EXP NO: 8 DATE: 02.04.25

GENERATE THREE ADDRESS CODE FOR A SIMPLE PROGRAM USING LEX AND YACC

AIM:

To design and implement a **LEX and YACC** program that generates **three-address code** (**TAC**) for a simple arithmetic expression or program. The program will:

- Recognize **expressions** like addition, subtraction, multiplication, and division.
- Generate **three-address code** that represents the operations in a way that could be directly translated into assembly code or intermediate code for a compiler.

ALGORITHM:

1. Lexical Analysis (LEX) Phase:

Input: A string containing an arithmetic expression (e.g., a = b + c * d;).

Output: A stream of tokens such as identifiers (variables), numbers (constants), operators, and special characters (like =, ;, (), etc.).

1. Define the Token Patterns:

- o **ID:** Identifiers (variables) are strings starting with a letter and followed by letters or digits (e.g., a, b, result).
- o **NUMBER:** Constants (e.g., 1, 5, 100).
- o **OPERATOR:** Arithmetic operators (+, -, *, /).
- o **ASSIGNMENT:** Assignment operator (=).
- o **PARENTHESIS:** Parentheses for grouping ((and)).
- WHITESPACE: Spaces, tabs, and newline characters (which should be ignored).

2. Write Regular Expressions for the Tokens:

- \circ ID -> [a-zA-Z][a-zA-Z0-9]*
- NUMBER -> [0-9]+
- \circ OPERATOR -> [\+\-*/]
- o ASSIGN -> "="
- \circ PAREN -> [\(\)]
- WHITESPACE -> [\t\n]+ (skip whitespace)

3. Action on Tokens:

- When a token is matched, pass it to YACC using yylval to store the token values.
- 2. Syntax Analysis and TAC Generation (YACC) Phase:

Input: Tokens provided by the **LEX** lexical analyzer.

Output: Three-address code for the given arithmetic expression.

1. **Define Grammar Rules:**

Assignment:

bash CopyEdit statement: ID '=' expr

This means an expression is assigned to a variable.

• Expressions:

bash CopyEdit

expr: expr OPERATOR expr

An expression can be another expression with an operator (+, -, *, /).

bash CopyEdit expr: NUMBER expr: ID expr: '(' expr ')'

2. Three-Address Code Generation:

- o For every arithmetic operation, generate a temporary variable (e.g., t1, t2, etc.) to hold intermediate results.
- For a = b + c, generate:

ini CopyEdit t1 = b + c a = t1

• For a = b * c + d, generate:

ini CopyEdit t1 = b * ct2 = t1 + da = t2

3. Temporary Variable Management:

- Keep a counter (temp_count) for generating unique temporary variable names (t0, t1, t2, ...).
- Each time a new operation is encountered, increment the temp_count to generate a new temporary variable.

4. Rule Actions:

• When a rule is matched (e.g., expr OPERATOR expr), generate the TAC and assign temporary variables for intermediate results.

Detailed Algorithm:

1. Initialize Lexical Analyzer:

o Define the token patterns for ID, NUMBER, OPERATOR, ASSIGN, PAREN, and WHITESPACE.

2. Define the Syntax Grammar:

- o Define grammar rules for:
 - Assignments: ID = expr
 - Expressions: expr -> expr OPERATOR expr, expr -> NUMBER, expr
 -> ID, expr -> (expr)

3. Token Matching:

- **LEX:** Match input characters against the defined regular expressions for tokens.
- YACC: Use the tokens to parse and apply grammar rules.

4. TAC Generation:

- o For Assignment:
 - Upon parsing ID = expr, generate a temporary variable for the result of expr and assign it to the variable ID.
- For Arithmetic Operations:
 - For each operator (e.g., +, -, *, /), generate temporary variables for intermediate calculations.

5. Output TAC:

o Print the generated three-address code, with each expression and its intermediate results represented by temporary variables.

PROGRAM:

LEX file (expr.l)

```
% {
#include "y.tab.h"
% }

%%
[0-9]+ { yylval.str = strdup(yytext); return NUMBER; }
[a-zA-Z_][a-zA-Z0-9_]* { yylval.str = strdup(yytext); return ID; }
[+\-*/=()] { return yytext[0]; }
[ \t\n] { /* Ignore whitespace */ }
. { printf("Unexpected character: %s\n", yytext); }
%%
```

YACC Program expr.y

```
% {
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
int temp_count = 0;
```

```
char* new_temp() {
  char* temp = (char*)malloc(8);
  sprintf(temp, "t%d", temp_count++);
  return temp;
}
void emit(char* result, char* op1, char op, char* op2) {
  printf("%s = %s %c %s\n", result, op1, op, op2);
}
void emit_assign(char* id, char* expr) {
  printf("%s = %s\n", id, expr);
}
%}
%union {
  char* str;
%token <str> ID NUMBER
%type <str> expr term factor
%left '+' '-'
%left '*' '/'
%%
statement : ID '=' expr { emit_assign($1, $3); }
         : expr'+' term { $$ = new_temp(); emit($$, $1, '+', $3); }
expr
       | expr '-' term { $$ = new_temp(); emit($$, $1, '-', $3); }
       | term
                   { $$ = $1; }
       ;
         : term '*' factor { $$ = new_temp(); emit($$, $1, '*', $3); }
       | term '/' factor { $$ = new_temp(); emit($$, $1, '/', $3); }
                    { $$ = $1; }
       factor
factor
       : '(' expr ')' { $$ = $2; }
       | NUMBER  { $$ = $1; }
                { $$ = $1; }
       | ID
%%
int main() {
  yyparse();
  return 0;
}
```

```
void yyerror(const char* s) {
    fprintf(stderr, "Error: %s\n", s);
}

OUTPUT:
yacc -d expr.y
lex expr.l
gcc y.tab.c lex.yy.c -o expr_parser
./expr_parser
a = b * c + d;
t0 = b * c
t1 = t0 + d
a = t1
```

RESULT:

Thus the process effectively tokenizes the input, parses it according to defined grammar rules, and generates the corresponding Three-Address Code, facilitating further compilation or interpretation stages.