarrow 0
// exterior tangents (d = R + r) ----> 1
// exterior (d > R + r) ----> 2
int crcCrcPosition(Crc c1, Crc c2) {
Tf d = len(c1.o - c2.o);
int in = dcmp(d - abs(c1.r - c2.r)), ex = dcmp(d - (c1.r)// returns the pts on tangents that touches the crc
      + c2.r));
return in < 0 ? -2 : in == 0 ? -1 : ex == 0 ? 1 : ex > 0 Pt u = p - c.o;
// compute the intersection pts between two crcs c1 && c2
vec_p crcCrcIntscn(Crc c1, Crc c2) {
vec p ret;
Tf d = len(c1.o - c2.o);
 if(dcmp(d) == 0) return ret;
if(dcmp(c1.r + c2.r - d) < 0) return ret;
if(dcmp(abs(c1.r - c2.r) - d) > 0) return ret;
Pt v = c2.0 - c1.0:
     .r * len(v));
Tf si = sqrt(abs(1.0 - co * co));
Pt p1 = scale(rotatePrecise(v, co, -si), c1.r) + c1.o;
Pt p2 = scale(rotatePrecise(v, co, si), c1.r) + c1.o;
ret.push_back(p1);
if(p1 != p2) ret.push_back(p2);
return ret;
// intersection area between two crcs c1, c2
Tf crcCrcIntscnArea(Crc c1, Crc c2) {
Pt AB = c2.0 - c1.0:
Tf d = len(AB);
if(d \ge c1.r + c2.r) return 0;
if(d + c1.r <= c2.r) return PI * c1.r * c1.r;</pre>
if(d + c2.r <= c1.r) return PI * c2.r * c2.r;</pre>
Tf alpha1 = acos((c1.r * c1.r + d * d - c2.r * c2.r) /
     (2.0 * c1.r * d));
Tf alpha2 = acos((c2.r * c2.r + d * d - c1.r * c1.r) /
     (2.0 * c2.r * d));
return c1.sector(2 * alpha1) + c2.sector(2 * alpha2);
// returns tangents from a pt p to crc c
vec_p ptCrcTngs(Pt p, Crc c) {
vec_p ret;
Pt u = c.o - p;
Tf d = len(u):
if(d < c.r);
```

```
else if(dcmp(d - c.r) == 0) {
                                                            ret = { rotate(u, PI / 2) };
                                                           else {
                                                            Tf ang = asin(c.r / d);
                                                            ret = { rotate(u, -ang), rotate(u, ang) };
                                                           return ret;
                                                          vec_p ptCrcTngPts(Pt p, Crc c) {
                                                           Tf d = len(u):
                                                           if(d < c.r) return {};</pre>
                                                           else if (dcmp(d - c.r) == 0) return \{c.o + u\};
                                                            Tf ang = acos(c.r / d);
                                                            u = u / len(u) * c.r;
                                                            return { c.o + rotate(u, -ang), c.o + rotate(u, ang) };
                                                           // for two crcs c1 && c2, returns two list of pts a && b
Tf co = (c1.r * c1.r + sqlen(v) - c2.r * c2.r) / (2 * c1 \frac{1}{3} such that a[i] is on c1 && b[i] is c2 && for every i
                                                          // Line(a[i], b[i]) is a tangent to both crcs
                                                          // CAUTION: a[i] = b[i] in case they touch | -1 for c1 =
                                                           int crcCrcTngPts(Crc c1, Crc c2, vec_p &a, vec_p &b) {
                                                           a.clear(), b.clear();
                                                           int cnt = 0;
                                                           if(dcmp(c1.r - c2.r) < 0) {
                                                            swap(c1, c2); swap(a, b);
                                                           Tf d2 = sqlen(c1.o - c2.o);
                                                           Tf rdif = c1.r - c2.r, rsum = c1.r + c2.r:
                                                           if(dcmp(d2 - rdif * rdif) < 0) return 0;</pre>
                                                           if(dcmp(d2) == 0 \&\& dcmp(c1.r - c2.r) == 0) return -1;
                                                           Tf base = angle(c2.o - c1.o);
                                                           if(dcmp(d2 - rdif * rdif) == 0) {
                                                            a.push_back(c1.pt(base));
                                                            b.push back(c2.pt(base));
                                                            cnt++:
                                                            return cnt;
                                                           Tf ang = acos((c1.r - c2.r) / sqrt(d2));
                                                           a.push_back(c1.pt(base + ang));
                                                           b.push_back(c2.pt(base + ang));
                                                           cnt++:
                                                           a.push_back(c1.pt(base - ang));
                                                           b.push_back(c2.pt(base - ang));
                                                           cnt++;
```

```
if(dcmp(d2 - rsum * rsum) == 0) {
 a.push_back(c1.pt(base));
 b.push_back(c2.pt(PI + base));
 cnt++;
else if (dcmp(d2 - rsum * rsum) > 0) {
Tf ang = acos((c1.r + c2.r) / sqrt(d2));
 a.push_back(c1.pt(base + ang));
 b.push_back(c2.pt(PI + base + ang));
 a.push_back(c1.pt(base - ang));
 b.push_back(c2.pt(PI + base - ang));
 cnt++:
return cnt;
```

### 4.6 Half Plane

```
using Linear::lineLineIntscn;
struct DirLine {
Pt p, v;
Tf ang;
DirLine() {}
/// Directed line containing pt P in the direction v
DirLine(Pt p, Pt v) : p(p), v(v) { ang = atan2(v.y, v.x)
bool operator<(const DirLine& u) const { return ang < u.
     ang; }
// returns true if pt p is on the ccw-left side of ray 1
bool onLeft(DirLine 1, Pt p) { return dcmp(crs(1.v, p-1.p struct edge {
    )) >= 0; }
// Given a set of directed lines returns a polygon such
// the polygon is the intersection by halfplanes created
// left side of the directed lines. MAY CONTAIN DUPLICATE
int halfPlaneIntscn(vector<DirLine> &li, Poly &poly) {
int n = li.size();
sort(li.begin(), li.end());
int first, last;
Pt* p = new Pt[n];
DirLine* q = new DirLine[n];
q[first = last = 0] = li[0];
for(int i = 1; i < n; i++) {</pre>
 while(first < last && !onLeft(li[i], p[last - 1])) last</pre>
  while(first < last && !onLeft(li[i], p[first])) first</pre>
```

```
q[++last] = li[i];
 if(dcmp(crs(q[last].v, q[last-1].v)) == 0) {
 if(onLeft(q[last], li[i].p)) q[last] = li[i];
 if(first < last)</pre>
 lineLineIntscn(q[last - 1].p, q[last - 1].v, q[last].p
      , q[last].v, p[last - 1]);
while(first < last && !onLeft(q[first], p[last - 1]))</pre>
    last--;
if(last - first <= 1) {</pre>
 delete∏ p:
 delete[] q;
 poly.clear();
 return 0;
lineLineIntscn(q[last].p, q[last].v, q[first].p, q[first
    ].v, p[last]);
int m = 0:
poly.resize(last - first + 1);
for(int i = first; i <= last; i++) poly[m++] = p[i];</pre>
delete[] p;
delete[] q;
return m;
```

# Graph

# 5.1 Graph Template

```
int u. v:
edge(int u = 0, int v = 0) : u(u), v(v) {}
int to(int node) { return u ^ v ^ node: }
struct graph {
int n;
vector<vector<int>> adj;
vector<edge> edges;
graph(int n = 0) : n(n), adj(n) {}
void addEdge(int u, int v, bool dir = 1) {
 adj[u].push_back(edges.size());
 if (dir) adj[v].push_back(edges.size());
 edges.emplace_back(u, v);
int addNode() {
 adj.emplace_back();
 return n++;
edge &operator()(int idx) { return edges[idx]; }
```

```
vector<int> &operator[](int u) { return adj[u]; }
```

```
5.2 Lifting, LCA, HLD
using Tree = vector<vector<int>>;
int anc[B][N], sz[N], lvl[N], st[N], en[N], nxt[N], t =
void initLifting(int n) {
for (int b = 1; b < B; b++) {</pre>
 for (int i = 0: i < n: i++) {
  anc[b][i] = anc[b - 1][anc[b - 1][i]];
}
int kthAncestor(int u, int k) {
for (int b = 0; b < B; b++) {</pre>
 if (k \gg b \& 1) u = anc[b][u]:
return u;
int lca(int u, int v) {
if (lvl[u] > lvl[v]) swap(u, v);
v = kthAncestor(v, lvl[v] - lvl[u]);
if (u == v) return u;
for (int b = B - 1; b \ge 0; b--) {
 if (anc[b][u] != anc[b][v]) u = anc[b][u], v = anc[b][v]
     ];
return anc[0][u]:
int dis(int u, int v) {
int g = lca(u, v);
return lvl[u] + lvl[v] - 2 * lvl[g];
bool isAncestor(int u, int v) { return st[v] <= st[u] and
     en[u] <= en[v]; }
void tour(int u, int p, Tree &T) {
st[u] = t++:
int idx = 0;
for (int v : T[u]) {
 if (v == p) continue;
 nxt[v] = (idx++?v : nxt[u]); // only for hld
 anc[0][v] = u, lvl[v] = lvl[u] + 1;
```

```
tour(v, u, T);
en[u] = t; // [st, en] contains subtree range
void hld(int u, int p, Tree &T) {
sz[u] = 1:
int eld = 0, mx = 0, idx = 0;
for (int i = 0; i < T[u].size(); i++) {</pre>
 int v = T[u][i];
 if (v == p) continue;
 hld(v, u, T);
 if (sz[v] > mx) mx = sz[v], eld = i;
 sz[u] += sz[v];
swap(T[u][0], T[u][eld]);
LL climbQuery(int u, int g) {
LL ans = -INF;
while (1) {
 int u = nxt[u];
 if (isAncestor(g, _u)) _u = g;
 ans = max(ans, rmq ::query(st[_u], st[u]));
 if (_u == g) break;
 u = anc[0][u];
return ans:
LL pathQuery(int u, int v) {
int g = lca(u, v);
return max(climbQuery(u, g), climbQuery(v, g));
void init(int u, Tree &T) {
int n = T.size();
anc[0][u] = nxt[u] = u;
lvl[u] = 0;
hld(u, u, T);
tour(u, u, T);
initLifting(n);
```

### 5.3 SCC

```
vector<int> order, comp, idx;
vector<bool> vis;
```

```
vector<vector<int>> comps;
Graph dag;
void dfs1(int u, Graph &G, string s = "") {
vis[u] = 1:
for (int e : G[u]) {
 int v = G(e).to(u):
 if (!vis[v]) dfs1(v, G, s);
order.push_back(u);
void dfs2(int u, Graph &R) {
comp.push_back(u);
idx[u] = comps.size();
for (int e : R[u]) {
 int v = R(e).to(u);
 if (idx[v] == -1) dfs2(v, R);
void init(Graph &G) {
int n = G.n;
vis.assign(n, 0);
idx.assign(n, -1);
for (int i = 0; i < n; i++) {</pre>
 if (!vis[i]) dfs1(i, G);
reverse(order.begin(), order.end());
Graph R(n);
for (auto &e : G.edges) R.addEdge(e.v, e.u, 0);
 for (int u : order) {
 if (idx[u] != -1) continue;
 comp.clear();
 dfs2(u, R);
 comps.push_back(comp);
Graph &createDAG(Graph &G) {
int sz = comps.size();
dag = Graph(sz);
vector<bool> taken(sz);
vector<int> cur;
for (int i = 0; i < sz; i++) {
```

```
cur.clear();
taken[i] = 1;
for (int u : comps[i]) {
  for (int e : G[u]) {
    int v = G(e).to(u);
    int j = idx[v];
    if (taken[j]) continue; // rejects multi-edge
    dag.addEdge(i, j, 0);
    taken[j] = 1;
    cur.push_back(j);
  }
}
for (int j : cur) taken[j] = 0;
}
return dag;
}
```

# 5.4 Centroid Decompose

```
namespace ct {
int par[N], cnt[N], cntp[N];
LL sum[N], sump[N];
void activate(int u) {
int v = u, u = u;
ans += sum[u];
cnt[u]++;
while (par[u] != -1) {
 u = par[u];
 LL d = ta :: dis(_u, u);
 ans += sum[u] - sump[v];
 ans += d * (cnt[u] - cntp[v]);
 sum[u] += d;
 cnt[u]++;
 sump[v] += d;
 cntp[v]++;
 v = u;
namespace ctrd {
int sz[N];
bool blk[N];
int szCalc(Tree &T, int u, int p = -1) {
sz[u] = 1;
for (int v : T[u]) {
 if (v == p or blk[v]) continue;
```

```
sz[u] += szCalc(T, v, u);
}
return sz[u];
}
int getCentroid(Tree &T, int u, int s, int p = -1) {
  for (int v : T[u]) {
    if (v == p or blk[v]) continue;
    if (sz[v] * 2 >= s) return getCentroid(T, v, s, u);
}
return u;
}

void decompose(Tree &T, int u, int p = -1) {
  szCalc(T, u);
  u = getCentroid(T, u, sz[u]);
  ct ::par[u] = p;

blk[u] = 1;
  for (int v : T[u]) {
    if (!blk[v]) decompose(T, v, u);
  }
}
}
```

# 5.5 Euler Tour on Edge

```
// for simplicity, G[idx] contains the adjacency list of
    a node
// while G(e) is a reference to the e-th edge.
const int N = 2e5 + 5;
int in[N], out[N], fwd[N], bck[N];
int t = 0:
void dfs(graph &G, int node, int par) {
out[node] = t;
for (int e : G[node]) {
 int v = G(e).to(node);
 if (v == par) continue;
 fwd[e] = t++:
 dfs(G, v, node);
 bck[e] = t++;
in[node] = t - 1;
void init(graph &G, int node) {
t = 0:
dfs(G, node, node);
```

### 5.6 Virtual Tree

```
namespace lca1 {
int st[N], lvl[N];
```

```
int tb1[B][2 * N];
int t = 0;
void dfs(int u, int p, Tree &T) {
 st[u] = t;
 tbl[0][t++] = u;
for(int v: T[u]) {
 if(v == p) continue;
 lvl[v] = lvl[u] + 1;
  dfs(v, u, T);
  tbl[0][t++] = u;
int low(int u, int v) {
return make_pair(lvl[u], u) < make_pair(lvl[v], v) ? u :</pre>
void makeTable(int n) {
 int m = 2 * n - 1:
for(int b = 1; b < B; b++) {</pre>
 for(int i = 0; i < m; i++) {</pre>
  tbl[b][i] = low(tbl[b - 1][i], tbl[b - 1][i + (1 << b]
       - 1)]);
 }
int lca(int u, int v) {
 int 1 = st[u], r = st[v];
 if(1 > r) swap(1, r);
 int k = _- lg(r - 1 + 1);
return low(tbl[k][1], tbl[k][r - (1 << k) + 1]);</pre>
void init(int root, Tree &T) {
lvl[root] = 0;
t = 0;
dfs(root, root, T);
 makeTable(T.size());
namespace vt {
int st[N], en[N], t;
vector <int> adj[N];
void dfs(int u, int p, Tree &T) {
 st[u] = t++:
for(int v: T[u]) if(v != p) dfs(v, u, T);
en[u] = t++;
```

```
bool comp(int u, int v) {
return st[u] < st[v];</pre>
bool isAncestor(int u, int p) {
return st[p] <= st[u] and en[u] <= en[p];</pre>
void construct(vector <int> &nodes) {
sort(nodes.begin(), nodes.end(), comp);
int n = nodes.size();
for(int i = 0; i + 1 < n; i++) {</pre>
 nodes.push_back(lca1 :: lca(nodes[i], nodes[i + 1]));
sort(nodes.begin(), nodes.end(), comp);
nodes.erase(unique(nodes.begin(), nodes.end()), nodes.
     end());
n = nodes.size();
stack <int> s;
s.push(nodes[0]);
for(int i = 1; i < n; i++) {</pre>
 int u = nodes[i];
 while(not isAncestor(u, s.top())) s.pop();
 adj[s.top()].push_back(u);
 s.push(u);
void clear(vector <int> &nodes) {
for(int u: nodes) {
  adj[u].clear();
void init(int root, Tree &T) {
lca1 :: init(root, T):
t = 0:
dfs(root, root, T);
```

#### 5.7 Dinic Max Flow

```
/// flow with demand(lower bound) only for DAG
// create new src and sink
// add_edge(new src, u, sum(in_demand[u]))
// add_edge(u, new sink, sum(out_demand[u]))
// add_edge(old sink, old src, inf)
// if (sum of lower bound == flow) then demand satisfied
// flow in every edge i = demand[i] + e.flow
using Ti = long long;
const Ti INF = 1LL << 60;</pre>
```

```
struct edge {
int v, u;
Ti cap, flow = 0;
edge(int v, int u, Ti cap) : v(v), u(u), cap(cap) {}
const int N = 1e5 + 50;
vector<edge> edges;
vector<int> adj[N];
int m = 0, n;
int level[N], ptr[N];
queue<int> q;
bool bfs(int s, int t) {
for (q.push(s), level[s] = 0; !q.empty(); q.pop()) {
 for (int id : adj[q.front()]) {
  auto &ed = edges[id];
  if (ed.cap - ed.flow > 0 and level[ed.u] == -1)
   level[ed.u] = level[ed.v] + 1, q.push(ed.u);
return level[t] != -1;
Ti dfs(int v, Ti pushed, int t) {
if (pushed == 0) return 0;
if (v == t) return pushed;
for (int &cid = ptr[v]; cid < adj[v].size(); cid++) {</pre>
 int id = adj[v][cid];
  auto &ed = edges[id];
 if (level[v] + 1 != level[ed.u] || ed.cap - ed.flow <</pre>
      1) continue;
 Ti tr = dfs(ed.u, min(pushed, ed.cap - ed.flow), t);
 if (tr == 0) continue;
  ed.flow += tr;
  edges[id ^ 1].flow -= tr;
 return tr:
return 0;
void init(int nodes) {
m = 0, n = nodes;
for (int i = 0; i < n; i++) level[i] = -1, ptr[i] = 0,
     adj[i].clear();
void addEdge(int v, int u, Ti cap) {
edges.emplace_back(v, u, cap), adj[v].push_back(m++);
edges.emplace_back(u, v, 0), adj[u].push_back(m++);
Ti maxFlow(int s, int t) {
Ti f = 0;
for (auto &ed : edges) ed.flow = 0;
for (; bfs(s, t); memset(level, -1, n * 4)) {
```

```
for (memset(ptr, 0, n * 4); Ti pushed = dfs(s, INF, t);
    f += pushed)
;
}
return f;
}
```

#### 5.8 Min Cost Max Flow

```
mt19937 rnd(chrono::steady_clock::now().time_since_epoch
    ().count());
const LL inf = 1e9:
struct edge {
int v, rev;
LL cap, cost, flow;
edge() {}
edge(int v, int rev, LL cap, LL cost)
  : v(v), rev(rev), cap(cap), cost(cost), flow(0) {}
struct mcmf {
int src, sink, n;
vector<int> par, idx, Q;
vector<bool> inq;
vector<LL> dis;
vector<vector<edge>> g;
mcmf() {}
mcmf(int src, int sink, int n)
  : src(src),
   sink(sink),
   n(n),
   par(n),
   idx(n),
   inq(n),
   dis(n),
   g(n),
   \mathbb{Q}(10000005) {} // use \mathbb{Q}(n) if not using random
void add_edge(int u, int v, LL cap, LL cost, bool
     directed = true) {
 edge _u = edge(v, g[v].size(), cap, cost);
 edge _v = edge(u, g[u].size(), 0, -cost);
 g[u].pb(_u);
 g[v].pb(_v);
 if (!directed) add_edge(v, u, cap, cost, true);
bool spfa() {
 for (int i = 0; i < n; i++) {</pre>
  dis[i] = inf, inq[i] = false;
 int f = 0, 1 = 0;
 dis[src] = 0, par[src] = -1, Q[1++] = src, inq[src] =
```

```
while (f < 1) {</pre>
  int u = Q[f++];
  for (int i = 0; i < g[u].size(); i++) {</pre>
   edge &e = g[u][i];
   if (e.cap <= e.flow) continue;</pre>
   if (dis[e.v] > dis[u] + e.cost) {
    dis[e.v] = dis[u] + e.cost:
    par[e.v] = u, idx[e.v] = i;
    if (!inq[e.v]) inq[e.v] = true, Q[1++] = e.v;
    // if (!inq[e.v]) {
    // ing[e.v] = true;
    // if (f \&\& rnd() \& 7) Q[--f] = e.v;
   // else Q[1++] = e.v;
   // }
   }
  }
  ing[u] = false;
 return (dis[sink] != inf);
pair<LL, LL> solve() {
 LL mincost = 0, maxflow = 0;
 while (spfa()) {
  LL bottleneck = inf;
  for (int u = par[sink], v = idx[sink]; u != -1; v =
      idx[u], u = par[u]) {
   edge &e = g[u][v];
   bottleneck = min(bottleneck, e.cap - e.flow);
  for (int u = par[sink], v = idx[sink]; u != -1; v =
      idx[u], u = par[u]) {
   edge &e = g[u][v];
   e.flow += bottleneck:
   g[e.v][e.rev].flow -= bottleneck;
  mincost += bottleneck * dis[sink], maxflow +=
      bottleneck:
 return make_pair(mincost, maxflow);
// want to minimize cost and don't care about flow
// add edge from sink to dummy sink (cap = inf, cost = 0)
// add edge from source to sink (cap = inf, cost = 0)
// run mcmf, cost returned is the minimum cost
```

#### 5.9 Block Cut Tree

```
vector<vector<int> > components;
vector<int> cutpoints, start, low;
vector<bool> is_cutpoint;
```

```
stack<int> st:
void find cutpoints(int node, graph &G, int par = -1, int | low[u] = depth[u] = d;
low[node] = start[node] = d++;
st.push(node);
int cnt = 0;
for (int e : G[node])
 if (int to = G(e).to(node); to != par) {
  if (start[to] == -1) {
   find_cutpoints(to, G, node, d + 1);
   cnt++;
   if (low[to] >= start[node]) {
    is_cutpoint[node] = par != -1 or cnt > 1;
    components.push_back({node}); // starting a new
        block with the point
    while (st.top() != node)
     components.back().push_back(st.top()), st.pop();
   }
  }
  low[node] = min(low[node], low[to]);
 }
graph tree;
vector<int> id;
void init(graph &G) {
int n = G.n:
start.assign(n, -1), low.resize(n), is_cutpoint.resize(n)
    ), id.assign(n, -1);
find_cutpoints(0, G);
for (int u = 0; u < n; ++u)
 if (is_cutpoint[u]) id[u] = tree.addNode();
for (auto &comp : components) {
 int node = tree.addNode();
 for (int u : comp)
  if (!is_cutpoint[u])
   id[u] = node;
  else
   tree.addEdge(node, id[u]);
if (id[0] == -1) // corner - 1
 id[0] = tree.addNode();
```

# 5.10 Bridge Tree

```
vector<vector<int>> comps;
vector<int> depth, low, id;
stack<int> st;
vector<Edge> bridges;
Graph tree;
```

```
void dfs(int u, Graph &G, int ed = -1, int d = 0) {
st.push(u);
for (int e : G[u]) {
 if (e == ed) continue:
 int v = G(e).to(u);
 if (depth[v] == -1) dfs(v, G, e, d + 1);
 low[u] = min(low[u], low[v]);
 if (low[v] <= depth[u]) continue;</pre>
 bridges.emplace back(u, v);
 comps.emplace_back();
 do {
  comps.back().push_back(st.top()), st.pop();
 } while (comps.back().back() != v);
 if (ed == -1) {
 comps.emplace_back();
 while (!st.empty()) comps.back().push back(st.top()),
      st.pop();
Graph &createTree() {
for (auto &comp : comps) {
 int idx = tree.addNode();
 for (auto &e : comp) id[e] = idx;
for (auto &[1, r]: bridges) tree.addEdge(id[1], id[r]); LL C(int n, int r) {
return tree:
void init(Graph &G) {
int n = G.n:
depth.assign(n, -1), id.assign(n, -1), low.resize(n);
for (int i = 0; i < n; i++) {</pre>
 if (depth[i] == -1) dfs(i, G);
```

# 5.11 Tree Isomorphism

```
mp["01"] = 1;
ind = 1:
int dfs(int u, int p) {
int cnt = 0;
vector<int> vs:
for (auto v : g1[u]) {
 if (v != p) {
  int got = dfs(v, u);
  vs.pb(got);
  cnt++;
```

```
if (!cnt) return 1;
sort(vs.begin(), vs.end());
string s = "0";
for (auto i : vs) s += to_string(i);
vs.clear();
s.pb('1');
if (mp.find(s) == mp.end()) mp[s] = ++ind;
int ret = mp[s];
return ret:
```

#### 6 Math

### 6.1 Combi

```
array<int, N + 1> fact, inv, inv_fact;
void init() {
fact[0] = inv_fact[0] = 1;
for (int i = 1; i <= N; i++) {</pre>
 inv[i] = i == 1 ? 1 : (LL)inv[i - mod % i] * (mod / i +
       1) % mod;
 fact[i] = (LL)fact[i - 1] * i % mod;
 inv fact[i] = (LL)inv fact[i - 1] * inv[i] % mod;
return (r < 0 \text{ or } r > n) ? 0 : (LL)fact[n] * inv_fact[r]
     % mod * inv_fact[n - r] % mod;
```

#### 6.2 Linear Sieve

```
const int N = 1e7:
vector<int> primes;
int spf[N + 5], phi[N + 5], NOD[N + 5], cnt[N + 5], POW[N]
     + 51:
bool prime [N + 5];
int SOD[N + 5];
void init() {
fill(prime + 2, prime + N + 1, 1);
SOD[1] = NOD[1] = phi[1] = spf[1] = 1;
for (LL i = 2; i <= N; i++) {</pre>
 if (prime[i]) {
  primes.push_back(i), spf[i] = i;
  phi[i] = i - 1;
  NOD[i] = 2, cnt[i] = 1;
  SOD[i] = i + 1, POW[i] = i;
 for (auto p : primes) {
```

### 6.3 Pollard Rho

LL mul(LL a, LL b, LL mod) {

```
return ( int128) a * b % mod;
// LL ans = a * b - mod * (LL) (1.L / mod * a * b);
// return ans + mod * (ans < 0) - mod * (ans >= (LL) mod
     );
LL bigmod(LL num, LL pow, LL mod) {
LL ans = 1;
for (; pow > 0; pow >>= 1, num = mul(num, num, mod)){
 if (pow & 1) ans = mul(ans, num, mod);
}
return ans;
bool is_prime(LL n) {
if (n < 2 or n % 6 % 4 != 1) return (n | 1) == 3;
LL a[] = \{2, 325, 9375, 28178, 450775, 9780504,
     1795265022};
LL s = \_builtin\_ctzll(n - 1), d = n >> s;
for (LL x : a) {
 LL p = bigmod(x \% n, d, n), i = s;
 for (; p != 1 and p != n - 1 and x % n and i--; p = mul |}
      (p, p, n));
 if (p != n - 1 and i != s) return false;
}
return true;
LL get_factor(LL n) {
auto f = [\&](LL x) \{ return mul(x, x, n) + 1; \};
LL x = 0, y = 0, t = 0, prod = 2, i = 2, q;
```

```
for (; t++ % 40 or gcd(prod, n) == 1; x = f(x), y = f(f())
 (x == y) ? x = i++, y = f(x) : 0;
 prod = (q = mul(prod, max(x, y) - min(x, y), n)) ? q :
      prod;
return gcd(prod, n);
map<LL, int> factorize(LL n) {
map<LL, int> res;
if (n < 2) return res;</pre>
LL small_primes[] = {2, 3, 5, 7, 11, 13, 17, 19, 23, 29,
      31, 37, 41, 43, 47, 53, 59, 61, 67, 71, 73, 79,
    83, 89, 97};
for (LL p : small_primes) {
 for (; n % p == 0; n /= p) res[p]++;
auto _factor = [&](LL n, auto &_factor) {
 if (n == 1) return;
 if (is_prime(n))
  res[n]++;
 else {
  LL x = get_factor(n);
  _factor(x, _factor);
  _factor(n / x, _factor);
 _factor(n, _factor);
return res;
```

#### 6.4 Chinese Remainder Theorem

```
// given a, b will find solutions for
// ax + bv = 1
tuple<LL, LL, LL> EGCD(LL a, LL b) {
if (b == 0)
 return {1, 0, a};
 else {
 auto [x, y, g] = EGCD(b, a \% b);
 return \{y, x - a / b * y, g\};
// given modulo equations, will apply CRT
PLL CRT(vector<PLL> &v) {
LL V = 0, M = 1;
for (auto &[v, m] : v) { // value % mod
 auto [x, y, g] = EGCD(M, m);
 if ((v - V) % g != 0) return {-1, 0};
 V += x * (v - V) / g % (m / g) * M, M *= m / g;
 V = (V \% M + M) \% M;
```

```
return make_pair(V, M);
}
```

#### 6.5 Mobius Function

```
const int N = 1e6 + 5;
int mob[N];
void mobius() {
  memset(mob, -1, sizeof mob);
  mob[1] = 1;
  for (int i = 2; i < N; i++)
    if (mob[i]) {
      for (int j = i + i; j < N; j += i) mob[j] -= mob[i];
    }
}</pre>
```

#### 6.6 FFT

```
using CD = complex<double>;
typedef long long LL;
const double PI = acos(-1.0L);
int N;
vector<int> perm;
vector<CD> wp[2];
void precalculate(int n) {
assert((n & (n - 1)) == 0), N = n;
perm = vector<int>(N, 0);
 for (int k = 1; k < N; k <<= 1) {</pre>
 for (int i = 0; i < k; i++) {
  perm[i] <<= 1;
  perm[i + k] = 1 + perm[i];
 }
wp[0] = wp[1] = vector < CD > (N);
for (int i = 0; i < N; i++) {</pre>
 wp[0][i] = CD(cos(2 * PI * i / N), sin(2 * PI * i / N))
 wp[1][i] = CD(cos(2 * PI * i / N), -sin(2 * PI * i / N))
     );
void fft(vector<CD> &v, bool invert = false) {
if (v.size() != perm.size()) precalculate(v.size());
for (int i = 0; i < N; i++)</pre>
 if (i < perm[i]) swap(v[i], v[perm[i]]);</pre>
for (int len = 2; len <= N; len *= 2) {</pre>
 for (int i = 0, d = N / len; i < N; i += len) {</pre>
  for (int j = 0, idx = 0; j < len / 2; j++, idx += d) {
   CD x = v[i + i];
   CD y = wp[invert][idx] * v[i + j + len / 2];
```

```
v[i + j] = x + y;
   v[i + j + len / 2] = x - y;
if (invert) {
 for (int i = 0; i < N; i++) v[i] /= N;
void pairfft(vector<CD> &a, vector<CD> &b, bool invert =
    false) {
int N = a.size();
vector<CD> p(N):
for (int i = 0; i < N; i++) p[i] = a[i] + b[i] * CD(0,
    1);
fft(p, invert);
p.push back(p[0]);
for (int i = 0; i < N; i++) {</pre>
 if (invert) {
  a[i] = CD(p[i].real(), 0);
  b[i] = CD(p[i].imag(), 0);
 } else {
  a[i] = (p[i] + conj(p[N - i])) * CD(0.5, 0);
  b[i] = (p[i] - conj(p[N - i])) * CD(0, -0.5);
 }
}
vector<LL> multiply(const vector<LL> &a, const vector<LL>
     &b) {
int n = 1;
while (n < a.size() + b.size()) n <<= 1;</pre>
vector<CD> fa(a.begin(), a.end()), fb(b.begin(), b.end()
    );
fa.resize(n):
fb.resize(n);
         fft(fa); fft(fb);
pairfft(fa, fb);
for (int i = 0; i < n; i++) fa[i] = fa[i] * fb[i];</pre>
fft(fa, true);
vector<LL> ans(n);
for (int i = 0; i < n; i++) ans[i] = round(fa[i].real())</pre>
return ans;
const int M = 1e9 + 7, B = sqrt(M) + 1;
vector<LL> anyMod(const vector<LL> &a, const vector<LL> &
    b) {
int n = 1;
while (n < a.size() + b.size()) n <<= 1;</pre>
vector<CD> al(n), ar(n), bl(n), br(n);
```

```
for (int i = 0; i < a.size(); i++) al[i] = a[i] % M / B,</pre>
      ar[i] = a[i] % M % B;
for (int i = 0; i < b.size(); i++) bl[i] = b[i] % M / B,</pre>
      br[i] = b[i] % M % B;
 pairfft(al, ar);
 pairfft(bl, br);
          fft(al); fft(ar); fft(bl); fft(br);
 for (int i = 0; i < n; i++) {</pre>
 CD ll = (al[i] * bl[i]), lr = (al[i] * br[i]);
  CD rl = (ar[i] * bl[i]), rr = (ar[i] * br[i]);
  al[i] = 11;
  ar[i] = lr:
 bl[i] = rl:
  br[i] = rr;
pairfft(al, ar, true);
pairfft(bl, br, true);
          fft(al, true); fft(ar, true); fft(bl, true);
     fft(br, true);
 vector<LL> ans(n);
for (int i = 0; i < n; i++) {</pre>
 LL right = round(br[i].real()), left = round(al[i].real
 LL mid = round(round(bl[i].real()) + round(ar[i].real() | for (LL i = 1; i < n - 1; i++)
  ans[i] = ((left \% M) * B * B + (mid \% M) * B + right) \%
       Μ:
return ans;
6.7 NTT
const LL N = 1 << 18:
const LL MOD = 786433;
vector<LL> P[N];
LL rev[N], w[N | 1], a[N], b[N], inv_n, g;
LL Pow(LL b, LL p) {
 LL ret = 1:
 while (p) {
 if (p & 1) ret = (ret * b) % MOD;
 b = (b * b) \% MOD;
 p >>= 1;
 return ret;
LL primitive_root(LL p) {
vector<LL> factor;
```

LL phi = p - 1, n = phi;

```
for (LL i = 2; i * i <= n; i++) {</pre>
 if (n % i) continue;
factor.emplace_back(i);
 while (n \% i == 0) n /= i;
if (n > 1) factor.emplace_back(n);
for (LL res = 2; res <= p; res++) {</pre>
 bool ok = true;
 for (LL i = 0; i < factor.size() && ok; i++)</pre>
  ok &= Pow(res, phi / factor[i]) != 1;
 if (ok) return res;
return -1;
void prepare(LL n) {
LL sz = abs(31 - __builtin_clz(n));
LL r = Pow(g, (MOD - 1) / n);
inv_n = Pow(n, MOD - 2);
w[0] = w[n] = 1;
for (LL i = 1; i < n; i++) w[i] = (w[i-1] * r) % MOD;
for (LL i = 1; i < n; i++)</pre>
rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << (sz - 1));
void NTT(LL *a, LL n, LL dir = 0) {
 if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
for (LL m = 2; m <= n; m <<= 1) {
 for (LL i = 0; i < n; i += m) {</pre>
  for (LL j = 0; j < (m >> 1); j++) {
   LL &u = a[i + j], &v = a[i + j + (m >> 1)];
   LL t = v * w[dir ? n - n / m * j : n / m * j] % MOD;
   v = u - t < 0 ? u - t + MOD : u - t;
   u = u + t >= MOD ? u + t - MOD : u + t;
  }
 }
if (dir)
 for (LL i = 0; i < n; i++) a[i] = (inv_n * a[i]) % MOD;</pre>
vector<LL> mul(vector<LL> p, vector<LL> q) {
LL n = p.size(), m = q.size();
LL t = n + m - 1, sz = 1;
while (sz < t) sz <<= 1;
prepare(sz);
for (LL i = 0; i < n; i++) a[i] = p[i];
for (LL i = 0; i < m; i++) b[i] = q[i];</pre>
for (LL i = n; i < sz; i++) a[i] = 0;</pre>
for (LL i = m; i < sz; i++) b[i] = 0;</pre>
```

```
NTT(a, sz);
NTT(b, sz);
for (LL i = 0; i < sz; i++) a[i] = (a[i] * b[i]) % MOD;
NTT(a, sz, 1);

vector<LL> c(a, a + sz);
while (c.size() && c.back() == 0) c.pop_back();
return c;
}
```

### 6.8 WalshHadamard

```
#include <bits/stdc++.h>
using namespace std;
typedef long long LL;
#define bitwiseXOR 1
// #define bitwiseAND 2
// #define bitwiseOR 3
const LL MOD = 30011;
LL BigMod(LL b, LL p) {
LL ret = 1;
while (p > 0) {
 if (p % 2 == 1) {
  ret = (ret * b) % MOD;
 p = p / 2;
 b = (b * b) \% MOD;
return ret % MOD;
void FWHT(vector<LL>& p, bool inverse) {
LL n = p.size();
assert((n & (n - 1)) == 0):
for (LL len = 1; 2 * len <= n; len <<= 1) {
 for (LL i = 0; i < n; i += len + len) {</pre>
  for (LL j = 0; j < len; j++) {
   LL u = p[i + j];
   LL v = p[i + len + j];
#ifdef bitwiseXOR
   p[i + j] = (u + v) \% MOD;
   p[i + len + j] = (u - v + MOD) \% MOD;
#endif // bitwiseXOR
#ifdef bitwiseAND
   if (!inverse) {
    p[i + j] = v \% MOD;
    p[i + len + j] = (u + v) \% MOD;
```

```
} else {
    p[i + j] = (-u + v) \% MOD;
    p[i + len + j] = u \% MOD;
#endif // bitwiseAND
#ifdef bitwiseOR
   if (!inverse) {
    p[i + j] = u + v;
    p[i + len + j] = u;
   } else {
    p[i + j] = v;
    p[i + len + j] = u - v;
#endif // bitwiseOR
 }
 }
#ifdef bitwiseXOR
if (inverse) {
 LL val = BigMod(n, MOD - 2); // Option 2: Exclude
 for (LL i = 0; i < n; i++) {</pre>
 // assert(p[i]%n==0); //Option 2: Include
  p[i] = (p[i] * val) % MOD; // Option 2: p[i]/=n;
#endif // bitwiseXOR
```

### 6.9 Berlekamp Massey

```
struct berlekamp_massey { // for linear recursion
typedef long long LL;
static const int SZ = 2e5 + 5;
static const int MOD = 1e9 + 7; /// mod must be a prime
LL m, a[SZ], h[SZ], t_{SZ}, s[SZ], t[SZ];
// bigmod goes here
inline vector <LL> BM( vector <LL> &x ) {
 LL lf , ld;
 vector <LL> ls , cur;
 for ( int i = 0; i < int(x.size()); ++i ) {</pre>
 LL t = 0:
  for ( int j = 0; j < int(cur.size()); ++j) t = (t + x)
      [i - j - 1] * cur[j]) % MOD;
  if ((t - x[i]) \% MOD == 0) continue;
  if (!cur.size()) {
   cur.resize( i + 1 ):
   lf = i; ld = (t - x[i]) % MOD;
   continue:
```

```
LL k = -(x[i] - t) * bigmod(ld, MOD - 2, MOD) %
 vector <LL> c(i - lf - 1);
 c.push back( k );
 for ( int j = 0; j < int(ls.size()); ++j ) c.push_back</pre>
      (-ls[i] * k % MOD);
 if ( c.size() < cur.size() ) c.resize( cur.size() );</pre>
 for ( int j = 0; j < int(cur.size()); ++j ) c[j] = (c[</pre>
      i] + cur[i]) % MOD;
 if (i - lf + (int)ls.size() >= (int)cur.size() ) ls =
      cur, lf = i, ld = (t - x[i]) % MOD;
 cur = c:
}
for ( int i = 0; i < int(cur.size()); ++i ) cur[i] = (</pre>
     cur[i] % MOD + MOD) % MOD;
return cur;
inline void mull( LL *p , LL *q ) {
for ( int i = 0; i < m + m; ++i ) t [i] = 0;
for ( int i = 0; i < m; ++i ) if ( p[i] )
  for ( int j = 0; j < m; ++j ) t_{i} = (t_{i} + j)
      + p[i] * q[i]) % MOD;
for ( int i = m + m - 1; i >= m; --i ) if ( t_[i] )
  for ( int j = m - 1; ~j; --j ) t_[i - j - 1] = (t_[i
       - j - 1] + t_{[i]} * h[j]) % MOD;
for ( int i = 0; i < m; ++i ) p[i] = t_[i];
}
inline LL calc( LL K ) {
for ( int i = m; ~i; --i ) s[i] = t[i] = 0;
s[0] = 1; if (m!=1) t[1] = 1; else t[0] = h[0];
 while (K) {
 if (K&1) mull(s,t);
 mull( t , t ); K >>= 1;
}
LL su = 0;
for ( int i = 0; i < m; ++i ) su = (su + s[i] * a[i]) %
      MOD:
return (su % MOD + MOD) % MOD;
/// already calculated upto k , now calculate upto n.
inline vector <LL> process( vector <LL> &x , int n , int
     k ) {
auto re = BM( x );
x.resize(n+1);
for ( int i = k + 1; i \le n; i++ ) {
 for ( int j = 0; j < re.size(); j++ ) {</pre>
  x[i] += 1LL * x[i - j - 1] % MOD * re[j] % MOD; x[i]
       %= MOD;
 }
}
```

### 6.10 Lagrange

```
// p is a polynomial with n points.
// p(0), p(1), p(2), \dots p(n-1) are given.
// Find p(x).
LL Lagrange(vector<LL> &p, LL x) {
LL n = p.size(), L, i, ret;
if (x < n) return p[x];</pre>
L = 1:
for (i = 1; i < n; i++) {
 L = (L * (x - i)) \% MOD:
 L = (L * bigmod(MOD - i, MOD - 2)) % MOD;
ret = (L * p[0]) % MOD;
 for (i = 1; i < n; i++) {</pre>
 L = (L * (x - i + 1)) \% MOD;
 L = (L * bigmod(x - i, MOD - 2)) % MOD;
  L = (L * bigmod(i, MOD - 2)) \% MOD;
 L = (L * (MOD + i - n)) \% MOD;
 ret = (ret + L * p[i]) % MOD;
return ret:
```

# 6.11 Shanks' Baby Step, Giant Step

```
// Finds a^x = b (mod p)

LL bigmod(LL b, LL p, LL m) {}

LL babyStepGiantStep(LL a, LL b, LL p) {
   LL i, j, c, sq = sqrt(p);
   map<LL, LL> babyTable;
```

### 6.12 Xor Basis

```
struct XorBasis {
static const int sz = 64;
array<ULL, sz> base = {0}, back;
array<int, sz> pos;
void insert(ULL x, int p) {
 ULL cur = 0:
 for (int i = sz - 1; ~i; i--)
  if (x >> i & 1) {
   if (!base[i]) {
    base[i] = x, back[i] = cur, pos[i] = p;
    break:
  } else x ^= base[i], cur |= 1ULL << i;</pre>
pair<ULL, vector<int>> construct(ULL mask) {
 ULL ok = 0, x = mask;
 for (int i = sz - 1; ~i; i--)
  if (mask >> i & 1 and base[i]) mask ^= base[i], ok |=
      1ULL << i:
 vector<int> ans;
 for (int i = 0; i < sz; i++)</pre>
  if (ok >> i & 1) {
   ans.push_back(pos[i]);
   ok ^= back[i];
 return {x ^ mask, ans};
```

# String

# 7.1 Aho Corasick

```
struct AC {
int N, P;
const int A = 26;
vector<vector<int>> next;
```

```
vector<int> link, out link;
vector<vector<int>> out;
AC() : N(0), P(0) { node(); }
int node() {
next.emplace_back(A, 0);
link.emplace_back(0);
out_link.emplace_back(0);
out.emplace_back(0);
return N++;
inline int get(char c) { return c - 'a'; }
int add_pattern(const string T) {
int u = 0:
for (auto c : T) {
 if (!next[u][get(c)]) next[u][get(c)] = node();
 u = next[u][get(c)];
out[u].push_back(P);
return P++;
void compute() {
queue<int> q;
 for (q.push(0); !q.empty();) {
 int u = q.front(); q.pop();
 for (int c = 0; c < A; ++c) {
  int v = next[u][c]:
  if (!v) next[u][c] = next[link[u]][c];
  else {
   link[v] = u ? next[link[u]][c] : 0;
   out_link[v] = out[link[v]].empty() ? out_link[link[v]
       ]] : link[v];
   q.push(v);
 }
}
int advance(int u, char c) {
while (u && !next[u][get(c)]) u = link[u];
u = next[u][get(c)];
return u;
void match(const string S) {
int u = 0:
for (auto c : S) {
 u = advance(u, c);
 for (int v = u; v; v = out_link[v]) {
  for (auto p : out[v]) cout << "match " << p << endl;</pre>
}
```

```
};
int main() {
AC aho; int n; cin >> n;
while (n--) {
 string s; cin >> s;
 aho.add_pattern(s);
aho.compute(); string text;
cin >> text; aho.match(text);
return 0:
```

#### 7.2 Double hash

```
// define +, -, * for (PLL, LL) and (PLL, PLL), % for (
    PLL, PLL);
PLL base(1949313259, 1997293877);
PLL mod(2091573227, 2117566807);
PLL power(PLL a, LL p) {
PLL ans = PLL(1, 1);
for(; p; p >>= 1, a = a * a % mod) {
  if(p \& 1) ans = ans * a % mod;
return ans:
PLL inverse(PLL a) { return power(a, (mod.ff - 1) * (mod.
    ss - 1) - 1); }
PLL inv_base = inverse(base);
PLL val;
vector<PLL> P;
void hash_init(int n) {
P.resize(n + 1):
P[0] = PLL(1, 1);
for (int i = 1; i \le n; i++) P[i] = (P[i-1] * base) %
PLL append(PLL cur, char c) { return (cur * base + c) %
/// prepends c to string with size k
PLL prepend(PLL cur, int k, char c) { return (P[k] * c +
    cur) % mod: }
/// replaces the i-th (0-indexed) character from right
    from a to b:
PLL replace(PLL cur, int i, char a, char b) {
cur = (cur + P[i] * (b - a)) \% mod:
return (cur + mod) % mod;
/// Erases c from the back of the string
```

```
PLL pop back(PLL hash, char c) {
return (((hash - c) * inv_base) % mod + mod) % mod;
/// Erases c from front of the string with size len
PLL pop_front(PLL hash, int len, char c) {
return ((hash - P[len - 1] * c) % mod + mod) % mod;
/// concatenates two strings where length of the right is |7.3\>\>\>\> Manacher's
PLL concat(PLL left, PLL right, int k) { return (left * P
    [k] + right) % mod; }
/// Calculates hash of string with size len repeated cnt
/// This is O(\log n). For O(1), pre-calculate inverses
PLL repeat(PLL hash, int len, LL cnt) {
PLL mul = (P[len * cnt] - 1) * inverse(P[len] - 1);
mul = (mul % mod + mod) % mod;
PLL ret = (hash * mul) % mod;
 if (P[len].ff == 1) ret.ff = hash.ff * cnt;
if (P[len].ss == 1) ret.ss = hash.ss * cnt;
return ret;
LL get(PLL hash) { return ((hash.ff << 32) ^ hash.ss); }
struct hashlist {
int len:
 vector<PLL> H, R;
 hashlist() {}
 hashlist(string& s) {
 len = (int)s.size();
 hash init(len);
  H.resize(len + 1, PLL(0, 0)), R.resize(len + 2, PLL(0,
 for (int i = 1; i <= len; i++) H[i] = append(H[i - 1],
      s[i - 1]):
 for (int i = len; i \ge 1; i--) R[i] = append(R[i + 1],
      s[i - 1]):
/// 1-indexed
 PLL range_hash(int 1, int r) {
 return ((H[r] - H[1 - 1] * P[r - 1 + 1]) % mod + mod) %
       mod:
PLL reverse hash(int 1, int r) {
 return ((R[1] - R[r + 1] * P[r - 1 + 1]) % mod + mod) %
       mod;
PLL concat_range_hash(int 11, int r1, int 12, int r2) {
 return concat(range_hash(11, r1), range_hash(12, r2),
      r2 - 12 + 1):
```

```
PLL concat_reverse_hash(int 11, int r1, int 12, int r2)
 return concat(reverse_hash(12, r2), reverse_hash(11, r1
     ), r1 - 11 + 1);
};
```

```
vector<int> d1(n):
// d[i] = number of palindromes taking s[i] as center
for (int i = 0, l = 0, r = -1; i < n; i++) {
int k = (i > r) ? 1 : min(d1[1 + r - i], r - i + 1);
while (0 \le i - k \&\& i + k \le n \&\& s[i - k] == s[i + k])
    k++;
d1[i] = k--;
if (i + k > r) l = i - k, r = i + k;
vector<int> d2(n);
// d[i] = number of palindromes taking s[i-1] and s[i] as
for (int i = 0, l = 0, r = -1; i < n; i++) {
int k = (i > r) ? 0 : min(d2[1 + r - i + 1], r - i + 1);
while (0 \le i - k - 1 \&\& i + k \le n \&\& s[i - k - 1] == s[
    i + kl) k++:
d2[i] = k--;
if (i + k > r) l = i - k - 1, r = i + k;
```

### 7.4 Suffix Array

```
vector<VI> c;
VI sort cyclic shifts(const string &s) {
int n = s.size():
const int alphabet = 256;
VI p(n), cnt(alphabet, 0);
c.clear();
c.emplace_back();
c[0].resize(n);
for (int i = 0; i < n; i++) cnt[s[i]]++;</pre>
for (int i = 1; i < alphabet; i++) cnt[i] += cnt[i - 1];</pre>
for (int i = 0; i < n; i++) p[--cnt[s[i]]] = i;</pre>
c[0][p[0]] = 0;
int classes = 1;
for (int i = 1; i < n; i++) {</pre>
 if (s[p[i]] != s[p[i - 1]]) classes++;
 c[0][p[i]] = classes - 1;
```

```
VI pn(n), cn(n);
cnt.resize(n);
for (int h = 0; (1 << h) < n; h++) {
 for (int i = 0; i < n; i++) {</pre>
  pn[i] = p[i] - (1 << h);
  if (pn[i] < 0) pn[i] += n;</pre>
 fill(cnt.begin(), cnt.end(), 0);
 /// radix sort
 for (int i = 0; i < n; i++) cnt[c[h][pn[i]]]++;</pre>
 for (int i = 1; i < classes; i++) cnt[i] += cnt[i - 1];</pre>
 for (int i = n - 1; i \ge 0; i--) p[--cnt[c[h][pn[i]]]]
      = pn[i];
  cn[p[0]] = 0;
  classes = 1;
  for (int i = 1; i < n; i++) {
  PII cur = {c[h][p[i]], c[h][(p[i] + (1 << h)) % n]};
  PII prev = \{c[h][p[i-1]], c[h][(p[i-1] + (1 << h))]\}
        % n]};
  if (cur != prev) ++classes;
  cn[p[i]] = classes - 1;
 c.push_back(cn);
return p;
VI suffix_array_construction(string s) {
s += "!";
VI sorted_shifts = sort_cyclic_shifts(s);
sorted_shifts.erase(sorted_shifts.begin());
return sorted_shifts;
/// LCP between the ith and jth (i != j) suffix of the
    STRING
int suffixLCP(int i, int j) {
assert(i != j);
int log_n = c.size() - 1;
int ans = 0;
for (int k = log_n; k >= 0; k--) {
 if (c[k][i] == c[k][j]) {
  ans += 1 << k;
  i += 1 << k;
  j += 1 << k;
return ans;
```

```
VI lcp_construction(const string &s, const VI &sa) {
   int n = s.size();
   VI rank(n, 0);
   VI lcp(n - 1, 0);

   for (int i = 0; i < n; i++) rank[sa[i]] = i;

   for (int i = 0, k = 0; i < n; i++, k -= (k != 0)) {
      if (rank[i] == n - 1) {
        k = 0;
        continue;
    }
   int j = sa[rank[i] + 1];
   while (i + k < n && j + k < n && s[i + k] == s[j + k])
        k++;
   lcp[rank[i]] = k;
}
   return lcp;
}</pre>
```

#### 7.5 Z Algo

```
vector<int> calcz(string s) {
int n = s.size();
vector<int> z(n);
int 1 = 0, r = 0;
for (int i = 1; i < n; i++) {</pre>
 if (i > r) {
  1 = r = i;
  while (r < n \&\& s[r] == s[r - 1]) r++;
  z[i] = r - 1, r--:
 } else {
  int k = i - 1:
  if (z[k] < r - i + 1) z[i] = z[k];
  else {
   1 = i:
   while (r < n \&\& s[r] == s[r - 1]) r++;
   z[i] = r - 1, r--;
  }
 }
return z;
```

# **Equations and Formulas**

### Catalan Numbers

$$C_n = \frac{1}{n+1} {2n \choose n}$$
  $C_0 = 1, C_1 = 1 \text{ and } C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}$ 

# Stirling Numbers First Kind

The Stirling numbers of the first kind count permutations according to their number of cycles (counting fixed  $\binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}$ points as cycles of length one).

$$S(n,k) = (n-1) \cdot S(n-1,k) + S(n-1,k-1),$$

$$where, S(0,0) = 1, S(n,0) = S(0,n) = 0 \sum_{k=0}^{n} S(n,k) = \begin{cases} i=0 & k \\ n,r \in N, n > r, \sum_{i=r}^{n} {i \choose r} = {n+1 \choose r+1} \end{cases}$$

$$n!$$
If  $D(n) = \sum_{k=0}^{n} {n \choose k} = \sum_{i=r}^{n} {n \choose i} = \sum_{i=r}^{n} {$ 

The unsigned Stirling numbers may also be defined algebraically, as the coefficient of the rising factorial:

$$x^{\bar{n}} = x(x+1)...(x+n-1) = \sum_{k=0}^{n} S(n,k)x^k$$

$$\begin{bmatrix} n & n \\ n & -k \end{bmatrix} = \sum_{0 \le i_1 < i_2 < i_k < n} i_1 i_2 \dots i_k.$$

# 8.3 Stirling Numbers Second Kind

Stirling number of the second kind is the number of ways to partition a set of n objects into k non-empty subsets.  $S(n,k) = k \cdot S(n-1,k) + S(n-1,k-1), where S(0,0) =$  $1, S(n,0) = S(0,n) = 0 \ S(n,2) = 2^{n-1} - 1 \ S(n,k) \cdot k! =$ number of ways to color n nodes using colors from 1 to k such that each color is used at least once.

An r-associated Stirling number of the second kind is the number of ways to partition a set of n objects into k subsets, with each subset containing at least r elements. It 8.6 GCD and LCM is denoted by  $S_r(n,k)$  and obeys the recurrence relation.

$$S_r(n+1,k) = kS_r(n,k) + \binom{n}{r-1}S_r(n-r+1,k-1)$$

Denote the n objects to partition by the integers 1, 2, ..., n.  $|\gcd(a_1 \cdot a_2, b) = \gcd(a_1, b) \cdot \gcd(a_2, b)$ . Define the reduced Stirling numbers of the second kind,  $\gcd(a, \operatorname{lcm}(b, c)) = \operatorname{lcm}(\gcd(a, b), \gcd(a, c))$ . denoted  $S^d(n,k)$ , to be the number of ways to partition  $\operatorname{lcm}(a,\gcd(b,c)) = \gcd(\operatorname{lcm}(a,b),\operatorname{lcm}(a,c))$ . elements in each subset have pairwise distance at least d. both zero,  $gcd(n^a - 1, n^b - 1) = n^{gcd(a,b)} - 1$ That is, for any integers i and j in a given subset, it is re- $|\gcd(a,b) = \sum \phi(k)$ quired that  $|i-j| \geq d$ . It has been shown that these numbers satisfy,  $S^d(n,k) = S(n-d+1,k-d+1), n \geq k \geq d$ 

# 8.4 Other Combinatorial Identities

tions according to their number of cycles (counting fixed points as cycles of length one). 
$$S(n,k) \text{ counts the number of permutations of } n \text{ elements}$$
with  $k$  disjoint cycles. 
$$S(n,k) = (n-1) \cdot S(n-1,k) + S(n-1,k-1),$$

$$where, S(0,0) = 1, S(n,0) = S(0,n) = 0 \sum_{k=0}^{n} S(n,k) = 0$$

$$n!$$

$$S(n,k) = (n-1) \cdot S(n-1,k) + S(n-1,k-1),$$

$$S(n,k) = (n+i) = \sum_{k=0}^{n} {n \choose k} = {n+i \choose k} =$$

If 
$$P(n) = \sum_{k=0}^{n} {n \choose k} \cdot Q(k)$$
, then,

$$Q(n) = \sum_{k=0}^{n} (-1)^{n-k} \binom{n}{k} \cdot P(k)$$

Lets 
$$[n, k]$$
 be the stirling number of the first kind, then  $\left| \text{If } P(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot Q(k) \right|$ , then,

$$Q(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot P(k)$$

# Different Math Formulas

Picks Theorem : A = i + b/2 - 1

**Deragements:**  $d(i) = (i-1) \times (d(i-1) + d(i-2))$ 

$$\frac{n}{ab}$$
 -  $\left\{\frac{b'n}{a}\right\}$  -  $\left\{\frac{a'n}{b}\right\}$  +

if m is any integer, then  $gcd(a + m \cdot b, b) = gcd(a, b)$ 

The gcd is a multiplicative function in the following sense: if  $a_1$  and  $a_2$  are relatively prime, then

the integers 1, 2, ..., n into k nonempty subsets such that all For non-negative integers a and b, where a and b are not  $\sum[\gcd(i,n) = k] = \phi\left(\frac{n}{k}\right)$  $\sum_{k=1}^{n} \gcd(k, n) = \sum_{d|n} d \cdot \phi\left(\frac{n}{d}\right)$  $\sum_{k=1}^{n} x^{\gcd(k,n)} = \sum_{d|n} x^d \cdot \phi\left(\frac{n}{d}\right)$  $\sum_{k=1}^{n} \frac{1}{\gcd(k,n)} = \sum_{d|n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{1}{n} \sum_{d|n} d \cdot \phi(d)$  $\sum_{k=1}^{n} \frac{k}{\gcd(k,n)} = \frac{n}{2} \cdot \sum_{d|n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{n}{2} \cdot \frac{1}{n} \cdot \sum_{d|n} d \cdot \phi(d)$ 

$$\sum_{k=1}^{n} \frac{\gcd(k,n)}{\gcd(k,n)} = 2 * \sum_{k=1}^{n} \frac{k}{\gcd(k,n)} - 1, \text{ for } n > 1$$

$$\sum_{i=1}^{n} \sum_{j=1}^{n} [\gcd(i,j) = 1] = \sum_{d=1}^{n} \mu(d) \left\lfloor \frac{n}{d} \right\rfloor^{2}$$

$$\sum_{i=1}^{n} \sum_{j=1}^{n} \gcd(i,j) = \sum_{d=1}^{n} \phi(d) \left\lfloor \frac{n}{d} \right\rfloor^{2}$$

$$\sum_{i=1}^{n} \sum_{j=1}^{n} i \cdot j[\gcd(i,j) = 1] = \sum_{i=1}^{n} \phi(i)i^{2}$$

$$F(n) = \sum_{i=1}^{n} \sum_{j=1}^{n} \operatorname{lcm}(i, j) = \sum_{l=1}^{n} \left( \frac{\left(1 + \lfloor \frac{n}{l} \rfloor\right) \left(\lfloor \frac{n}{l} \rfloor\right)}{2} \right)^{2} \sum_{d \mid l} \mu(d)$$