Assignment 4

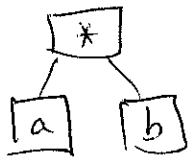
1. Infix to expression tree: $a * b / (c + d - e * (f + g) / h)$

After 3 tokens: opstack: [*] exptree stack: [b] [a]

/ has same precedence as *, so pop * from opstack, 2 expressions from expstack, merge them with * as op:

opstack: [ ] expstack: [*] [a] [b]

T_1:

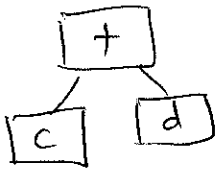


After 5 tokens: op stack: [+] [(] [/] expstack: [d] [c] [T_1]

- has same precedence as +, pop + from stack, 2 expressions from exp stack and merge them with + as operator.

Opstack: [(] [/] expstack: [T_2] [T_1]

T_2:

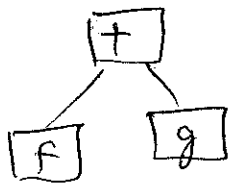


After 7 tokens: opstack: [+] [(] [*] [-] [(] [/] expstack: [g] [f] [e] [T_2] [T_1]

When) is processed, pop opstack up to (:

opstack: [*] [-] [(] [/] exp: [T_3] [e] [T_2] [T_1]

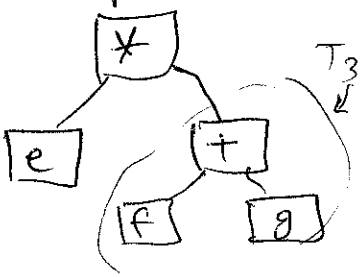
T_3:



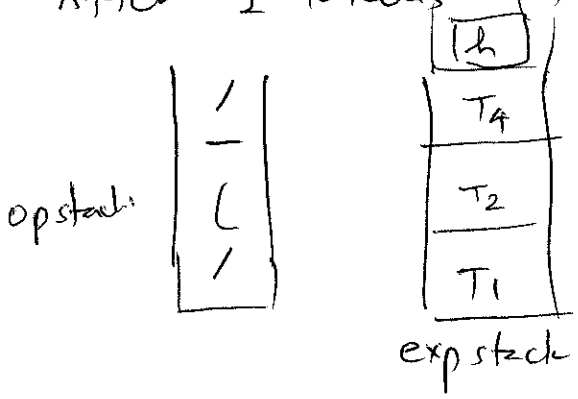
/ has same precedence as *, pop * and 2 exp:

[-] [(] [/] exp: [T_4] [T_2] [T_1]

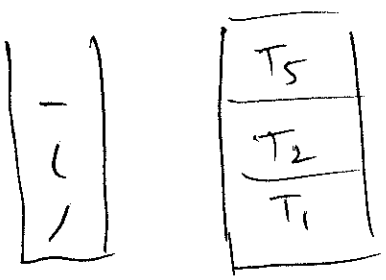
T_4:

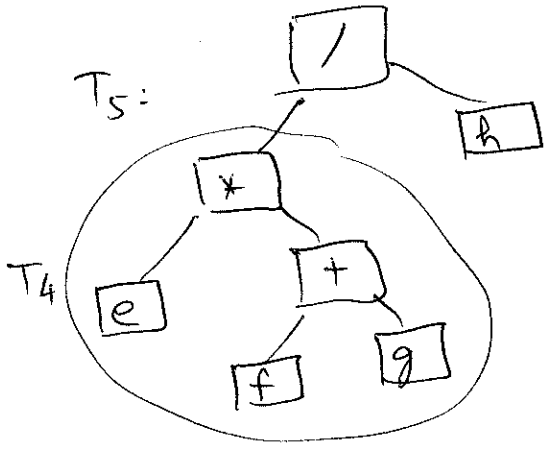


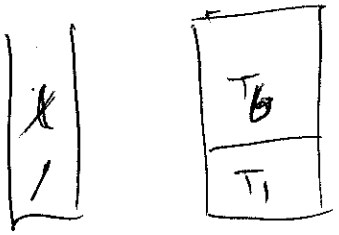
After 2 tokens (/ h :

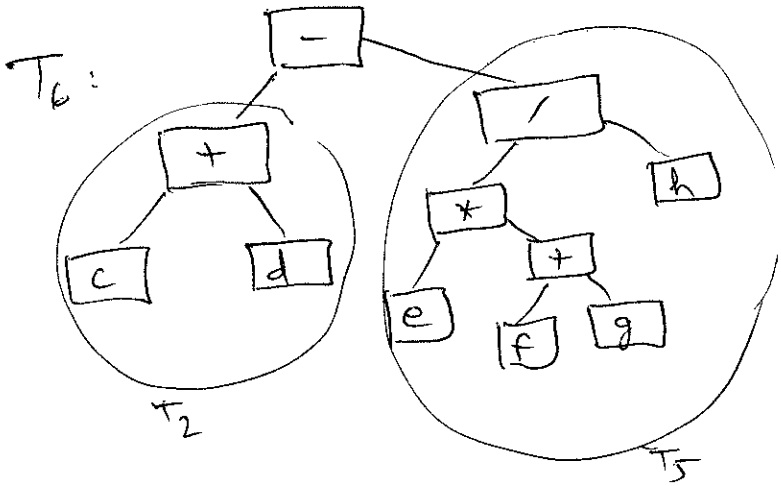


When processing), pop operators from stack and process them until (is removed:

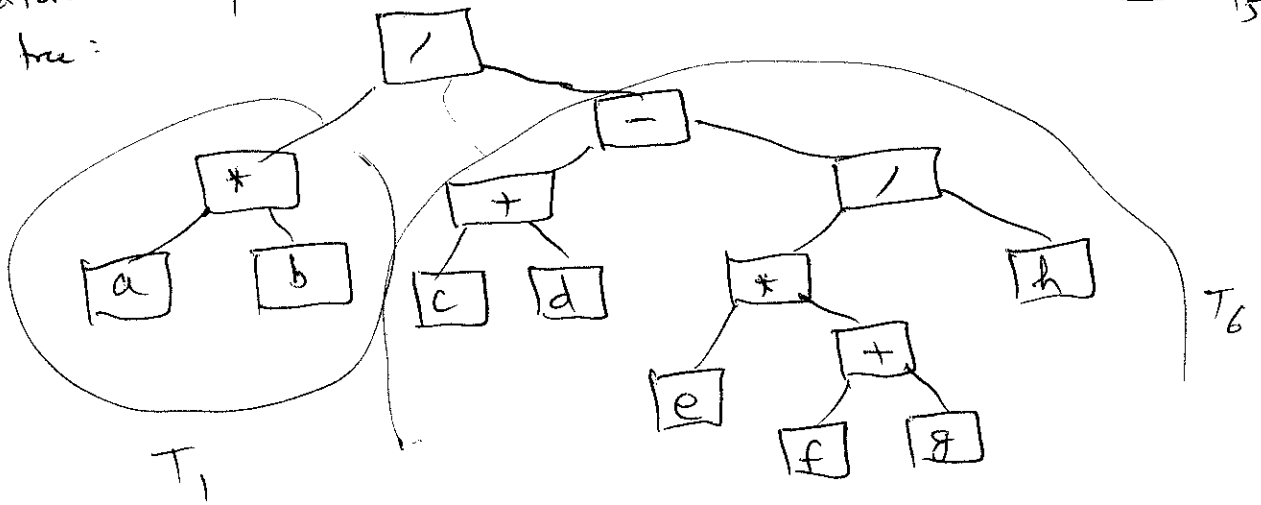








Input is done: log remaing operators and process them.
Final tree:



Solution to Q2 on Assignment 4

```
// return elements at index s, s+k, s+2k, s+3k, ...
List filter( List l, int s, int k ): // RT = O(n).
    List result ← empty list of int
    if l.size() >= 1+s then // result's size > 0
            Iterator it ← l.iterator();
            for i ← 0 to s do
                x ← it.next()
            result.add(x)
            count ← 0
            // LI: s+q*k+count elements processed
            while it.hasNext() do
                count ← count + 1
                x ← it.next()
                if count = k then
                        result.add(x)
                        count ← 0
    return result
```

Binary Search Trees in Java Library

TreeSet, TreeMap — Red Black trees (extension of BST's)

TreeSet: implementation of interface Set.

Usage:
```
Set<Integer> s = new TreeSet<>();
s.add(x);   // add new integer x to set
            // duplicates are rejected
```

TreeSet<Vertex> for user defined class Vertex:

(i) Vertex class defines: (a) Comparable<Vertex>:

public int compareTo(Vertex other)

in Vertex class — natural ordering

or (b) Comparator<Vertex> passed as a parameter

Set<Vertex> s = new TreeSet(comp)

comp = object that implements Comparator<Vertex>

- user defined ordering:

public int compare(Vertex one, Vertex two)

Convention: $a \leq b$? a.compareTo(b) or compare(a,b)

$a < b$: return -1 (any negative no) $a = b$: 0 $a > b$: return $+1$ (any positive no)

Note: It is customary to provide:

boolean equals (Object other), — a.equals(b)?

int hashCode( ); — used in ~~Lists~~ hash tables.

```
s.remove(x);      //remove x from s
s.contains(x);    //does x appear in set?
s.iterator();     ← create an iterator to go
                    through elements of s
                    in sorted order of keys.
```

Usage of iterator:

explicit:

```
Iterator<Integer> it
          = s.iterator();
//s = TreeSet<Integer>
while (it.hasNext()) {
    Integer x = it.next();
        :
    it.remove(); //if needed
}
```

implicit

```
//for classes that are Iterable.
for (Integer x : s) {
        //process x
        :
}
```

Implicit iterator does not allow you to remove current object.

TreeMap: Map of tuples K, V (Key-value pairs)

K = Keys that have no duplicates, no null keys. $K \to V$

V = no restrictions.

Usage: Map<Vertex, Integer> m = new TreeMap<>(

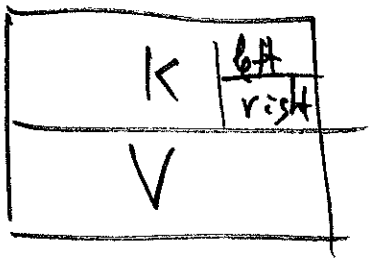
// map using natural ordering of Vertex class.

Can also create TreeMap<>(comp)

↑ Comparator<Vertex>

for user defined ordering.

BST: Each Entry is



Operations:

m.get(key): value stored with ~~key~~ K key.
null if there is no such key.

m.put(key, value): (a) if key exists already:
replace its old value by value.
old value is returned by put.

(b) if key is new,
add a new entry (key, value)
return null.

Iterate over map:

```
for(Map.Entry<Vertex, Integer> e : m.entrySet()) {
    .
    .
    .
    e.getKey()
    e.getValue()
```

Iterate over objects in order of keys.

Additional ops: `m.containsKey(key)`

`m.containsValue(value)` ← $O(n)$ don't use.

Key based operations of TreeMap/TreeSet:

$O(\log n)$ per operation (n = # of keys in tree)

Iteration: $O(\log n)$ per op — worst case

$O(1)$ per op — amortized.