

CSE460/560 DATA MODELS AND QUERY LANGUAGES

Relational Database Design

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(Slides Adopted from Jan Chomicki and Ning Deng)



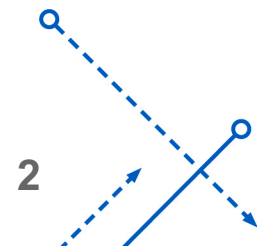
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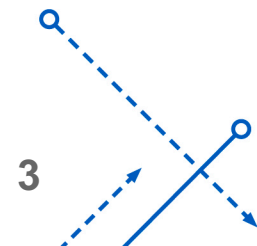
Outline

1. Functional Dependencies
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 2. Inference of FDs
2. Normal Forms
 1. BCNF and 3NF



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1. Functional Dependencies
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Functional Dependencies (FDs)

- Keys of Relations
 - Given a set of attributes $\{A_1, A_2, \dots, A_n\}$
 - Impossible for two distinct tuples of R
 - Agree on all $\{A_1, A_2, \dots, A_n\}$
 - No proper subset of $\{A_1, A_2, \dots, A_n\}$ determines all other attributes
 - Key must be **minimal**
 - Key {title, year, starName}
 - Is {title, year} a key?

<i>title</i>	<i>year</i>	<i>length</i>	<i>genre</i>	<i>studioName</i>	<i>starName</i>
Star Wars	1977	124	SciFi	Fox	Carrie Fisher
Star Wars	1977	124	SciFi	Fox	Mark Hamill
Star Wars	1977	124	SciFi	Fox	Harrison Ford
Gone With the Wind	1939	231	drama	MGM	Vivien Leigh
Wayne's World	1992	95	comedy	Paramount	Dana Carvey
Wayne's World	1992	95	comedy	Paramount	Mike Meyers

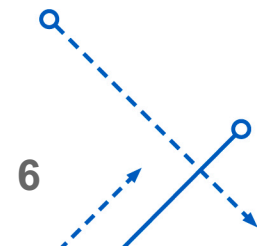
Functional Dependencies (FDs)

- Superkeys
 - Superset of a key
 - Every key is a superkey
 - Some superkeys are not keys
 - Why?
 - 📍 Example of superkey?

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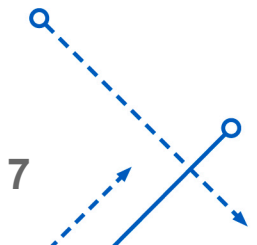
Functional Dependencies (FDs)

- Dependency Implication
 - A set of FDs F implies an FD $X \rightarrow Y$
 - If every relations instance that satisfies all dependencies in F
 - Also satisfies $X \rightarrow Y$
 - Notation
 - $F \models X \rightarrow Y$
 - Closure of a dependency set F
 - $F^+ = \{X \rightarrow Y : F \models X \rightarrow Y\}$
 - Suppose $\{A_1, \dots, A_n\}$ is a set of attributes and F is a set of FDs
 - The closure of $\{A_1, \dots, A_n\}$ under F is set of attribute B
 - such that every relation that satisfies all the FDs in F also satisfies $A_1, \dots, A_n \rightarrow B$
 - Notation: $\{A_1, \dots, A_n\}^+$



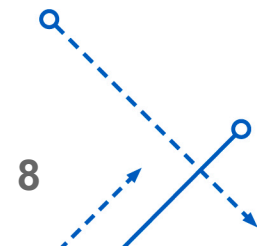
Functional Dependencies (FDs)

- Formal Specification of Key
 - $X \subseteq \{A_1, \dots, A_n\}$ is a key of R if
 - The dependency $X \rightarrow A_1, \dots, A_n$ is in F^+
 - For all proper subsets Y of X
 - The dependency $Y \rightarrow A_1, \dots, A_n$ is not in F^+
- Notations
 - Superkey: superset of a key
 - Primary key: one designated key
 - Candidate key: one of the keys



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 - 1. BCNF and 3NF**



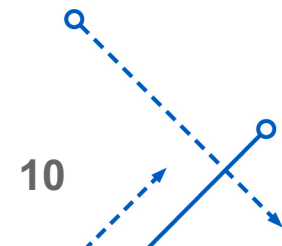
Inference of FDs

- How to tell **whether** $X \rightarrow Y \in F^+$
- Inference Rules (Armstrong Axioms)
 - Reflexivity: infer $X \rightarrow Y$ if $Y \subseteq X \subseteq attr(R)$ (trivial dependency)
 - Augmentations: from $X \rightarrow Y$ infer $XZ \rightarrow YZ$ if $Z \subseteq attr(R)$
 - Transitivity: from $X \rightarrow Y$ and $Y \rightarrow Z$, infer $X \rightarrow Z$
- Armstrong axioms are
 - Sound: if $X \rightarrow Y$ is derived from F , then $X \rightarrow Y \in F^+$
 - Complete: if $X \rightarrow Y \in F^+$, then $X \rightarrow Y$ is derived from F
- Additional (implied) inference rules
 - Union: from $X \rightarrow Y$ and $X \rightarrow Z$, infer $X \rightarrow YZ$
 - Decomposition: from $X \rightarrow Y$ infer $X \rightarrow Z$, if $Z \subseteq Y$

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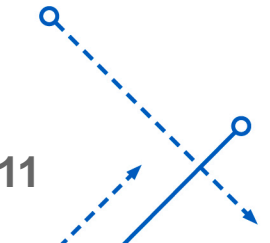
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Normal Forms

- To eliminate or reduce redundancy in database relations
 - 1NF: a relation is in 1NF
 - Every attribute in that relation is atomic
 - Reduce the need for restructuring the collection of relations, as new types of data are introduced
 - Make the relational model more informative to users
 - 2NF: a relation is in 2NF
 - iff it's in 1NF
 - And has no partial dependency
 - No non-prime attribute (attributes which are not part of any candidate key) is dependent on any proper subset of any candidate key of the relation



Normal Forms (3NF, BCNF)

- A relation R is in 3NF if for every non-trivial FD $X \rightarrow A \in F$
 - *either*
 - X contains a key of R
 - A is part of some key of R
- A relation R is in BCNF if
 - For every non-trivial FD $X \rightarrow A \in F$
 - X contains a key of R
 - BCNF does not contain a redundancy
- 3NF vs. BCNF
 - If R is in BCNF, it is also in 3NF
 - A relations the are in 3NF but not in BCNF



Normal Forms (3NF, BCNF)

- 3NF vs. BCNF
 - A relations the are in 3NF but not in BCNF
 - $R(A,B,C)$
 - FD: $\{A,B\} \rightarrow \{C\}$, $\{C\} \rightarrow \{B\}$
 - Candidate Keys: $\{A,B\}$, $\{A,C\}$
 - In 3NF
 - $\{A,B\} \rightarrow \{C\}$ (A contains a key of R)
 - $\{C\} \rightarrow \{B\}$ (B is part of some key of R)
 - Not BCNF
 - $\{C\} \rightarrow \{B\}$ (C does NOT contain a key of R)

Recommended Reading

Database Systems: The Complete Book
Chapter 3.1 – 3.6